1 Infiniteness of delay-1, arity-2 unary deterministic oritatami system

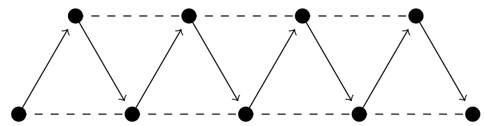


Fig. 1. zig-zag conformation

1.1 Introduction

In this section, we consider on finiteness of structures produced deterministically at delay 1. In result, cases of arity 1 and 3 can only yield finite structures which is size of $\mathcal{O}(n)$, and cases of arity 4 and more can only yield finite structures which is size of $mathcalO(n^2)$, and a case of arity 2 can yield infinite structures but they are only the zig-zag conformation shown in Fig.1.

Let Ξ be a deterministic oritatami system of delay 1 and arity 2. Assume its seed σ consists of n beads. For $i \geq 0$ let C_i be the unique elongation of σ by $w[1..i]$ that is foldable by Ξ. Hence $C_0 = \sigma$.

Let us consider the stabilization of the i-th bead a_i upon C_{i-1}. The bead cannot collaborate with any succeeding bead $w[i+1], w[i+2], \cdots$ at delay 1. There are just two ways to get stabilized at delay 1. One way is to be bound to another bead. The other way is through a *tunnel section*. A tunnel section consists of four beads that occupy four neighbors of a point (Fig.2). Accordingly, how they are stabilized can be described by a binary sequence S of b's (bound) and t's (tunnel section); priority is given to t, that is, $S[i] = t$ if the i-th bead w_i is stabilized not only by being bound but also through a tunnel section.

Assume that four of the six neighbors of a point p are occupied by beads $a_{j_1}, a_{j_2}, a_{j_3}, a_{j_4}$ while the other two are free. We call such a bead p as $inside\ of\ a\ tunnel$ and such beads p' as $entrance\ of\ a\ tunnel$ without a case that p' is inside of a tunnel. If the beads $w[i-2]$ and $w[i-1]$ are stabilized respectively at one of the two free neighbors and at p one after another, then the next bead $w[i]$ cannot help but be stabilized at the other free neighbor. In this way, $w[i]$ can get stabilized without being bound.

If a bead is stabilized through a tunnel section, then it can provide some bonds. Let us consider on bonds of C_i. C_i is represented $C = (W, P, H)$ where $|W| = i+n$ C_i contains $\alpha \cdot (i+n) - 2|H|$ bonds because C_i consists of $i+n$ beads and a bead has just α bonds and then $2|H|$ of the those bonds are already used. However, even if a bead has an available bond, $w[j]$ $(j > i)$ might not be able to use this bond because the bond has possibility that it is blocked by transcripts

$w[i+1..j-1]$. Number of *binding capabilities* does not contain that case so that it is at most $\alpha \cdot (i+n) - 2|H|$.

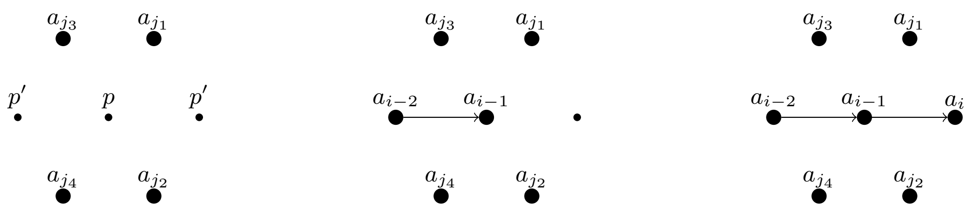


Fig. 2. Through a tunnel section

Theorem 1 (Tunnel Troll Theorem). *Let Ξ be a unary oritatami system of $\delta = 1, \alpha \geq 2$. If there is a part of transcript $w[i..j+1] = bttt...tb$, then C_i's binding capabilities is more than C_j's binding capabilities.*

1.2 Proof of Tunnel Troll Theorem

Assume Ξ is deterministic. Each of beads in transcript are bound either inside of tunnel or outside (Fig. 4). If a bead is stabilized at inside of tunnel, then its successor is already decided the position either inside of a tunnel or outside. Moreover, If a bead is stabilized at outside of a tunnel, then its position is an entrance of a tunnel or otherwise.

Tunnel sections have three possible shape with considering symmetry such as straight, acute turn and obtuse turn (Fig. 3).

Let us consider each of cases of tunnel A, B, and C. Accordingly, We use lemma 1

Lemma 1. *If a bead does not have any bond, then neighbors of it must be occupied by $\alpha + 2$ beads at $\delta = 1$ and unary except first and last beads.*

A bead in transcript needs predecessor and successor except first and last beads. If the bead does not have ant bond, then it use hand with α neighbors. Thus, lemma 1 is clearly true.

Let us consider tunnel sections only tunnel A and B. See Fig. 4. From appendix (Entrance of Tunnel A, B and Exit of Tunnel), we add maximum number which each transitions provide binding capabilities for each edge where a is number of consuming binding capabilities when the bead is stabilized at position of *successor in outside*.

Next, we consider on tunnel C section. If $w[i]$ is stabilized by tunnel C and $S[i+1]$ is t, then $w[i+1]$ is stabilized by tunnel A or B because if $w[i+1]$ is stabilized by tunnel C, then C_{i+1} is a terminal. Hence, tunnel C section is divided cases such as Figure 5. Cases of $S[i...] = bt^l(l \geq 2)$ are already considered (Upper). According to appendix (Tunnel C), cases of $S[i..i+2] = btb$ also consume some binding capabilities (Lower).

Thus, if a bead is stabilized through a tunnel section, then it consume some binging capabilities.

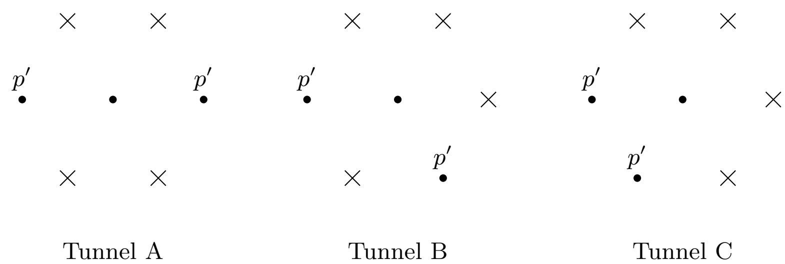


Fig. 3. All possible tunnel sections: straight, acute turn, and obtuse turn

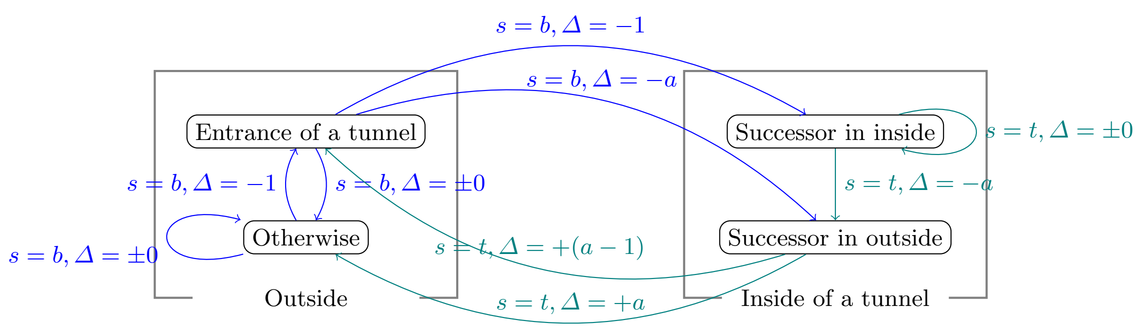


Fig. 4. Increment on Tunnel A,B

1.3 Appendix of Tunnel Troll Theorem

Entrance of TunnelA,B Fig.6 exhibits all the three kinds of entrance of tunnel A, B. Any cases in $\delta = 1, \alpha = 2$ consume some binding capabilities into the follows.

– Case of t_0

  Let us consider points of n_3, n_4 either occupied or not. A point n_3 or n_4 is free because if both of them are occupied, p' is inside of tunnel. If n_3 is free, then p' has to be bound to a bead except n_1 due to deterministically

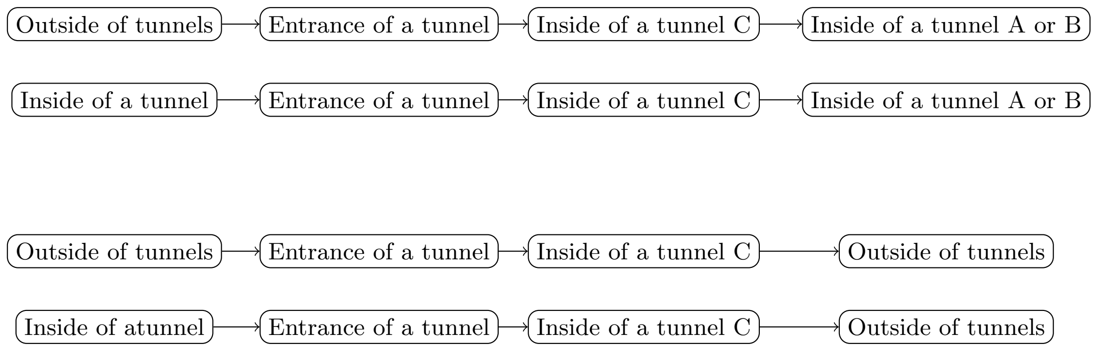


Fig. 5. Case of Tunnel C

stabilize. In this situation, at least three neighbors of n_1 are free that is n_1 leave at least one bond from lemma 1. Hence, p' must be bound to n_1. Thus, a case of t_0 consumes two binding capabilities and it does not supply any binding capabilities.

– Case of $t_{\pm 60}$
  In this case, a point n_4 or n_5 is free, too. If n_5 is free, p' has to be bound to n_1 or n_2. If n_5 is occupied, then n_4 is free. This time, n_2 has some binding capabilities so p' has to be bound to n_2.
  In this situation, p' is able to supply a binding capability. if this capability is active, p' bind a bead into n_4 or n_5. However, n_2 and n_3 are exist in back bone. According to Jordan curve theorem, any successors of p' cannot reach a point n_4 or n_5 so this capability is inactive. Thus, a case of $t_{\pm 60}$ consumes some binding capabilities.
– Case of $t_{\pm 120}$
  Binding capabilities that p' supply are inactive according to Jordan curve theorem on n_1 and n_2. Moreover, p' has to be bound to one of n_3, n_4, n_5 in order to deterministically stabilize. Thus, a case of $t_{\pm 120}$ consumes some binging capabilities.

Exit of Tunnel Fig.7 exhibits all the two kings of exit of tunnel. At least one of point n_1 or n_2 is free because if both of them are occupied, p' is inside of tunnel.

$\delta = 1, \alpha = 2$ Any cases of $\delta = 1, \alpha = 2$ supply at most a binding capabilities into follows where a is number of predecessor of p' consumes binging capabilities.

– Case of n_1 and n_2 are free
  This case can be regarded same situation as entrance. See Fig.7 (Left). Predecessor n_5 has to be bound n_4 and n_5 because each of n_3 and n_4 leave

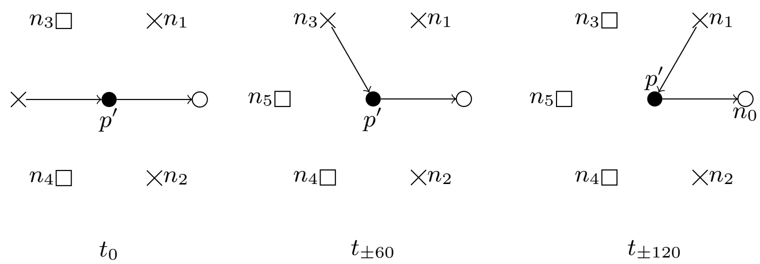


Fig. 6. Direction into a entrance

binding capabilities. Hence, at least $a = 2$. This time, $\alpha = 2$ that is this case supply at most only a binding capabilities.

- Case of c is occupied
  See Fig.7 (Right). If n_1 is occupied, then n_2 is free so that n_5 has to be bound n_4. Hence, at least $a = 1$. This case can supply two binding capabilities but p' can bind to only one of n_0 or n_2 because n_0 or n_2 will be occupied a successor of p'. Therefore, this case supply at most only $a = 1$ active binding capability.

$\delta = 1, \alpha \geq 3$ Any cases of $\delta = 1, \alpha \geq 3$ consume some binding capabilities into follows.

- Case of n_1 and n_2 are free
  In $\alpha \geq 3$, if three neighbors of a bead leave, then it can supply two binding capabilities. Therefore Predecessor n_5 has to be bound n_3 and n_4, and p', too. In this case, at least consumes four bindings and supplies at most two bindings. Thus, it consumes some binding capabilities, totally.
- Case of c is occupied
  In this case, n_4 leave at least two bindings and n_3, n_1 also leave at least one binding. Therefore n_5 has to be bound n_3 and n_4, and p' also has to be bound n_1 and n_4. In this case, at least consumes four bindings and supplies at most one binding. Thus, it consumes some binding capabilities, totally.

Point n_1 and n_2 are free　　Point n_1 is occupied

Fig. 7. Exit of Tunnel

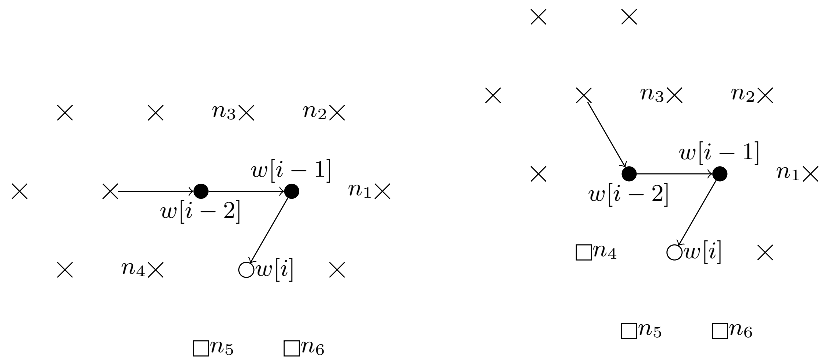
Tunnel C Assume $w[i]$ is a bead which stabilized by tunnel C. Let us consider kinds of stabilization $S[i-2..i] = tbt$ or $S[i-2..i] = bbt$ except cases of $w[i]$ is inside of tunnel A, B.

Case of $S[i-2..i] = tbt$ Fig.8 exhibits all the two kinds of stabilization depending on structures of tunnel C.

- Left of Fig.8
  In this figure, Bead n_4 has at least one binding so that $w[i-2]$ has to bound n_4. Moreover, $w[i-1]$ has to bound one of n_1, n_2, n_3 in order to stabilize deterministically. On the other hand, $w[i]$ can supply two bindings but free neighbors of $w[i]$ are two points. One of them is occupied by a successor. Therefore $w[i]$ can only bind one of n_5, n_6 that is $w[i]$ supplies at most one binding. Thus, this case consumes some binding capabilities.
- Right of Fig.8
  This cases are divided on number of capabilities that $w[i-2]$ consumes.
  - $w[i]$ does not consume any bindings
    $w[i-1]$ has to bound one of n_1, n_2, n_3 in order to stabilize deterministically. $w[i]$ has to be bound to $w[i-2]$ because $w[i-2]$ has bindings. This time, let us consider either n_4 is occupied or not. If n_4 is occupied, then $w[i-2]$ has no active bindings that is this situation consumes some binding capabilities. If n_4 is free and $w[i+1]$ is stabilized in n_4, then $w[i-2]$ has to bind $w[i+1]$. Therefore, In this case, stabilization of $w[i-2..i+1]$ consumes some bindings. If n_4 is free and $w[i+1]$ is stabilized except n_4, then this oritatami system has to use two binding capabilities in order to bind $w[i+1]$. Therefore, in this case consumes some bindings. Thus in this cases consume some binding capabilities.
  - w_i consumes one binding
    In this case, w_{i-1} has to be bound one of n_1, n_2, n_3. In addition, $w[i-2]$ and $w[i]$ are not supply any bindings. Thus, in this cases consume some binding capabilities.
  - w_i consumes two bindings
    In this case, $w[i-2]$ already consumes two binding. $w[i-1]$ has to be bound. $w[i]$ supplies two bindings. Thus, in this cases consume some binding capabilities.

Case of $S[i-2..i] = bbt$ Let us consider number of consumed by $w[i-2]$ (Fig.9).

- $w[i-2]$ consumes one binding
  In this situation, $w[i-2]$ supplies one active binding whereas $w[i]$ consumes this binding. In addition, $w[i-1]$ has to bound to one of n_1, n_2, n_3. Thus, in this cases consume some binding capabilities.



Pattern 1 Pattern 2

Fig. 8. Case of $S[i-2..i] = tbt$

- $w[i-2]$ consumes two bindings
  In this case, $w[i-2]$ already consumes two binding. $w[i-1]$ has to be bound. $w[i]$ supplies at most two bindings. Thus, in this cases consume some binding capabilities.

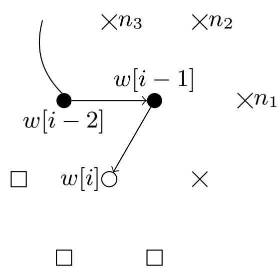


Fig. 9. Case of $S[i-2..i] = bbt$

By Tunnel Troll Theorem, any tunnel sections which represented in bbt^+ or bt^+bt^+ consume binding capabilities. If the sequence S is free from any subsequence of the form bt^+bt^+, then it can factorize as $S = u_1u_2u_3\cdots$ for some $u_1, u_2, u_3, \cdots \in \{b\} \cup bbt^+$. Assume the length of σ is n, seed supplies at most $2n$ binding capabilities. Therefore formula 1 hold.

$$\exists i \in \mathbb{N} \quad s.t. \quad u_{i-1}, u_i, u_{i+1}, u_{i+2}, \cdots \in \{b\} \tag{1}$$

Let us represent S as $S[i.i+1...] = v_iv_{i+1}v_{i+2}\cdots$ for some $v_i, v_{i+1}, v_{i+2}, \cdots \in \{a, o\}$ where if v_k is a, then v_{k+1} is bound to v_{k-1}, if v_k is o, then v_{k+1} is NOT bound to v_{k-1}.

Let us consider the case of v_k is o. See Fig.10. v_{i-1} consumes some binding capabilities because $S[i-1]$ is b. If the number of v_{i-1}'s bindings is one binding, then v_{i+1} has to be bound except A or B so that v_{i+1} must consumes two bindings except the case of A and B are occupied and v_i consumes at least one binding. If A and B are occupied, then v_{i-1}'s binginds are inactive that is v_{i-1} consumes two binding capabilities. Therefore, this case consumes binding capabilities. If v_{i-1} dose Not have any bindings, then v_{i-1} already consumes two bindings. In addition, v_i and v_{i+1} consume at least one binding. Therefore this case consumes binding capabilities. Thus, the formula 2 hold and according to the formula 1 and the formula 2, the formula 3 is hold.

$$\exists j \in \mathbb{N} \quad s.t. \quad u_j, u_{j+1}, u_{j+2}, \cdots \in \{a\} \tag{2}$$

$$|S| > \forall m \in \mathbb{N} \quad \rightarrow \quad \exists n \in \mathbb{N} \quad s.t. \quad S[n], S[n+1], \cdots \in \{a\} \tag{3}$$

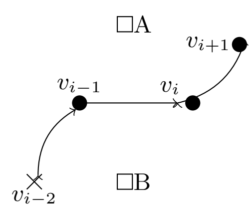


Fig. 10. Case of $S[i]$