

Math 240: Midterm 2 Q6bc

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Given positive integers a and b , we want to compute some integers s and t such that

$$\gcd(a, b) = sa + tb$$

Consider the following iterative program LIN_COMB (a, b) which is supposed to accomplish this:

Initialize variables $c = a, d = b, s_0 = 1, t_0 = 0, s_1 = 0, t_1 = 1$

While $c \neq d$, do the following:

 If $c < d$, then decrement d by c , decrement s_1 by s_0 , decrement t_1 by t_0

 Else if $c > d$, then decrement c by d , decrement s_0 by s_1 , decrement t_0 by t_1

Return s_0 and t_0

§1 6b

Assume, in addition to (a), that $\gcd(a, b) = \gcd(c, d)$ is a loop invariant for the while loop in LIN_COMB. Prove that LIN_COMB satisfies partial correctness.

Proof. Suppose that LIN_COMB halts. Fix $n \in \mathbb{N}$ such that exactly n iterations of the while loop are executed. Consider the values of c, d, s_0, t_0 after the n th iteration. Since the $(n+1)$ th iteration is not executed, the loop condition fails, which means $c = d$. This must mean that $\gcd(c, d) = c = d$. Moreover, $\gcd(a, b) = c$ by the given loop invariant. From (a), we know $(c = s_0a + t_0b) \wedge (d = s_1a + t_1b)$ is a loop invariant. Since $\gcd(a, b) = c = s_0a + t_0b$, we get the correct output. \square

§2 6c

Prove that $(c > 0) \wedge (d > 0)$ is a loop invariant for the while loop in LIN_COMB.

Proof. Let $P(n)$ be the predicate asserting that if the while loop has run for n iterations, then $(c > 0) \wedge (d > 0)$. Domain: \mathbb{N} . We prove by induction that for all k , $P(k)$ holds.

Base case: Before any iteration of the loop: $c = a$ and $d = b$. a and b are positive integers by definition. Hence, $P(0)$ holds as desired.

Inductive step: Suppose that $n \in \mathbb{N}$ such that $P(n)$ holds. We prove that $P(n+1)$ holds.

Suppose the loop has run for $n+1$ iterations. Let $c', d', s'_0, t'_0, s'_1, t'_1$ be the value of the variables after the loop has run for n iterations. Now we consider what happens in the $(n+1)^{th}$ iteration.

i. Case 1: $c' < d'$

$$\begin{aligned} c &= c' \\ d &= d' - c' \end{aligned}$$

By our induction hypothesis, we know that $c' > 0$, so $c > 0$ holds. Since $c' < d'$ and $c' > 0$, so $d' - c' > 0$. Because $d = d' - c'$, this means that $d > 0$ holds. Thus, $P(n+1)$ holds for when $c' < d'$.

ii. Case 2: $c' > d'$

$$c = c' - d'$$

$$d = d'$$

By our induction hypothesis, we know that $d' > 0$, so $d > 0$ holds. Since $c' > d'$ and $d' > 0$, so $c' - d' > 0$. Because $c = c' - d'$, this means that $c > 0$ holds. Thus, $P(n+1)$ holds for when $c' < d'$.

This proves the inductive step. By induction, we conclude that $P(n)$ holds for all $n \in \mathbb{N}$ which makes it an loop invariant for LIN_COMB.

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