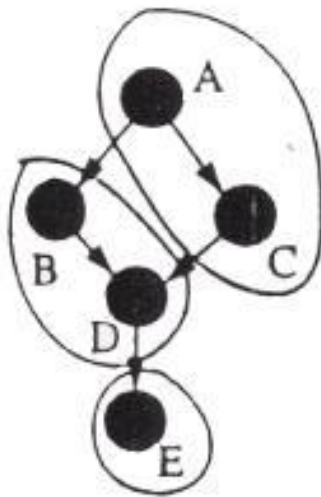


Clustering for Delay Minimization

- Allow gate duplication.
- Gate duplication may help reduce delay.

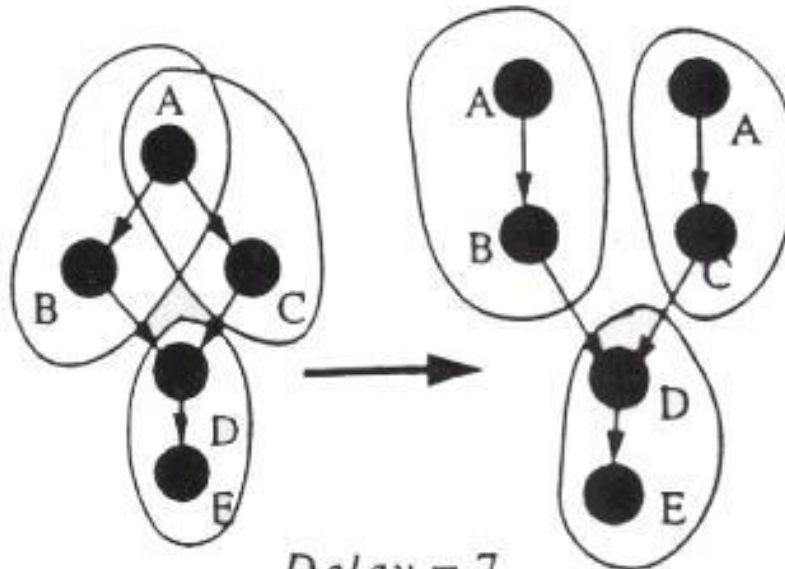
$D=3; M=2; \delta(v)=1, w(v)=1$, for each v .

Without gate duplication



Delay = 10

With gate duplication

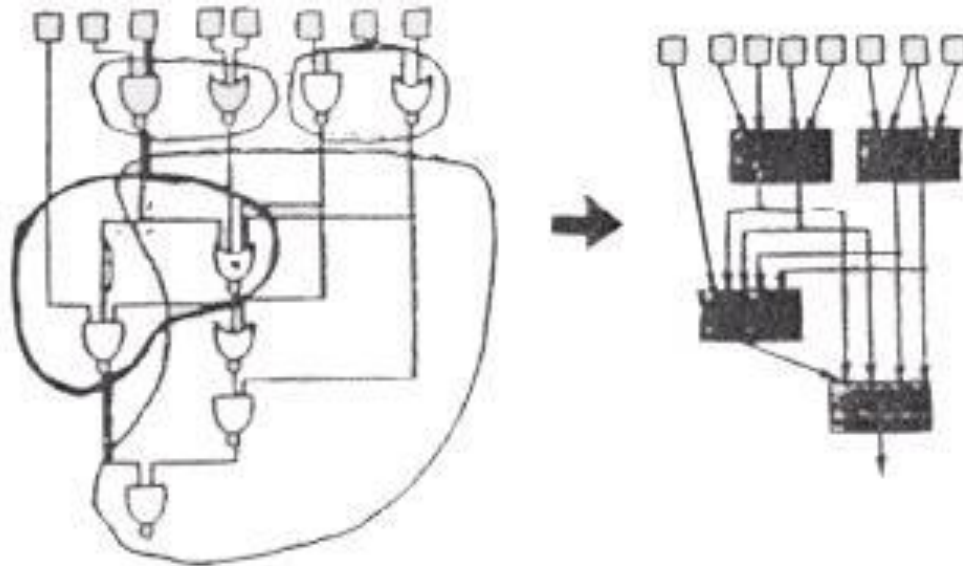


Delay = 7

Unit Delay Model

- No gate delay.
- No interconnection delay within a cluster.
- Delay of 1 unit for an interconnection between 2 clusters.
- An optimal algorithm for area constraint only (Lawler, Levitt and Turner, IEEE TC, 1966).
- An optimal algorithm for pin constraint only (Cong and Ding, ICCAD, 1992).

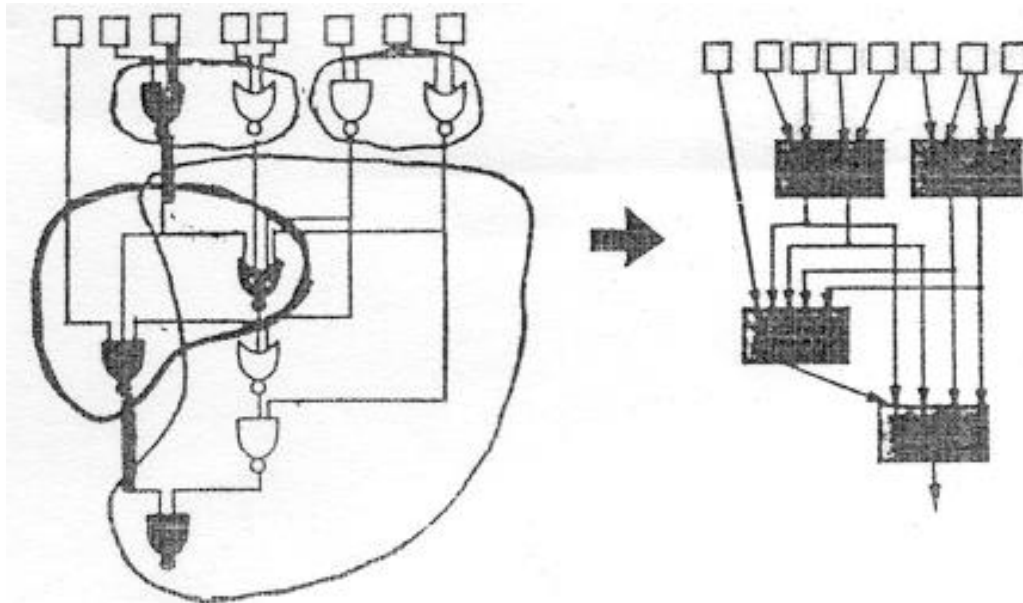
Circuit delay = 3.



General Delay Model

- Each gate has a delay.
- No interconnection delay within a cluster.
- Delay of D units for an interconnection between 2 clusters.
- A heuristic algorithm for area constraint only (Murgai, Brayton and Sangivanni-Vincentelli, ICCAD, 1991).

$D = 2$, $\delta(v) = 1$, circuit delay = $6 + 4 = 10$.



General Delay Model (Cont'd)

- Rajaraman and Wong, “Optimal clustering for delay minimization,” DAC, 1993.
 - Optimal algorithm: $O(n^2 \log n + nm)$, where n is # of gates, m is # of interconnections.
 - Definitions:
 - M : the area constraint on a cluster.
 - $W(C)$: the total area of the gates in cluster C .
 - N : a given combinational circuit.
 - N_v : v and all its *predecessors* in N .
 - $\delta(v)$: the delay of v .
 - $\Delta(u, v)$: maximum delay along any path from the output of u to the output of v , ignoring delays on interconnections.
 - $w(v)$: the area of v .
 - $l(v)$: the delay at v in an optimal clustering of N_v .
- For each *primary input* v , $l(v) = \delta(v)$.
- $l'(u) = l(u) + \Delta(u, v)$, for each u in $N_v - \{v\}$.

General Delay Model (Cont'd)

- Algorithm: labeling phase + clustering phase.
- **Labeling phase:** compute $l(v)$ for each v in a topological order.
 - P : the set of nodes in $N_v - \{v\}$ sorted in non-increasing order in the value of l' .

```

Algorithm Labeling( $v$ );
begin
  done  $\leftarrow$  false;
  cluster( $v$ )  $\leftarrow$  { $v$ };
  while (not done)
    Remove the first node  $u$  in  $P$ ;
    if ( $W(\text{cluster}(v)) + w(u) \leq M$ )
      cluster( $v$ )  $\leftarrow$  cluster( $v$ )  $\cup$  { $u$ };
      if  $P$  is empty
        done  $\leftarrow$  true;
      endif
    else
      done  $\leftarrow$  true;
    endif
  endwhile
   $l_1(v) \leftarrow \max\{l'(x) \mid x \in \text{cluster}(v) \cap PI\}$ ;
   $l_2(v) \leftarrow \cancel{l'(u) + D}, \max\{l'(u) + D \mid u \in N_v - \text{cluster}(v)\}$ ;
   $l(v) \leftarrow \max\{l_1(v), l_2(v)\}$ ;
end
    
```

General Delay Model (Cont'd)

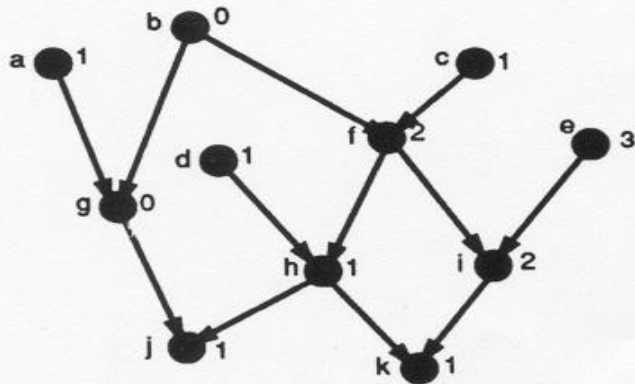
- **Clustering phase:** generate the clusters based on the information obtained in the labeling phase.
- Overall algorithm:

```

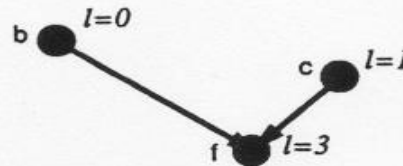
begin
  Compute the maximum delay matrix  $\Delta$ .  $\Delta(i, j)$  is the
  maximum delay along any path from the output of  $i$ 
  to the output of  $j$ ;
  for each PI  $i$ , do  $l(i) \leftarrow \delta(i)$ ;
  Sort the non-PI nodes of  $N$  in topological order
  to obtain list  $T$ ;
  while  $T$  is non-empty
    Remove the first node  $v$  from  $T$ ;
    Compute  $N_v$ ;
    for each node  $u \in N_v \setminus \{v\}$  do
       $l'(u) \leftarrow l(u) + \Delta(u, v)$ ;
    Sort the nodes in  $N_v \setminus \{v\}$  in order of
    decreasing value of  $l'$  to form list  $P$ ;
    Call Labeling( $v$ );
  endwhile
   $L \leftarrow \mathcal{PO}$ ;
   $S \leftarrow \phi$ ;
  while  $L$  is not empty
    Remove a node  $v$  from  $L$ ;  $N$ -cluster( $v$ )
     $S \leftarrow S \cup \{\text{cluster}(v)\}$ ;
    for all nodes  $x$  in  $N$ , such that  $x$  is adjacent
    to  $v$ , for some  $y \in \text{cluster}(v)$ ,  $L \leftarrow L \cup \{x\}$ ;
  endwhile
end
  
```

General Delay Model (Cont'd)

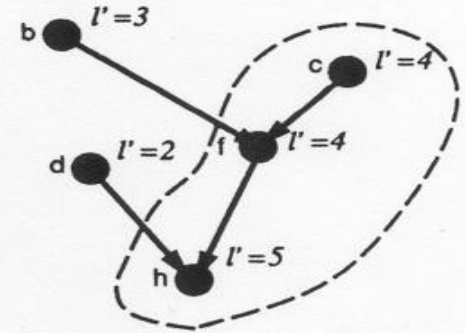
- An example: $M=3, D=3$



(a)



(b)



(c)

