On the power of oritatami cotranscriptional folding with unary bead sequence ⁰

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What is oritatami system?

Oritatami system is a mathematical model for co-transcriptional folding(CF). (Geary, Meunier, Schabanel and Seki. MFCS 2016.)

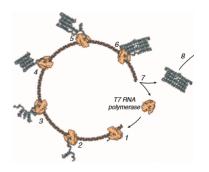
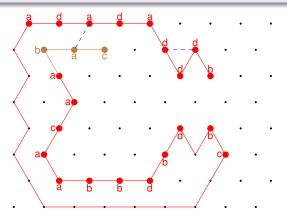


Figure: RNA tile is self-assembled (RNA Origami)

(Geary, Rothemund and Andersen. Science 345(6198), 2014)

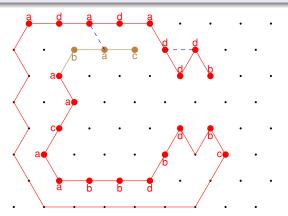
An example

 $\Sigma = \{a, b, c, d\}, R = \{(a, a), (b, b), (c, c), (d, d)\},$ arity $\alpha = 2$, delay $\delta = 3$, $w = \underline{bac}$ bcadbcbab



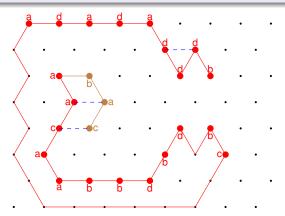
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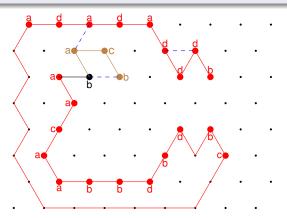
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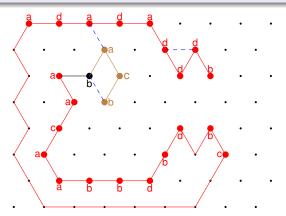
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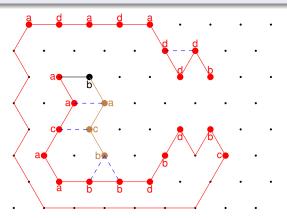
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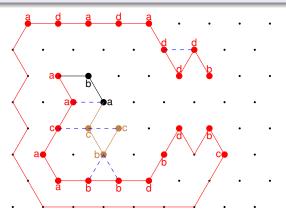
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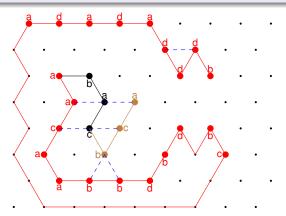
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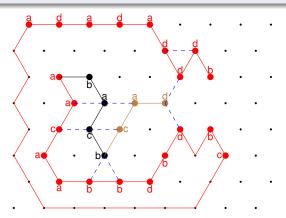
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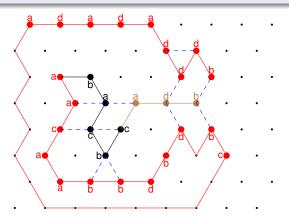
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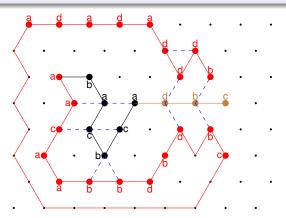
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An example

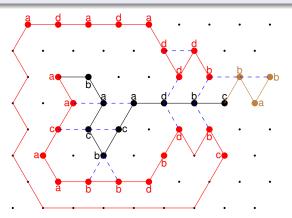
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Oritatami system and determinism

An example

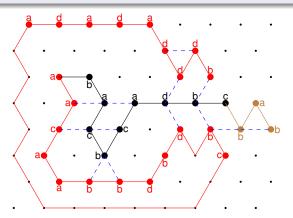
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Oritatami system and determinism

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Turing universality

Theorem (Geary, Meunier, Schabanel and Seki. ISAAC 2018.)

Deterministic oritatami system at *delay* $\delta=3$ are Turing complete. (the simulation needs 542 bead types)

- A smaller Turing universal system
- The characterization of non-Turing-universal systems

Theorem

Polynomial upper bounds on the size of structures

→ Not Turing universal

The system at *delay* $\delta=1$ and *arity* $\alpha=1$ is not Turing universal. (Demaine et al. DNA24, 2018.)

Problem

Problem

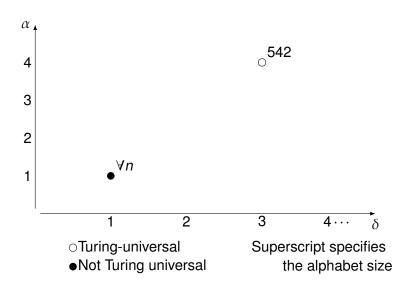
Give an upper bound on the length of a transcript of a *delay* δ , *arity* α deterministic oritatami system by a function in δ , α , and seed length n.

- input : delay δ , arity α and seed of length n
- output : an upper bound

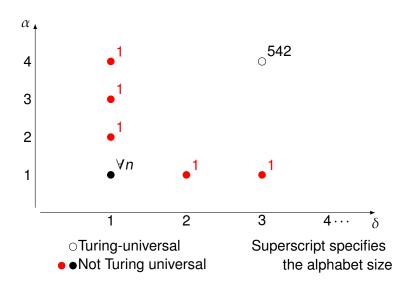
Unary variant

Let's focus on unary.

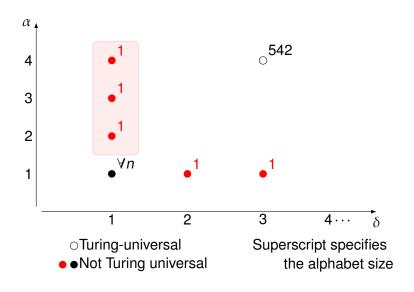
Turing universality and oritatami systems



Turing universality and oritatami systems



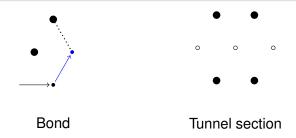
Turing universality and oritatami systems



Oritatami systems at delay 1

The two ways to stabilize a bead at delay 1

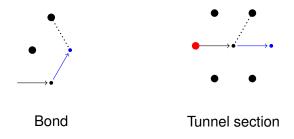
- To be bound to another bead.
- Through a 1-in-1-out structure called the tunnel section.



Oritatami systems at delay 1

The two ways to stabilize a bead at delay 1

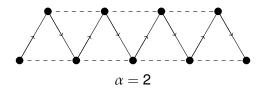
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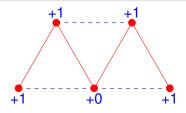


Results $(\delta = 1)$

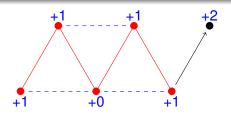
$$\alpha=4$$
 $3n^2+3n+1$
 $\alpha=3$ $4n+14$
 $\alpha=2$ ∞ but zigzag after $(27n^2+9n+1)$

^ac.f. $\alpha = 1$: 10*n* (Demaine et al. 2018 DNA24)

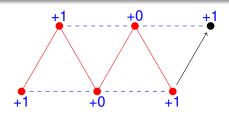




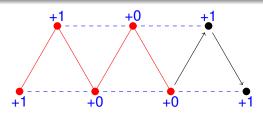
$$\alpha = 2$$



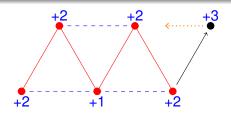
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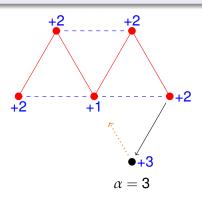
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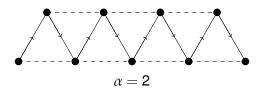
$$\alpha = 3$$



Results $(\delta = 1)$

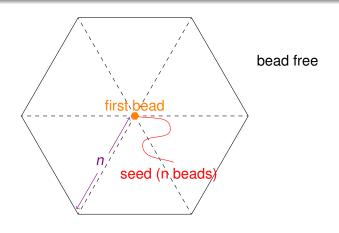
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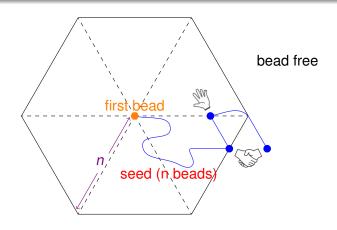
$$\alpha = 4$$

The terminal conformation at $\alpha = 4$ is of length at most $3n^2 + 3n + 1(\bigcirc_O^n)$.



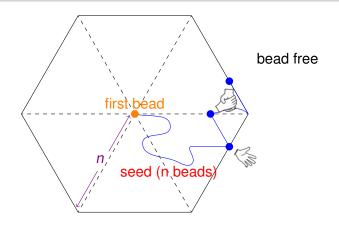
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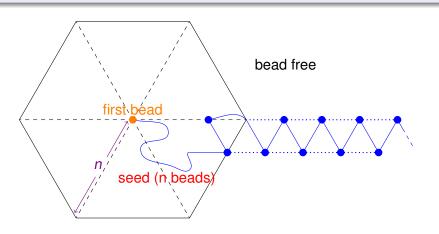
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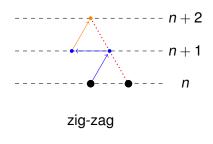
$$\alpha = 2 (\delta = 1)$$

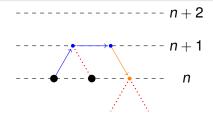
A transcript folds into the zig-zag conformation after its $(27n^2 + 9n + 1)$ -th bead (\bigcirc_O^{3n}) .



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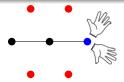
free hands: -2

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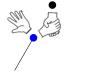
free hands: ±0

free hands: -2

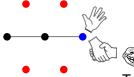
free hands: $\leq +2$

$\alpha = 2 (\delta = 1)$

A transcript folds into the zig-zag conformation after its $(27n^2 + 9n + 1)$ -th bead (\bigcirc_O^{3n}) .







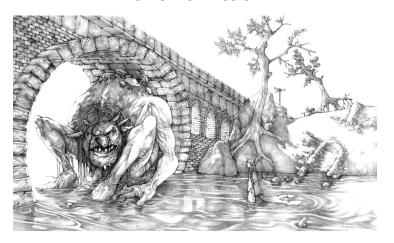
Troll

free hands: ±0

free hands: -2

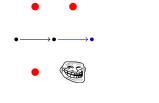
free hands: ± 0

Tunnel Troll Theorem

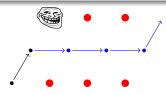


Tunnel Troll Theorem

- $\alpha \ge 4$ # of free hands does not increase / tunnel section.
- $\alpha = 3$ Troll consumes bonds / tunnel section.
- $\alpha = 2$ Troll consumes bonds / tunnel.



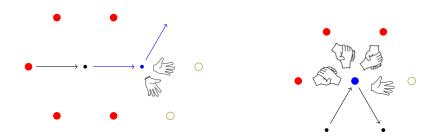




Tunnel

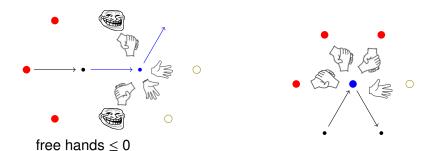
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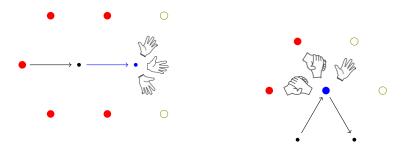


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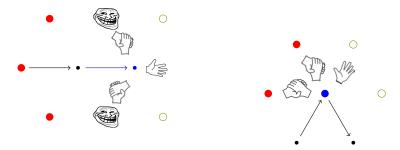


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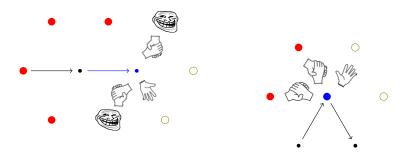
 $\alpha = 3$ At least one free hand is consumed / tunnel section.



free hands ≤ -1

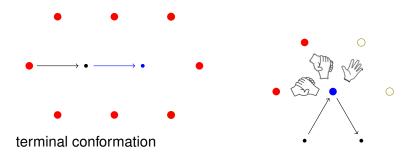
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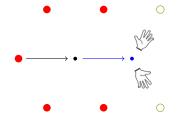


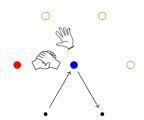
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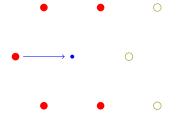


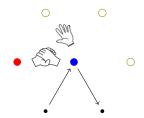
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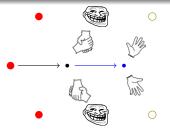


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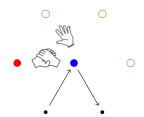




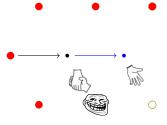
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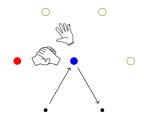




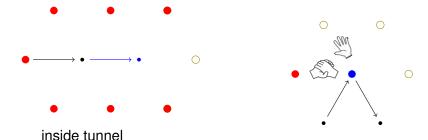
Tunnel Troll Theorem





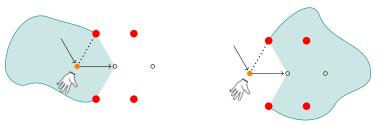


Tunnel Troll Theorem



Jordan curve theorem

A closed curve which is non-self-intersecting, divides the plane into inside and outside.



At $\alpha = 2$, Troll consumes free hands at an entrance of tunnel, too.

ご清聴ありがとうございました。



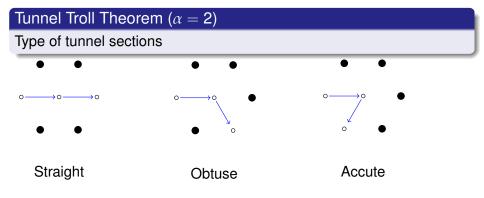
Szilárd Zsolt Fazekas Collaborators

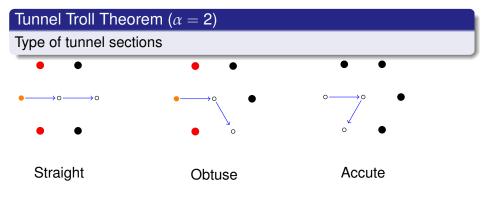


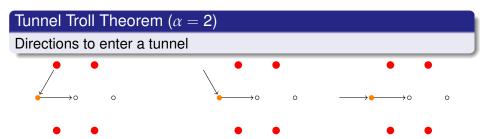
Shinnosuke Seki

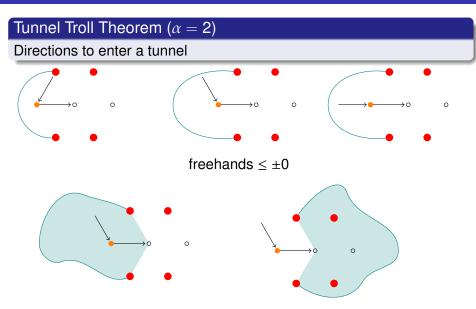
Future works

- Case of $\delta \geq 4$
- Non unary case
- Other ways of simulating Turing machine



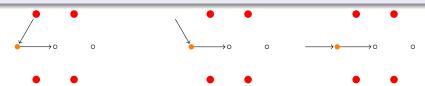


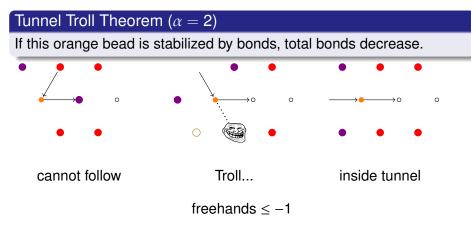




Tunnel Troll Theorem ($\alpha = 2$)

If this orange bead is stabilized by bonds, total bonds decrease.





References I