

# Propositional Equivalences

- Discrete Mathematics (Kenneth Rosen)
  - 8<sup>th</sup> edition – Section 1.3

# Motivation

- How do we simplify complex logical statements?
- Can we replace a proposition with an equivalent one?
- Why it matters: circuits, automating proofs, solving puzzles, checking correctness of a program, checking feasibility etc.

# Definition of Logical Equivalence

- $p \equiv q$  means  $p \leftrightarrow q$  is a tautology.
- Example:  $\neg(p \wedge q) \equiv \neg p \vee \neg q$

# Tautology, Contradiction, Contingency

- Tautology: always true (e.g.,  $p \vee \neg p$ )
- Contradiction: always false (e.g.,  $p \wedge \neg p$ )
- Contingency: sometimes true, sometimes false (e.g.,  $p \wedge q$ )

# Common Logical Equivalences

- Identity:  $p \vee F \equiv p$ ,  $p \wedge T \equiv p$
- Domination:  $p \vee T \equiv T$ ,  $p \wedge F \equiv F$
- Idempotent:  $p \vee p \equiv p$ ,  $p \wedge p \equiv p$
- Double Negation:  $\neg(\neg p) \equiv p$

# More Logical Equivalences

- Commutative:  $p \vee q \equiv q \vee p, p \wedge q \equiv q \wedge p$
- Associative:  $(p \vee q) \vee r \equiv p \vee (q \vee r)$
- Distributive:  $p \wedge (q \vee r) \equiv (p \wedge q) \vee (p \wedge r)$

# De Morgan's Laws

- $\neg(p \wedge q) \equiv \neg p \vee \neg q$
- $\neg(p \vee q) \equiv \neg p \wedge \neg q$

# Implication Laws

- $p \rightarrow q \equiv \neg p \vee q$
- Contrapositive:  $p \rightarrow q \equiv \neg q \rightarrow \neg p$
- Note: Converse ( $q \rightarrow p$ ) and inverse ( $\neg p \rightarrow \neg q$ ) are not equivalent to  $p \rightarrow q$ . Are they equivalent to each other?

# Biconditional Equivalences

- $p \leftrightarrow q \equiv (p \wedge q) \vee (\neg p \wedge \neg q)$
- $p \leftrightarrow q \equiv (p \rightarrow q) \wedge (q \rightarrow p)$

# Puzzle: Which Are Equivalent?

- A.  $\neg(p \wedge q)$
- B.  $\neg p \wedge \neg q$
- C.  $\neg p \vee \neg q$
- D.  $\neg(p \vee q)$

# Puzzle: Simplify to p

- 1.  $(p \vee q) \wedge (p \vee \neg q)$
- 2.  $\neg(\neg p \wedge \neg q)$
- 3.  $p \vee (p \wedge q)$
- 4.  $p \wedge (p \vee q)$

# Truth Table Challenge

- Simplify:  $\neg(p \rightarrow q) \vee (\neg q \rightarrow \neg p)$
- Hint: Use implication identity and De Morgan's laws

# Application: Simplifying Circuits

- Use equivalence to reduce logic gates.
- Example:  $(p \wedge T) \vee (p \wedge F) \rightarrow p$

# Satisfiability Problem (SAT)

- A formula is satisfiable if some assignment makes it true.
- Unsatisfiable if no assignment makes it true.
- SAT is the first known NP-complete problem.

# SAT vs Tautology

- Satisfiable: at least one assignment is true.
- Tautology: all assignments are true.
- Example:  $p \vee \neg p$  is both;  $p \wedge \neg p$  is neither.

# Solving SAT – Brute Force

- 1. List all assignments
- 2. Evaluate the expression
- 3. Check if any assignment makes it true

# Solving SAT – Smart Techniques

- Use equivalence rules and simplification.
- Convert to CNF for SAT solvers.
- Apply DPLL, resolution, etc.

# Satisfiability in Daily Life

- Scheduling constraints
- Smart home automation logic
- Games like Sudoku, Minesweeper

# Knights and Knaves – Setup

- Island where:
  - - Knights always tell the truth
  - - Knaves always lie
- You meet A and B.
- A says: 'B is a knave'
- B says: 'A and I are opposites'

# Knights and Knaves – Case 1: A is a knight

- If A is a knight → B is a knave (truth)
- B says 'A and I are opposites'
- But knaves lie → statement must be false
- So A and B are not opposites ⇒ contradiction
- → A cannot be a knight

## Knights and Knaves – Case 2: A is a knave

- If A is a knave  $\rightarrow$  'B is a knave' is false  $\Rightarrow$  B is a knight
- B says 'A and I are opposites'  $\rightarrow$  true (B is a knight)
- Consistent with all statements
- Conclusion: A is a knave, B is a knight

# Truth-Teller, Liar & Random – Setup

- Three people:
  - One always tells the truth
  - One always lies
  - One answers randomly
- **Three Doors**
  - One leads to **freedom**
  - The other two lead to **doom** (or nothing)
- You can ask **one yes/no question** to one person
- Goal: Find the correct door to freedom

# Truth-Teller, Liar & Random – Hint

- Can't rely on random person
- Need a question that works for both truth-teller and liar
- Use a self-referential question:
- Something similar to :  
'If I asked you whether door A is correct, would you say yes?'

# Truth-Teller, Liar & Random – Logic

- Truth-teller answers truthfully about what they'd say
- Liar lies about what they'd say (double negation)
- → Both give consistent answer!
- Go to door A if answer is YES

# Logic Grid Puzzle – Setup

- People: Alice, Bob, Charlie
- Drinks: Tea, Coffee, Juice
- Information Given as follows:
  - Alice didn't order coffee
  - The person who ordered juice is not Bob
  - Charlie didn't order tea
- Who Ordered what?

# Logic Grid Puzzle – Deduction

- From the information:
  - Alice → tea or juice
  - Bob → tea or coffee
  - Charlie → coffee or juice
- Try assigning:
- Alice → Tea, Bob → Coffee, Charlie → Juice
- ✓ All the constraints are satisfied.

Take Home Assignment:

Solve all this puzzle using propositional logic.

# Summary

- Equivalence: Replace and simplify logic expressions
- SAT: Determine if some assignment satisfies a formula
- Logic applies to daily life, puzzles, AI, and computing.