

5/3/21.

## Modifiers

### Access Modifiers

- ① public
- ② private
- ③ protected
- ④ default

### Non Access Modifiers

- ⑤ static
- ⑥ final
- ⑦ ~~strictFP~~
- ⑧ abstract
- ⑨ synchronized
- ⑩ transient
- ⑪ native
- ⑫ volatile

1) public → same class, same package need you  
different package need you.

2) private → same class start

3) protected → same package you will get all  
package of child class

4) default → same package you.

1) static → ① data member → memory address storage  
class and static

② final method → related with non static class  
class regular final, if object or static method  
can not change.

\* Local static. static block.

2) final → ① data member → to create constant variable

② final methods → prevent method overriding

③ final classes → prevent inheritance

\* we must initialize final variable.

3) synchronized

- multithreading help multiple threads same  
resource at a time access.

- help about thread can access the  
resource.

- just a way to do it by using synchronized  
keyword like it execute block by  
using "Monitors".

④ transient → serialization  
object in conversion file of a program or  
read serialization.  
transient keyword serialization elidan..  
security purpose nilai.

⑤ Native → java di hardware & interaction  
c/c++ & mod.  
c/c++ help function will interact with Java  
Native

⑥ abstract  
function declaration doesn't have keyword  
but it has  
multiple function definition under one child class  
needed multiple file. (one class has multiple definitions)

## ⑦ strictfp.

java FPU standard

① GPU → graphical processing unit → six CPU  
are used which one of distributed send work  
each

GPU all support rendering will elid.

② FPU → floating processing unit & standard  
float standard number down to be keyword  
will process or float & action of float  
will strictfp interface of float some of floatosis.

## ⑧ volatile.

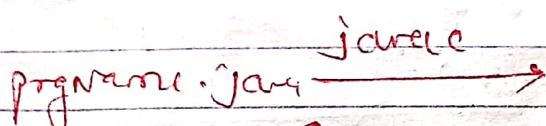
Volatile int \*p = (int \*) mem\_alloc → don't perform any  
optimization on p.

anything which can change beyond compiler scope  
you should declare its volatile.

## Computer Architecture

## JVM Architecture

Java's class loader takes input from command line arguments or environment variables. It uses the classpath variable to find classes. The classpath variable is a list of directories and JAR files where Java looks for classes. The classpath can be set in the environment or specified on the command line using the -cp or -classpath option.



→ Rem as draft (process)

javac Rem as our class will inherit  
from a main file itself.

- Q. \* :-

  - pdf
  - ~~exe~~ exe.
  - pptx

:- —

—

ই সিনেট ওস ঢাক্টেলি কো  
মন এবন গুড়ি.

ਮੁਹ ਏਵੇਂ ਨਾਈ,

$\Rightarrow$  Market has two main applications.

All written work on the first two lines

But each state selects its own method of calculating its

31401 અને એટાનિમ ફોર્મ ક્રીડા પ્રેરણ

i. st à affric type ii st ð m n s ð z ð.

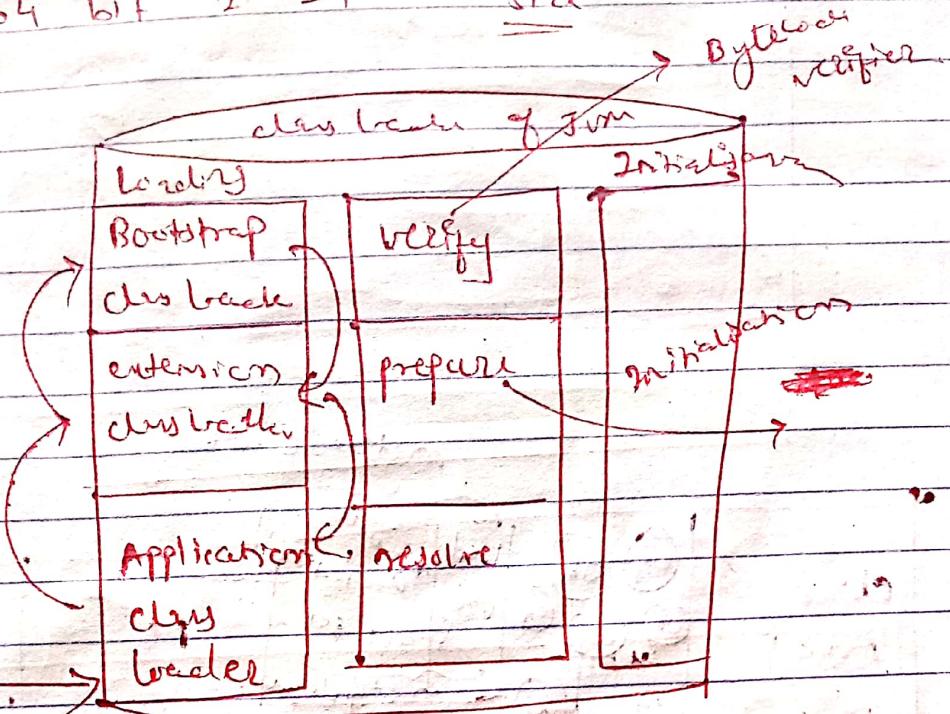
Device driver → software company printer model (ex)

Third device driver quick. And o.s mismatch

211 device 27 interaction and control.

\* Virtual address       $2^{32} - 1$        $2^{48} - 1$        $\rightarrow$  ~~generate from CPU~~  
 32 bit      64 bit      physical address to RAM

Size



- \* A.C.L dield Cell E.C.L mi dili fogni cell  
 $B.C.L$  mi noi eniute elnra  $B.C.L$  greti of 12  
 extension day loader ohe Application.
- \* 8120). 20100111 Heratil predefined classeun

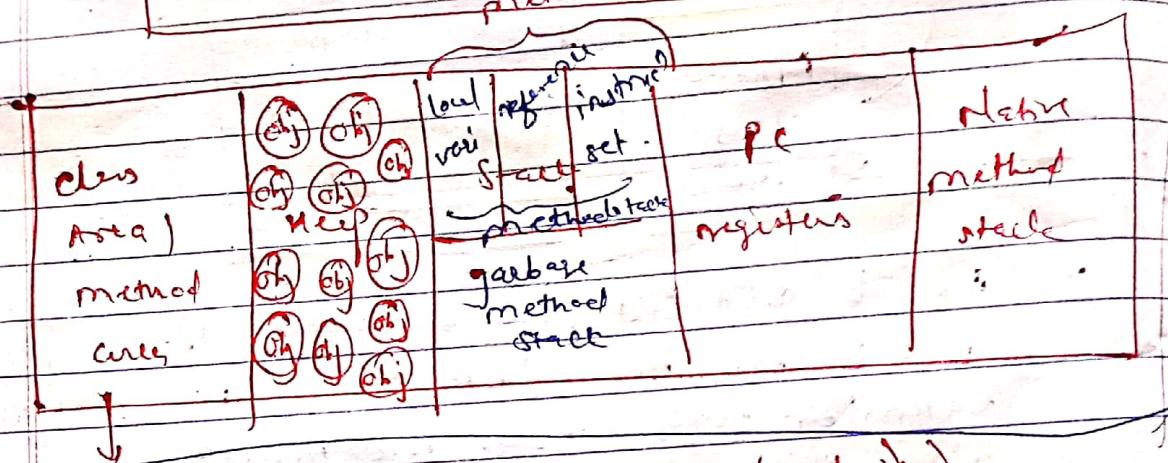
prepare  $\rightarrow$  static field mi memory aifni initialised  
 $\rightarrow$  0.

Initializaton  $\rightarrow$  static field initialised  $\rightarrow$  user defined  
 value

static int a=10      static aifni miel aifni initialise  
 aifni nra 10 aifni 8120)

1111

static int a=fun() e for mi cell aifni aifber  
 aifber eniuet ali 8120 0 mi initialise  
 data : datad.



class area method = ein class local eln.  
 area = global line instruction block  
 stack eln, variables uo  
 instruction block ~~stack~~ block. mit memory den  
 geln

global function cell und local d' func' zog  
 stack stack as also execution mit mit  
 initial variables mit value geln.

Heap is global object heap as well as d'.

stack → function execution.

- A) ① method stack → method in cell ~~instruction~~  
 zgi d' mit execution  
 mit 3 part.

- ① local variable → func<sup>n</sup> frame local variables in memory stack.
- ② reference & heap are same smart objects with null reference etc. main.
- ③ function sets → execution return instructions.

### B) Garbage method & stack →

~~garbage~~ method need object address allocation, then memory address along with heap going delete scenario in garbage method stack work.

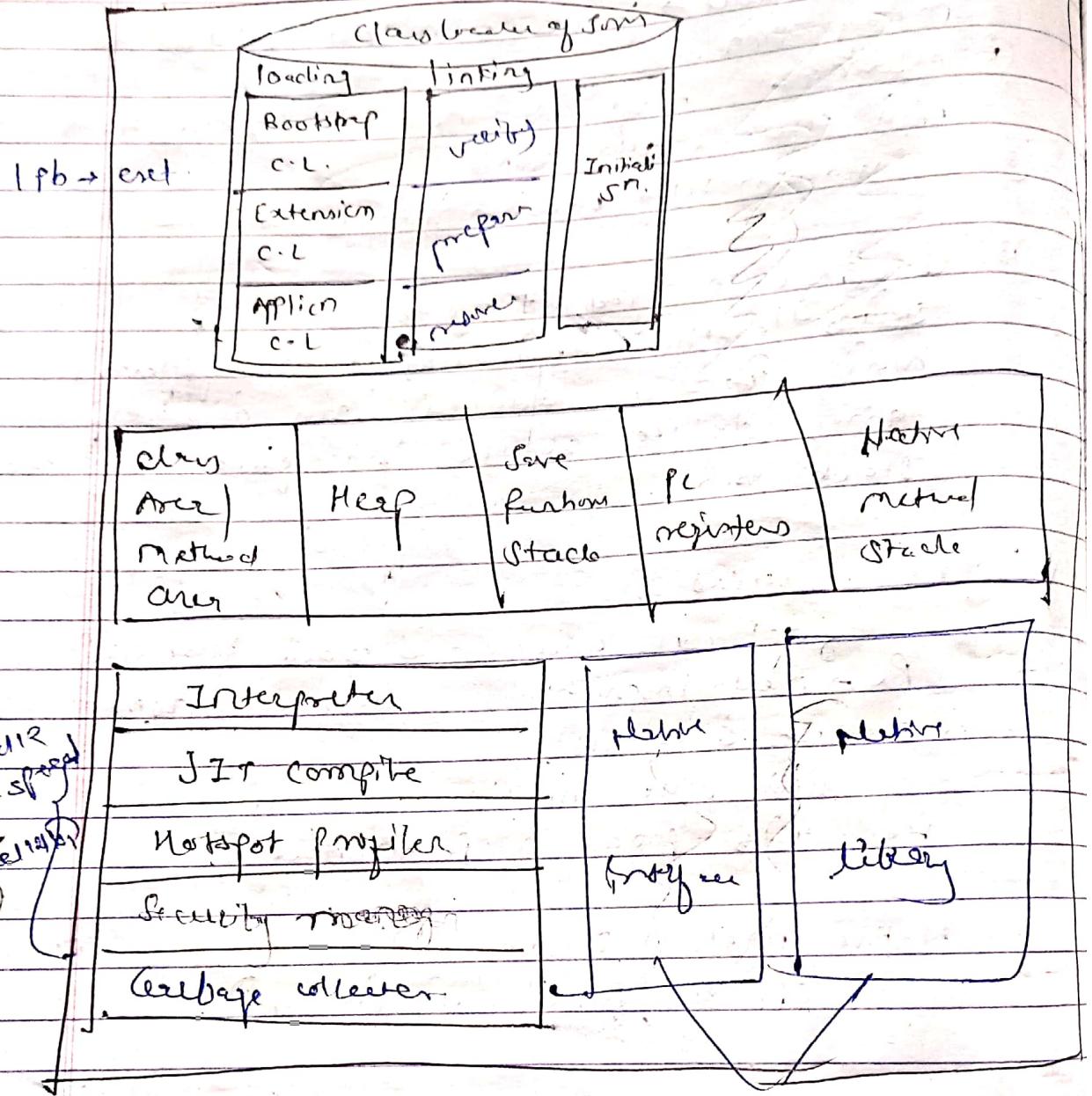
### C) PC registers →

function execution is done sequentially.  
PC → next instruction, current instruction is trailing.

also in multithreading use same PC register process in PC registers.

### D) Native method stack

c, c++ → func<sup>n</sup> run simultaneously.

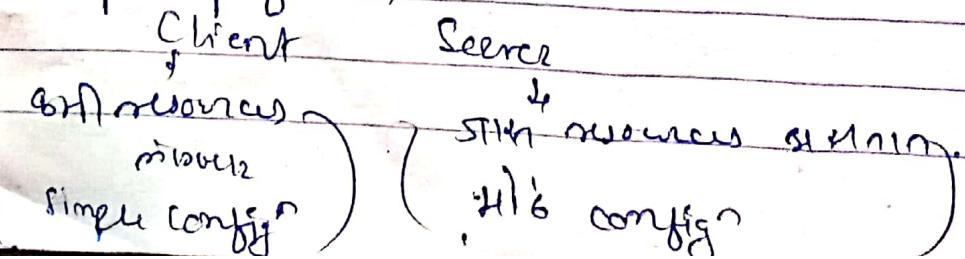


JIT → Just in time

C/C++ →

Interpreter → line by line code execute *slow*  
 Interpreter → line by line code execute *slow*  
 Compiler → step by step fast *fast* *slow* *slow*  
*slow* *fast* *slow* interpreter as in *slow* *slow*  
 compiler *slow*.

Hotspot profiler →



572 Seeser all setup done, no to setup  
72112 0910am | 1011 am 0910am  
and user in 0910am 0910am 0910am  
start | setup done 0910am

by default client is the size  
clients select the size of switch size  
java -server demo ← ← 240 kb size  
the server file must be small.

## Security manager

- Multithreading → Synchronized keyword  
of of security in Java for key word security  
and aligned in environment like class  
program execution & (b)  
application security provide security.

given the code runs free till to a point  
very (byte code verifier error)

multiple threads can run at a time.

## Garbage collector

Garbage collector

Garbage collector determining garbage  
memory  
state

Stack or LIFO area shared by all

Object will be deleted when it goes out.

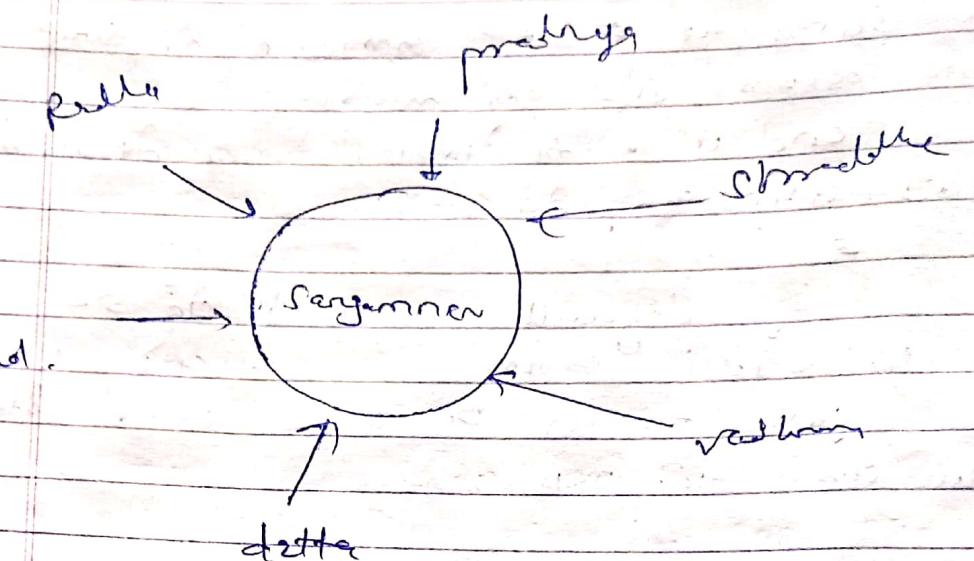
if memory can't be collected then  
dynamically allocated objects will

reference Count → `45 laptop` has a reference count = 3

Local resource manager 234 का क्या करता है

मार्ग से का reference count बढ़ाता है

31



Predatory

प्रिया जो Senganner को किसे कहती है

इसका memory location

Predatory → Senganner 102

∴ dynamically object को किसी भी reference variable की तरह  
किसी भी object को किसी reference variable की तरह

demol = new demol();

30

garbage collector का किसी भी object का किसी  
reference count को 0 के रूप में delete करने की क्षमता  
जो stages की है उसके लिए किसी भी विद्युत  
का reference count को 0 के रूप में delete  
करने की क्षमता है। इसके लिए किसी भी stages  
की direct destroy की है। किसी भी stages  
की है। (जो की क्षमता है)

⇒

Heap memory का value का किसी garbage  
collector का किसी garbage

## \* Native method interface

Java to communicate C/C++ code

through XML files define Java

data communication mechanism

interface.

## \* Native method library

C/C++ code communicate

library

resolved (Symbolic  
link)

O.S. OS

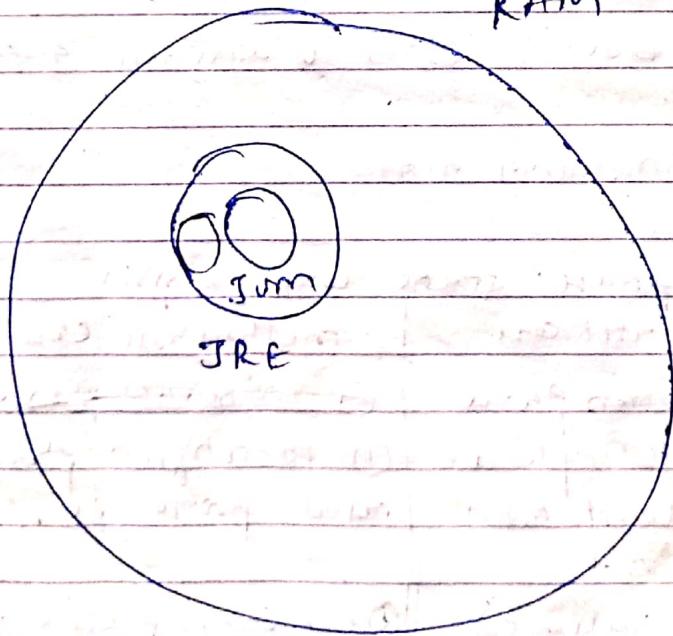
JVM < Virtual < MM < RAM

address address

(Virtual)

than mapping.

RAM



JRE word

JVM program static method. ∵ JVM program will be called from native  
program or native program will be called from JVM.

## Static & non static

- (1) Static variable / field
- (2) static method
- (3) static block → loading will do at execute time.

### \* Static field

\* static local within obj  
→ static method uses local level obj  
↳ a) class level object can have static obj.

### \* static field

512 memory object having 512 bits  
field will same value across the objects  
↳ common static field of object will be shared  
among static fields are called static variables  
↳ static field is  
eq → Arcot

→ static block is class loading will do at execute time.  
main memory obj.

### \* 4 phases static local memory

field will come	method will come
① Initialization phase	② Initialization phase
③ Identification phase var. will load obj memory location default value for.	④ Identification phase method prototype assigned
④ Execution phase → user defined value assign to it.	④ Execution phase → function in cell structure will prototype method by algorithm execution acci.

## static method

method mi cell aro 2021 mali object nahi re  
call by class name.

object nahi abhi static function  
call kرنi se hi non static function hi available  
hi nahi na function object nahi den  
memory varage hain. static and method

.. static mi .. object ka place nahi.

## static block.

Static

{  
}  
}

static block return

① static field mi value hain.

② static block he, private function mi hao  
~~static block return~~ aao ~~function~~ share ho  
hi call kar nahi.

Abhi return execute ho rahi hao code hao  
den main & main method ka return.

new load / execute private method static  
den hao.

\* Prg

Q. → static block ka return function static function  
mi call kar nahi ~~return~~

→ ct.

Class ka nahi call hao.

execution mi a method static ka return.

a method si aur static block hain.

Return value ki form or direct ho leti.

## Encapsulation

(methods)

Binding of fields & member functions

~~object orientation~~ ~~with~~

- ① object
- ② class
- ③ encapsulation
- ④ abstraction
- ⑤ inheritance
- ⑥ polymorphism

Bank

clerk

{

~~private float balance;~~

{

  |

  |

  |

  |

  |

  |

}

j.

Clerk of balance will make valid changes

will use a valid rule on itself & bank rule

itself. no setter method

Bank will balance clerk will tell him

\* abstraction is no paradigm of O.O -



data abstraction

field (var) number  
hide access.

method abstraction

→ implement hide access

C++ → pure virtual func

Java → abstract func

→ interface implements

access modifier method

abstraction

## \* final Keyword \*

C/C++ Has const.

const int a = 10;

a = 20 reassign does not work.

final keyword →

- ① final field (datamember) } constant
- ② final local variable } cannot immutable
- ③ final method (function) } member data types
- ④ final class. } final.

- non final method will inherit final local variable else it will.

### ① final field:

class Demo

{

final int a;

p. S. v. m (String arr [ ] )

{

Demo d = new Demo();

} final class will initial static address & initialized

reference variable & final class all initial parent class will

↓

final, return parent class will

initial.

String s = new String ("Sanginner");

Stack reference s = 100

variable new & dynamically memory allocate etc.

Heap de

Heap

s sanginner  
→ 100

s.concat ("422605") s →

8anginner ← s  
100 Hash code

sanginner 422605

sanginner  
422605  
100 pointer add.

String is immutable

∴ s will not change concat

immutable → cannot change

immutable data types

(1) primitive

(2) String (final class)

final

String Buffer - Mutable class.

mutable object state ~~will~~ can be changed.

StringBuffer s = new StringBuffer ("Engguru").  
s.append ("922605")

new

S=1100

StringBuffer

दि value को  
नहीं नया object  
implan change  
होता है।

हर immutable हो सकते हैं  
जिसका state बदलना कठिन हो।

① प्रारंभिक स्टेटमें  
class Demo

{  
final int a = 10;

Not final

object

p. J. v. m (String args[]) {  
value

final

object

value

final

object

value

final

object

value

constructor of class having default value for all fields.  
Constructor will have default values for all final variables cannot be assigned  
than initialise

(2) Constructor with static class Demo

```
class Demo {  
    final int a;  
  
    Demo()  
    {  
        a = 10;  
    }  
}
```

(3) Non static block will run | instance block  
↓ | block constructor will run  
with object creation will other  
block will run.

```
class Demo {  
    static { "ANCOB" }  
    final int a; { Demo.collegenme }  
    { static final int a; }  
    a = 10; }  
}
```

↓ |  
↓ ↓  
↓ ↓ ↓ ↓ ↓ ↓  
Static static block will execute first and  
static will share static block  
which has static block  
will value get shared among them  
variable static will compare  
e.g. collegename final + static both  
↓ & non static same

field mi se rani chal hiya value (10) oka  
nhi error hoga. constructor me se pani oka or  
default value final hoga.

- ② final local variable  
method ka hoga

class Demo.

{

void fun()

{

final int a; ekko defult value hoga  
na hoga blank final  
variable  
21(m) error hoga.

}

blank final var. nahi

value dhya kro se hoga.  
assign chalo kro na hoga.

value of ahi direktly yahan hoga  
hera error hoga.

- ③ final method.

class Demo.

{

final void fun()

{

void fun()

{

}

}

~~and a class extends another class~~

class test extends Demo

{

void fun ()

error also find methods

{



cannot compile

}

void fun ()

body has changes

{

}

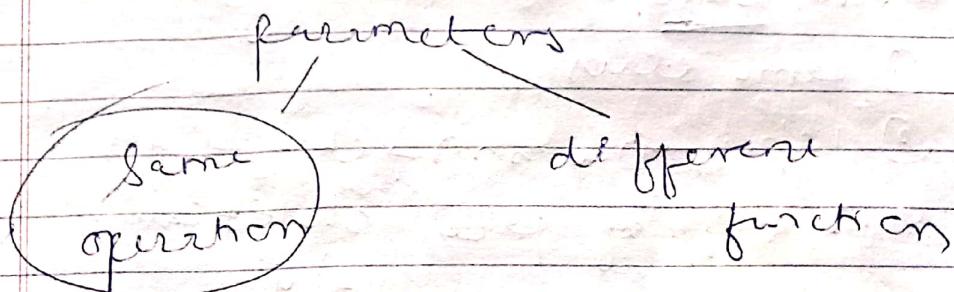
4) final day → cannot be inherited.

inheritance prevention.

## Method overloading

Same scope

Single name multiple behaviour



prototype

compile time

polyorphism

early binding

~~late bind.~~

dynamic  
binding

parent w/ reference

in object child w/

operator overloading → } implicitly  
obv. in w.

String s1 s2 + concat

+ add?

↳ ~~char string print etc.~~  
↓ + + ~~as m. concat~~  
println( ) etc.

PH 12/2021 scope Heir. overriding

- (1) same class
- (2) inheritance  $\rightarrow$  child overloads  
that (i) type of func. child + parent  
(2) No. of func. ~~not~~
- (3) sequence of func.

C++ used global function

overload as in 9/19/21

5/2 parent will be static

func child will be non static

2/12/21 d2/20/21 static func will be static

6/2 in { (1) type } method overriding

(2) no.

(3) sequence

method overriding will give

new modifiers will not apply

2/11/2021

method overriding

Access specifier (public, private)

Size child is same as parent

Access - Public child is access

modifier + public

class demo

{

void add (float f)

{

}

void add (double d)

{

}

{

demo d = new demo();

d.add (10);

}

int

fun (value).

fun (type) value);

fun (a+b);

fun (fun());

↳ non void.

① exact match

② Java type need convert

int → var args

3.

widening

var args

byte(10)

var args

fun(int)

widening

③ autoboxing

④ autoboxing widening  
int a=10;

a. 

10
----

 method by

100 104 assign

new  
box(a)

Integer a=10; implicitly

new Integer(10);

a. 

100
-----

10

obj

a. method();

wrapper classes

(class demo)

{

wid fun (float x)

{

}

wid fun (Integer x)

{

}

main

Demod = new Democ();

d.fun(10);

float x = 10, widening

Integer x = 10; Autoboring

Still widening & no SIR

① return no autoboring & ②

exact autoboring & no SIR

autoboring widening certain.

overloading here ① reference anyway

↳ 2nd 3rd ② parameters

Q2

① copy

② call by value reference Q2

# Arrays

`Demol d = new Demol();`

ते या तीन नोंद value.

तुका एक अलग category  
value है १०२।८।६।७।

`int a[10]; c/c++ def`

\$1201 memory size.

\* memory का क्षेत्र के नाहीं runtime  
में लगता.

~~ज्ञ~~ variable

primitive reference

int Integer

float Object

float

Array

`int a[10] नोंद आउट गरी गया।`

क्या ? 

type and value of



global variable

\* i) declarations

\* ii) definitions

object → runtime → new

optional syntax

{ <access modifier> <non-access modifier> }

datatype variable-name

[ ], ,

declarat<sup>n</sup>

int a[];

int array      memory allocation  
                  chart type

a will allot memory in stack.

Java has ~~address~~ pointer

→ concept of pointers

int a, [ ] b; error

declare int[] a;  
int a[];  
int []a;

1D array

int a[][];

int []a[];

int [][]a;

2D array

int [] []a;

int [][]a;

int [] []a;

int [] a, b; → a[], b[].

int [] a, b; → a[], b[]

which means at int. allocation

int a[], b → a[], b is integer  
error var

int a, b; a-variable } Compile error

int a, b[]; → a int var | b array

comma separated elements

the subscript gives different value than a)

multiple variables with same name

different data type struct and

\*. decl with size - dim 7/8  
size must same as declared.  
`int a[10]; X`

defn → value initialisation  
→ using new keyword,  
explicit

i) decl + defn; initialisation

`int a[] = {10, 20, 30, 40};`

implicitly new requires  
alloc object new arr.

a is reference.

② `int a[]; a, [ ]`

`a[0] = 10 X memory is 10 bytes`

11/12

a → reference (here ~~new~~ will save)

`a = new int[10];`

∴ Array is reference type

-11

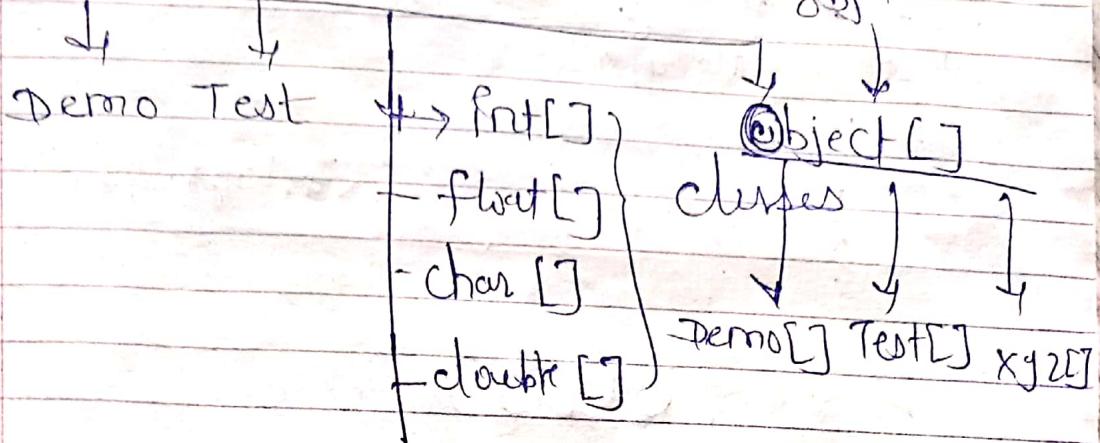
object array type -11

Object is parent class and array & function are its objects  
object[] classes are simple array type classes  
(length) ~~length~~ inheritance etc.

Object

→ ~~single~~ parent class

fixed array.



10, 20, 30 primitive

Integer → wrap for classes.

primitive data object ref.

Ex int[] → class

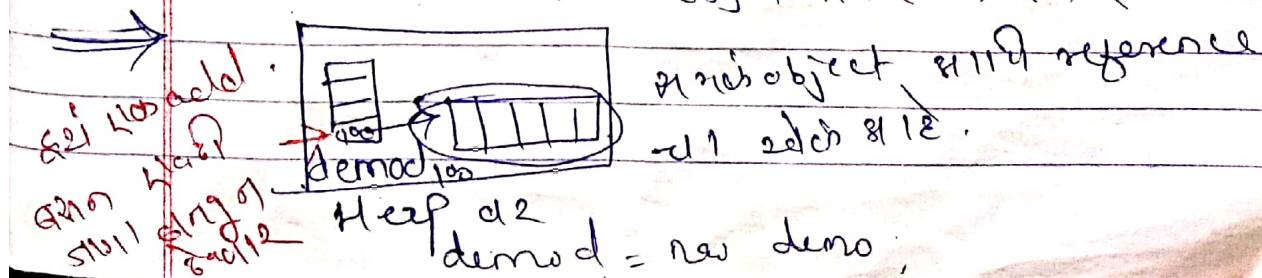
primitive data primitive regular  
variable.

int[] result contains address of memory  
allocated by OS.

element

int a

Q. array is an object or not? Ans: Yes  
and answer is object have memory



\* (Hed)

n initiall igitate  
n = size of array  
total size.

for (int i=0; i<a.length; i++)  
{

s-o. P(a[i]),

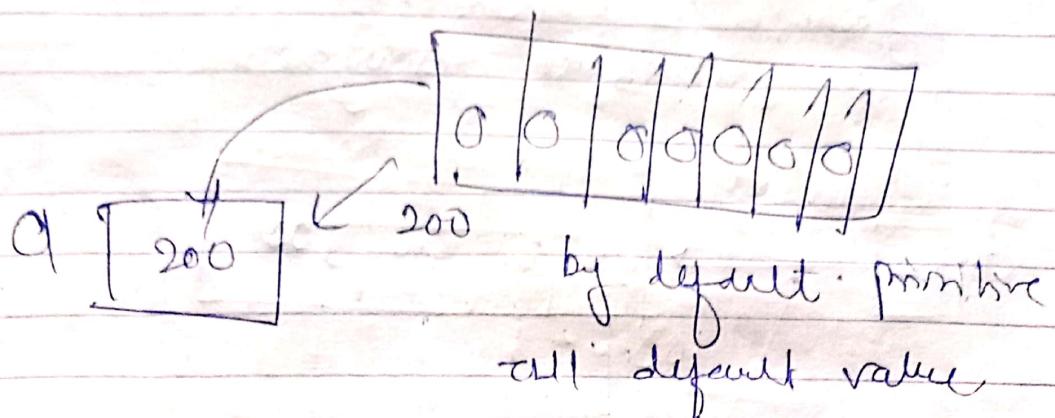
j

\* gets array total length → Norathi  
→ jfirst

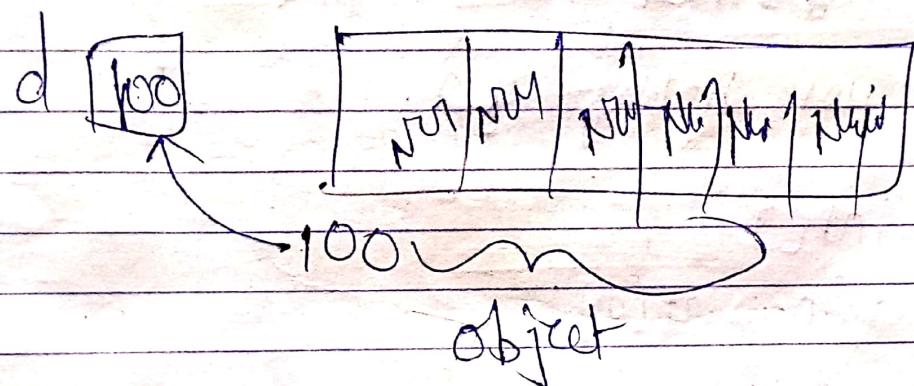
otherwise impudely

2) variable object [] array ->

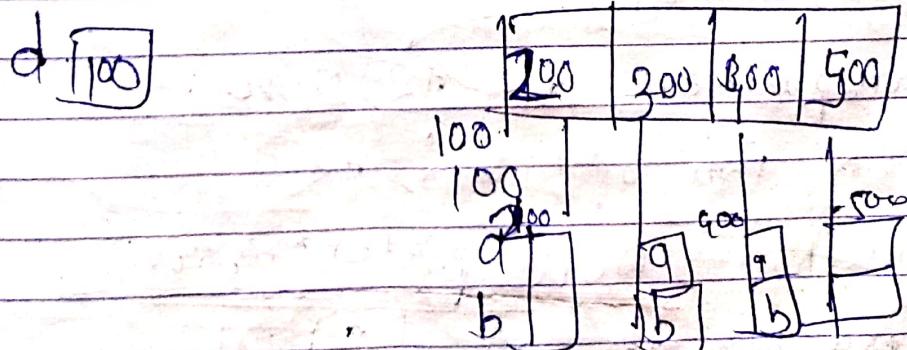
`int [] a = new int[10];`



`Demo [] d = new Demo[10];`



`Demo [] d = new Demo, new Demo, new  
Demo, new Demo,  
new`



## Var - arg

fun(10); fixed parameter

fun(10, 20)

std progs have func def.

multiple values will multiply  
para intell array will be created

int a[] = {10, 20, 30};

fun(int a[]). fun(a);

var-arg →

void fun(int... a)

{

a.length =

for loop it can go beyond

}

fun(25)

fun(10)

fun(10, 20),

10 20  
25

fun(new int[3]{10, 20, 30})

Program

\* Multidimensional array

abstraction

Concrete Class →

abstract class

concrete method

abstract method

interface

abstraction

data  
abstraction

using  
private  
keyword

Access

specifier

method abstraction  
(implementation  
abstraction)

method w/ private  
dim no return  
syntax error

## { Interface

interface Demo.

abstract method,

abstract void fun();

abstract void gen();

}

Java 8 till y6 static final

default method interface till

min implement abstract class  
done default.

static final default keyword

done.

## { interface Demo

void fun();

}

} till now.

min till now abstract method

✓ class  $\xleftarrow{\text{extends}}$  class

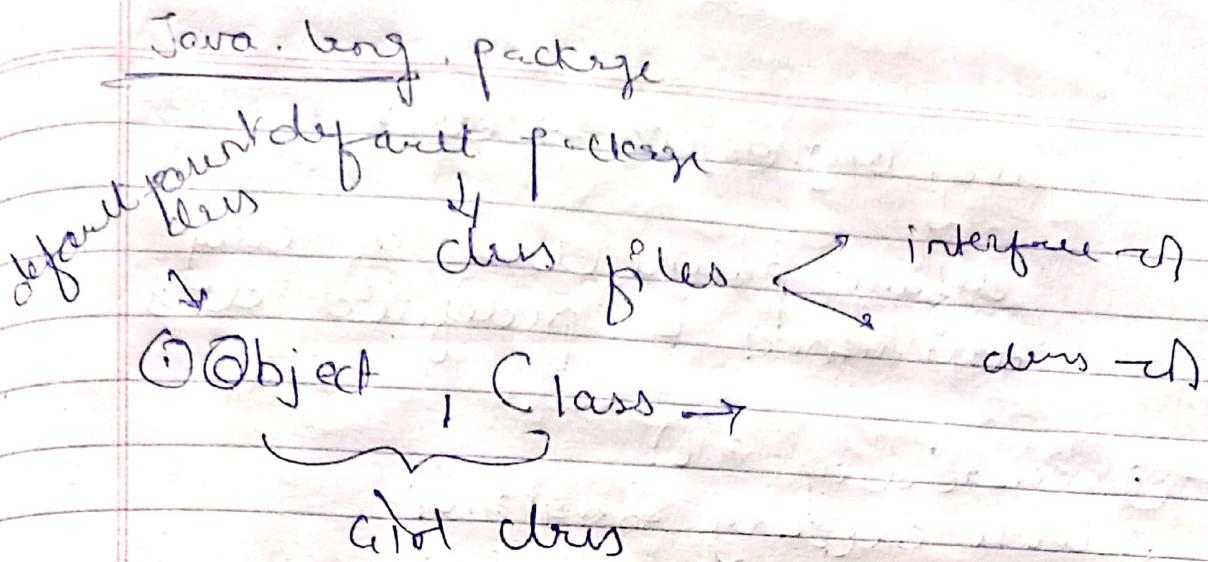
✓ class  $\xleftarrow{\text{Implements}}$  Interface

✓ I  $\xleftarrow{\text{extends}}$  Interface

Int  $\xleftarrow{\text{extends}}$  ~~CCay~~

I  $\cancel{\xleftarrow{\text{Imp}}}$  C

C  $\cancel{\xleftarrow{\text{Imp}}} \quad C$



② String, StringBuffer, StringBuilder

③ Wrapper classes

④ Throwable, Exception, Error

⑤ System, Random, Process, ClassLoader  
SecurityManager

⑥ Thread, ThreadGroup  $\rightarrow$  multithreading

\* Class  $\rightarrow$  Super class / interface class method can  
be overridden in sub class / interface class

↳ class cell / object can

JVM (22)

java. lang. Object → class

↓  
default parent class of every  
user defined & predefined class

- ① accessibility 21/16)
- ② loose coupling 20/24/21/16)

parental reference 21/21/21

↳ *function* 21 4/8d7/17) direct print  
→ 21 reference test,

parental reference 21/11 object 21/11/16)  
and 21/1 class 21/1

① public native int hashCode();

② public boolean equals ( Object obj );

③ public String toString();

④ public final native Class getClass();  
Class getclass();

⑤ protected native Object clone();  
throws CloneNotSupportedException

⑦ protected void finalize() throws  
Throwable;

⑧ wait();

⑨ wait(-);

⑩ wait(-, -);

⑪ notify();

⑫ notify();

clone → its object's copy

→ all own changes will not reflect.

original & compare opertions etc etc

Shallow copy

deep copy

\* public String toString();

class Object {

{

public String toString()

{

return

}

{

S. op ( d. tostring() )

St. p( d );

→ class name @ hashcode is hexadecimal  
format.

3114011 calls → member print 3111

411851.

Hot override ob 31 . 31P1.

public String toString()

{

return this.x + " " + this.y;

}

String

at data member print 3111.

\* public native int hashCode();

→ 3114011 calls hashCode current

31110112

∴ it overrides.

hashCode() → virtual address ~~3111~~

C/C++ 311011 31121

∴ it method native 3111.

return at address → int.