# Features of (Un)decidable Logics

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February 28, 2025

Student Logic Colloquium

# Plan for the talk

- (Un)decidability: what and why?
- Propositional team logics and their decidability
- Exploring boundaries between the decidable and the undecidable
    - Solving problems and obtaining insights along the way
    - Using insights to solve one last problem

# (Un)decidability: what and why?

What?

A decision problem is a collection of inputs $I$, with a yes-or-no question for each $i \in I$.

A decision problem is decidable if there is an algorithm that, given any $i \in I$, accurately answers the question. Otherwise, it is undecidable.

A logic $\mathbf{L}$, in a language $\mathcal{L}$, is decidable if there is an algorithm that, given any $\varphi \in \mathcal{L}$, determines whether $\varphi \in \mathbf{L}$. Otherwise, it is undecidable.

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**Why?** *Because it is a profound conceptual distinction.*

# Propositional team logics and their decidability

Traditionally (in, e.g., CPL), formulas $\varphi$ are evaluated at single valuations $v : \mathbf{Prop} \to \{0, 1\}$,

$$v \models \varphi.$$

In team semantics, formulas $\varphi$ are evaluated at sets ('teams') of valuations $s \subseteq \{v \mid v : \mathbf{Prop} \to \{0, 1\}\}$,

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**Definition (some team-semantic clauses)**

Let $X := \{v \mid v : \mathbf{Prop} \to \{0, 1\}\}$. For $s \in \mathcal{P}(X)$, we define

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| $s \models p$ | iff | $\forall v \in s : v(p) = 1,$ |
| $s \models \varphi \wedge \psi$ | iff | $s \models \varphi$ *and* $s \models \psi,$ |
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**Observation.** All propositional team logics are decidable: given $\varphi$, simply check whether $s \models \varphi$ for all $s \subseteq \{v \mid v : \mathbf{Prop}(\varphi) \to \{0, 1\}\}$.

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What is it that makes propositional team logics decidable, *and others not?*

# Team semantics as relational semantics

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$$V(p) := \{s \in \mathcal{P}(X) \mid \forall v \in s : v(p) = 1\} = \downarrow\{v \in X \mid v(p) = 1\}.$$

**In fact**, if we take all powerset frames $(\mathcal{P}(X), \cup)$, redefine the base clause

$$(\mathcal{P}(X), \cup, V), s \Vdash p \qquad \text{iff} \qquad s \in V(p),$$

and only allow principal valuations $V : \mathbf{Prop} \to \{\downarrow s \mid s \in \mathcal{P}(X)\}$, we get sound and complete relational semantics for team logics.

*Proof.* A simple p-morphism argument.

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s \vDash p & \text{iff} & \forall v \in s : v(p) = 1, \\
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s \vDash \varphi \vee \psi & \text{iff} & s \vDash \varphi \text{ or } s \vDash \psi, \\
s \vDash \neg\varphi & \text{iff} & s \nvDash \varphi, \\
s \vDash \varphi \circ \psi & \text{iff} & \text{there exist } s', s'' \in \mathcal{P}(X) \text{ such that } s' \vDash \varphi; \\
 & & s'' \vDash \psi; \text{ and } s = s' \cup s''.
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This induces a powerset frame $\mathbb{F} = (\mathcal{P}(X), \cup)$, where '$\circ$' is a binary modality referring to the ternary $\cup$-relation: $s = s' \cup s''$; and a model $\mathbb{M} = (\mathcal{P}(X), \cup, V)$ with a 'principal valuation', i.e.,

$$V(p) := \{s \in \mathcal{P}(X) \mid \forall v \in s : v(p) = 1\} = {\downarrow}\{v \in X \mid v(p) = 1\}.$$

**In fact**, if we take all powerset frames $(\mathcal{P}(X), \cup)$, redefine the base clause

$$(\mathcal{P}(X), \cup, V), s \Vdash p \qquad \text{iff} \qquad s \in V(p),$$

and only allow principal valuations $V : \mathbf{Prop} \to \{{\downarrow}s \mid s \in \mathcal{P}(X)\}$, we get sound and complete relational semantics for team logics.

*Proof.* A simple p-morphism argument.

# Team semantics as relational semantics

Recall our semantic clauses: For $X := \{v \mid v : \mathbf{Prop} \to \{0,1\}\}$ and $s \in \mathcal{P}(X)$, we had

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Summarizing, (i) team logics are decidable, and (ii) relational semantics for team logics are given by powerset frames $(\mathcal{P}(X), \cup)$ with principal valuations $V : \mathbf{Prop} \to \{\downarrow s \mid s \in \mathcal{P}(X)\}$.

Question: *Sticking with the signature $\{\wedge, \vee, \neg, \circ\}$, what happens if we allow for arbitrary valuations $V : \mathbf{Prop} \to \mathcal{P}\mathcal{P}(X)$? Does the logic remain decidable?*

In fact, this question is intimately related with an open problem: Goranko and Vakarelov (1999) consider the logic of Boolean frames – instead of a powerset $\mathcal{P}(X)$, the carrier is a Boolean algebra $B$ – and raises the problem of its decidability.[1]

**Theorem**

The logic of powerset frames, in the signature $\{\wedge, \vee, \neg, \circ\}$, with arbitrary valuations is **undecidable**. And so is the hyperboolean modal logic of Goranko and Vakarelov (1999).

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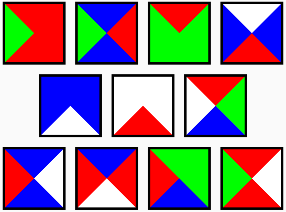
# Proof method: tiling

- A (Wang) tile is a square with colors on each side.
- The tiling problem: given any finite set of tiles $\mathcal{W}$, determine whether each point in the quadrant $\mathbb{N}^2$ can be assigned a tile from $\mathcal{W}$ such that neighboring tiles share matching colors on connecting sides.
- The tiling problem was introduced by Wang (1963) and proven undecidable by Berger (1966).



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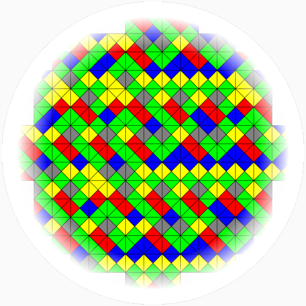

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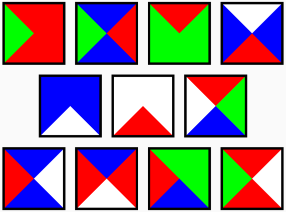

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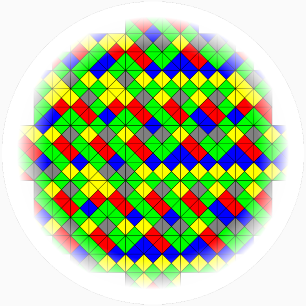

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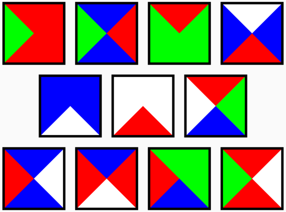

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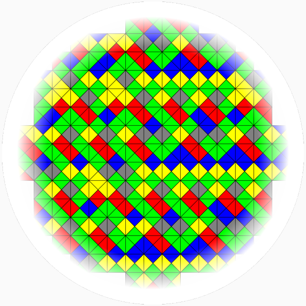

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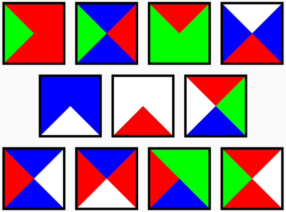

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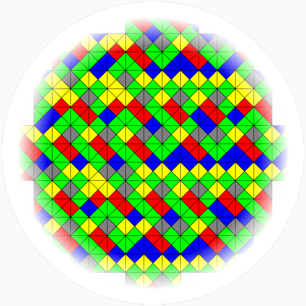

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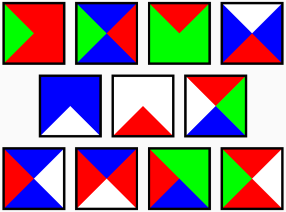

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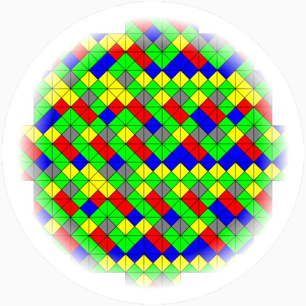

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**Proof idea.**

For each finite set of tiles $\mathcal{W}$, we construct a formula $\phi_{\mathcal{W}}$ such that $\mathcal{W}$ tiles the quadrant if and only if $\phi_{\mathcal{W}}$ is satisfiable. $\square$

Dividing the proof into two lemmas, corresponding to a direction each, we can prove both results in one go:

**Lemma**

If $\phi_{\mathcal{W}}$ is satisfiable (in a Boolean frame), then $\mathcal{W}$ tiles $\mathbb{N}^2$.

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# Insight 1: valuations matter

# Semilattice frames, associativity and negation

Question: *Since we can weaken from powersets to Boolean algebras and stay undecidable, how much further can we go while remaining undecidable?*

Weakening from powersets $(\mathcal{P}(X), \cup)$ to general (join-)semilattices $(S, \sqcup)$, we get a problem posed by Bergman (2018) and Jipsen et al. (2021) (and SBK (2023a)).

**Theorem**

For any class of semilattices containing $(\mathcal{P}(\mathbb{N}), \cup)$, its logic in the signature $\{\wedge, \vee, \neg, \circ\}$, is undecidable.

*Proof.* *Formulas in handout (manuscript with proof available on request)*

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**Answer 2:** As semilattices are associative, commutative, idempotent functions, we could also consider the logic of all associative ternary relations. This is undecidable (Kurucz et al. 1995).

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# Insight 2: associativity matters

Insight 3: negation matters

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$\mathbf{S}$ is the logic of semilattice frames $(S, \sqcup, \mathbf{0})$ with a bottom element $\mathbf{0}$, with arbitrary valuations, in the signature $\{\wedge, \vee, \rightarrow\}$. '$\rightarrow$' is closely connected to '$\circ$' (it is its residual).

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**Problem of concern:** *Is relevant logic $\mathbf{S}$ decidable?*

$\mathbf{S}$ is the logic of semilattice frames $(S, \sqcup, \mathbf{0})$ with a bottom element $\mathbf{0}$, with arbitrary valuations, in the signature $\{\wedge, \vee, \rightarrow\}$. '$\rightarrow$' is closely connected to '$\circ$' (it is its residual).

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[2] See SBK (2024)

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We cover the no-FMP proof instead, since it is considerably simpler than the undecidability proof, yet effectively illustrates some of the same key ideas.[3]

**Theorem: $\mathbb{S}$ lacks the FMP.**

*Proof.* We show that the formula $\psi_\infty$ from the handout only is refuted by infinite models.

[3] Additionally, it addresses an open problem (as recently raised in Weiss 2021)

**Refuting model**

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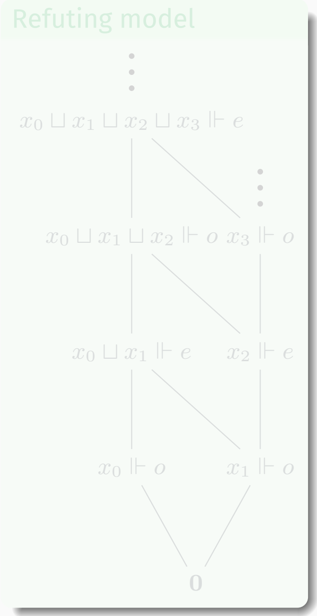

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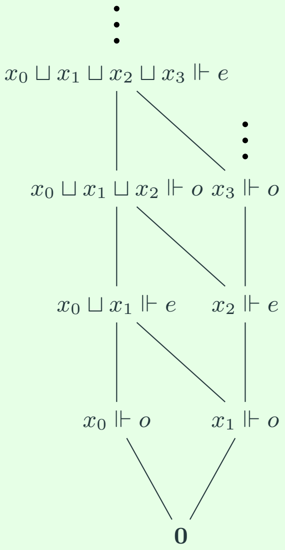

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**Question:** *While $\mathbf{S}$ escaped those techniques, can we extend the present proof to include $\mathbf{R}$ (and $\mathbf{E}$ and $\mathbf{T}$) as well?*

**Answer:** Yes! It extends to $\mathbf{R}^+$, hence $\mathbf{R}$, as well ($\mathbf{E}$ and $\mathbf{T}$ to be checked).

**Two comments on the proof:**

1. We use the $n \times n$, for all $n \in \mathbb{N}$, tiling problem instead.
2. On associativity and tiling (modulo commutativity):
    - Associativity for ternary relations: $R(ab)cd \Rightarrow Ra(bc)d$.[4]
    - Write $aR_b\,c$ for $Rabc$. Then $R(ab)cd$ means $\exists e{:}aR_b\,eR_c\,d$; and $R(ac)bd$ means $\exists f{:}aR_c\,fR_b\,d$.
    - So sp. $(m,n)R_x(m+1,n)R_y(m+1,n+1)$. From associativity, we get that there is a point $(m,n+1)$ s.t. $(m,n)R_y(m,n+1)R_x(m+1,n+1)$. I.e.:

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[4] Mod comm., it is $R(ab)cd \Rightarrow R(ac)bd$. $R(ab)cd$ means $\exists e{:}Rabe \wedge Recd$; and $R(ac)bd$ means $\exists f{:}Racf \wedge Rfbd$.

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Question: Is Bunched Implication Logic (BI) decidable?

Theorem: BI is undecidable

## Conclusion

We obtained new (undecidability) results, including:

- Hyperboolean modal logic is undecidable.[5]
- Modal logic of semilattices is undecidable.[6]
- $\mathbf{S}$ is undecidable [cf. SBK 2024].[7]
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- BI is undecidable.[8]

We compared them with known decidability results:

- Propositional team logics are decidable.
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Core messages:

- Valuations matter.
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We compared them with known decidability results:

- Propositional team logics are decidable.
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Core messages:

- Valuations matter.
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# Conclusion

**We obtained new (undecidability) results, including:**

- Hyperboolean modal logic is undecidable.[5]
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Thank you!

How can we think of this algebraically?

# From relations to algebras

Given any set $A$ with a ternary relation $R$, we can form the complex algebra:

$$(\mathcal{P}A, \cap, \cup, ^c, \circ),$$

where

$$B \circ C := \{a \in A \mid Rabc, b \in B, c \in C\}.$$

The result is a Boolean algebra with an operator $\circ$.

In our case, $A$ is the powerset $\mathcal{P}X$ and $R$ the union relation $\cup$, so we get

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$$B \circ C := \{a \in A \mid a = b \cup c, b \in B, c \in C\}.$$

Finally,

- Let *Pow*$^+$ denote the class of complex algebras of powersets with union.
- Team logic is the theory of *Pow*$^+$ where homomorphisms send variables to principal downsets.
- If *arbitrary* homomorphisms, then we get $\mathbf{V}(Pow^+)$.

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