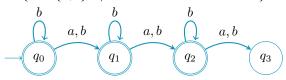
Automata

 $\begin{array}{c} \text{Course Work 1.3} \\ \text{F29LP} \end{array}$

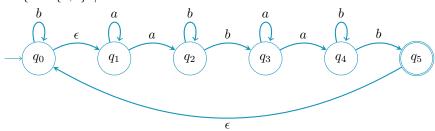
SUBMITTED BY

 $\mathop{\rm Yoav}_{\it H00347035} \mathop{\rm Levi}_{\it H00347035}$

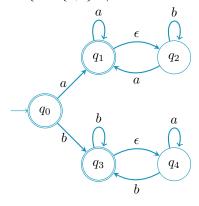
- 1 /(ab)*/
- 2 /(b*a) + a + b[ab]*/
- 3 NFA
 - 1. $L = \{w \in \{a,b\} * | \text{w contains at most two a's} \}$



2. $L = \{w \in \{a,b\}*| \text{w contains an even number of occurrences of ab as a subword}\}$



3. $L = \{w \in \{a,b\} * | \text{the first and the last letter of w are identical} \}$



4
$$/a*(ba\{2,\})*/$$

5

$$S \rightarrow aA$$

$$A \rightarrow aB$$

$$B \rightarrow aS|aC$$

$$C \rightarrow aS|\epsilon$$

6 Unmarked, N/A

7

1.

$$S \to aA|bB$$

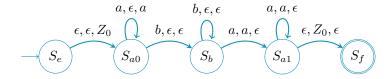
$$A \to aA|bS|aB|\epsilon$$

$$B \to aS$$

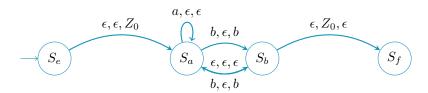
2. Is ambiguous as "aaaa" can be constructed in two ways

$$(I) \begin{tabular}{lll} Rule & Result & Rule & Result \\ $S \to aA$ & a \\ $A \to aA$ & aa \\ $A \to aA$ & aaa \\ $A \to aA$ & aaaa \\ $A \to \epsilon$ & \underline{aaaa} & A \to \epsilon$ & \underline{aaaa} \\ \end{tabular}$$

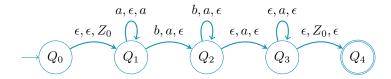
- 8 The CFG is used to create a number of a's with an equivalent number of b's, in any order.
- 9 $L = \{a^m b^n a^m | m, n \ge 0\}$, alphabet $= \{a, b, Z_0\}$, Start stack symbol $= Z_0$



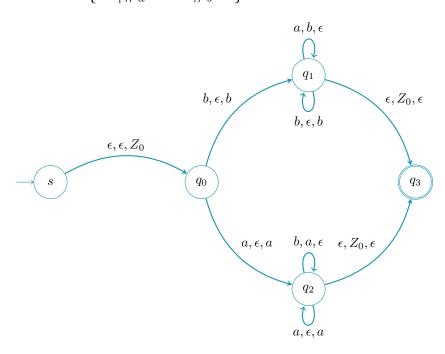
- Not possible as DFA express regular languages, and $a^mb^na^m$ is a only expressible as a context-free language.
- 11 $L = \{a^m b^{2n} | m, n \ge 0\}$



12 L={ $a^mb^n|m>n>0$ }



$L = \{W | \#_a W = \#_b W\}$



14 L={ $W|\#_aW = 2\#_bW$ }

