## Does Logic-based Reasoning Work for Dutch?

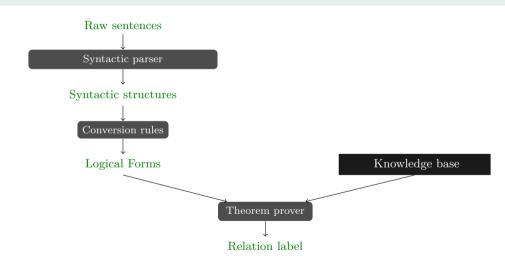
Lasha Abzianidze Konstantinos Kogkalidis

Utrecht Institute of Linguistics OTS, Utrecht University

<del>06-07-2021 @ CLIN</del> 21-06-2021 @ NLP RG

TIOS/FLIAGT-NETHERLANDS.PIRGESUITS

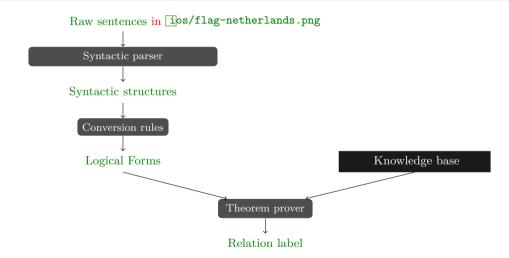
# Logic-based approach to NLI



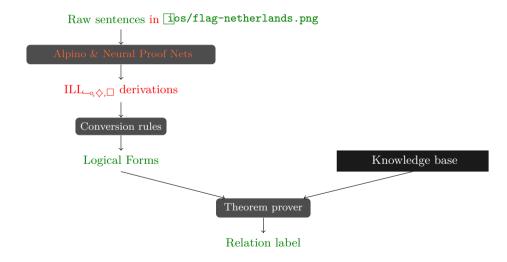
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Overview

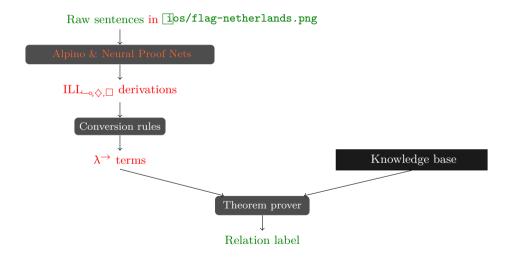
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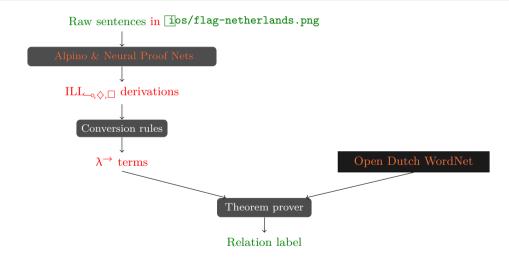
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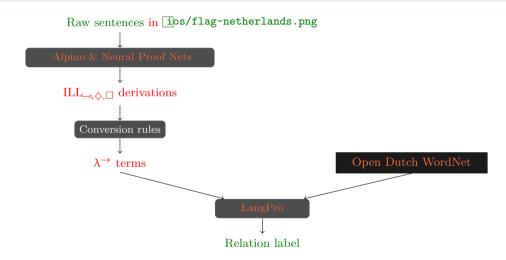


# Logic-based approach to NLI



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# Logic-based approach to NLI



TIONS/PHIAGENETHERLANDS.PIRGESUITS

### The linear $\lambda$ -calculus

Words are assigned ILL types, inductively defined as:  $\mathcal{T} := a \mid t_1 \multimap t_2 \mid \diamondsuit^{\delta} t \mid \Box^{\alpha} t$  where

- ► a an atom, from a finite set A:

  np,  $s_{main}$ ,  $s_{sub}$ , pron, ...
- ▶  $t_1 \multimap t_2$  a linear function that consumes  $t_1$  to produce  $t_2$  $np \multimap s_{main}$ ,  $np \multimap (np \multimap s_{main})$ ,  $(np \multimap s_{sub}) \multimap (np \multimap np)$ , ...

Syntactic derivations  $\equiv$  proofs  $\stackrel{\text{chc}}{\equiv}$  functional programs

$$\tau := \operatorname{c}^{\operatorname{t}} \mid \left(\tau_1^{\operatorname{t}_1 - \operatorname{ot}_2} \tau_2^{\operatorname{t}_1}\right)^{\operatorname{t}_2} \mid \left(\lambda \operatorname{x}^{\operatorname{t}_1} . \tau^{\operatorname{t}_2}\right)^{\operatorname{t}_1 - \operatorname{ot}_2} \mid \delta_{\square} \tau \mid \delta^{\square} \tau \mid \delta_{\diamondsuit} \tau \mid \delta^{\diamondsuit} \tau$$

A boy plays

$$\operatorname{play}^{\diamond \operatorname{sunp} - \circ \mathbf{s}_{\operatorname{main}}} \left( \operatorname{su}^{\diamond} \left( \left( \operatorname{det}_{\square} \ \operatorname{a}^{\square \operatorname{det} (\mathbf{n} - \circ \mathbf{np})} \right) \operatorname{boy}^{\mathbf{n}} \right) \right)$$

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A boy plays:

$$\operatorname{play}^{\diamondsuit^{\operatorname{su}} np - \circ s_{\operatorname{main}}} \left( \operatorname{su}^{\diamondsuit} \left( \left( \operatorname{det}_{\square} \ \operatorname{a}^{\square^{\operatorname{det}} (n - \circ np)} \right) \operatorname{boy}^n \right) \right)$$

I Distribut NETHERLANDS . PIResults

Overview

Types are trees, with nodes inductively polarized: + we have - we seek



Conclusion

## **Proof Nets**

Overview

Types are trees, with nodes inductively polarized: + we have – we seek

> npn example a very easy

I Distribut NETHERLANDS . PIResults

### **Proof Nets**

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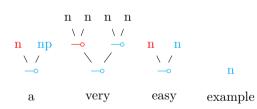
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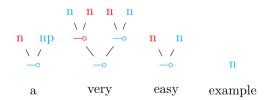
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Overview

Proof net

0000

a polarized forest (proof frame) and a bijection between + and - (axiom links)

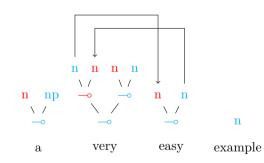


### Proof Nets

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very easy

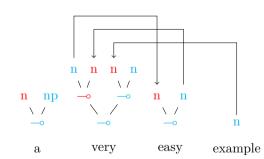
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very easy example

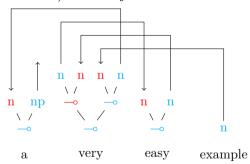
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0000

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a (very easy example)

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Neural Proof Nets: from sentences to  $\lambda$ -terms

Supertagging

Overview

From sentences to proof frames with seq2seq transduction

example a very easy

Conclusion

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Neural Proof Nets: from sentences to  $\lambda$ -terms

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n<sub>1</sub>

a very easy example

TIOS/FLIAGT-NETHERLANDS.PIRGESUITS

Conclusion

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From sentences to proof frames with seq2seq transduction

very

 $\begin{array}{ccc}
\mathbf{n_1} & \mathbf{np_2} \\
& & \\
& & \\
& & \\
\mathbf{a}
\end{array}$ 

easy example

TIOS/FLIAGT-NETHERLANDS.PIRGESUITS

Conclusion

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 $\begin{array}{ccc} n_1 & np_2 \\ & & & \\ & & - \\ & & & \end{array}$   $\begin{array}{ccc} a & & \text{very} \end{array}$ 

easy example

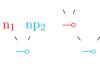
TIOS/FLIAGT-NETHERLANDS.PIRGESUITS

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a

0000

very

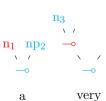
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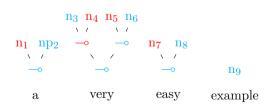


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### Neural Proof Nets: from sentences to $\lambda$ -terms

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Proving

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From proof frames to axiom links with Sinkhorn-Knopp

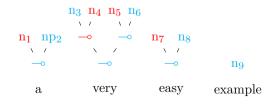
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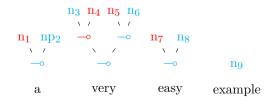
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# Other ingredients: Alpino & ODWN

### Alpino:

Overview

- ► Stochastic Attribute Value Grammar (HPSG) for Dutch
- ► Builds dependency structures
- ► Used for pre-processing Dutch treebanks

#### Open Dutch WordNet

- ► 52K synsets (vs 117K Princeton WN)
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- ► We converted ODWN into the prolog forma
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- ► Natural logic + semantic tableau
- Uniform to Ent./Cont.
- ▶ Prove with refutation
- $\triangleright \approx \text{Syntactic trees}$

- ► Native higher-order logic

#### ios/flag-united-kingdom.png Natural Tableau: proving

- 1 a hedgehog (be  $(\lambda x. a \text{ bov } (\lambda y. \text{ by } y \text{ cradle } x))) : \mathbb{T}$ 2 a person ( $\lambda x$ . an animal ( $\lambda y$ . hold y x):  $\mathbb{F}$
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  - 6 by b cradle: [h]:  $\mathbb{T}$

### Natural Tableau: proving

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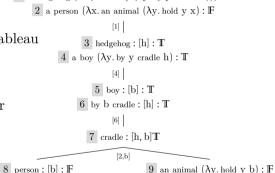
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person: [b]: F

[5,8]× 9 an animal (λy. hold y b): F

### Natural Tableau: proving

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```
2 a person (\lambda x. an animal (\lambda y. hold y x): \mathbb{F}
             3 hedgehog: [h]: T
       4 a boy (\lambda y. by y cradle h): T
                      [4]
                5 boy: [b]: T
            6 by b cradle: [h]: T
                      [6]
               7 cradle: [h, b]T
                       [2,b]
```

X

[9,h]

### Natural Tableau: proving

Overview

#### ios/flag-united-kingdom.png

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```

[6]

7 cradle: [h, b]T

[2,b] person: [b]: F 9 an animal (λy. hold y b): F [5,8]

10 animal : [h] : F

11 hold: [h, b]: F

person: [b]: F

×

[5,8]

Overview

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X

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[3,10]

person: [b]: F

×

[5,8]

Conclusion

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6 by b cradle: [h]: T

[6] 7 cradle: [h, b]T

[2,b]

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[9,h]

10 animal : [h] : F [3,10]

11 hold: [h, b]: F [7,11]

×

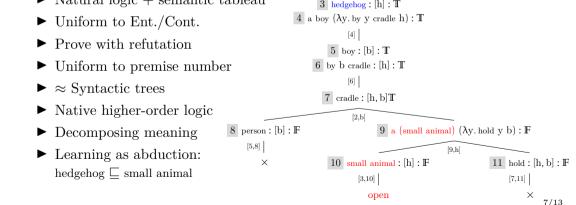
[3,10]open 11 hold: [h, b]: F

7/13

[7,11]

4 a boy ( $\lambda y$ . by y cradle h):  $\mathbb{T}$ Uniform to Ent./Cont. Prove with refutation 5 boy: [b]: T 6 by b cradle: [h]: T ► Uniform to premise number [6]  $\triangleright \approx \text{Syntactic trees}$ 7 cradle: [h, b]T Native higher-order logic [2,b] 8 person : [b] : **F** 9 a (small animal) (λy. hold y b): F Decomposing meaning [5,8][9,h]10 small animal: [h]: F ×





IDES/FILLAGENETHERLANDS.PRGesults

Conclusion

Overview

### Syntactic $\lambda$ -terms to $\lambda$ -logical forms

$$\begin{array}{lll} \operatorname{play}^{\diamondsuit \operatorname{su}} \operatorname{np} - \circ \operatorname{s_{\operatorname{main}}} \left( \operatorname{su}^{\diamondsuit} \left( \left( \operatorname{det}_{\square} \operatorname{a}^{\square \operatorname{det}(\mathbf{n} - \circ \mathbf{np})} \right) \operatorname{boy}^{\mathbf{n}} \right) \right) & \longrightarrow \operatorname{play}^{\mathbf{np}, \mathbf{s}} \left( \operatorname{a}^{\mathbf{n}, \mathbf{np}} \operatorname{boy}^{\mathbf{n}} \right) \\ \operatorname{large}^{\operatorname{np}, \operatorname{np}} \left( \operatorname{brown}^{\operatorname{np}, \operatorname{np}} \left( \operatorname{a}^{\mathbf{n}, \operatorname{np}} \operatorname{dog}^{\mathbf{n}} \right) \right) & \longrightarrow \operatorname{a}^{\mathbf{n}, \operatorname{np}} \left( \operatorname{large}^{\mathbf{n}, \mathbf{n}} \left( \operatorname{brown}^{\mathbf{n}, \mathbf{n}} \operatorname{dog}^{\mathbf{n}} \right) \right) \\ \operatorname{and} \left( \lambda \operatorname{x.} \operatorname{brown} (\operatorname{x} \operatorname{dog}) \right) \left( \lambda \operatorname{y.} \operatorname{black} (\operatorname{y} \operatorname{dog}) \right) \operatorname{no} & \longrightarrow \operatorname{and}^{\operatorname{np}, \operatorname{np}, \operatorname{np}} \left( \operatorname{no} \left( \operatorname{brown} \operatorname{dog} \right) \right) \left( \operatorname{no} \left( \operatorname{black} \operatorname{dog} \right) \right) \end{array}$$

- ▶ Map POS tags and shift to slightly Generalized POS tags: UPOS & Penn
- ► Use only these syntactic categories: n, np., s<sub>x</sub>, pp, pr
- ightharpoonup Function words  $\mapsto$  canonical terms (excl. prepositions)

Conclusion

### Syntactic $\lambda$ -terms to $\lambda$ -logical forms

```
\operatorname{play}^{\diamond \operatorname{su} \operatorname{np} - \circ \operatorname{s}_{\operatorname{main}}} \left( \operatorname{su}^{\diamond} \left( \left( \operatorname{det}_{\square} \ \operatorname{a}^{\square^{\operatorname{det}} (\operatorname{n} - \circ \operatorname{np})} \right) \operatorname{boy}^{\operatorname{n}} \right) \right)
                                                                                                                                                                                                                                                                          \rightsquigarrow play<sup>np,s</sup> (a<sup>n,np</sup> boy<sup>n</sup>)
large^{np,np} (brown<sup>np,np</sup> (a<sup>n,np</sup> dog<sup>n</sup>))
                                                                                                                                                                                                                            \rightsquigarrow a^{n,np} \left( large^{n,n} \left( brown^{n,n} dog^{n} \right) \right)
```

Natural Tableau

- ▶ Map POS tags and shift to slightly Generalized POS tags: UPOS & Penn
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```
\begin{array}{lll} \operatorname{play}^{\diamondsuit \operatorname{su}\operatorname{np} \multimap \operatorname{s}_{\operatorname{main}}} \left( \operatorname{su}^{\diamondsuit} \left( \left( \operatorname{det}_{\square} \operatorname{a}^{\square \operatorname{det}(\operatorname{n} \multimap \operatorname{np})} \right) \operatorname{boy}^{\operatorname{n}} \right) \right) & \longrightarrow \operatorname{play}^{\operatorname{np},\operatorname{s}} \left( \operatorname{a}^{\operatorname{n},\operatorname{np}} \operatorname{boy}^{\operatorname{n}} \right) \\ \operatorname{large}^{\operatorname{np},\operatorname{np}} \left( \operatorname{brown}^{\operatorname{np},\operatorname{np}} \left( \operatorname{a}^{\operatorname{n},\operatorname{np}} \operatorname{dog}^{\operatorname{n}} \right) \right) & \longrightarrow \operatorname{a}^{\operatorname{n},\operatorname{np}} \left( \operatorname{large}^{\operatorname{n},\operatorname{n}} \left( \operatorname{brown}^{\operatorname{n},\operatorname{n}} \operatorname{dog}^{\operatorname{n}} \right) \right) \\ \operatorname{and} \left( \lambda \operatorname{x.} \operatorname{brown}(\operatorname{x} \operatorname{dog}) \right) \left( \lambda \operatorname{y.} \operatorname{black}(\operatorname{y} \operatorname{dog}) \right) \operatorname{no} \\ & \longrightarrow \operatorname{and}^{\operatorname{np},\operatorname{np},\operatorname{np}} \left( \operatorname{no} \left( \operatorname{brown} \operatorname{dog} \right) \right) \left( \operatorname{no} \left( \operatorname{black} \operatorname{dog} \right) \right) \\ \operatorname{cut}^{\operatorname{pp},\operatorname{n},\operatorname{np},\operatorname{s}} \left( \operatorname{in}^{\operatorname{np},\operatorname{pp}} \operatorname{slice}_{\operatorname{n}} \right) \operatorname{meat}_{\operatorname{n}} & \longrightarrow \operatorname{cut}^{\operatorname{pp},\operatorname{np},\operatorname{np},\operatorname{s}} \left( \operatorname{in}^{\operatorname{np},\operatorname{pp}} \left( \operatorname{slice}^{\operatorname{n},\operatorname{np}} \right) \right) \\ \operatorname{cut}^{\operatorname{pp},\operatorname{np},\operatorname{np},\operatorname{s}} \left( \operatorname{in}^{\operatorname{np},\operatorname{pp}} \operatorname{slice}_{\operatorname{n}} \right) \operatorname{meat}_{\operatorname{n}} & \longrightarrow \operatorname{cut}^{\operatorname{pp},\operatorname{np},\operatorname{np},\operatorname{s}} \left( \operatorname{in}^{\operatorname{np},\operatorname{pp}} \left( \operatorname{slice}^{\operatorname{n},\operatorname{np}} \right) \right) \\ \operatorname{cut}^{\operatorname{pp},\operatorname{np},\operatorname{np},\operatorname{s}} \left( \operatorname{in}^{\operatorname{np},\operatorname{pp}} \left( \operatorname{slice}^{\operatorname{n},\operatorname{np}} \right) \right) \\ \operatorname{and} \left( \operatorname{a}^{\operatorname{np},\operatorname{np},\operatorname{np}} \left( \operatorname{no} \left( \operatorname{brown} \operatorname{np} \right) \right) \right) \\ \operatorname{cut}^{\operatorname{pp},\operatorname{np},\operatorname{np},\operatorname{s}} \left( \operatorname{in}^{\operatorname{np},\operatorname{pp}} \left( \operatorname{slice} \right) \right) \\ \operatorname{cut}^{\operatorname{np},\operatorname{np},\operatorname{np}} \left( \operatorname{np} \left( \operatorname{np} \right) \right) \\ \operatorname{cut}^{\operatorname{np},\operatorname{np},\operatorname{np},\operatorname{s}} \left( \operatorname{np} \left( \operatorname{np} \right) \right) \\ \operatorname{np}^{\operatorname{np},\operatorname{np},\operatorname{np},\operatorname{np},\operatorname{s}} \left( \operatorname{np}^{\operatorname{np},\operatorname{np},\operatorname{np},\operatorname{np},\operatorname{s}} \left( \operatorname{np}^{\operatorname{np},\operatorname{np},\operatorname{np},\operatorname{np},\operatorname{s}} \right) \right) \\ 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```

- ▶ Map POS tags and shift to slightly Generalized POS tags: UPOS & Penn
- ► Use only these syntactic categories: n, np., sx, pp, pr
- ightharpoonup Function words  $\mapsto$  canonical terms (excl. prepositions)

### Syntactic $\lambda$ -terms to $\lambda$ -logical forms

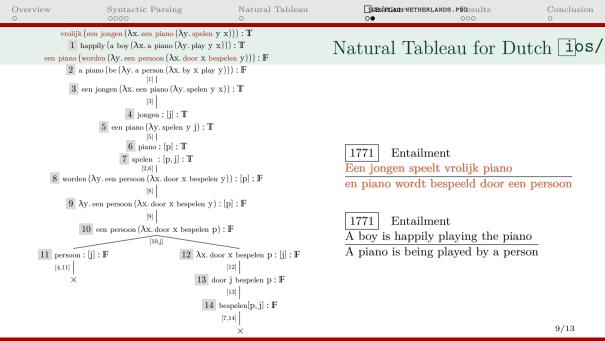
```
\begin{array}{lll} \operatorname{play}^{\diamondsuit \operatorname{su} \operatorname{np} \multimap \operatorname{s}_{\operatorname{main}}} \left( \operatorname{su}^{\diamondsuit} \left( \left( \operatorname{det}_{\square} \operatorname{a}^{\square \operatorname{det} (\operatorname{n} \multimap \operatorname{np})} \right) \operatorname{boy}^{\operatorname{n}} \right) \right) & \longrightarrow \operatorname{play}^{\operatorname{np}, \operatorname{s}} \left( \operatorname{a}^{\operatorname{n}, \operatorname{np}} \operatorname{boy}^{\operatorname{n}} \right) \\ \operatorname{large}^{\operatorname{np}, \operatorname{np}} \left( \operatorname{brown}^{\operatorname{np}, \operatorname{np}} \left( \operatorname{a}^{\operatorname{n}, \operatorname{np}} \operatorname{dog}^{\operatorname{n}} \right) \right) & \longrightarrow \operatorname{a}^{\operatorname{n}, \operatorname{np}} \left( \operatorname{large}^{\operatorname{n}, \operatorname{n}} \left( \operatorname{brown}^{\operatorname{n}, \operatorname{n}} \operatorname{dog}^{\operatorname{n}} \right) \right) \\ \operatorname{and} \left( \lambda \operatorname{x.} \operatorname{brown} (\operatorname{x} \operatorname{dog}) \right) \left( \lambda \operatorname{y.} \operatorname{black} (\operatorname{y} \operatorname{dog}) \right) \operatorname{no} \\ & \longrightarrow \operatorname{and}^{\operatorname{np}, \operatorname{np}, \operatorname{np}} \left( \operatorname{no} \left( \operatorname{brown} \operatorname{dog} \right) \right) \left( \operatorname{no} \left( \operatorname{black} \operatorname{dog} \right) \right) \\ \operatorname{cut}^{\operatorname{pp}, \operatorname{n}, \operatorname{np}, \operatorname{s}} \left( \operatorname{in}^{\operatorname{n}, \operatorname{pp}} \operatorname{slice}_{\operatorname{n}} \right) \operatorname{meat}_{\operatorname{n}} & \longrightarrow \operatorname{cut}^{\operatorname{pp}, \operatorname{np}, \operatorname{np}, \operatorname{s}} \left( \operatorname{in}^{\operatorname{np}, \operatorname{pp}} \left\{ \operatorname{slice}^{\operatorname{n}} \right\}^{\operatorname{np}} \right) \left\{ \operatorname{meat}^{\operatorname{n}} \right\}^{\operatorname{np}} \\ \end{array}
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IDES/FILLAGENETHERLANDS.PRGesults

```
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```
Parser / pos

Alpino / Alpino
Alpino / spaCy
NPN / Alpino
NPN / spaCy

LangPro 2×2
```

- ▶ POS tagging: spaCy is 1.7< better than Alpino:
- ▶ Parsers: Alpino > NPN? Sentences parsed 98.6% vs 94.9%;
- ▶ The ensemble is  $1.2 \le$  better than the best monoLP;
- The number of false proofs increases by 4% after αbductive learning.

Parser / pos	$\mathrm{T} \varepsilon$	
Alpino / Alpino	72.7	
Alpino / spaCy	74.8	
NPN / Alpino	72.0	
NPN / spaCy	74.3	
LangPro 2×2	76.0	

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Parser / pos	$\mathrm{T}\epsilon$	$T\alpha:T\epsilon$	
Alpino / Alpino	72.7	(+9.3)	
Alpino / spaCy	74.8	(+9.5)	
NPN / Alpino	72.0	(+8.6)	
NPN / spaCy	74.3	(+9.1)	
$\overline{\mathrm{LangPro}2{\times}2}$	76.0	(+9.8)	

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Parser / pos	$\mathrm{T}\epsilon$	$T\alpha:T\varepsilon$	$\mathrm{E}\epsilon$	
Alpino / Alpino	72.7	(+9.3)	74.1	
Alpino / spaCy	74.8	(+9.5)	75.9	
NPN / Alpino			72.8	
NPN / spaCy	74.3	(+9.1)	75.0	
LangPro 2×2	76.0	(+9.8)	77.1	

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Conclusion

Parser / pos	$\mathrm{T}\varepsilon$	$T\alpha$ : $T\epsilon$	$\mathrm{E}\varepsilon$	$T\alpha$ : $E\epsilon$
Alpino / Alpino	72.7	(+9.3)	74.1	(+1.8) 75.9
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NPN / Alpino	72.0	(+8.6)	72.8	(+1.5) 74.3
NPN / spaCy	74.3	(+9.1)	75.0	(+1.4) 76.4
$\operatorname{LangPro} 2 \times 2$	76.0	(+9.8)	77.1	(+1.6) 78.7

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### Experiments & Results

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Models	All	Ent	Cont
$\overline{\mathrm{LangPro}2{\times}2}$	78.7	50.6	66.3
BERTje	82.0	86.2	86.7
mBERT	79.9	79.0	81.9
RobBert	81.7	76.9	85.3

Models	All ±Δ	Ent	Cont
$\overline{\mathrm{LangPro}2{\times}2}$	78.7	50.6	66.3
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Conclusion

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Problems failed by all three DL models but solved by LangPro:

1556 Entailment A man is carrying a tree	[175] Entailment A family is watching a little boy who is hitting a baseball
A man is carrying a plant	A boy is hitting a baseball

#### Comparison to the transformer-based NLI models

Models	All $\pm \Delta$	Ent $\pm \Delta$	Cont $\pm \Delta$
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RobBert	$81.7 \pm 0.9$	$76.9 \pm 6.4$	$85.3 \pm 1.1$

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v	v G
A man is carrying a tree A man is carrying a plant	Entailment A family is watching a little boy who is hitting a baseball A boy is hitting a baseball

[4092] Contradiction

The person is not drawing

A woman in a red dress is putting away an instrument

A man is drawing a picture

A woman in a red dress is playing an instrument

11/13

TOS/FLIAGT-NETHERLANDS.PResults

#### Findings: SICK NL & Open Dutch WN

SICK NL is more challenging that the original dataset due to MT:

► Transferred gold labels are not gold:

3181 Neutral?

A man is trekking in the woods

The man is not hiking in the woods

■ State of the man is an het wandelen in het bos

■ De man is niet aan het wandelen in het bos

► Extra reasoning due to translation shifts: drawing a picture → een tekening maakt | tekent een foto dirt bike race → crossmotorwedstrijd | crossmotorrace.

Lexical relations learned from the training set:

 $\begin{array}{ll} lopen \equiv rennen & halter/dumbbell \sqsubseteq gewicht/weight \\ pizza \sqsubseteq voedsel/food & leeg/empty | vol/full \end{array}$ 

TIONS/PHIAGENETHERLANDS.PIRESUITS

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#### Conclusion & Future work

- ► First Dutch NLI system based on logic
- ► YES! Logic-based Reasoning Works for Dutch: promising results
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#### Conclusion & Future work

Overview

- ► First Dutch NLI system based on logic
- ► YES! Logic-based Reasoning Works for Dutch: promising results
- ► Automatic translation makes the NLI data more challenging

- ► Employ the Dutch CCG parser
- ► Qualitative comparison of the Dutch syntactic parsers
- ▶ Multilingual LangPro: SICK ioiof/f/files gollsardiforms

Conclusion