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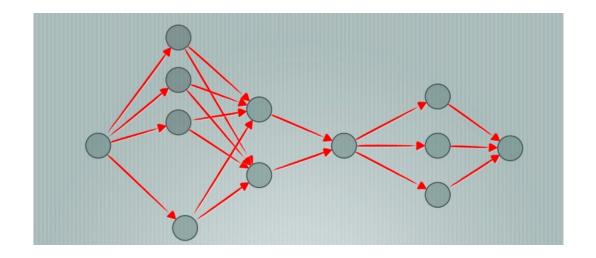
Κωνσταντίνος Αλεξόπουλος ΣΗΜΜΥ ΕΜΠ CSLab

MOTIVATION

- HPC, Multi-node & Heterogeneous Systems
- Communication with low latency
- Reliability
- Provide an instant employment method
- The CERN use-case

ZEROMQ

- Messaging Library
- Does not employ a messaging broker
 - Low Latency
- Asynchronous I/O
- Easy deployment of complex topologies



ZEROMQ

- Supports many platforms and language bindings
- Open source project with active development
- IPC, UDP, TCP/IP
- Port to an RDMA-enabled interconnect

COMMUNICATION PARADIGMS

TCP Socket Semantics

- High Overhead
- Decent Bandwidth
- Low implementation effort

RDMA Semantics

- Low Overhead
- Very Low Latency
- Very High Throughput
- High implementation effort

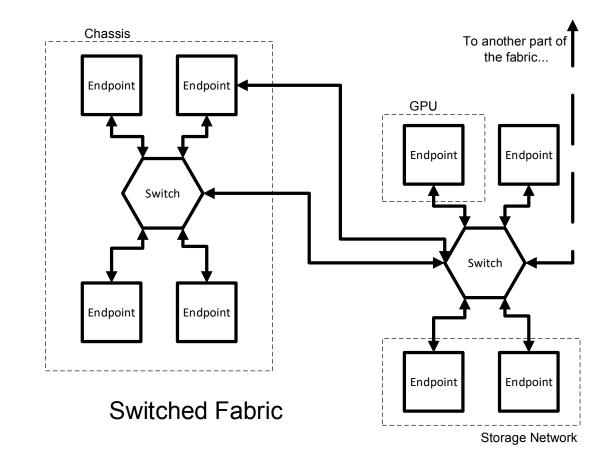
SWITCHED FABRICS

Switched Fabric Architectures

- Allow for any topology
- Reliable
- "Inside-the-box" & "Outside-the-box"

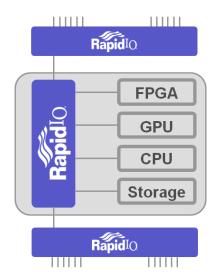
Shared Bus Architectures

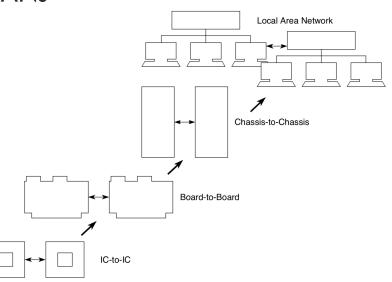
- Topology restrictions
- Bottlenecks
- "Inside-the-box" only



RAPIDIO

- System-level interconnect originally
- Independent from Physical Implementation
- Lately oriented towards chassis-to-chassis and SANs
- Protocol stack processed in HW
- Destination Based Routing





RIO OPERATIONS — MESSAGING INTERFACE

Channelized Messages

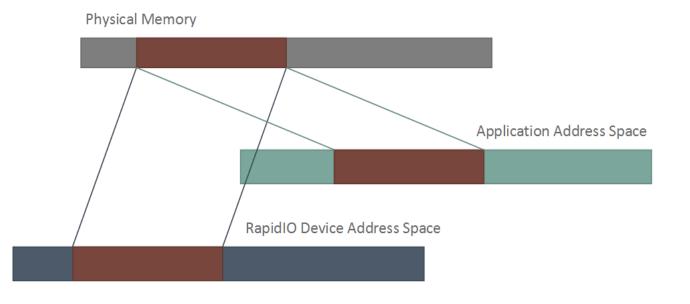
- Maximum 4K
- Socket-like interface
- Sent to a channel

Doorbells

- Hardware Signals
- 8B software-defined payload
- Have to allocate exclusive range to receive doorbells

RIO OPERATIONS — MMIO

- Remote Direct Memory Access
 - Read/Write
- Zero-copy
- "One-way communication"
- Device memory mapped to physical memory
 - Needs to be done at boot time
 - Kernel boot parameter
- Physical memory mapped to process address space
 - Done through a library call in the application layer
- Supports multi-cast

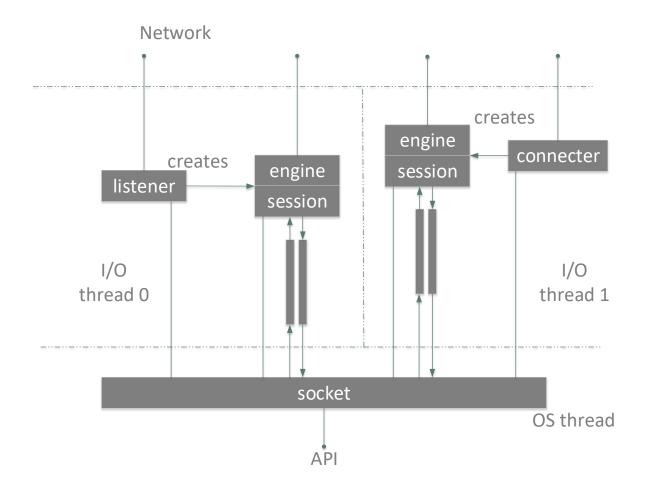


DESIGN & IMPLEMENTATION

RIOZMQ

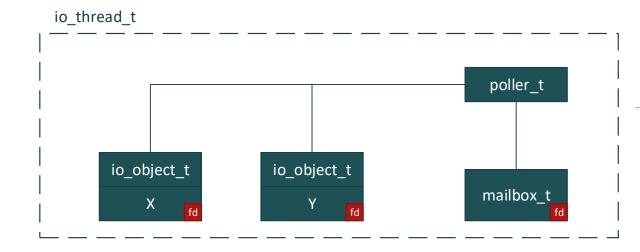
ZEROMQ INTERNAL ARCHITECTURE

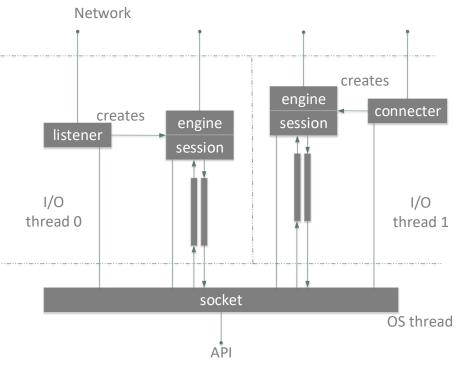
- User creates (Omq!) Socket
- Socket binds/connects
- Listener/Connecter create Session
- Session/Engine object for any new connection



ZEROMQ INTERNAL ARCHITECTURE

- •"I/O threads" for asynchronous operations
- in_event() and out_event() for every io_object_t



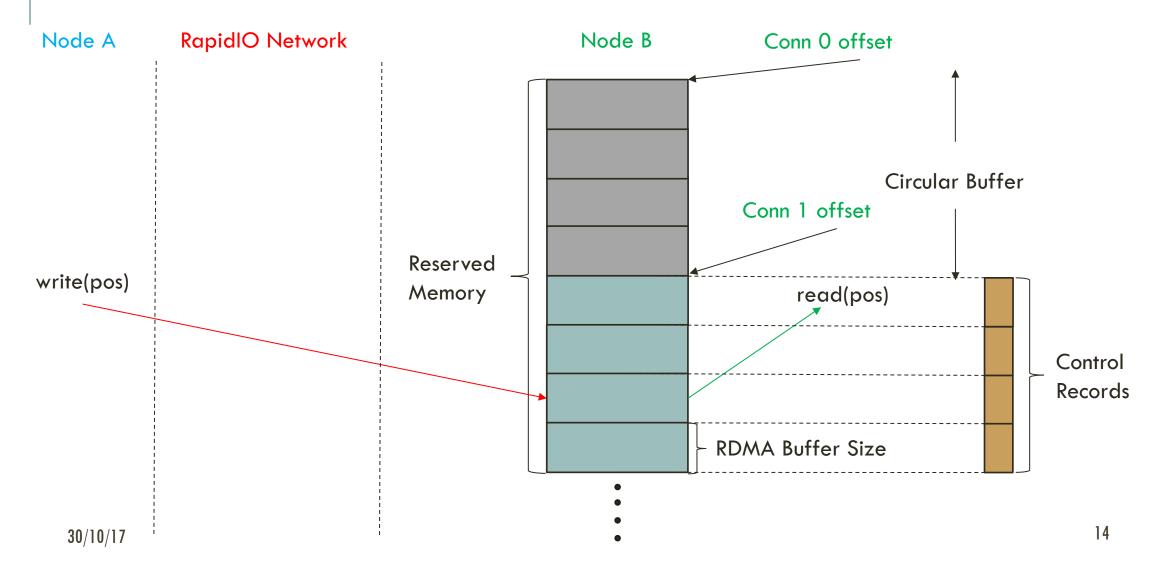


RIOZMQ EXTENSION

- rio_address
 - resolves address of rio://[destID]:[channel]format
- o rio_connecter / rio_listener
 - o connect()/bind()
 - RDMA target addresses exchange
 - doorbell range allotment

- orio_engine
 - RDMA write/recv
- o rio_mailbox
 - doorbell operations
 - FD registered with io_object_t
- glue code in socket/session

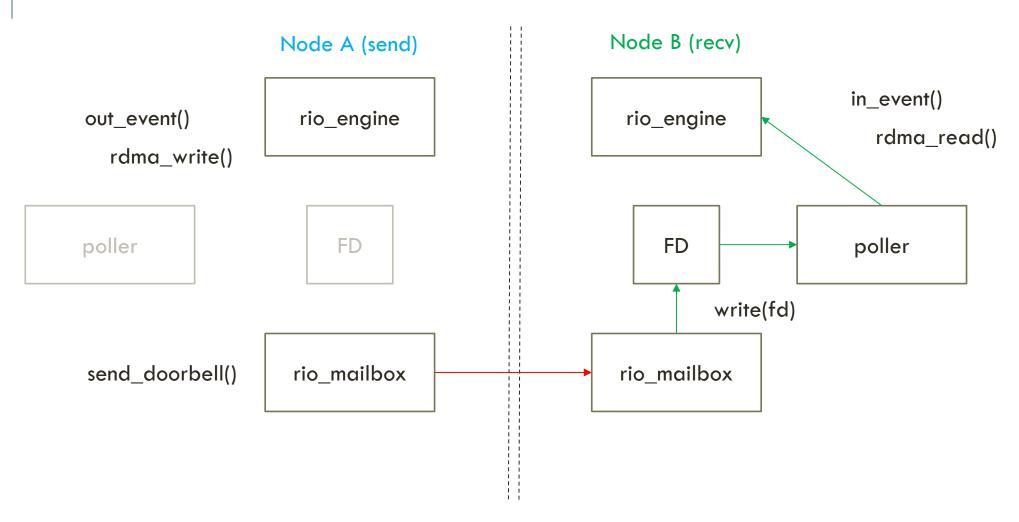
MEMORY SCHEME



DOORBELLS AS NOTIFICATIONS

```
Node A (send)
                                                  Node B (recv)
check_write_idx(pos)
update_write_idx(pos)
write(pos)
send_doorbell(pos)
                                             update_read_idx(pos)
                                              check_read_idx(pos)
                                              update_read_idx(pos)
                                              read(pos)
                                              send_doorbell(pos)
update_write_idx(pos)
```

DOORBELLS - FILE DESCRIPTORS - POLLER



EVALUATION

HARDWARE SETUP

- 4x 2U Quad Units
- 4 Nodes per Unit
- Intel Xeon L5640 @ 2.27Ghz
- 48GB of DDR3 1333MHz RAM
- IDT Tsi721 RapidIO to PCle bridge cards
- QSFP+ cables
- 38-port Top of Rack (ToR) RapidIO Gen2 switch
- CERN CentOS



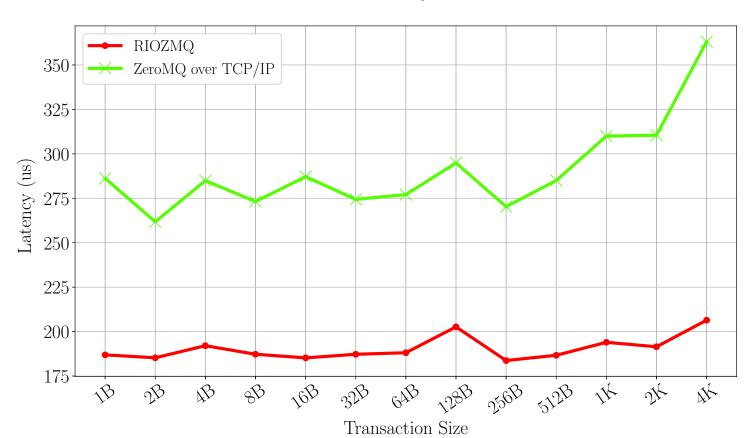


BENCHMARKS

- Standard ZeroMQ Benchmarks
- Measures round trip time (RTT) between 2 nodes

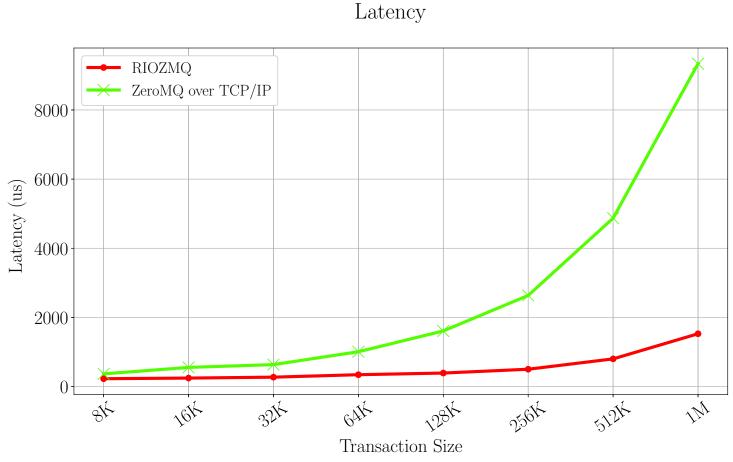
EVALUATION — LATENCY SMALL





- RDMA Buffer Size : 4K
- O Circular Buffer Length: 1
- RIOZMQ 65% faster

EVALUATION — LATENCY LARGE

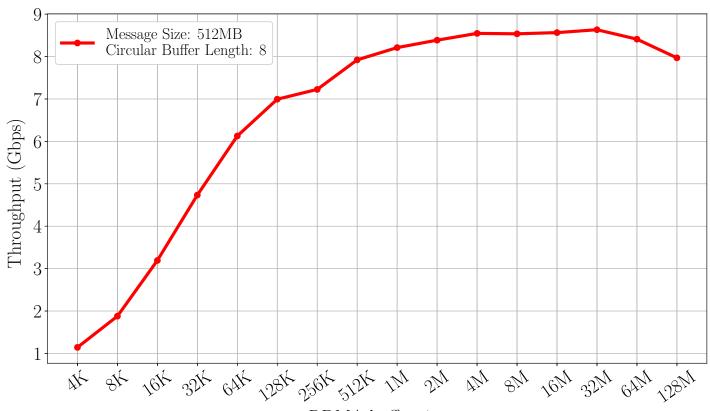


O RDMA Buffer Size: 1M

Circular Buffer Length: 1

EVALUATION — RDMA BUFFER SIZE SCAN

RDMA Buffer Scan

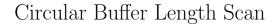


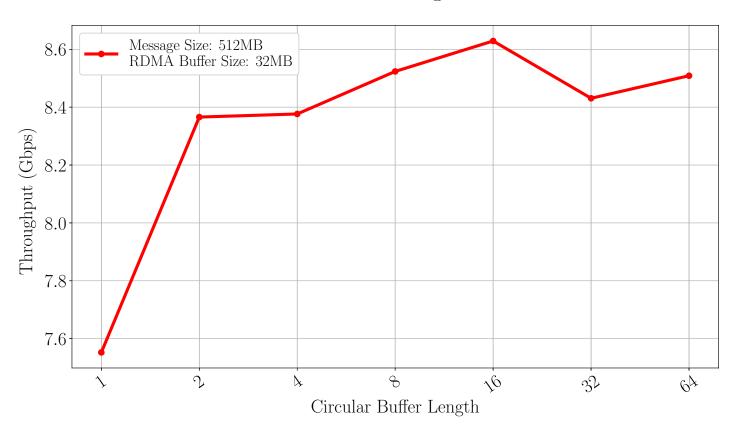
An **RDMA Buffer** is the smallest possible data size to transmit

The **Circular Buffer** consists of the number of RDMA-enabled blocks assigned to a connection

RDMA buffer size

EVALUATION — CIRCULAR BUFFER LENGTH SCAN

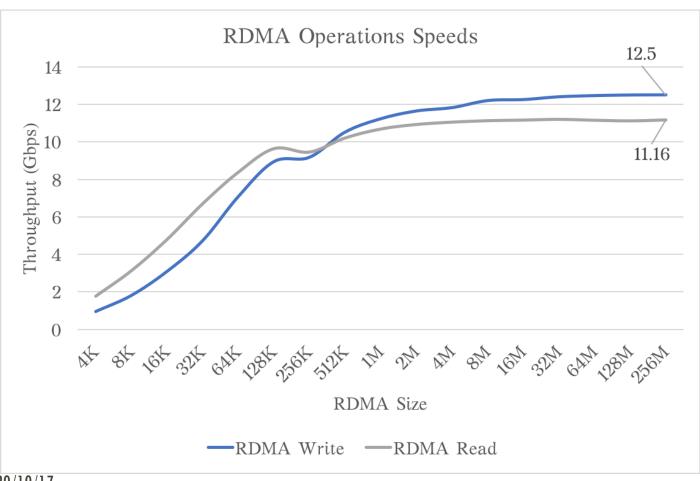




An **RDMA Buffer** is the smallest possible data size to transmit

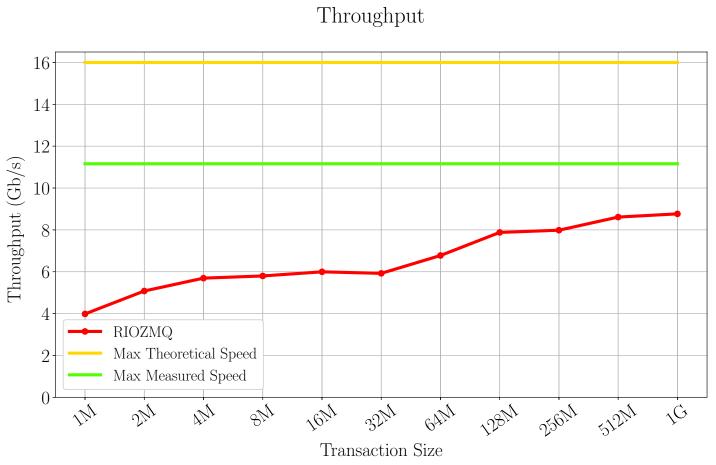
The **Circular Buffer** consists of the number of RDMA-enabled blocks assigned to a connection

EVALUATION — RDMA SPEEDS



- Maximum Measured Speeds
- Around rdma_write() call
- RapidIO <--> PCleTranslations
- Library Overhead

EVALUATION — THROUGHPUT

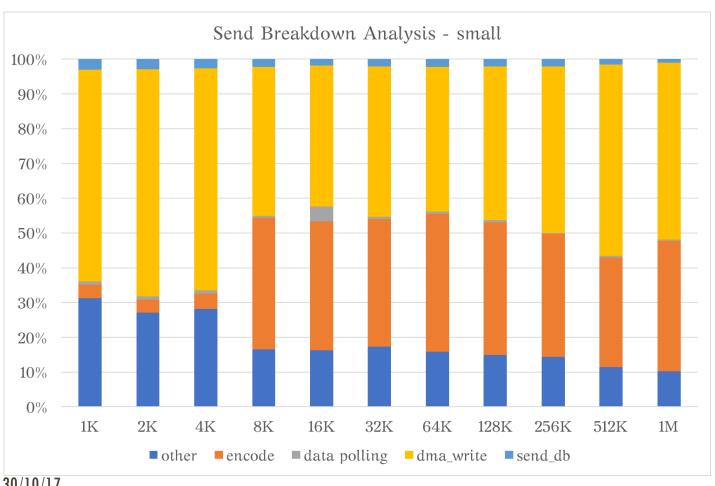


RDMA Buffer Size : 32M

O Circular Buffer Length: 16

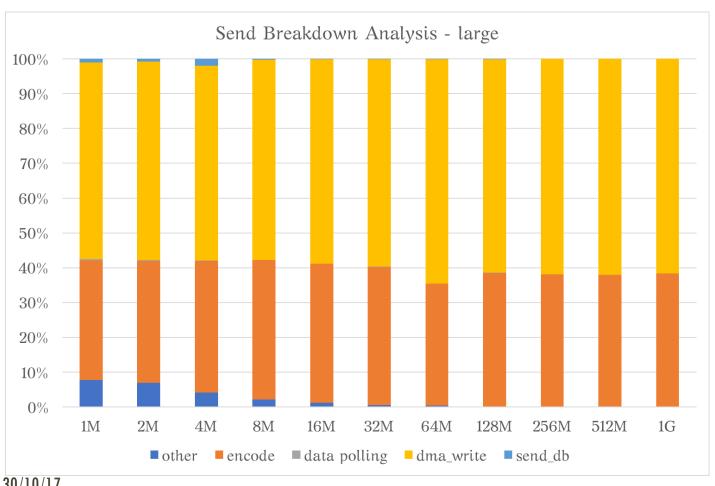
Achieved ~75% saturation

EVALUATION — BREAKDOWN ANALYSIS SEND **SMALL**



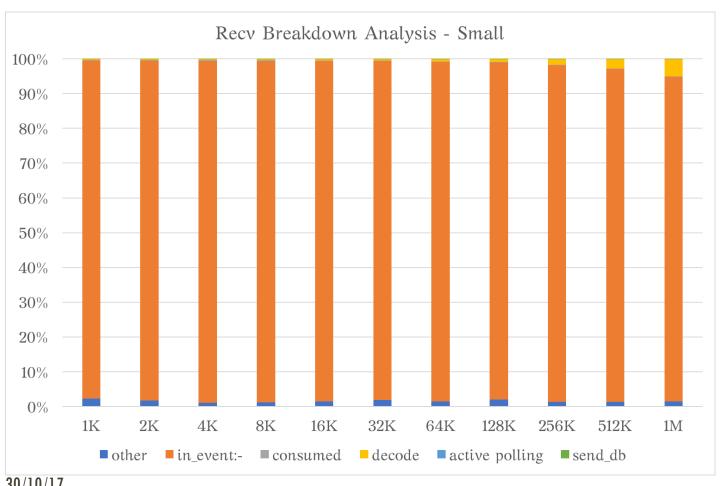
- RDMA Buffer Size: 32M
- Circular Buffer Length: 16
- Main bottleneck dma_write
- Encode also consumes time

EVALUATION — BREAKDOWN ANALYSIS SEND LARGE



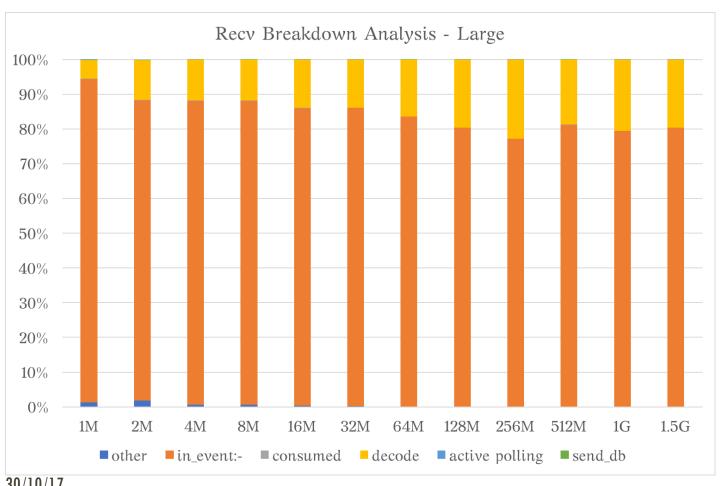
- RDMA Buffer Size: 32M
- Circular Buffer Length: 16
- Main bottleneck dma_write
- Encode also consumes time

EVALUATION — BREAKDOWN ANALYSIS RECV **SMALL**



- RDMA Buffer Size: 32M
- Circular Buffer Length: 16
- in event: between calls
- poll operations
- fd operations
- doorbell handling
- Can't measure asynchronous operations

EVALUATION — BREAKDOWN ANALYSIS RECV LARGE



- RDMA Buffer Size: 32M
- Circular Buffer Length: 16
- in event: between calls
- poll operations
- fd operations
- doorbell handling
- Can't measure asynchronous operations
- Decode for larger transactions

CONCLUSIONS

- ZeroMQ extended to use the RapidIO transport
- Achieved better latency for small messages compared to TCP/IP
- Designed for use with arbitrary number of nodes
- Use in existing setups by changing the address from tcp* to rio*
- Used RDMA semantics within ZeroMQ
- Same scheme for other RDMA-enabled interconnects
- Work is open-source can be found at <u>github.com/kostorr/libzmq</u> (soon...)

FUTURE WORK

Implementation

- Optimize circular buffer
- Zero-Copy in the critical path
- Possible removal of file descriptor use

Evaluation

- More extensive breakdown on recv() performance
- Employ on a system with more than 16 nodes
- Run a real-life benchmark on a distributed system

THANK YOU!

THE PROBLEM (1)

Moore's Law

Semiconductor performance increases at an exponential rate

The conjunction of these laws leads to an imbalance, limiting performance

Amdahl's Law – Law of diminishing returns

The performance of a system can only be assessed as the balance between:

- CPU
- Memory Bandwidth
- I/O Performance

New Interconnects Technologies!