

function for	G is not unique	etwork $G = (V, E, s, t, c)$ with $ V  \le 4$ and $c : V \times V \to 0, 1$ such that max–flow is not unique. On the flow network you design, show at least two different ions and state the value of the max–flow.			







3.	capa	pose that we are given a flow network $G$ and max flow $f$ on $G$ . We know that all the acities are integers in $G$ , and all flow values are also integers. Suppose that the capacity of dge $(u,v)$ is increased by 1.					
(8	(a)	) How can you decide if the increase in the capacity affects the max flow value? Find an efficient algorithm to decide this.					
	(b)	If f is not max-flow any more, what is the new max-flow value?					
	(c)	How do you find the new max-flow function in an efficient way?					





## 4. (a) Graph and Johnson's Algorithm Problem

Consider the following directed graph G = (V, E), where:

- $V = \{1, 2, 3, 4, 5\}$
- $E = \{(1,2), (2,5), (3,2), (3,4), (4,1), (4,5), (5,1), (5,3)\}$

Suppose that we are given the following weight function  $w: E \to \mathbb{R}$  for the edges:

$$w(1,2) = -5$$

$$w(2,5) = 8$$

$$w(3,2) = 1$$

$$w(3,4) = -4$$

$$w(4,1) = 7$$

$$w(4,5) = 6$$

$$w(5,1) = 0$$

$$w(5,3) = 2$$

For all pairs shortest path problem, if we would like to apply Dijkstra's algorithm, we have to reweight the edges, as suggested by Johnson's algorithm.

Johnson's algorithm is based on first finding some suitable weights for the vertices of the graph. Let us use  $x_1, x_2, x_3, x_4, x_5$  as the weight of the nodes 1, 2, 3, 4, 5 respectively. After we find a suitable weight for each node, we will define a new weight function  $\tilde{w}: E \to \mathbb{R}$  for the edges, such that all edges will have non-negative weights when we use

$$\tilde{w}(i,j) = w(i,j) - x_i + x_j \quad \forall (i,j) \in E$$

For each edge (i, j) in E, write down the constraint on the difference of  $x_i$  and  $x_j$  based on the way  $\tilde{w}$  is defined in Equation (1) above.





