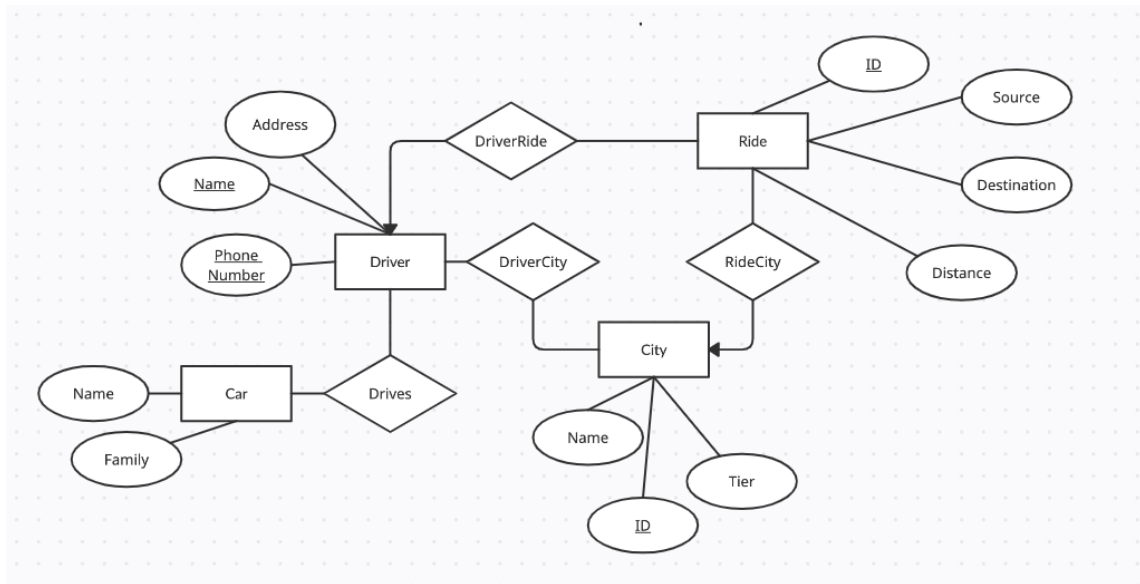


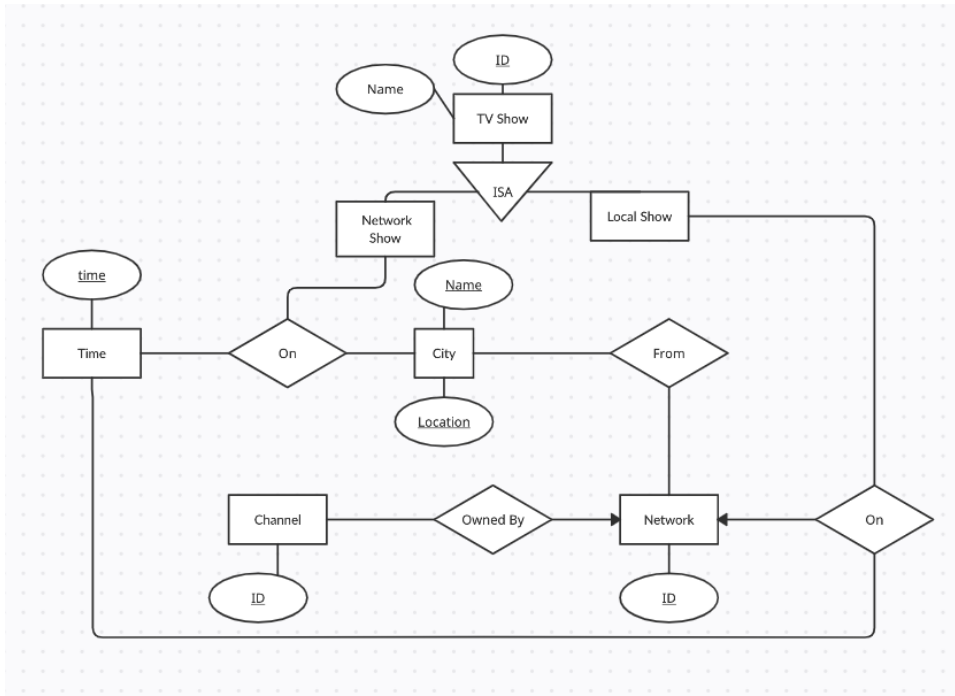
1)



a)

b) Relational Schema:

- i) CREATE TABLE Driver (Address VARCHAR(100), Name VARCHAR(100), PhoneNumber INT, PRIMARY KEY (Name, PhoneNumber))
- ii) CREATE TABLE Car (Name VARCHAR(100), Family VARCHAR(100))
- iii) CREATE TABLE City (Name VARCHAR(100), ID INT, Tier INT, PRIMARY KEY (ID))
- iv) CREATE TABLE Ride (ID INT, Source VARCHAR(100), Destination VARCHAR(100), Distance FLOAT, PRIMARY KEY (ID))
- v) CREATE TABLE Drives (DriverName VARCHAR(100), PhoneNumber INT, CarName VARCHAR(100), CarFamily VARCHAR(100), PRIMARY KEY (DriverName, PhoneNumber, CarFamily, CarName))
- vi) CREATE TABLE DriverRide (DriverName VARCHAR(100), PhoneNumber INT, RideID INT)
- vii) CREATE TABLE DriverCity (Name VARCHAR(100), PhoneNumber INT, CityID INT)
- viii) CREATE TABLE RideCity (RideID INT, CityID INT, PRIMARY KEY (RideID))



- 2)
- 3) Programmer(name), TeamLeader(name,team_name), WorksWith(team_name,name,project)
- 4) Yes, decomposition is lossless. A is the shared unique key and able to get (A,D,E) from R1(A,D,E)
- 5) $(C \rightarrow A)$ and $(A \rightarrow B)$
- 6)
- One to many: sid-> dept,cnum; dept,cnum -> sid
 - Many to one: sid -> dept,cnum
- 7)
- Yes, $\{A\}^+ = R$
 - Yes, $\{BC\}^+ = R$
- 8) No not in BCNF
- $\{A\}^+ = \{A,B,C,D,E\}$
- $\{C\}^+ = \{C,E\}$
- $\{B\}^+ = \{B,D\}$
- All of the above violate BCNF

Key for relation: A,F

$R(A,B,C,D,E,F) \rightarrow$ using $(C \rightarrow E)$

$R_1(C,E)$

$R_2(A,B,C,D,F) \rightarrow$ using $(B \rightarrow D)$

$R_3(B,D)$

$R_4(A,B,C,F) \rightarrow$ using $A \rightarrow BC$

$R_5(A,B,C)$

$R_6(A,F)$

Final BCNF tables:

$R_1(C,E), R_3(B,D), R_5(A,B,C), R_6(A,F)$