

# Ch1 Informed Search

1. State space graph  $G(V, E)$   $V = \text{state}$   
 $E = \text{arcs} = \text{successors}$  with length cost
2. Search tree: ~~expand~~ expand successor as child
3. maintain a partial plan as expanding

```
function TREE-SEARCH(problem, strategy) returns a solution, or failure
  initialize the search tree using the initial state of problem
  loop do
    if there are no candidates for expansion then return failure
    choose a leaf node for expansion according to strategy
    if the node contains a goal state then return the corresponding solution
    else expand the node and add the resulting nodes to the search tree
  end
```

fringe, expansion  
exploration

2. DFS/BFS
3. fringe: LIFO stack / FIFO queue  $\rightarrow$  shallowest solution
4. time complexity:  $O(b^m)$  /  $O(b^d)$  / space complexity:  $O(bm)$  /  $O(b^d)$
5. Uniform Cost Search (UCS)  $\rightarrow$  if solution cost  $C^*$   
最少 cost 的 擴展 fringe  $\rightarrow$  single arc at least cost  $\epsilon \rightarrow$  effective depth  $\rightarrow C^*/\epsilon$   
 $\hookrightarrow$  time complexity  $O(b^{C^*/\epsilon}) = \text{space complexity}$

```
function GRAPH-SEARCH(problem, fringe) return a solution, or failure
  closed  $\leftarrow$  an empty set
  fringe  $\leftarrow$  INSERT(MAKE-NODE(INITIAL-STATE[problem]), fringe)
  loop do
    if fringe is empty then return failure
    node  $\leftarrow$  REMOVE-FRONT(fringe)
    if GOAL-TEST(problem, STATE[node]) then return node
    if STATE[node] is not in closed then
      add STATE[node] to closed
      for child-node in EXPAND(STATE[node], problem) do
        fringe  $\leftarrow$  INSERT(child-node, fringe)
    end
  end
```

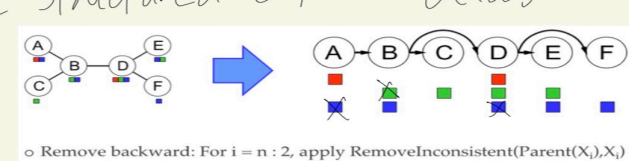
# Ch2 Uninformed Search

1. Search heuristics = a function that estimates how close goal to
2. Greedy Search: "expand the node (you think) nearest to goal"  
might lead to some worst case (最少 hcn)
3.  $A^*$  search: "expand by 最少  $f(n) = g(n) + h(n)$ "  
implementation: should stop when  $\neq$  dequeue the goal
4. what might went wrong: estimated heuristics  $\rightarrow$  評估觀

5. Heuristics (1) admissible (optimistic)  $\rightarrow$  估計  $h(n) \leq \text{actual cost}$   
 $\Rightarrow 0 \leq h(n) \leq h^*(n)$  (true cost)
2. optimality:  $f(n) \leq f(A) < f(B)$   
 $\begin{cases} f(n) \leq f(A) \\ f(A) \leq f(B) \\ n \text{ before } B \end{cases}$

3. characteristics: dominance  $h_a \geq h_b$  if  $\forall n, h_a(n) \geq h_b(n)$   
exact - non-dominant  $\rightarrow$  可以 construct semi-lattice

4. Graph Search: (1) idea: never expand a node twice  
(2)  $\Rightarrow$  tree search + set of expanded states (closed set)  
(3) 避免誤判是 same set (node 不同但 path diff)  $\rightarrow$  針對 heuristics 修正
5. consistency: arc heuristic cost (兩 node 相減)  $\leq$  actual cost  
 $\Rightarrow h(A) - h(C) \leq \text{cost}(A \text{ to } C)$
6. consistency  $\Rightarrow$  f along a path never decrease

7. Tree structured CSP:  $O(nd^3)$  can solve
8. a/g: 

9. nearly-tree structure: 1. 挑 1, assign, 修其他 constraint, 找解
10. cut set: 把 1 擴大至 挑 1 set  $\rightarrow O(d^c (n-c)d^2)$
11. tree decomposition: mega node 來表示

12. Iterative algorithms
13. Local search
14. no fringe, 是由完整的 option 去修改 (3) simulated annealing
15. hill climbing  $\rightarrow$  might stuck local optimum
16. local beam search
17. 同時跑 k 組 local search, 產生 successor, 找最好 k 個

Algorithm: While not solved,  
Variable selection: randomly select any conflicted variable  
Value selection: min-conflicts heuristic:  
Choose a value that violates the fewest constraints  
i.e., hill climb with  $h(x) = \text{total number of violated constraints}$

```
function SIMULATED-ANNEALING(problem, schedule) returns a state
  current  $\leftarrow$  problem.initial-state
  for t = 1 to  $\infty$  do
    T  $\leftarrow$  schedule(t)
    if T = 0 then return current
    next  $\leftarrow$  a randomly selected successor of current
     $\Delta E \leftarrow$  next.value - current.value
    if  $\Delta E > 0$  then current  $\leftarrow$  next
    else current  $\leftarrow$  next only with probability  $e^{\Delta E}$ 
```

# Ch3 Constraint Satisfaction problem

1. Formulation: state  $X_i$  with values from  $D_i$ ,  
goal test by satisfying all constraint
2. Binary CSP graph: nodes = var, arc = constraints  
one variable at one time + DFS
3. Backtracking search: check constraint as you go
4. Filtering
5. forward checking  
當 assign 時同時確認其他相關的 var 值可能值
6. constraint propagation:  $X \rightarrow Y$  is constraint: 對於每個箭頭 + 發端的 every x 都應該被指端有
7. check all arc consistent
8. if X loses a value  $\Rightarrow$  neighbor of X should recheck 對其
9. 可能 no sol.

```
function BACKTRACKING-SEARCH(csp) returns solution/failure
  return RECURSIVE-BACKTRACKING({}, csp)
function RECURSIVE-BACKTRACKING(assignment, csp) returns soln/failure
  if assignment is complete then return assignment
  var  $\leftarrow$  SELECT-UNASSIGNED-VARIABLE(VARIABLES[csp], assignment, csp)
  for each value in ORDER-DOMAIN-VALUES(var, assignment, csp) do
    if value is consistent with assignment given CONSTRAINTS[csp] then
      add {var = value} to assignment
      result  $\leftarrow$  RECURSIVE-BACKTRACKING(assignment, csp)
      if result  $\neq$  failure then return result
      remove {var = value} from assignment
  return failure
```

```
function AC-3(csp) returns the CSP, possibly with reduced domains
  inputs: csp, a binary CSP with variables  $\{X_1, X_2, \dots, X_n\}$ 
  local variables: queue, queue of arcs, initially all the arcs in csp
  while queue is not empty do
     $(X_i, X_j) \leftarrow$  REMOVE-FIRST(queue)
    if REMOVE-INCONSISTENT-VALUES( $X_i, X_j$ ) then
      for each  $X_k$  in NEIGHBORS( $X_j$ ) do
        add  $(X_i, X_k)$  to queue
  function REMOVE-INCONSISTENT-VALUES( $X_i, X_j$ ) returns true iff succeeds
  removed  $\leftarrow$  false
  for each x in DOMAIN( $X_i$ ) do
    if no value y in DOMAIN( $X_j$ ) allows  $(x, y)$  to satisfy the constraint  $X_i \leftrightarrow X_j$ 
      then delete x from DOMAIN( $X_i$ ); removed  $\leftarrow$  true
  return removed
```

10. k-consistency:  $\forall$  node 都有可 assign
11. 從 1-consistency:  $\forall$  node 都有可 assign
12. k-consistency:  $\forall$  挑 k node  $\rightarrow$  k<sup>th</sup> node 有 consistent assign
13. Ordering: (1) minimum remaining values: 選最少可能性的 node  
(2) least constraining value: 最不被限制 (fail fast)
14. Structure: 挑解  $O(d^n) \rightarrow O(Vd^d)$

# Chapter 4 adversarial search

1. Formulation: States S (start sol), players  $P_i$ , actions A, ...
2. 零和 game  $\rightarrow$  maximize utility

3. Minimax computation
4. characteristics  
 $\rightarrow$  很像 DFS  
 $\rightarrow$  time  $O(b^m)$ , space  $O(bm)$

5. alpha beta pruning  $O(b^{m/2})$
6. General configuration (MIN version)
  - o We're computing the MIN-VALUE at some node n
  - o We're looping over n's children
  - o n's estimate of the children's min is dropping
  - o Who cares about n's value? MAX
  - o Let  $\alpha$  be the best value that MAX can get at any choice point along the current path from the root
  - o If n becomes worse than  $\alpha$ , MAX will avoid it, so we can stop considering n's other children (it's already bad enough that it won't be played)

7. if depth limited, 用 evaluation function 找
8. 有可能 state, 如果  $f(n)$  不好  $\rightarrow$  可能 真實值
9. Expectimax: 對 outcome by chance

10. pruning: 因為取 EC 了 無法 因一部分就捨棄

```
def value(state):
  if the state is a terminal state: return the state's utility
  if the next agent is MAX: return max-value(state)
  if the next agent is EXP: return exp-value(state)
def max-value(state):
  initialize v = - $\infty$ 
  for each successor of state:
    v = max(v, value(successor))
  return v
def exp-value(state):
  initialize v = 0
  for each successor of state:
    p = probability(successor)
    v += p * value(successor)
  return v
```

```
def max-value(state,  $\alpha, \beta$ ):
  initialize v = - $\infty$ 
  for each successor of state:
    v = max(v, value(successor,  $\alpha, \beta$ ))
    if v  $\geq \beta$  return v
     $\alpha = \max(\alpha, v)$ 
  return v
def min-value(state,  $\alpha, \beta$ ):
  initialize v = + $\infty$ 
  for each successor of state:
    v = min(v, value(successor,  $\alpha, \beta$ ))
    if v  $\leq \alpha$  return v
     $\beta = \min(\beta, v)$ 
  return v
```

