

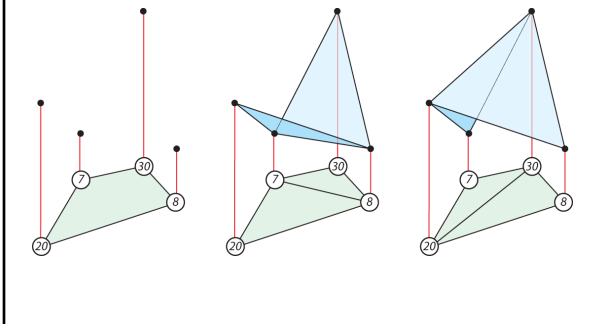
Announcements

- Lab today will be office hours instead
- No lecture on Thursday 2/28
- Switching lab from Tuesday 3/5 to Thursday 3/7 – next week only

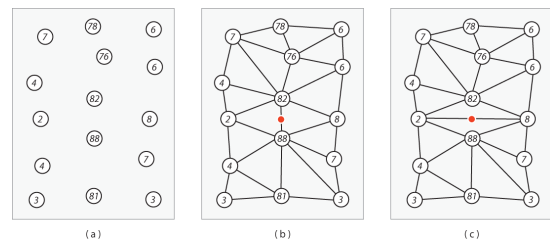
Computational Geometry

Delaunay and Other Special Triangulations

Lifting the Triangles



Skinny is Bad

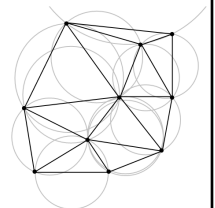


Angle Sequence

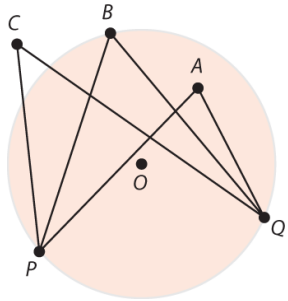
- Let T be a triangulation of a point set S , and suppose T has n triangles. The angle sequence $\{a_1, a_2, \dots, a_n\}$ lists all $3n$ angles of T in sorted order.
- A triangulation T_1 is fatter than T_2 ($T_1 > T_2$) if the angle sequence of T_1 is lexicographically greater than T_2 's.
 $\{30^\circ, 45^\circ, 65^\circ, 120^\circ\} > \{30^\circ, 45^\circ, 60^\circ, 120^\circ\}$

Delaunay Triangulation

- For each convex quad in a triangulation T_1 with diagonal e , if a diagonal flip results in a triangulation T_2 , s.t. $T_1 \geq T_2$, then e is legal.
- A Delaunay triangulation is a triangulation with all legal edges.

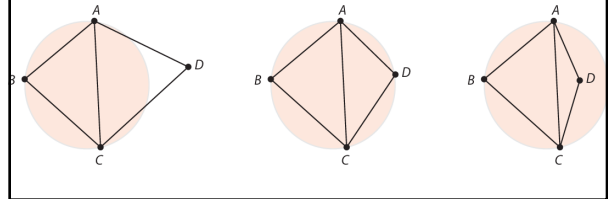


Thales' Theorem

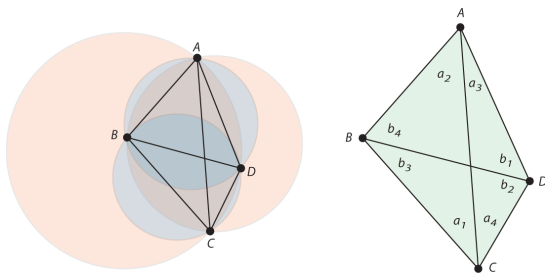


Thales' and Illegal Edges

- Let e be an edge of a triangulation, where $e = AC$ belongs to the two triangles ABC and ACD . Then e is legal if D is outside of the circumcircle of ABC and illegal if D is inside.

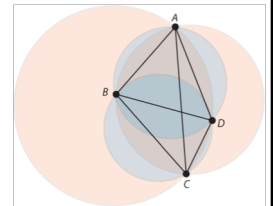


Proof

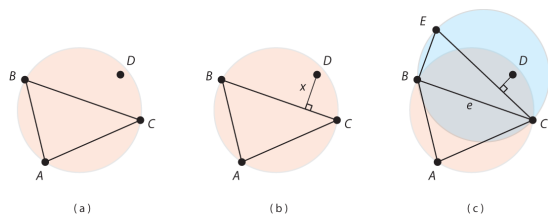


Empty Circle Property

- Let S be a point set in general position. A triangulation T is Delaunay if and only if no point from S is in the interior of any circumcircle of any triangle of T .



Proof



Naïve Delaunay

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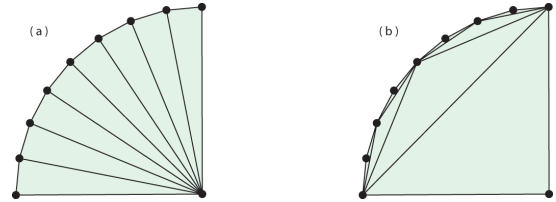
Delaunay( $\{p_1, \dots, p_n\}$ )
  for i in  $[1, n]$ 
    for j in  $[1, i]$ 
      for k in  $[1, j]$ 
        ( $c, r$ ) = Circumcircle( $p_i, p_j, p_k$ )
        isValid = true
        for l in  $[1, k]$ 
          if ( $|p_l - c| < r$ )
            isValid = false
        if (isValid) Output( $p_i, p_j, p_k$ )

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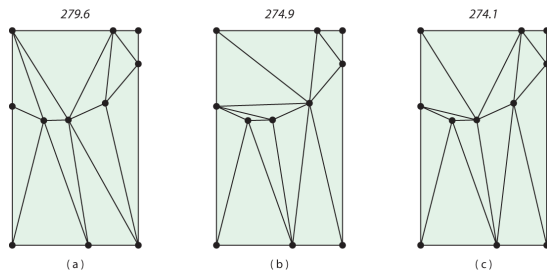
Delaunay via Edge Flipping

- Create any triangulation (incremental, say)
- Flip every illegal edge until done
- Complexity?

Delaunay is not MWT



Delaunay vs. Greedy vs. MWT



Theorem

- For point set S , a minimum spanning tree of S is a subset of the Delaunay triangulation of S .
- Proof by contradiction.

Compatible Triangulations

