

## Networking / Video Conferencing

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## Admin

- PA2
  - ❖ Write-up due now
- HW 5
  - ❖ Due now
- HW 6
  - ❖ Will be assigned Wed.
  - ❖ Due Wed. 12/2/2015

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# Video Conferencing

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## Video Bit Rate Scaling

Temporal scaling - simply reduce frame rate

- Easy to do @ encoding
- Harder to do in network or other layer

I P P P P P I P P P P ..

Interactive streaming usually has no B-frame

Best to drop  $\Rightarrow$  no dependency on it

This one you would like to be

next but: m.r. get screwed up

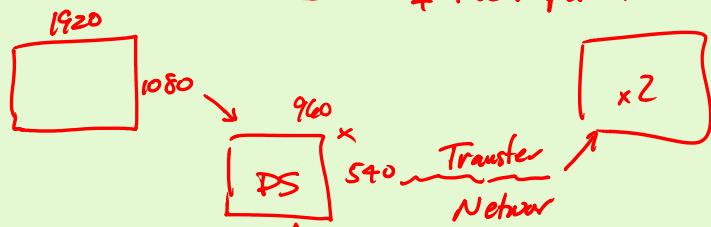
drift errors  $\Leftarrow$  m.r. that are reconstructed but off a little

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## Video Bit Rate Scaling

Spatial scaling - reducing the number of pixels

Cut resolution in  $\frac{1}{2}$   $\rightarrow \frac{1}{4}$  the # pixels!



You would expect  $\uparrow$  to be  $\frac{1}{4}$  the bit rate

In practice DS needs higher <sup>lower</sup> Quantization quality to make up details in image

Shoot for  $\frac{1}{2}$  bit rate.

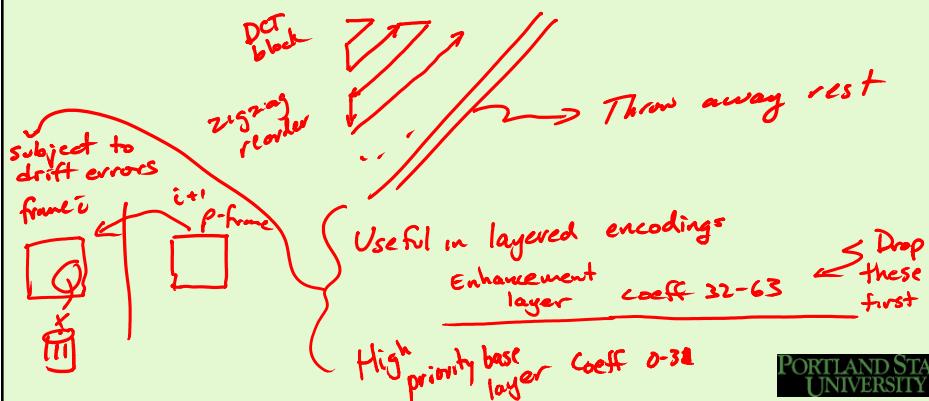
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## Video Bit Rate Scaling

Freq scaling - reduce DCT coefficients

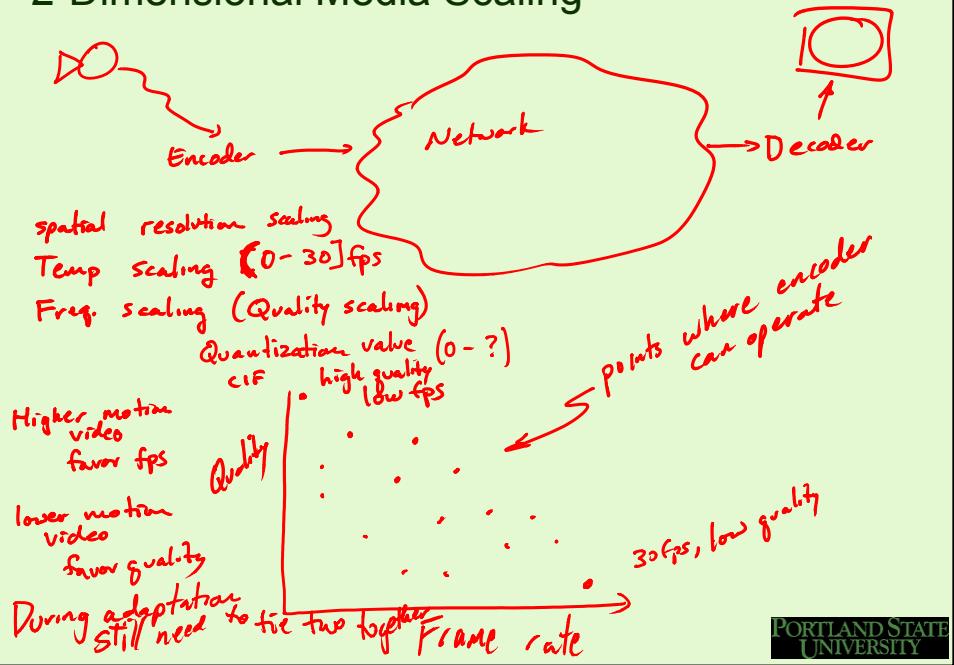
•) Traditionally accomplished by increasing quantization factor

•) Limit the number of coefficients



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## 2-Dimensional Media Scaling



## ~~Video Bit Rate Scaling~~

### 2 D Scaling

For videos operating points Q1, Q2, Q3, Q4  
fps 5, 10, 15, 24, 30 fps

Highest quality = 30fps, Q1

30fps, Q2

24fps, Q2

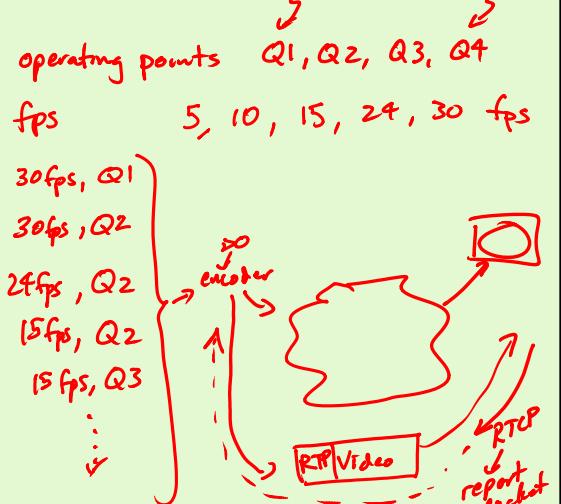
15fps, Q2

15fps, Q3

⋮

Lowest quality

5fps, Q4



## Live Video over Best Effort

Principles / characteristics of best-effort networking

Ideal: no delay jitter, no packet loss

Delay jitter is variation in end-to-end delay

loss in network due to congestion

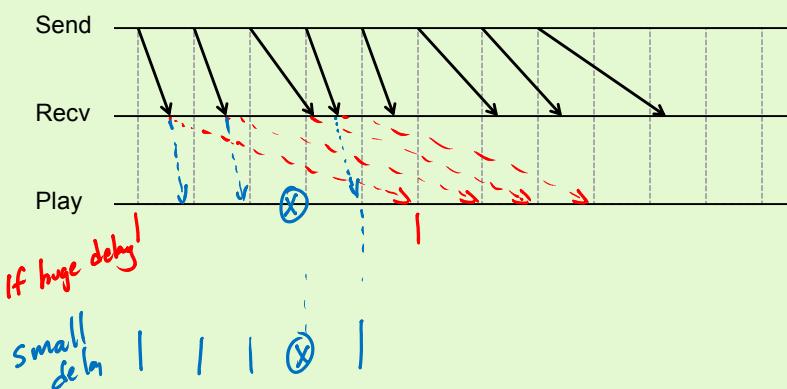
- show goes on
- manage through FEC - forward error correction
- Retransmit (increases delay)
- show goes on but conceal loss

Out-of-order delivery



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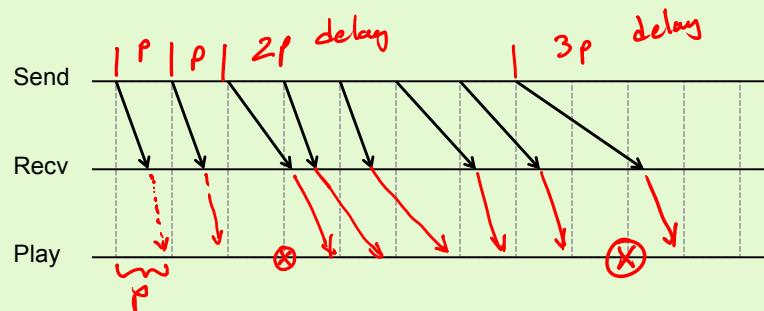
## Taking Care of Delay Jitter



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## Fundamental Initial Playout Strategies

Let naturally occurring network delay determine playout latency



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## Fundamental Initial Playout Strategies

Avoid initial sequence of playout gaps by estimating network delay and setting it accordingly



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## Estimating Network Delay

Send a packet with a timestamp

Enqueued @ other side timestamp + d

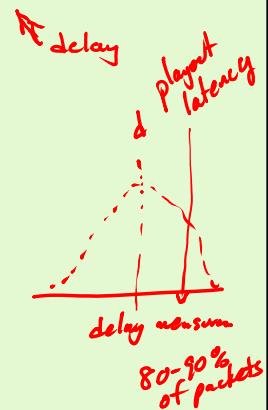
Basic algorithm

$$\text{playout latency} = d + k * v$$

d = estimated average delay

v = estimated variation in delay

k = congestion estimator

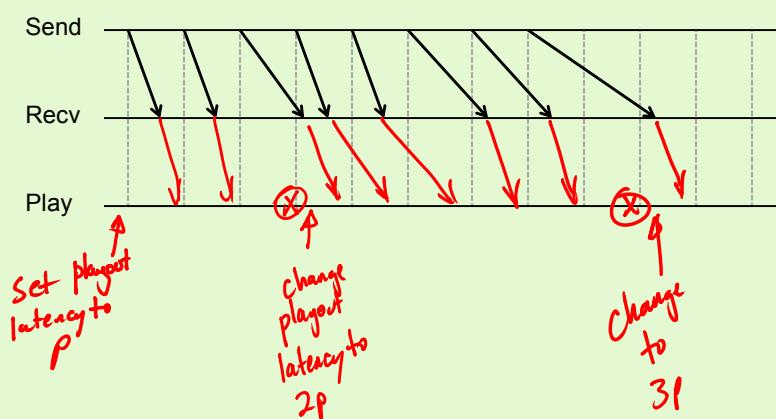


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## Fundamental Late Arrival Strategies

Play media as they arrive

Essentially change playout latency

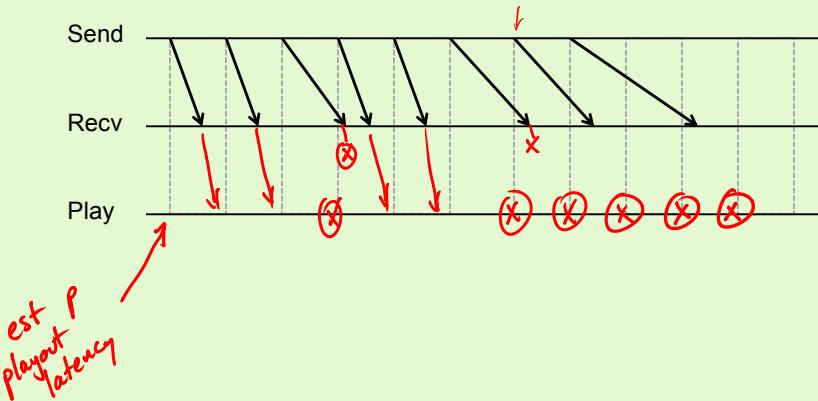


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## Fundamental Late Arrival Strategies

Discard "late" samples

Keep playout latency the same



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## Fundamental Tension

Large delay = continuous playback  
poorer interactivity

Small delay = gaps in playback  
better interactivity

Carefully balance both

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## Queue Monitoring

Rather than compute network delay, infer it from length of display queue

growing display queue - network delay is decreasing

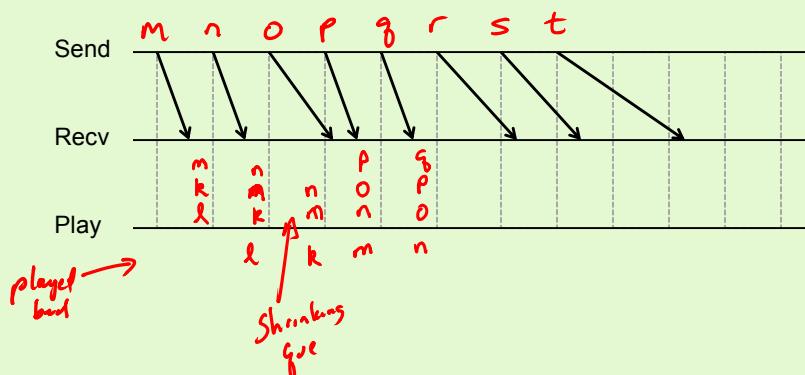
shrinking display queue - network delay is increasing

constant = stable over time

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## Queue Monitoring

Example



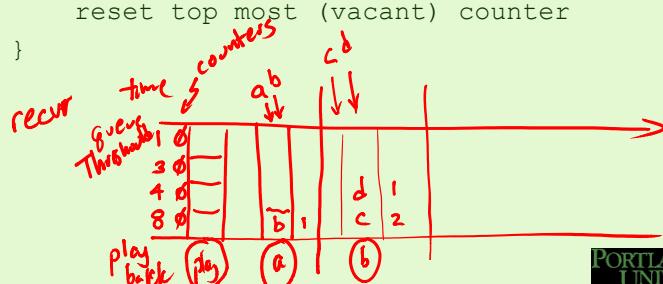
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## Queue Monitoring Algorithm Pseudo-code

```

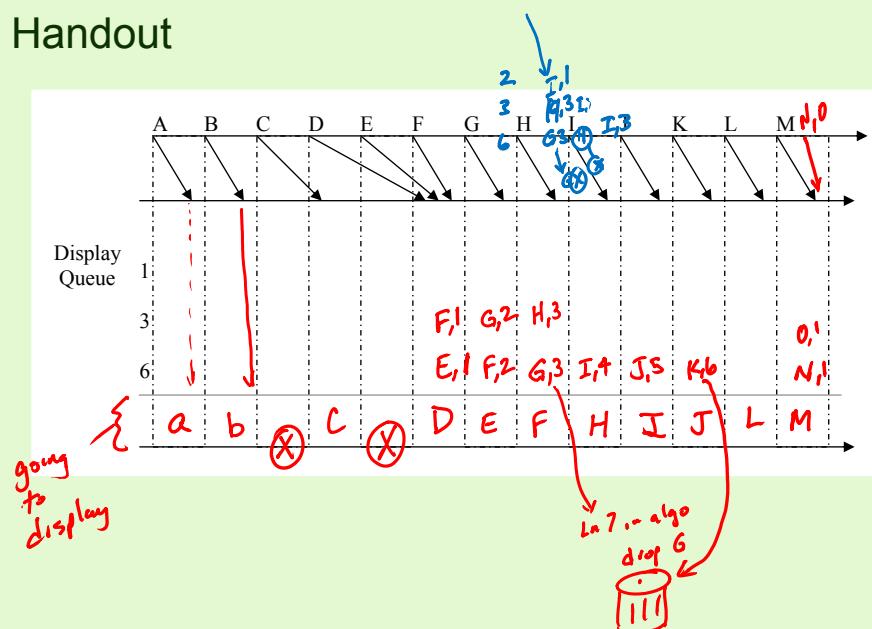
1. Set counters for all levels = 0;
2. while (1){
3.     Move oldest data to playback;
4.     Enqueue all arriving samples in period
5.     Increment all counters for levels with data
6.     Reset all counters for levels w/ no data
7.     If any counter >= to its threshold
8.         drop the oldest sample in queue
9.         reset top most (vacant) counter
10.    }

```



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## Handout



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