

## Algorithms Analysis - Handout 2

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### Exercise 1

1.  $x^2 = O(x^3)$
2.  $x^3 = O(x^2)$
3.  $5x^3 + 3x^2 = \Omega(x^4)$
4.  $3x^2 + 5x + 2 = \Theta(x^2)$
5.  $2^{n+1} = O(2^n)$
6.  $3^n = O(2^n)$
7.  $3^n = \Omega(2^n)$

#### 1. True.

Big-O notation represents an upper bound on the growth of a function. Since  $x^3$  grows faster than  $x^2$ ,  $x^2$  is  $O(x^3)$ .

$$x^2 \leq x^3 \cdot c$$

divide by  $x^3$

$$\frac{1}{x} \leq c$$

As  $x \rightarrow \infty$ ,  $\frac{1}{x} \rightarrow 0$ , so the equation is true for any  $c \geq 0$  and  $x_0 \geq 1$ .

#### 2. False.

$x^3$  grows faster than  $x^2$ , so  $x^3$  cannot be bounded above by  $x^2$ .

$$x^3 \leq x^2 \cdot c$$

divide by  $x^2$

$$x \leq c$$

This inequality is false for large  $x$  because  $x$  grows without bound.

**3. False.**

The term  $\Omega(x^4)$  suggests a lower bound, but  $5x^3 + 3x^2$  grows slower than  $x^4$ . Hence, it is not bounded below by  $x^4$ .

$$5x^3 + 3x^2 \geq c \cdot x^4$$

divide by  $x^4$

$$\frac{5}{x} + \frac{3}{x^2} \geq c$$

As  $x \rightarrow \infty$ ,  $\frac{5}{x} + \frac{3}{x^2} \rightarrow 0$ , so there is no such  $c$  that inequality would be true.

**4. True.**

The polynomial  $3x^2 + 5x + 2$  has a dominant term of  $x^2$  as  $x \rightarrow \infty$ . Therefore, it has both upper and lower bounds that are proportional to  $x^2$ , making it  $\Theta(x^2)$ .

$$c_1 \cdot x^2 \leq 3x^2 + 5x + 2 \leq c_2 \cdot x^2$$

We can ignore terms  $5x + 2$  for large  $x$ , for which both lower and upper bounds can be satisfied by choosing constants.

For example  $c_1 = 1, c_2 = 4$

**5. False.**

$2^{n+1} = 2 \cdot 2^n$ , which grows faster than  $2^n$ . Thus,  $2^{n+1}$  is not  $O(2^n)$ .

$$2 \cdot 2^n \leq c \cdot 2^n$$

divide by  $2^n$

$$2 \leq c$$

Which implies that the statement is false.

**6. False.**

$3^n$  grows exponentially faster than  $2^n$ , so it cannot be bounded above by  $2^n$ . Therefore, the statement is false.

$$3^n \leq c \cdot 2^n$$

divide by  $2^n$

$$\left(\frac{3}{2}\right)^n \leq c$$

$\left(\frac{3}{2}\right)^n$  grows exponentially, so this inequality would be false for large  $n$

### 7. True.

Since  $3^n$  grows faster than  $2^n$ ,  $3^n$  is indeed bounded below by  $2^n$ , making this statement true.

$$3^n \geq c \cdot 2^n$$

divide by  $2^n$

$$\left(\frac{3}{2}\right)^n \geq c$$

$\left(\frac{3}{2}\right)^n$  grows exponentially, which makes this inequality true for large  $n$  and any  $c \geq 1$ .

## Exercise 2

```
MAX-HEAPIFY( $A, i$ )
1   $l \leftarrow \text{LEFT}(i)$ 
2   $r \leftarrow \text{RIGHT}(i)$ 
3  if  $l \leq \text{heap-size}[A]$  and  $A[l] > A[i]$ 
4      then  $\text{largest} \leftarrow l$ 
5      else  $\text{largest} \leftarrow i$ 
6  if  $r \leq \text{heap-size}[A]$  and  $A[r] > A[\text{largest}]$ 
7      then  $\text{largest} \leftarrow r$ 
8  if  $\text{largest} \neq i$ 
9      then exchange  $A[i] \leftrightarrow A[\text{largest}]$ 
10     MAX-HEAPIFY( $A, \text{largest}$ )
```

```
MAX-HEAPIFY( $A, i$ )
while (true)
     $l = \text{LEFT}(i)$ 
     $r = \text{RIGHT}(i)$ 
    if  $l \leq \text{heap-size}[A]$  and  $A[l] > A[i]$ 
         $\text{largest} = l$ 
    else
         $\text{largest} = i$ 
```

```

    if r <= heap-size[A] and A[r] > A[largest]
        largest = r
    if largest != i
        exchange A[i] with A[largest]
        i = largest
    else
        break
end

```

### Exercise 3

The following pictures show sorting an array using heap sort.









