

CSE 230: Data Structures

Lecture 4-2: Application of Stacks

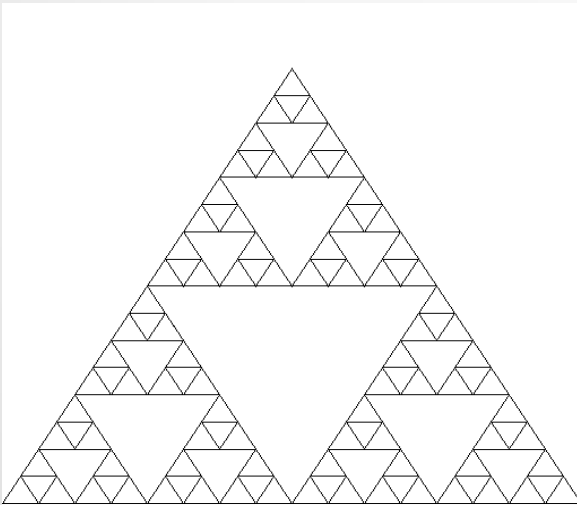
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Application of Stacks

- Recursion
- Tower of Hanoi
- Evaluation of Expressions

Recursion

- Concept of defining a function that calls itself as a sub-routine
 - Allows us to take advantage of the repeated structure in many problems
 - e.g finding the factorial of a number



<http://introcs.cs.princeton.edu/java/23recursion/images/sierpinski5.png>

CSE 201: Data Structures and Algorithms



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<http://www.spektyr.com/Gallery/Recursion.jpg>

Linear Recursion

- Function is defined so that it makes at most one recursive call at each time it is invoked
- Useful when an algorithmic problem
 - Can be viewed in terms of a first or last element, plus a remaining set with same structure

```
{factorial 6}  
(* 6 {factorial 5})  
(* 6 (* 5 {factorial 4}))  
(* 6 (* 5 (* 4 {factorial 3})))  
(* 6 (* 5 (* 4 (* 3 {factorial 2}))))  
(* 6 (* 5 (* 4 (* 3 (* 2 {factorial 1})))))  
(* 6 (* 5 (* 4 (* 3 (* 2 1)))))  
(* 6 (* 5 (* 4 (* 3 2))))  
(* 6 (* 5 (* 4 6)))  
(* 6 (* 5 24))  
(* 6 120)  
720
```

Src:mitpress.mit.edu

Algorithm LinearSum(A, n):

Input: Integer array A and element n

Output: Sum of first n elements of A

if $n=1$ **then**

return $A[0]$

else

return LinearSum($A, n-1$) + $A[n-1]$

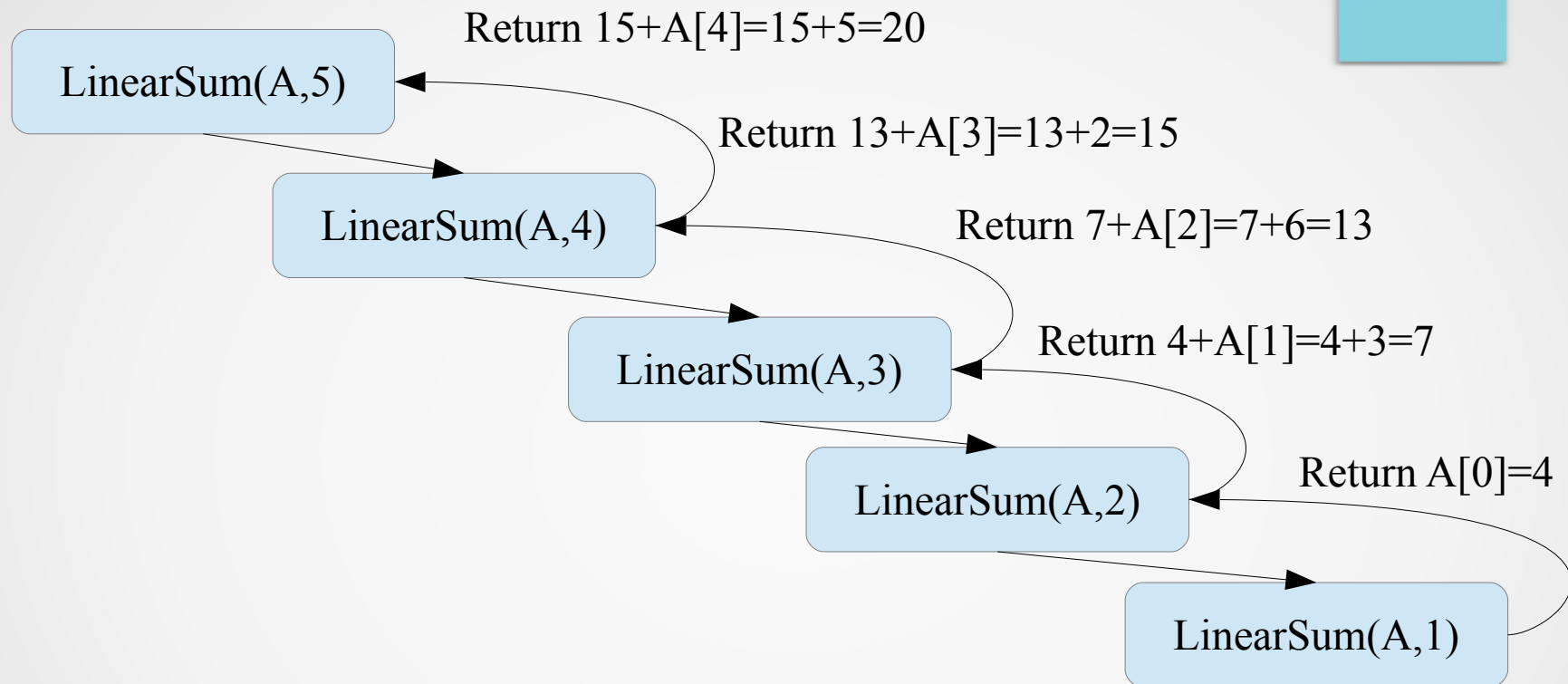
Linear Recursion

- Algorithm using linear recursion uses the following
 - Test for base cases
 - Base cases defined so that every possible chain of recursive call reaches base case
 - Helps in termination
 - Recurse
 - May decide on one of multiple recursive calls to make
 - Recursive calls must progress to base case

Analyzing Recursive Algorithms

- Use visual tool – recursion trace
 - Contains box for each recursive call
 - Contains parameters of the call
 - Links showing the return value
- Use recurrence relation
 - Mathematical formulation to capture the recursion process

Analyzing Recursive Algorithms



- Example `LinearSum`
 - Running time is $O(n)$
 - Space Complexity: $O(n)$

Stack Trace

- Example stack trace

LinearSum(A,1)	Return $A[0]=4$
LinearSum(A,1)+A[1]	Return $4+A[1]=4+3=7$
LinearSum(A,2)+A[2]	Return $7+A[2]=7+6=13$
LinearSum(A,3)+A[3]	
LinearSum(A,4)+A[4]	

Problem 1

- Reverse an array using linear recursion

- Solution

- **Algorithm** ReverseArray(A, i, n):

- Input:*** Integer array A and integers i, n

- Output:*** Reversal of n integers in A starting from i

- if** $n \leq 1$ **then**

- return**

- else**

- Swap $A[i]$ and $A[i+n-1]$

- Call ReverseArray($A, i+1, n-2$)

- return**

Problem 2:

- Computing Powers via Linear Recursion

$$power(x, n) = \begin{cases} 1 & \text{if } n=0 \\ x \cdot power(x, n-1) & \text{otherwise} \end{cases}$$

$$power(x, n) = \begin{cases} 1 & \text{if } n=0 \\ x \cdot power(x, (n-1)/2)^2 & \text{if } n>0 \text{ is odd} \\ power(x, n/2)^2 & \text{if } n>0 \text{ is even} \end{cases}$$

Problem 3

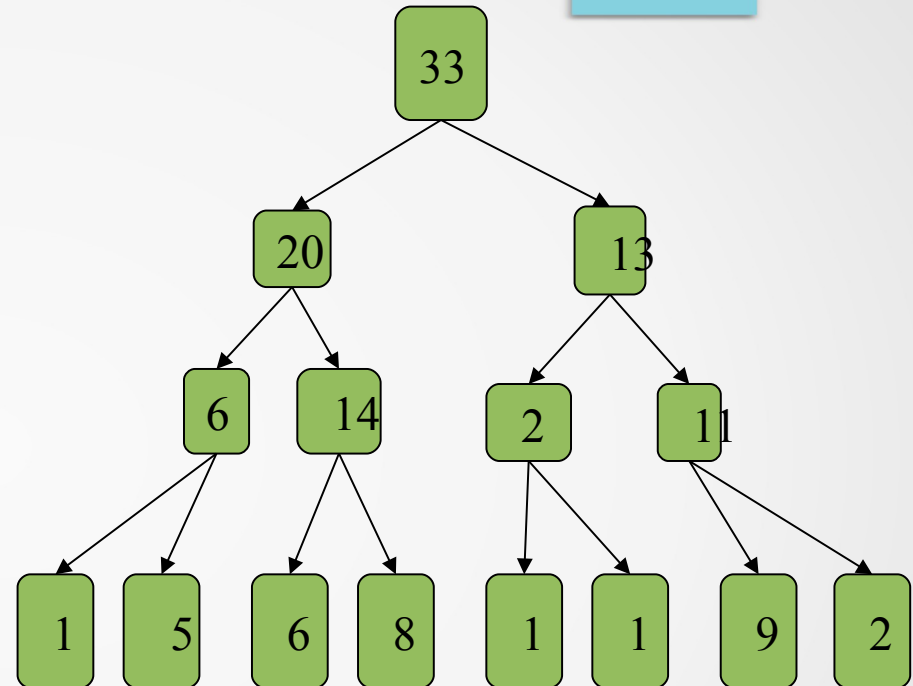
- Describe a linear recursive algorithm for finding the minimum element in an n-element array
- Write a function using recursion to print numbers from n to 0.
- Given n write a recursive function to find the factorial of n
- Given a non-negative integer n, return the sum of its digits
- Given a string, compute recursively a new string where all the adjacent chars are now separated by a “*”. e.g abc- → a*b*c

Higher-Order Recursion

- Uses more than one recursion call
 - e.g 3-way merge sort
- Binary recursion
 - Two recursion calls
 - e.g BinarySum
 - **Algorithm** BinarySum(A,i,n):
Input: Integer array A and integers i, n
Output: Sum of first n elements of A starting at index i
if n=1 then
 return A[i]
else
 return BinarySum(A,i,[n/2])+BinarySum(A,i+[n/2],[n/2])

Analysis

- Recursion trace is a tree
- Depth of recursion
 - $O(\log n)$
 - Lesser additional space needed
- Running time
 - $O(n)$
 - Have to visit every element in the array



Problem

- Generate the kth Fibonacci number using Binary Recursion
- Solution (This method not recommended !! Exponential complexity)

– **Algorithm** BinaryFib(k):

Input: k

Output: k^{th} Fibonacci Number

if $k \leq 1$ **then**

return k

else

return BinaryFib($k-1$)+BinaryFib($k-2$)

Problem

- Describe a binary recursive method for searching an element x in an n -element unsorted array A .
 - Compute the running time and space complexity of your algorithm

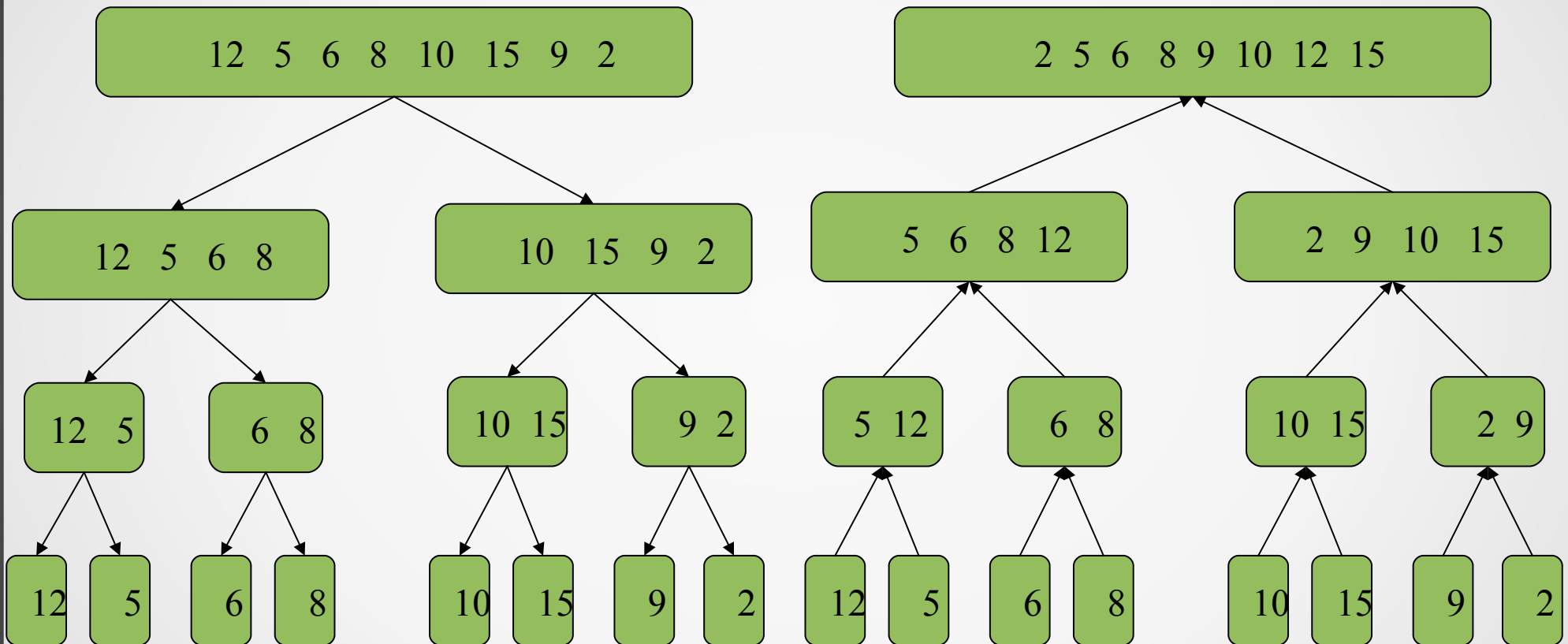
Merge Sort

- Uses divide and conquer strategy (multiple recursion) to sort a set of numbers
- Divide:
 - If S has zero or one element, return S
 - Divide S into two sequences S_1 and S_2 each containing half of the elements of S
- Recur
 - Recursively apply merge sort to S_1 and S_2
- Conquer
 - Merge S_1 and S_2 into a sorted sequence

Merging of Sorted Sequences

- Iteratively remove smallest element from one of the two sequences S_1 and S_2 and add it to end of output sequence S

Merge Sort



Merge Sort

- mergesort(S, low, high) {
 if (low < high) {
 middle = (low+high)/2;
 mergesort(s,low,middle);
 mergesort(s,middle+1,high);
 merge(s, low, middle, high);
 }
}

Src: Skiena, Algorithm Design Manual, Chapter 4