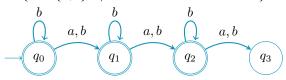
## Automata

 $\begin{array}{c} \text{Course Work 1.3} \\ \text{F29LP} \end{array}$ 

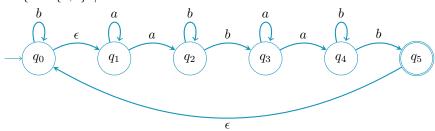
SUBMITTED BY

 $\mathop{\rm Yoav}_{\it H00347035} \mathop{\rm Levi}_{\it H00347035}$ 

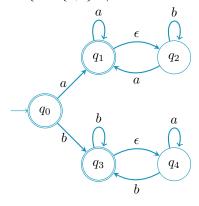
- 1 /(ab)\*/
- 2 /(b\*a) + a + b[ab]\*/
- 3 NFA
  - 1.  $L = \{w \in \{a,b\} * | \text{w contains at most two a's} \}$



2.  $L = \{w \in \{a,b\}*| \text{w contains an even number of occurrences of ab as a subword}\}$ 



3.  $L = \{w \in \{a,b\} * | \text{the first and the last letter of w are identical} \}$ 



4 
$$/a*(ba\{2,\})*/$$

5

$$S \rightarrow aA$$

$$A \rightarrow aB$$

$$B \rightarrow aS|aC$$

$$C \rightarrow aS|\epsilon$$

## 6 Unmarked, N/A

7

1.

$$S \to aA|bB$$

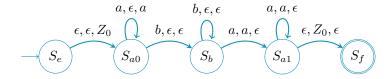
$$A \to aA|bS|aB|\epsilon$$

$$B \to aS$$

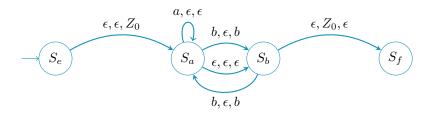
2. Is ambiguous as "aaaa" can be constructed in two ways

$$(I) \begin{tabular}{lll} Rule & Result & Rule & Result \\ $S \to aA$ & a \\ $A \to aA$ & aa \\ $A \to aA$ & aaa \\ $A \to aA$ & aaaa \\ $A \to \epsilon$ & $\underline{aaaa}$ & A \to \epsilon$ & $\underline{aaaa}$ \\ \end{tabular}$$

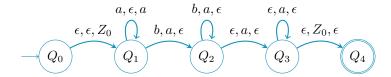
- 8 The CFG is used to create a number of a's with an equivalent number of b's, in any order.
- 9  $L = \{a^m b^n a^m | m, n \ge 0\}$ , alphabet  $= \{a, b, Z_0\}$ , Start stack symbol  $= Z_0$



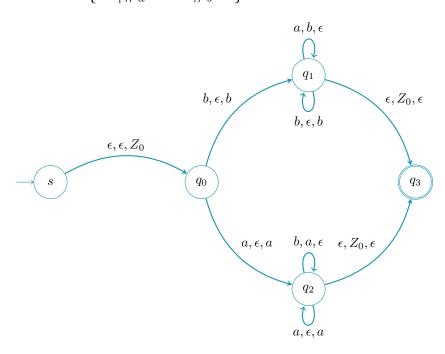
- Not possible as DFA express regular languages, and  $a^mb^na^m$  is a only expressible as a context-free language.
- 11  $\mathbf{L} = \{a^m b^{2n} | m, n \ge 0\}$



12 L={ $a^m b^n | m > n > 0$ }



## $L = \{W | \#_a W = \#_b W\}$



## 14 L={ $W|\#_aW = 2\#_bW$ }

