

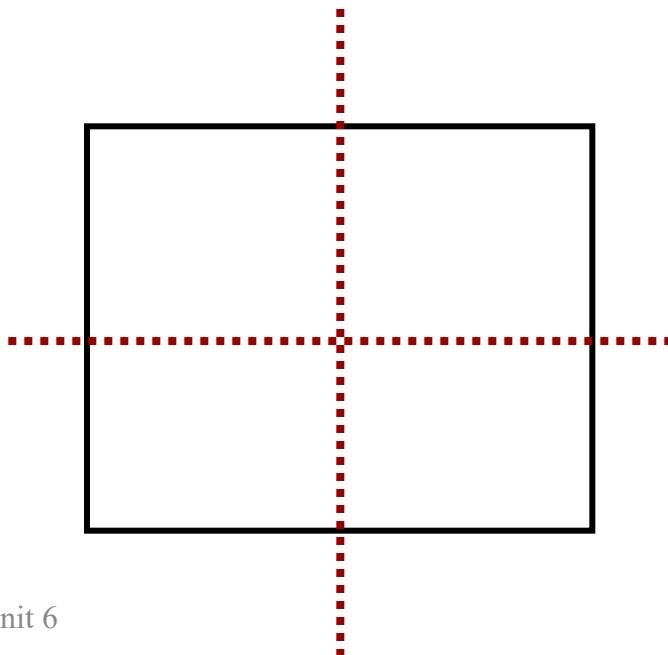
Hierarchical Routing Framework

- The hierarchical approach recursively divides a routing region into a set of subregions and solve those subproblems independently.
- Drawbacks: lack the global information for the interaction among subregions.



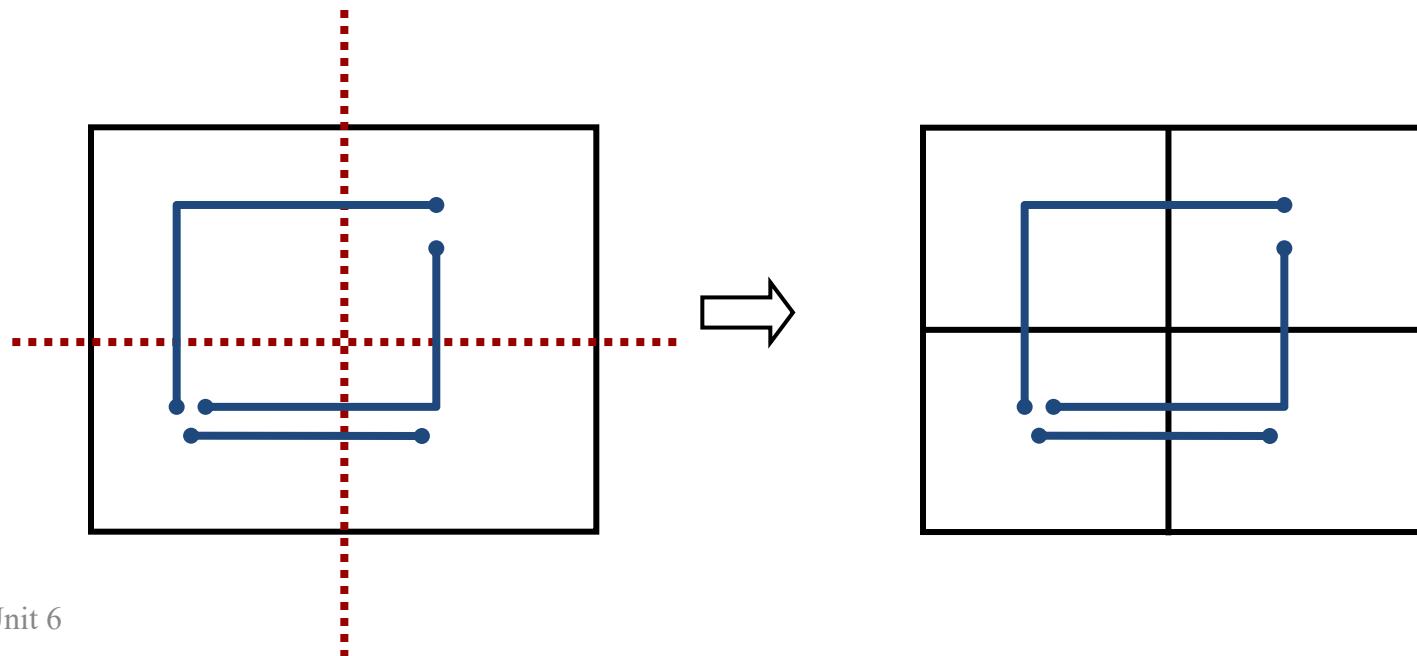
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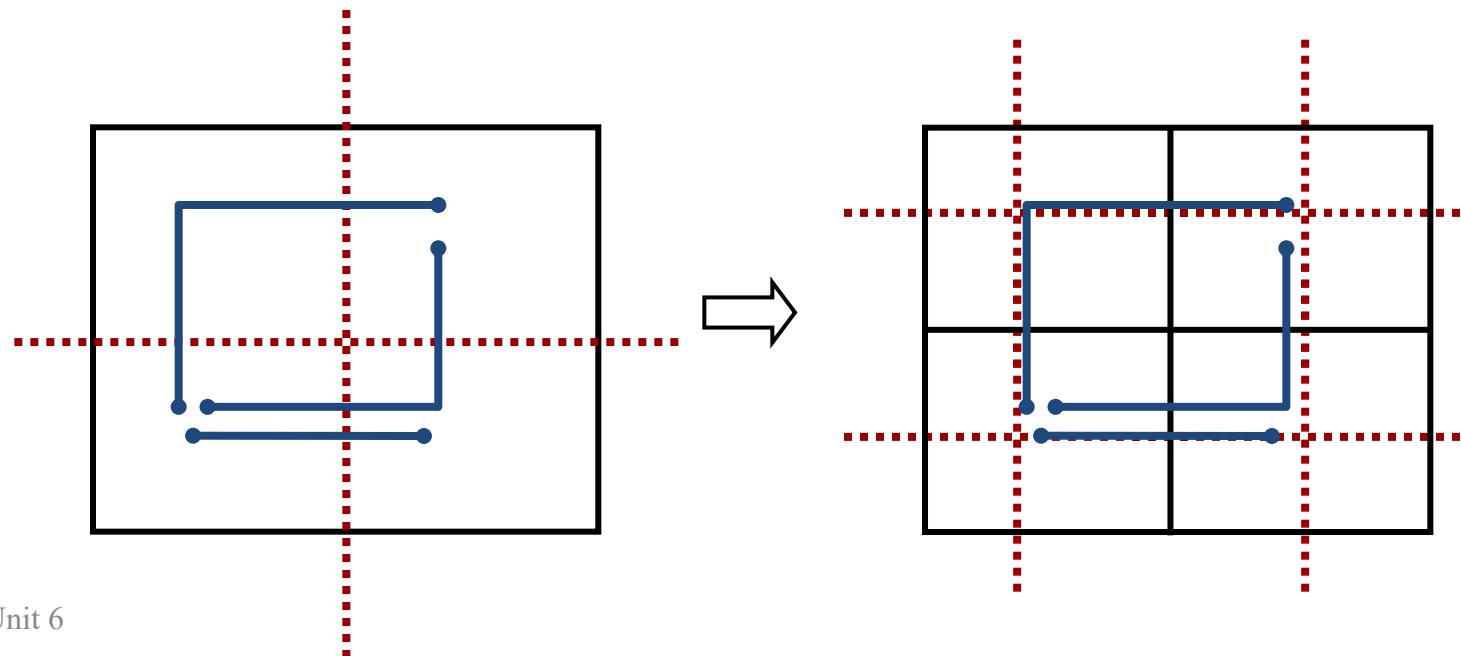
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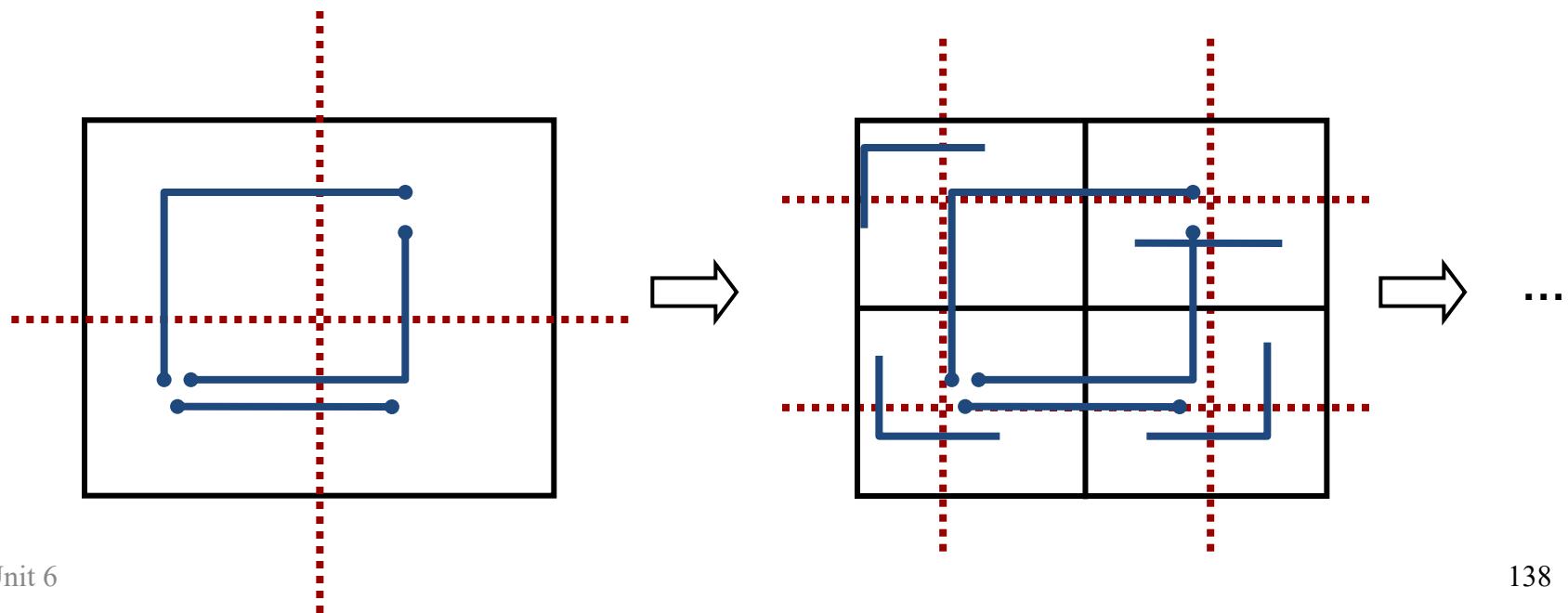
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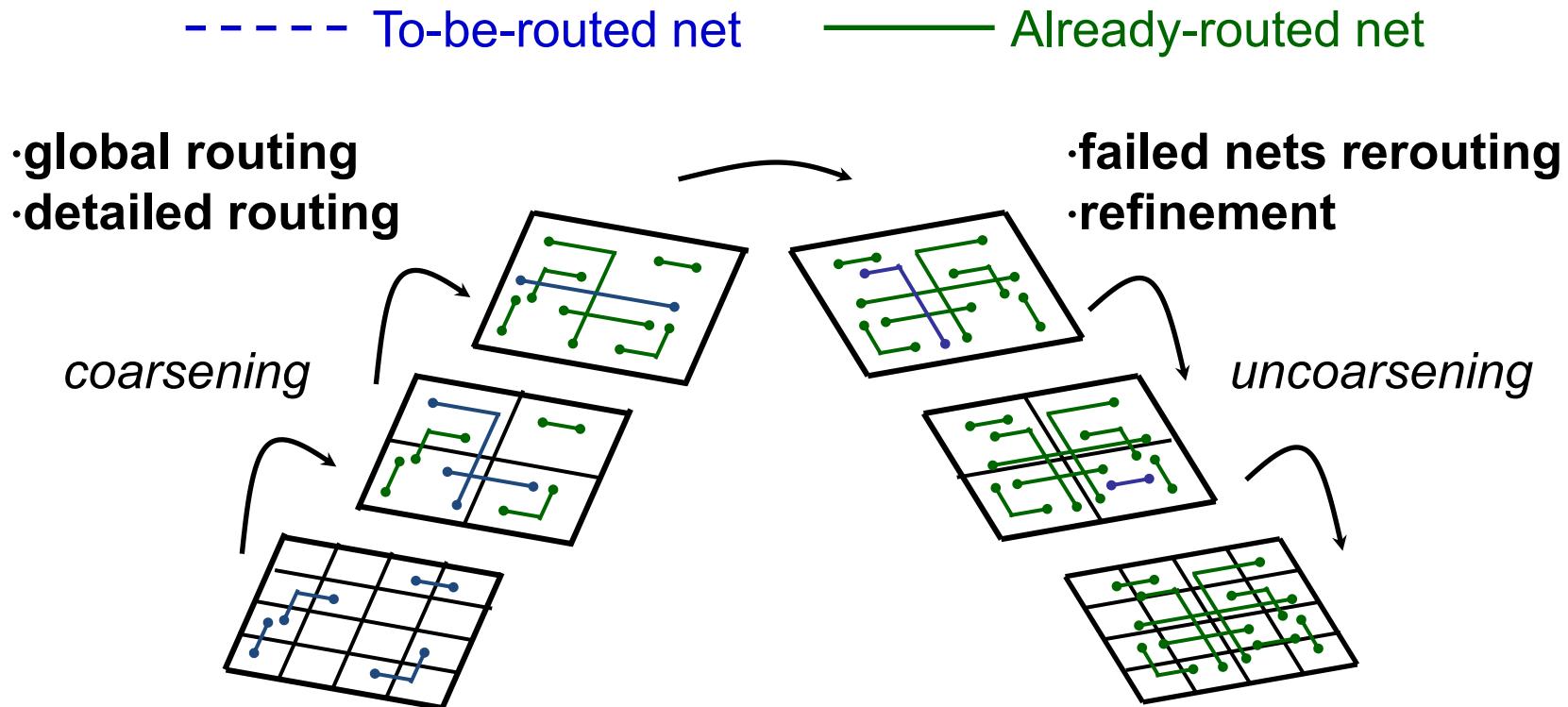
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Multilevel Full-Chip Routing Framework

- Lin and Chang, “A novel framework for multilevel routing considering routability and performance,” ICCAD-02 (TCAD-03).
- Multilevel framework: coarsening followed by uncoarsening.
- Coarsening (bottom-up) stage:
 - Constructs the net topology based on the minimum spanning tree.
 - Processes routing tiles one by one at each level, and only local nets (connections) are routed.
 - Applies two-stage routing of global routing followed by detailed routing.
 - Uses the L-shaped & Z-shaped pattern routing.
 - Performs resource estimation after detailed routing to guide the routing at the next level.
- Uncoarsening (top-down) stage
 - Completes the failed nets (connections) from the coarsening stage.
 - Uses a global and a detailed maze routers to refine the solution.

Λ -Shaped Multilevel Routing Framework

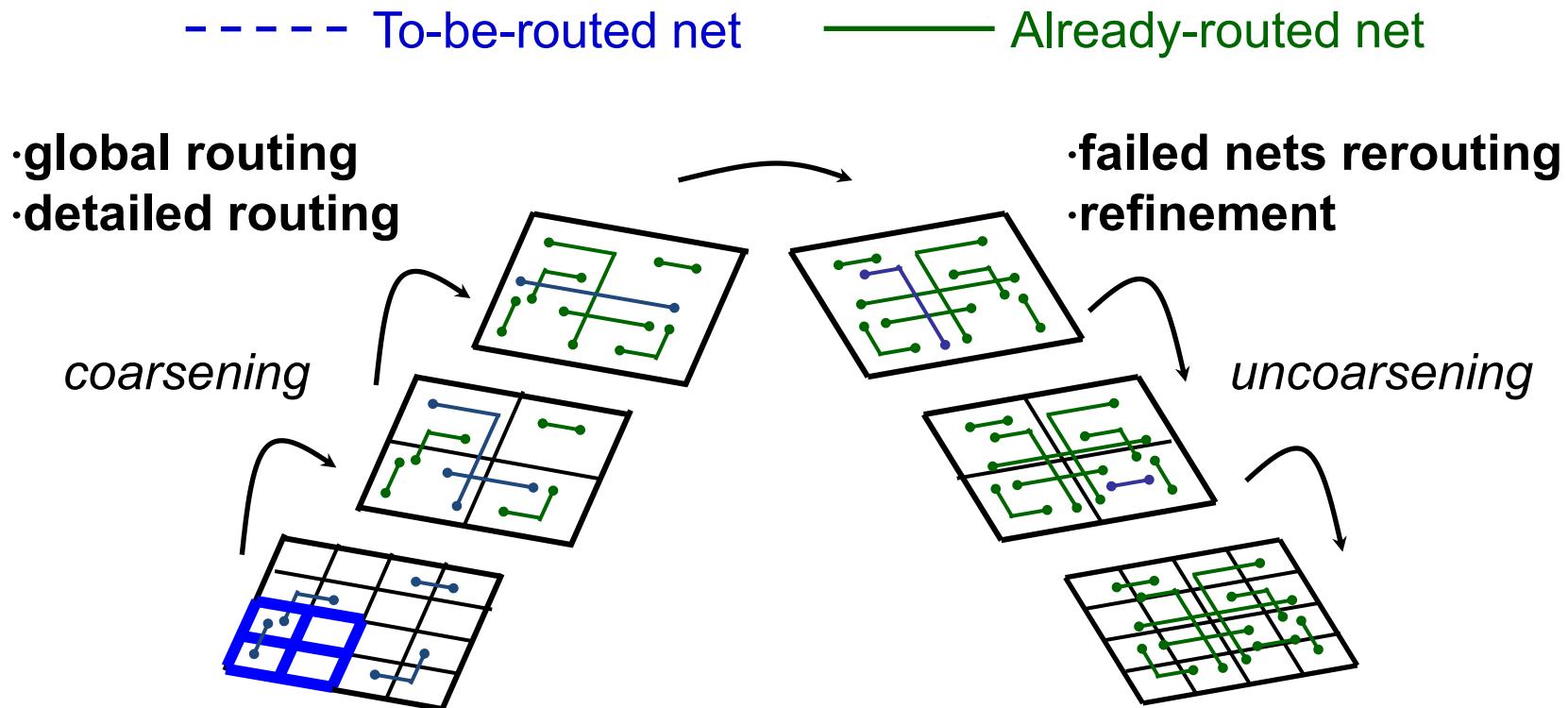


Perform global pattern & detailed routing for local connections and then estimate routing congestion for the next level

Use maze routing to reroute failed connections and refine the solution

Cost: congestion + net length

Λ-Shaped Multilevel Routing Framework

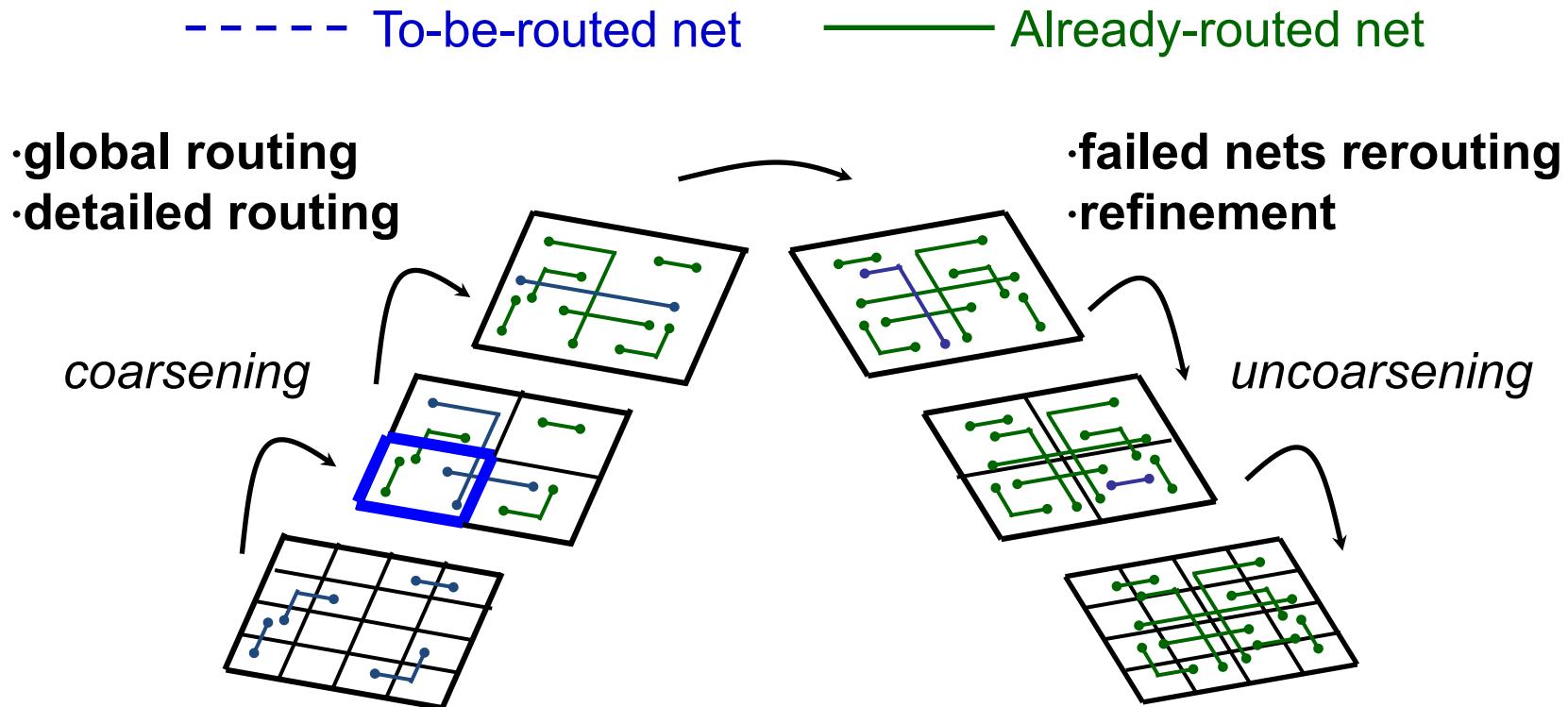


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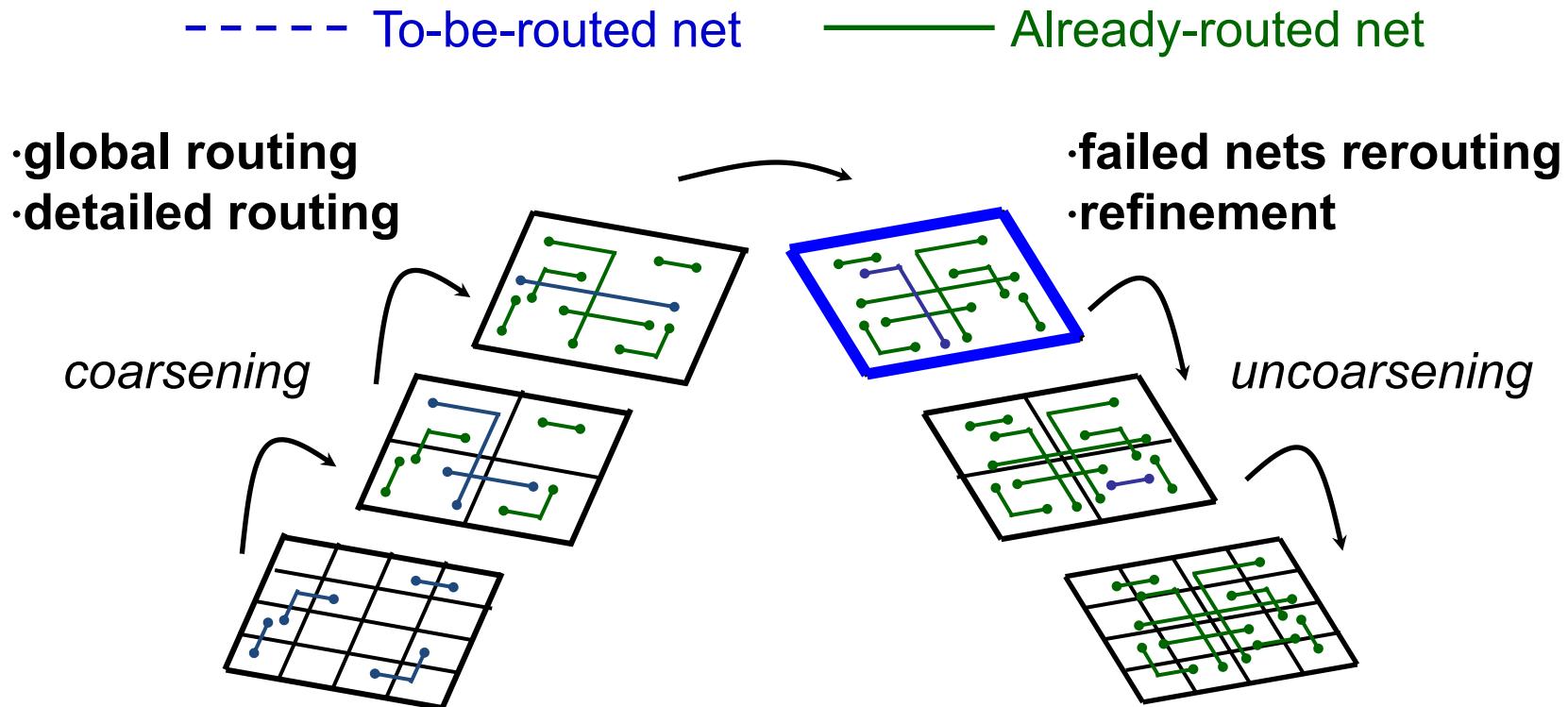


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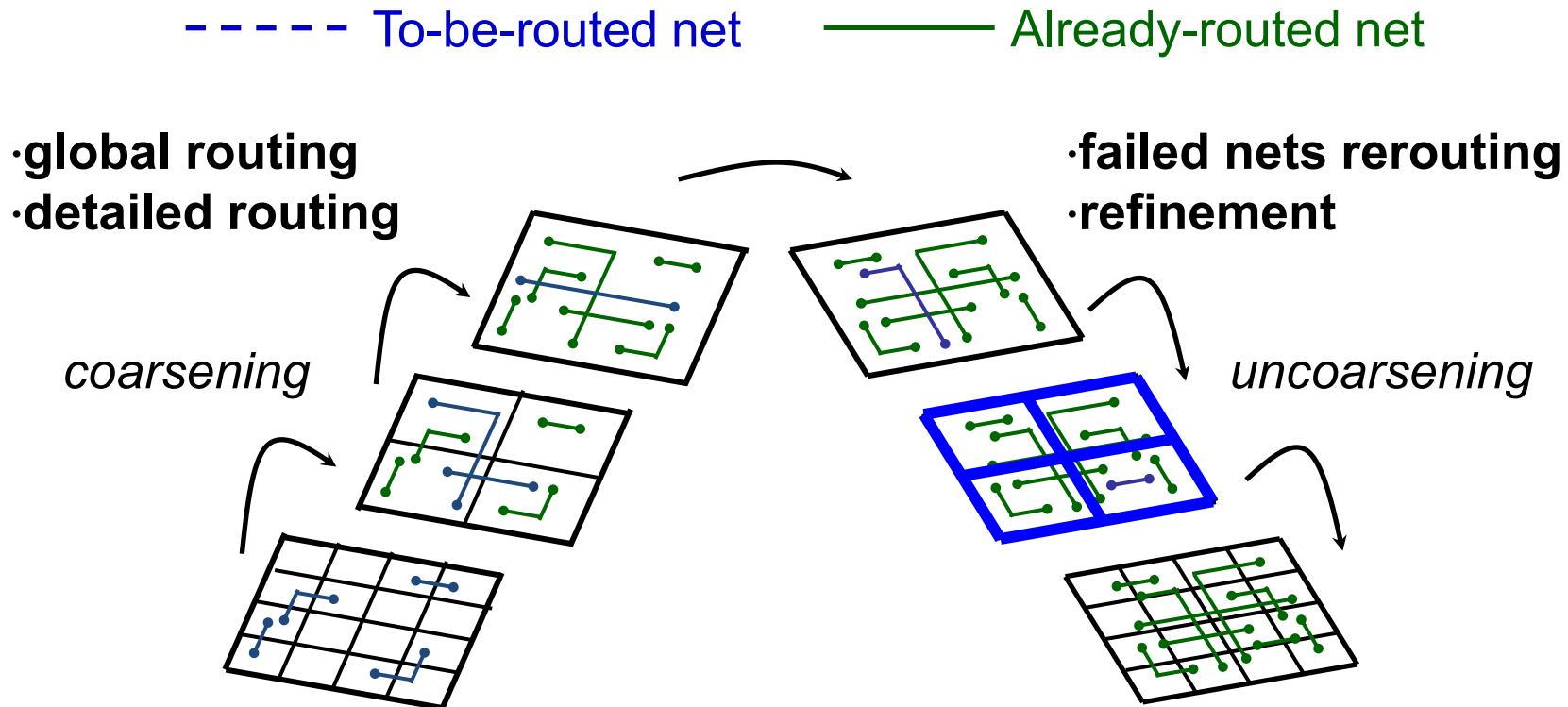


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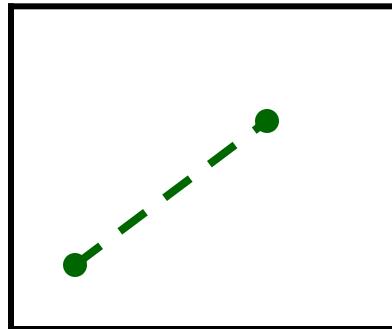
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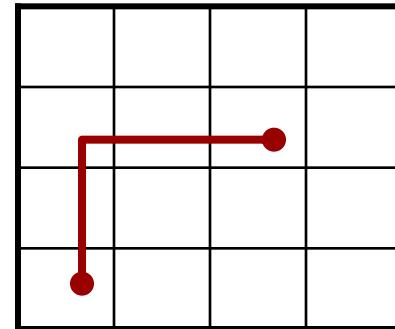
Coarsening Stage

- Build MSTs for all nets and decompose them into two-pin connections.
- Route **local nets (connections)** from level 0.
 - Two-stage routing (global + detailed routing) for a local net.

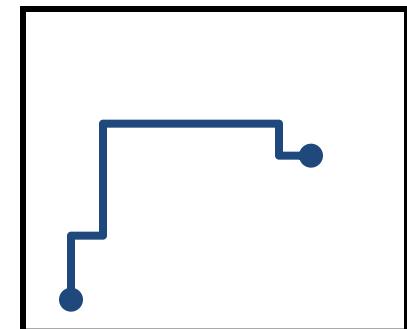
— — — an MST edge



— global route

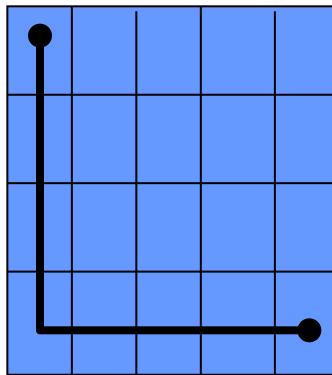


— detailed route

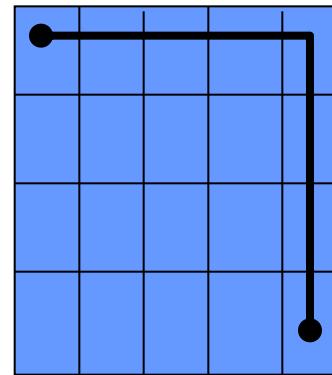


Global Routing

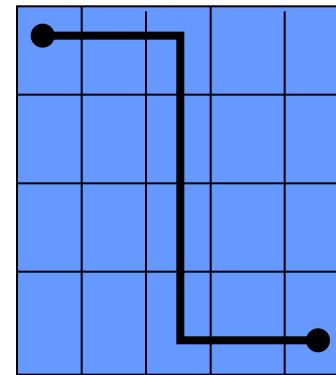
- Apply pattern routing for global routing
 - Use L-shaped and Z-shaped connections to route nets.
 - Has lower time complexity than maze routing.



**Lower L-Shaped
connection**



**Upper L-Shaped
connection**



**Z-Shaped
connection**

Cost Function in Global Pattern Routing

$G_i = (V_i, E_i)$: multilevel routing graph at level i .

Define $R_e = \{e \in E_i \mid e \text{ is the edge chosen to route}\}$

Then

$$\text{cost}(R_e) = \sum_{e \in E} C_e,$$

where C_e is the congestion of edge e

and

$$C_e = \frac{1}{2^{(p_e - d_e)}},$$

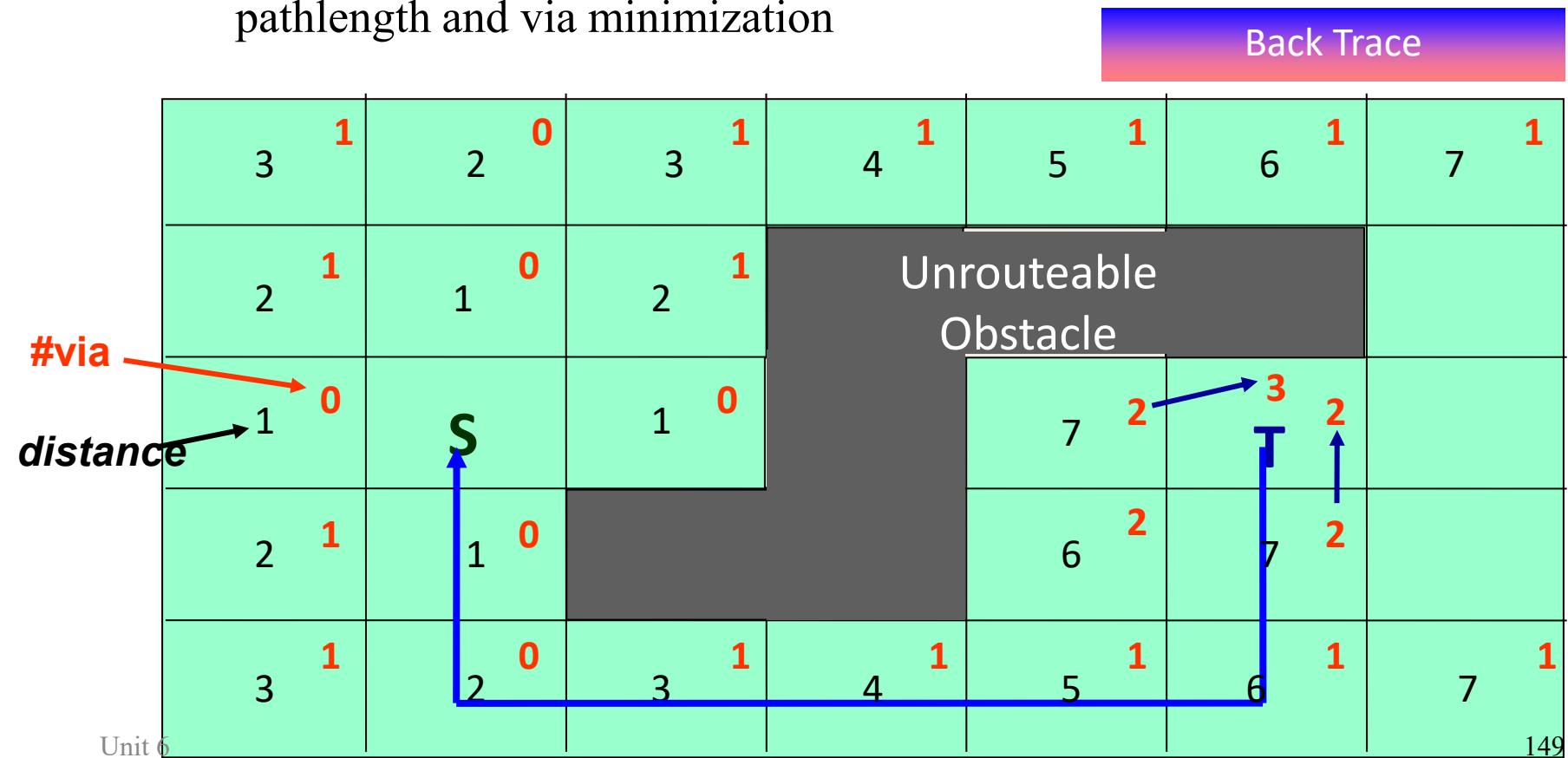
where p_e and d_e are the capacity and density associated with e , respectively.

Detailed Routing

- Via minimization
 - Modify the maze router to minimize the number of bends.
- Local refinement
 - Apply general maze routing to improve the detailed routing results.
- Resource estimation
 - Update the edge weights of the routing graph after detailed routing.

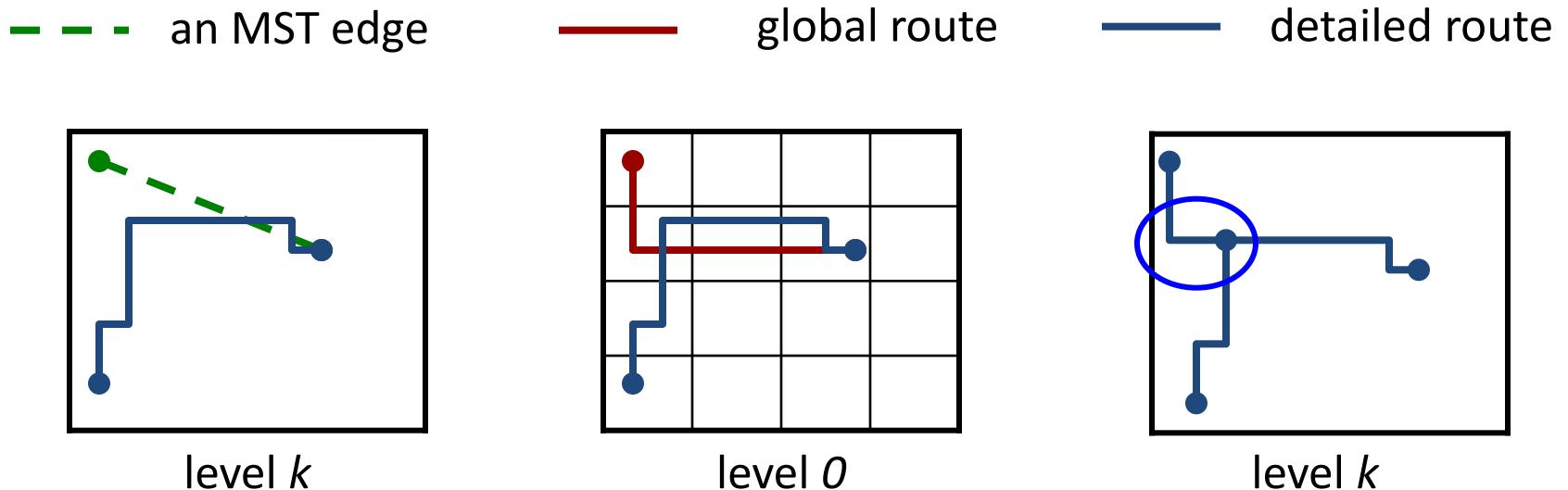
Via Minimization

- Simultaneous pathlength and via minimization (SPVM)
 - perform modified maze routing that simultaneously considers the pathlength and via minimization



Local Refinement

- Local refinement improves detailed routing results by merging two connections which are decomposed from the same net.



Resource Estimation

- Global routing cost is the summation of congestions of all routed edges.
- Define the congestion, C_e , of an edge e by

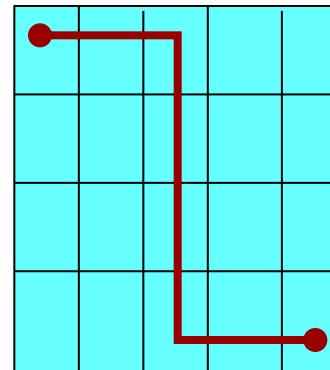
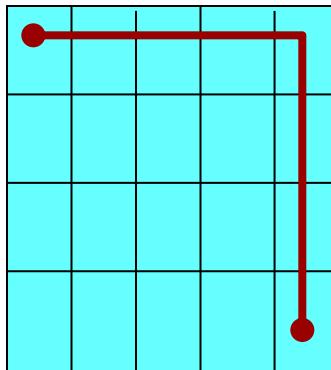
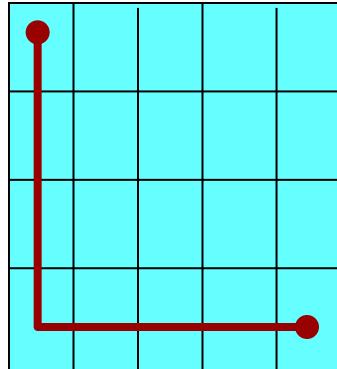
$$C_e = \frac{1}{2^{(p_e - d_e)}},$$

where p_e and d_e are the capacity and density, respectively.

- Update the congestion of routed edges to guide the subsequent global routing.

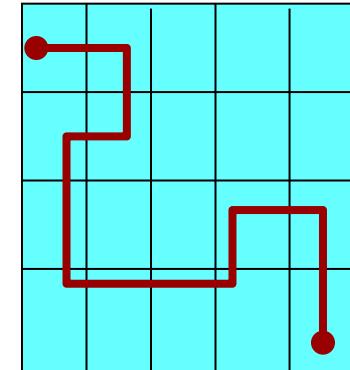
Uncoarsening Global Routing

- Use maze routing.
- Iterative refinement of a failed net is stopped when a route is found or several tries have been made.



Coarsening stage

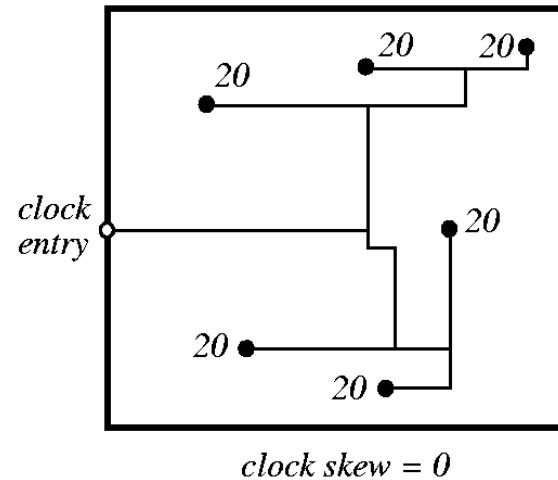
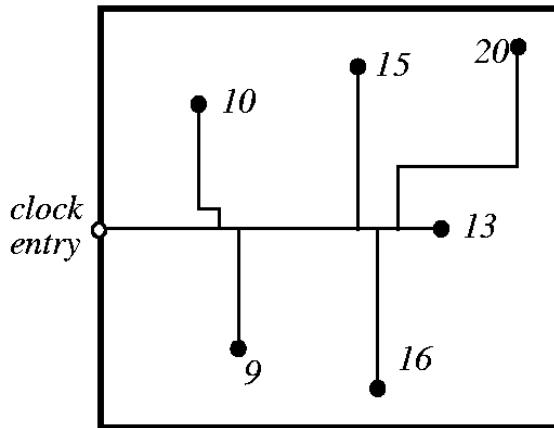
⋮



Uncoarsening stage

Clock Routing

- **Clock skew** is defined as the difference in the minimum and the maximum arrival time of the clock.



- Route clock nets such that
 - Clock signals arrive simultaneously
 - Clock delay is minimized
 - * Other issues: total wirelength, power consumption.

Clock Routing Problem

- Given the routing plane and a set of points $P = \{p_1, p_2, \dots, p_n\}$ within the plane and clock entry point p_0 on the boundary of the plane, the **Clock Routing Problem (CRP)** is to interconnect each $p_i \in P$ such that $\max_{i,j \in P} |t(0,i) - t(0,j)|$ (skew) and $\max_{i \in P} t(0,i)$ (delay) are both minimized.
- Pathlength-based approaches
 - H-tree: Dhar, Franklin & Wang, *ICCD*, 1984.
 - Methods of means & medians (MMM): Jackson, Srinivasan & Kuh, *DAC*, 1990.
 - Geometric matching: Kahng, Cong & Robins, *IEEE TCAD*, 1993.
- RC-delay based approaches:
 - Exact zero skew: Tasy, *ICCAD*, 1991.
 - Deferred merge embedding (DME) algorithm: Boese & Kahng, *ASICON-92*; Chao, Hsu & Ho, *DAC-92*; Edahiro, NEC R&D, 1991.