

TTK4145 – Real-time Programming

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Innhold

1 Fault model and software fault masking

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This chapter ...

Fault model The one used in this chapter involves three entities: *processes, messages and storage*.

Unexpected faults Faults that are not tolerated by the design. Two categorizations:

- *Dense faults*: The algorithms will be n -fault tolerant. If there are more than n faults within a repair period, the service may be interrupted (in this case the system should be designed to stop).
- *Byzantine faults*: The fault model postulates certain behavior – for example it may postulate that programs are failfast. Faults in which the system does not conform to the model behavior are called Byzantine.

Underlying progression

- *Failfast*: They either execute the next step, or they fail and reset to the null state.

- *Available*: Failfast + repairability
- *Reliable*: Continuous operation

Checkpoint-Restart Write state to storage after each acceptance test, but before each event. This approach can be somewhat slow (hours/days).

Process pairs OS generates a backup process for each new primary process. The primary sends *I'm Alive* messages to the backup on a regular basis.

The backup can take over in three different ways:

- *Checkpoint-restart*. The primary records its state on a duplexed storage module. At takeover, the backup starts by reading these duplexed storage pages. (quick repair)
- *Checkpoint message*. The primary sends its state changes as messages to the backup. At takeover, the backup gets its current state from the most recent message. (basic pairs must checkpoint)
- *Persistent*. The backup restarts in the null state and lets the transaction mechanism clean up (undo) any recent uncommitted state changes. (simple)

