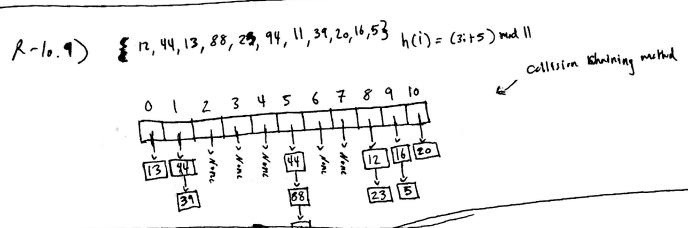
R-10.6) (open harding)

Separate chaining odlision-handling statemes can tolerate a land factor above 1 (although they shouldn't for effectioney), while open addressing schemes regular that the land factor is never greater than 1, and connect tolerate greater land factors.



R -10.10)

0 1 2 3 4 5 6 7 8 9 10

Linear probing method

R-10.12)

0 1 2 3 4 5 6 7 8 9 10

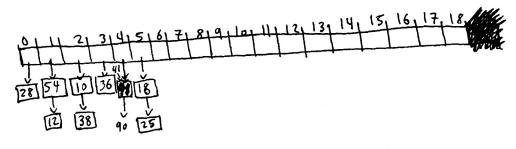
Donble - hoshing method

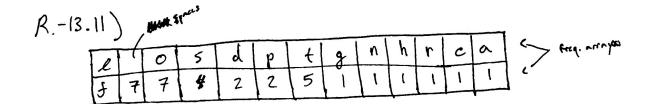
E1: h(i) = (3i+5) mid 11
E2: h'(K) = 7 - (K mid 7)

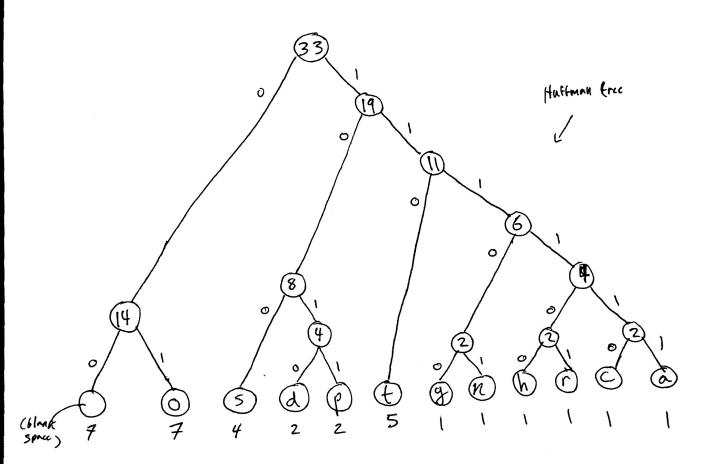
R-10.14)

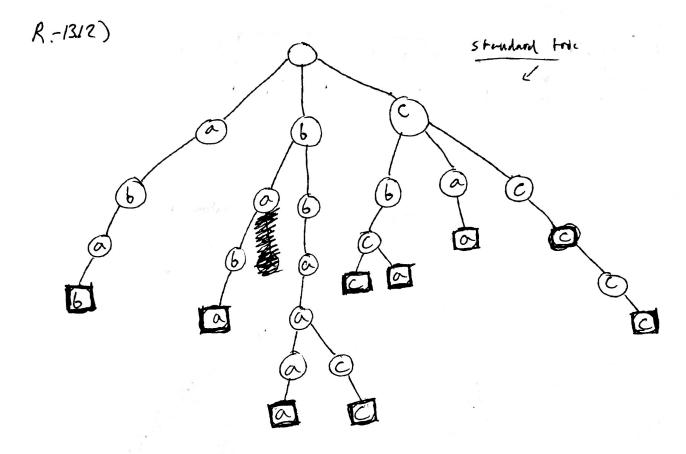
£ 54, 18, 10, 25, 28, 36, 38, 41, 12, 90 }

h(K) = 3K mid 7

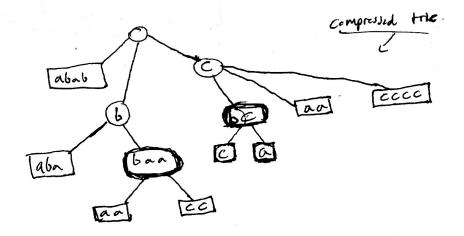






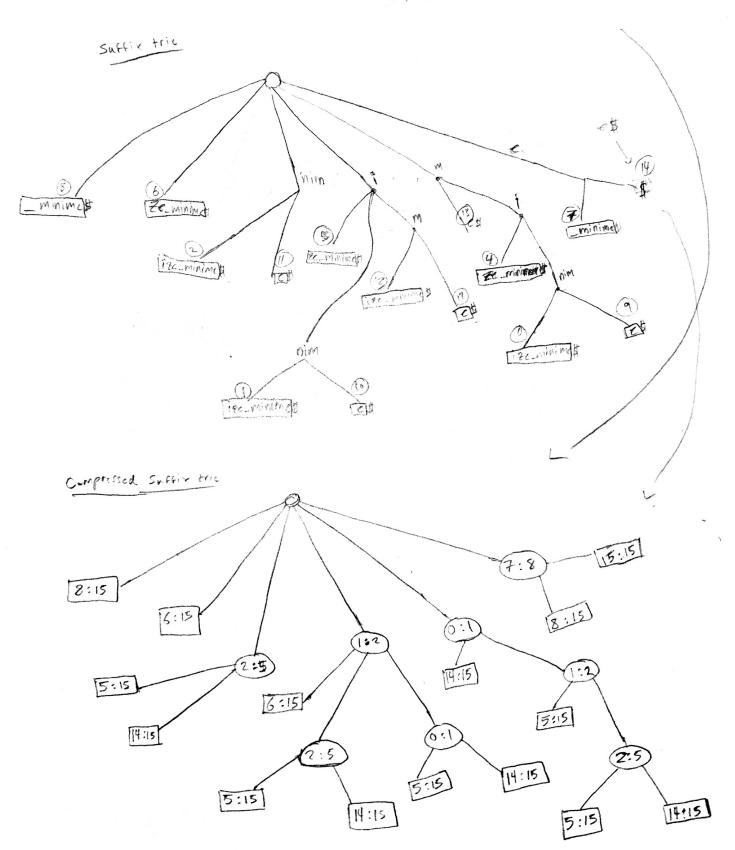


R-13.13)



R.-13.14)

0 1 2 3 4 5 6 7 8 9 9 1 1 12 13 14 \$ minimizer minimizer



Given a standard tric T and a strong s to delete from the tric, search the tric. for the given strong (traverse paths of the tric checking for char's in order of given strong).

If the search were unsuccessful, return the original tric.

Else, but in be node where s was bound.

- check if s is not equal to strong andry at us or has a child.

- if so return tries else :

Bos set ov = u's pront.

- · delete a from tric
- if v has a child
 - · Set v's string to concentration of v's and child string (vestiring terchild string)
 - · then delete V5 child node
- Return the new tric but strong removed.

Time complexity: O(N.M), where n = alphotet size and 102 = depth of leng representing the string to be removed.

R-15.4)

- · Suppose T 15 a maltiply fore of each informal ande having 54 € 58 children.
- · Ench internal mode other than root most have

at limit d-1 kys and contrin at most 2d-1 kys.

Thou, to guarantee a light B- tree, need to guarantee that:

d-1≤5 mm 2d-1≥8,

50 de cm be 5 = d = 8; { 5, 6, 7, 8 3.

