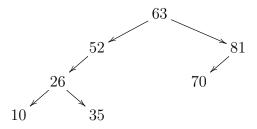
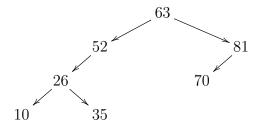
This exam has 10 questions, for a total of 100 points.

1. 8 points For each of the nodes 35, 52, 63, and 81, list the successor and predecessors of the node (if they exist).

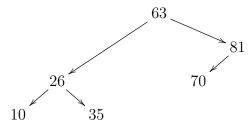


Solution: (2 points each)

- For 35, 52 is the successor and 26 is the predecessor.
- For 52, 63 is the successor and 35 is the predecessor.
- For 63, 70 is the successor and 52 is the predecessor.
- For 81, there is no successor and 70 is the predecessor.
- 2. 10 points Draw the result of deleting the node with key 52, followed by deleting the node with key 63, from the below Binary Search Tree.

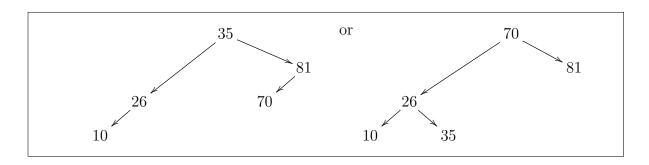


Solution: The node with key 52 has only one child, so we can attach its child to its parent. (5 points)



The node with key 63 has two children, so we can replace with with either its predeccessor 25 or its successor 60. (5 points) Both solutions are shown below.

Name:



3. 10 points Implement the post_order methods of the Tree and Node classes. They should return a list of the keys in the tree, ordered by the "post-order" traversal. For example:

```
class Tree:
    def __init__(self, r):
        self.root = r

    def post_order(self):
        return self.root.post_order()

class Node:
    def __init__(self, k, l=None, r=None):
        self.key = k
        self.left = l
        self.right = r

def post_order(self):
```

Name: _____

```
L = self.left.post_order() if self.left else []
R = self.right.post_order() if self.right else []
return L + R + [self.key]
```

4. 10 points What is a best DNA sequence alignment for the sequences TAG and ATG? When computing the score, use +2 for a match, -2 for a mismatch, and -1 for the gap penalty. Show your work by filling in the below dynamic programming table.

	T	A	G
A			
Т			
G			

Solution: One of several solutions is

TA_G _ATG

		Т	A	G
	0	←- 1	←-2	← -3
A	↑-1	△ -2	<u>\</u>	$\leftarrow 0$
\overline{T}	↑-2	<u>\</u>	↑0	_ -1
G	↑- 3	† 0	←-1	$\nwarrow 2$

(2 points for correct initialization of the table. 5 points for the rest of the table. 3 points for correct result from traceback.)

- 5. 2 points Suppose that L is a Python list (array) of length n. Categorize the worst-case execution time of the below expressions as either
 - 1. constant time (takes the same amount of time no matter what n is).
 - 2. logarithmic time (takes time proportional to $\lg n$).
 - 3. linear time (takes time proportional to n)
 - 4. quadratic (takes time proportional to n^2)

Label each operation with the above item number.

Name:

- L[n]
- L.remove(n/2)
- L + L
- 0 in L

Solution: (3 points each)

- (1), L[n] is constant time,
- (3), L.remove(n/2) is linear time,
- (3), L + L is linear time.
- (3), 0 in L is linear time,
- 6. 6 points What is the big-O worst-case time complexity of the following function, which raises x to the power n, in terms of the input n? Assume that all the primitive number operations (multiplication, division, etc.) are constant time. Explain why your answer is correct.

```
def exponent(x, n):
    if n == 0:
        return 1
    elif n == 1:
        return x
    elif n % 2 == 0:
        return exponent(x * x, n / 2)
    else:
        return x * exponent(x, n - 1)
```

Solution: In the worst-case, each time we divide in half we get an odd number, and have to make an extra call to exponent. So we have $2 \lg n$ calls, but the factor of 2 is ignored for the big-O, so the time complexity is $O(\lg n)$.

7. 10 points Give the definition of Θ . Prove that if $f(n) \in \Theta(g(n))$ then $g(n) \in \Theta(f(n))$.

Solution: The definition of Θ is: (2 points) $\Theta(g(n)) = \{f(n) \mid \exists n_0. \forall n \geq n_0. \exists c_1 c_2. \ 0 \leq c_1 g(n) \leq f(n) \leq c_2 g(n)\}$

Name:

So we can assume

$$\forall n \geq n_0. \ 0 \leq c_1 g(n) \leq f(n) \leq c_2 g(n)$$

and we need to show that

$$\forall n \ge n'_0. \ 0 \le c'_1 f(n) \le g(n) \le c'_2 f(n)$$

for some suitable choice for n'_0, c'_1, c'_2 . From $f(n) \leq c_2 g(n)$ we have

$$\forall n \ge n_0. \frac{1}{c_2} f(n) \le g(n)$$
 (2 points)

and from $c_1g(n) \leq f(n)$ we have

$$\forall n \ge n_0. \ g(n) \le \frac{1}{c_2} f(n)$$
 (2 points)

Also, from $0 \le f(n)$ we have

$$\forall n \ge n_0. \ 0 \le \frac{1}{c_2} f(n)$$
 (1 point)

Putting these together we have

$$\forall n \ge n_0. \ 0 \le \frac{1}{c_2} f(n) \le g(n) \le \frac{1}{c_1} f(n)$$

So we choose $n'_0 = n_0$ (1 point), $c'_1 = \frac{1}{c_2}$ (1 point), and $c'_2 = \frac{1}{c_1}$ (1 point).

8. 9 points What is the output of the following Python program?

import math

```
def abs(n):
    return n
class Point1D:
    def __init__(self, x):
        self.x = x
    def distance(self, other):
        return abs(self.x - other.x)
    def abs(self, n):
        return n if n >= 0 else -n
def transform(L, f):
    return [f(x) for x in L]

p1 = Point1D(3)
p2 = Point1D(5)
print(p1.distance(p2))
L = [Point1D(3), Point1D(5)]
```

NT		
- Name:		

```
print(transform(L, lambda p: p.x * p.x))
D = {3: 'x', 'x': 3}
print(D['x'], D[3])
```

```
Solution: (3 points each)

-2
[9, 25]
(3, 'x')
```

9. 10 points Define a function in Python named swap_range with four parameters. The function should swap the elements of two sub-ranges within an array given by parameter A. The beginning of the first subrange is given by parameter begin1 and the beginning of the second subrange is given by parameter begin2. The two sub-ranges have the same length, given by parameter n. You may assume that begin1 + n <= begin2.

Example use of swap_range:

```
A = [1,2,3,4]
swap_range(A, 0, 2, 2)
assert A == [3,4,1,2]
```

```
Solution:
    def swap(A, i, j):
        tmp = A[i]
        A[i] = A[j]
        A[j] = tmp

    def swap_range(A, begin1, begin2, n):
        for i in range(0, n):
            swap(A, begin1 + i, begin2 + i)

4 points for correct logic for swapping
4 points for looping over the array correctly
2 points for swapping the right elements
```

10. 15 points The following code defines a linked-list data structure comprised of a List and Node class. Fill in the blanks of the implementation so that the unit test passes.

```
class Node:
    def __init__(self, d, n = None):
        self.data = d
        self.next = n
class List:
    def __init__(self, array = []):
        self.head = ___(a)___
```

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Name:

```
Solution: (3 points each)
```

- (a) None
- (b) other.head
- (c) n1.data
- (d) n1.next
- (e) True