

Blatt 6 NB

$$1.1) T_{\text{worst}}(n) = T_{\text{best}}(n) = C(1) = 1 \\ \Rightarrow O(1)$$

$$1.3) T_{\text{best}}(n, m) = C(1) + C(1) + C(1) + \text{Konstruktor } n = 4$$

$$T_{\text{worst}}(n, m) = C(1) + C(1) + C(1) + T_{\text{worst}}(n-m, m) \\ = 3 \cdot \frac{n}{m} + 1$$

$$\Rightarrow O(n)$$

$$1.4) T_{\text{best}}(n) = T_{\text{best}} \text{ fun3}(n, 2) + \text{"True"/"False"} \\ = ~~3 \cdot \frac{n}{2} + 1~~ \\ = ~~3 \cdot \frac{n}{2} + 1~~ 4 + 1 = 5 \\ \Rightarrow O(n)$$

$$T_{\text{worst}}(n) = T_{\text{worst}} \text{ fun3}(n, 2) + \text{"True"/"False"} \\ = 3 \cdot \frac{n}{2} + 1 + 1 \\ = 3 \cdot \frac{n}{2} + 2 \\ \Rightarrow O(n)$$

$$1.5) T_{\text{best}}(0, n) = T_{\text{worst}}(0, n) = \text{Konstruktor } [] = 1$$

$$T_{\text{worst}}(n, m) = C(1) + T_{\text{worst}}(n-1, m-1) \\ = C(1) + \dots + T_{\text{worst}}(0, m-n) \\ = n + 1 \\ \Rightarrow O(n)$$

$$1.6) T_{\text{best}}(n, 0) = T_{\text{worst}}(n, 0) = \text{Konstruktor } [] = 1$$

$$T_{\text{best}}(n, m) = C(1) + T_{\text{best}}(n-1, m-1) \\ = C(1) + \dots + T_{\text{best}}(0, m-n) / T_{\text{best}}(n-m, 0) \\ = (n+1), \quad n \geq m \Rightarrow O(n+1) = T_{\text{worst}}(n, m) \\ \Rightarrow O(n)$$

$$1.7) T_{\text{best}}(1) = T_{\text{worst}}(1) = \text{Konstante für } 1 = 1$$

$$T_{\text{best}}(n) = C(1) + T_{\text{best}}(1)$$

$$= 1 + 1$$

$$= 2$$

$$T_{\text{worst}}(n) = C(1) + T_{\text{worst}}(n-1)$$

$$= C(1) + \dots + T_{\text{worst}}(1)$$

$$= n \cdot 1 + 1$$

$$= n$$

$$\Rightarrow O(n)$$

$$1.8) T_{\text{best}}(n) = C(1) + \text{Konstante für } 1 = 1$$

$$= n + 1 - 1 = n$$

$$T_{\text{worst}}(n) = C(1) + C(1) + T_{\text{worst}}(n-1)$$

$$= C(1) + C(1) + T_{\text{worst}}(0)$$

$$= 2 \cdot n + 1$$

$$\Rightarrow O(n)$$

$$1.9) T_{\text{worst}}(0) = 1$$

$$T_{\text{worst}}(n) = C(n) + T_{\text{worst}}(n-1)$$

$$T_{\text{worst}}(n) = \frac{1}{2} n \cdot (n-1)$$

$$= \frac{1}{2} n^2 - \frac{1}{2} n$$

$$T_{\text{best}}(n) = 3(n-1)$$

$$\Rightarrow O(n^2)$$

$$2c) O(3n \cdot \log n + 7n) = O(3n \cdot \log n) + O(7n)$$

$$= O(n \cdot \log n) + O(n)$$

$$= O(n \cdot \log n) \text{, da } O(n) \leq O(n \cdot \log n)$$

$$\text{also } O(3n \cdot \log n + 7n) \in O(n \log n) \checkmark$$



$$2. f) 3^n \in 2^{\alpha n}$$

$$\log(3)^n \in \log(2)^{\alpha n}$$

$$n \cdot \log(3) \in O(n \cdot \log(2))$$

$$n \cdot \log(3) \in O(n)$$

$$O(n \log 3) = O(n) \cdot O(\log 3)$$

$$= O(n)$$

$$\Rightarrow 3^n \in 2^{\alpha n} \checkmark$$