## Automata

 $\begin{array}{c} \text{Course Work 1.3} \\ \text{F29LP} \end{array}$ 

SUBMITTED BY

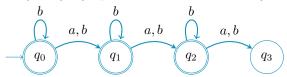
 $\mathop{\rm Yoav}_{\it H00347035} \mathop{\rm Levi}_{\it H00347035}$ 

1 
$$/(ab)*/$$

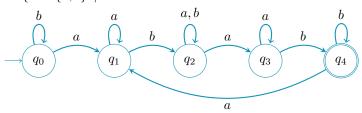
2 
$$/(b*a) + a + b[ab]*/$$

## 3 NFA

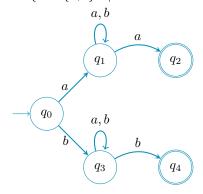
1.  $L = \{w \in \{a,b\} * | \text{w contains at most two a's} \}$ 



2.  $L = \{w \in \{a, b\}* | w \text{ contains an even number of occurrences of ab as a subword}\}$ 



3.  $L = \{w \in \{a,b\} * | \text{the first and the last letter of w are identical} \}$ 



4 
$$/a * (ba\{2,\})*/$$

**5** 

$$S \rightarrow aA$$
 
$$A \rightarrow aB$$
 
$$B \rightarrow aS|aC$$
 
$$C \rightarrow aS|\epsilon$$

6 Unmarked, N/A

7

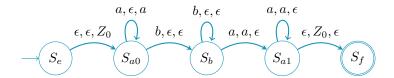
1.

$$S \rightarrow aA|bB$$
 
$$A \rightarrow aA|bS|aB|\epsilon$$
 
$$B \rightarrow aS$$

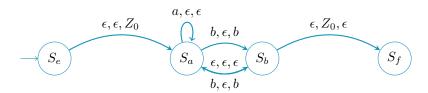
2. Is ambiguous as "aaaa" can be constructed in two ways

$$(I) \begin{tabular}{lll} Rule & Result & Rule & Result \\ $S \rightarrow aA$ & a & $S \rightarrow aA$ & a \\ $A \rightarrow aA$ & aaa & (II) & $A \rightarrow aB$ & aa \\ $A \rightarrow aA$ & aaaa & $B \rightarrow aS$ & aaa \\ $A \rightarrow \epsilon$ & $aaaa$ & $A \rightarrow \epsilon$ & $\underline{aaaa}$ \\ \end{tabular}$$

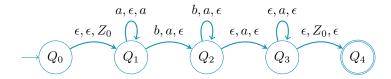
- 8 The CFG is used to create a number of a's with an equivalent number of b's, in any order.
- 9  $L = \{a^m b^n a^m | m, n \geq 0\}$ , alphabet  $= \{a, b, Z_0\}$ , Start stack symbol  $= Z_0$



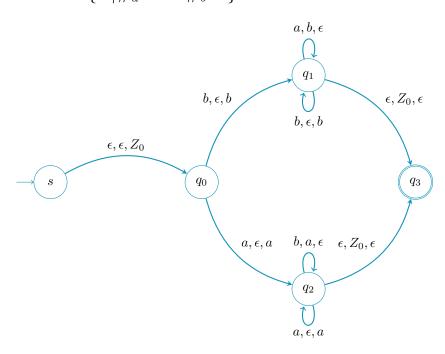
- Not possible as DFA express regular languages, and  $a^mb^na^m$  is a only expressible as a context-free language.
- 11  $L = \{a^m b^{2n} | m, n \ge 0\}$



12 L={ $a^mb^n|m>n>0$ }



## $L = \{w | \#_a W = \#_b W\}$



## 

