Automata

 $\begin{array}{c} \text{Course Work 1.3} \\ \text{F29LP} \end{array}$

SUBMITTED BY

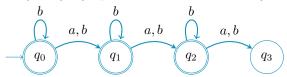
 $\mathop{\rm Yoav}_{\it H00347035} \mathop{\rm Levi}_{\it H00347035}$

1
$$/(ab)*/$$

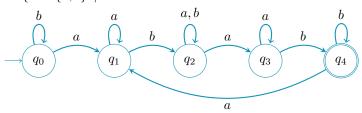
2
$$/(b*a) + a + b[ab]*/$$

3 NFA

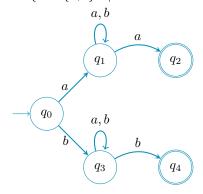
1. $L = \{w \in \{a,b\} * | \text{w contains at most two a's} \}$



2. $L = \{w \in \{a, b\}* | w \text{ contains an even number of occurrences of ab as a subword}\}$



3. $L = \{w \in \{a,b\} * | \text{the first and the last letter of w are identical} \}$



4
$$/a * (ba\{2,\})*/$$

5

$$S \rightarrow aA$$

$$A \rightarrow aB$$

$$B \rightarrow aS|aC$$

$$C \rightarrow aS|\epsilon$$

6 Unmarked, N/A

7

1.

$$S \rightarrow aA|bB$$

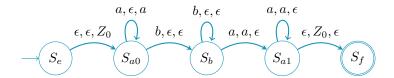
$$A \rightarrow aA|bS|aB|\epsilon$$

$$B \rightarrow aS$$

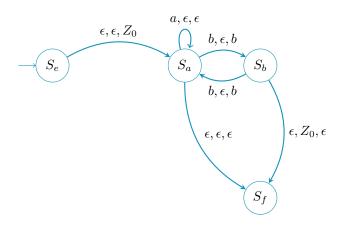
2. Is ambiguous as "aaaa" can be constructed in two ways

$$(I) \begin{tabular}{llll} Rule & Result & Rule & Result \\ $S \to aA$ & a & $S \to aA$ & a \\ $A \to aA$ & aaa & (II) & $A \to aB$ & aa \\ $A \to aA$ & aaaa & $B \to aS$ & aaa \\ $A \to e$ & $aaaa$ & $A \to \epsilon$ & $\underline{a}\underline{a}\underline{a}\underline{a}\underline{a}$ \\ \end{tabular}$$

- 8 The CFG is used to create a number of a's with an equivalent number of b's, in any order.
- 9 $L = \{a^m b^n a^m | m, n \geq 0\}$, alphabet $= \{a, b, Z_0\}$, Start stack symbol $= Z_0$



- Not possible as DFA express regular languages, and $a^mb^na^m$ is a only expressable as a context-free language.
- 11 L= $\{a^mb^{2n}|m,n\geq 0\}$



13
$$L = \{w | \#_a = \#_b\}$$

