

- Möbius Function

$$\mu(n) = \begin{cases} 0 & \exists d \in \mathbb{N}, d > 1 \text{ s.t. } d^2 | n \\ (-1)^k, & n = p_1 \cdots p_k \end{cases}$$

- Dirichlet Convolution

$$(f * g)(n) = \sum_{d|n} f(d) g\left(\frac{n}{d}\right)$$

$$\text{ex) } (\mu * 1)(n) = \sum_{d|n} \mu(d) = \varepsilon \begin{cases} 0 & (n \neq 1) \\ 1 & (n = 1) \end{cases}$$

$$\begin{aligned} (f * \varepsilon)(n) &= \sum_{d|n} f(d) \varepsilon\left(\frac{n}{d}\right) = f(n) \\ &= (\varepsilon * f)(n) \end{aligned}$$

• Xudyh's Sieve

$$S_f(n) := \sum_{i=1}^n f(i) \rightarrow \text{빠르게 계산원칙}$$

$$\begin{aligned} S_{f * g}(n) &= \sum_{i=1}^n \sum_{d|i} g(d) f\left(\frac{i}{d}\right) \\ &= \sum_{u=1}^n \sum_{v=1}^{\lfloor n/u \rfloor} f(v) g(u) \quad (i=uv) \end{aligned}$$

$$= g(1) S_f(n) + \sum_{u=2}^n g(u) S_f\left(\left\lfloor \frac{n}{u} \right\rfloor\right)$$

$$\rightarrow S_f(n) = \frac{1}{g(1)} \left\{ \underbrace{S_{f * g}(n)}_{\text{쉽게 계산 가능 가정}} - \sum_{u=2}^n g(u) S_f\left(\left\lfloor \frac{n}{u} \right\rfloor\right) \right\}$$

• Time Complexity 계산

$\left\lfloor \frac{n}{u} \right\rfloor$  은 최대  $2\sqrt{n}$  개 ( $\sqrt{n}$  이하 큰 것 :  $\sqrt{n}$  개,  $1 \sim \sqrt{n}$ )

$$\rightarrow T(n) = O(\sqrt{n}) + \sum_{i=1}^{\sqrt{n}} \sqrt{i} + \sum_{i=1}^{\sqrt{n}} \sqrt{\left\lfloor \frac{n}{i} \right\rfloor}$$

$$\approx O(\sqrt{n}) + \int_0^{\sqrt{n}} \sqrt{x} dx + \int_0^{\sqrt{n}} \sqrt{\frac{n}{x}} dx$$

$$= O(\sqrt{n}) + \frac{2}{3} n^{\frac{3}{4}} + 2 n^{\frac{3}{4}} = O(n^{\frac{3}{4}})$$

• Optimization

D(2) k 개를 전처리

$$\rightarrow T(n) = O(\sqrt{n}) + O(k) + \int_0^k \sqrt{\frac{n}{x}} dx$$

$$= O(\sqrt{n}) + O(k) + \frac{2n}{\sqrt{k}}$$

$$O(k) = O\left(\frac{n}{\sqrt{k}}\right) \rightarrow k = n^{\frac{2}{3}} \rightarrow n^{\frac{2}{3}} \text{ 개수로 전처리 후 계산}$$

$$\text{ex) } \sum_{i=1}^n \phi(i) = ?$$

$$(\phi * 1)(n) = \sum_{d|n} \phi(d) = n$$

$$S_{\phi * 1}(n) = S_{\phi}(n) + \sum_{u=2}^n S_{\phi}\left(\left\lfloor \frac{n}{u} \right\rfloor\right)$$

$$\rightarrow S_{\phi}(n) = \frac{n(n+1)}{2} - \sum_{u=2}^n S_{\phi}\left(\left\lfloor \frac{n}{u} \right\rfloor\right)$$

$$\therefore \sum_{i=1}^n \phi(i) : O(n^{\frac{2}{3}}) \text{에 계산 가능}$$

다른 방법들도 있지만 되도록 하기

$$S_{\phi}(n) : |\{ (a, b) \mid \gcd(a, b) = 1, 1 \leq a \leq b \leq n \}|$$

$$|\{ (a, b) \mid \gcd(a, b) = g, 1 \leq a \leq b \leq n \}|$$

$$= |\{ (a, b) \mid \gcd(a, b) = 1, 1 \leq a \leq b \leq \left\lfloor \frac{n}{g} \right\rfloor \}|$$

$$= S_{\phi}\left(\left\lfloor \frac{n}{g} \right\rfloor\right)$$

$$\rightarrow \sum_{i=1}^n S_{\phi}\left(\left\lfloor \frac{n}{i} \right\rfloor\right) = \frac{n(n+1)}{2}$$

$$\therefore S_{\phi}(n) = \frac{n(n+1)}{2} - \sum_{i=2}^n S_{\phi}\left(\left\lfloor \frac{n}{i} \right\rfloor\right)$$

역시  $O(n^{\frac{2}{3}})$ 에 계산 가능