

The Lattice-Theoretic Essence of Property Directed Reachability Analysis

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PDR is ...

Knaster
- Tarski



Kleene

Property Directed Reachability (PDR)

Model checking Algorithm

To verify / refute safety problems. of state transition systems.

- Known for its performance

- hardware verification [Bradley, VMCAI'11]

- used in Spacer [Hader+, CAV'11] as part of Z3 [MouraB, TACAS'08]

- many variants:

1. IC3/PDR [Bradley, VMCAI'11], [Een+, FMCAD'11] :

original one.

2. Reverse PDR, fbPDR [Seufert & Scholl, DATE'18, '19] :

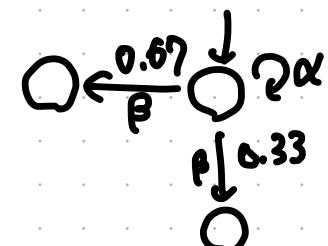
Search direction is reversed

3. PrIC3 [Barat, CAV'20] :

probabilistic extension.

) Model checking
MDPs.

) Model checking
Kripke structures



Our Contribution: Mathematical Abstraction of PDR

①

Lattice-Theoretical abstraction

Our algorithm

LT-PDR

for $\mu F \leq d$ in L

Essence: Combination of
Knaster-Tarski & Kleene

②

Instantiate the theory

IC3/PDR

Instance 1

Instance 2

Instance 3

for Kripke structures



for Markov Decision Processes

for Markov Reward Models

Outline

1. LT-PDR — lattice-theoretic essences

- Positive & Negative
- LT-PDR — let these help each other

2. Instances of LT-PDR for quantitative model checking

3. Categorical classification of PDR variants

4. From mathematical abstraction To programming abstraction.

5. Experiments

Overview of LT-PDR

Positive Side

Thm. (Knaster-Tarski)
Prefixed points form a complete lattice.

$\mu F \leq \alpha$
 \Leftrightarrow There exists X s.t. $Fx \leq X \leq \alpha$

$X_0 \leq X_1 \leq \dots \leq X_{n-2} \leq X_{n-1}$
s.t. $Fx_i \leq x_{i+1}$ and $x_{n-1} \leq \alpha$

verifies

$\mu F \leq ? \alpha$

Negative Side

Thm. (Kleene)
 $\mu F = \bigvee_{\text{near}} F^n \perp$

$\mu F \not\leq \alpha$
 \Leftrightarrow There exists n . s.t. $F^n \perp \not\leq \alpha$

(C_1, \dots, C_{n-1})
s.t. $C_{i+1} \leq F C_i$. $C_n \not\leq \alpha$

refutes

help
each other



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Positive LT-PDR

- ... aims to verify $\mu F \leq \alpha$.
- comes from KT thm.

Thm. (Knaster - Tarski')

$$\mu F \leq \alpha \Leftrightarrow \text{There exists } X \text{ s.t. } FX \leq X \leq \alpha$$

invariant

How to search for X ?

- (Naively) Because $FX \leq X \Rightarrow \forall i. F^i \perp \leq X$,

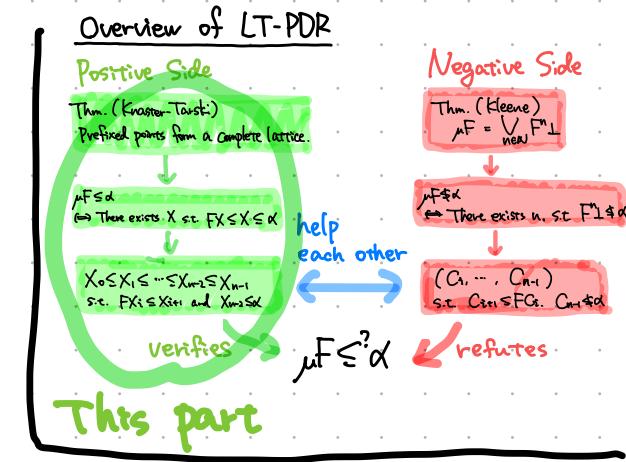
construct $\perp \leq F\perp \leq F^2\perp \leq \dots$ and check $F^{i+1}\perp \leq F^i\perp \leq \alpha$.
↑ Small candidates of X .

- (Positive LT-PDR)

• Construct $X_0 \leq X_1 \leq X_2 \leq \dots (\leq \alpha)$ and check $FX_i \leq X_{i+1}$.

↑ preemptively inflate $\perp \leq F\perp \leq \dots$. It accelerates the search

• Preemptively inflate and shrink if necessary.



It's difficult
to find this ...

stabilized

easier to find!

Negative LT-PDR

- Aims to refute $\mu F \leq \alpha$.
- Comes from Kleene thm.

Thm. (Kleene')

$\mu F \not\leq \alpha \Leftrightarrow$ There exists n s.t. $F^n \perp \not\leq \alpha$.

\Leftrightarrow There exists n, C s.t. $C \leq F^n \perp$ and $C \not\leq \alpha$.

How to search for C ? (in Negative LT-PDR)

1. Pick $C_n \not\leq \alpha$. unsafe

2. Construct C_0, C_1, \dots, C_n

s.t. $C_i \leq F^i \perp \Rightarrow C_{i+1} \leq F^{i+1} \perp \Rightarrow \dots \Rightarrow C_n \leq F^n \perp$

3. If $C_0 = \perp$ then C_n is the one!

(Let's hope this is the case ...)

Overview of LT-PDR

Positive Side

Thm. (Knaster-Tarski)
Prefixed points form a complete lattice.

$\mu F \text{sd}$

\Leftrightarrow There exists X s.t. $FX \leq X \leq \alpha$

$X_0 \leq X_1 \leq \dots \leq X_{n-1} \leq X_n$

s.t. $FX_i \leq X_{i+1}$ and $X_n = \alpha$

verifies $\mu F \leq \alpha$

Negative Side

Thm. (Kleene)
 $\mu F = \bigvee F^n \perp$ new

$\mu F \text{sk}$

\Leftrightarrow There exists n s.t. $F^n \not\leq \alpha$

(C_0, \dots, C_n)
s.t. $C_n \not\leq \alpha$, $C_{i+1} \leq F C_i$

refutes $\mu F \leq \alpha$

This part

Counter example

{

$C \leq F^n \perp$ and $C \not\leq \alpha$

try to prove
 $C_n \leq F^n \perp$
 (reachable)
 by backtracking

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LT-PDR

Combination of Positive & Negative ones,
helping each other.

Overview of LT-PDR

Positive Side

Thm. (Knaster-Tarski)
Prefixed points form a complete lattice.

$\mu F \leq d$
 \Leftrightarrow There exists X s.t. $FX \leq X \leq d$

$X_0 \leq X_1 \leq \dots \leq X_{n-1} \leq d$
s.t. $FX_i \leq X_{i+1}$ and $X_{n-1} \leq d$

verifies

$\mu F \leq ? \leq d$ refutes

Negative Side

Thm. (Kleene)
 $\mu F = \bigvee_{n \in \omega} F^n \perp$

$\mu F \leq x$
 \Leftrightarrow There exists n , s.t. $F^n \leq x$

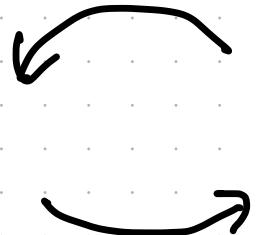
(C_0, \dots, C_{n-1})
s.t. $C_{n-1} \leq F C_n$, $C_n \leq d$

This part

hints for shrinking overly-inflated X

Positive Side

$$X_0 \leq \dots \leq X_{n-1}$$



Negative Side

$$(C_0, \dots, C_{n-1})$$

tells $C_0 = \perp$ will never happen

by a lattice-theoretic proposition:

Proposition 3.17. Let $C = (C_0, \dots, C_{n-1})$ be a Kleene sequence ($2 \leq n, 0 < i \leq n-1$) and $X = (X_0 \leq \dots \leq X_{n-1})$ be a KT sequence. Then

1. $C_i \not\leq X_i$ implies that C cannot be extended to a conclusive one, that is, there does not exist C_0, \dots, C_{i-1} such that (C_0, \dots, C_{n-1}) is conclusive.
2. $C_i \not\leq FX_{i-1}$ implies that C cannot be extended to a conclusive one.
3. There is no conclusive Kleene sequence with length $n-1$. □

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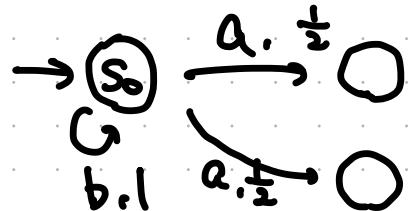
Instances of LT-PDR

- Instance for Markov Decision Processes

$$L := [0,1]^S$$

$$F(d: \delta \rightarrow [0,1])$$

$$= \left(S \mapsto \begin{cases} \max_a \sum_{s'} d(s') \cdot \delta(s)(a)(s') & \text{if } s \in a \\ 1 & \text{otherwise} \end{cases} \right)$$



Similar to PrIC3 [Batz et al., CAV '20]

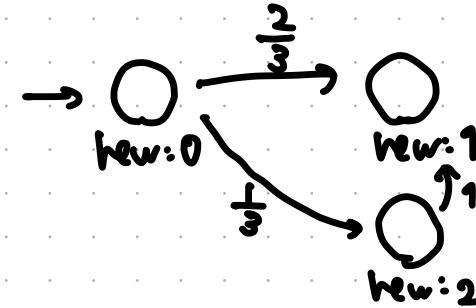
PrIC3 needs additional check before returning False, but our instance does not.

- Instance for Markov Reward Models.

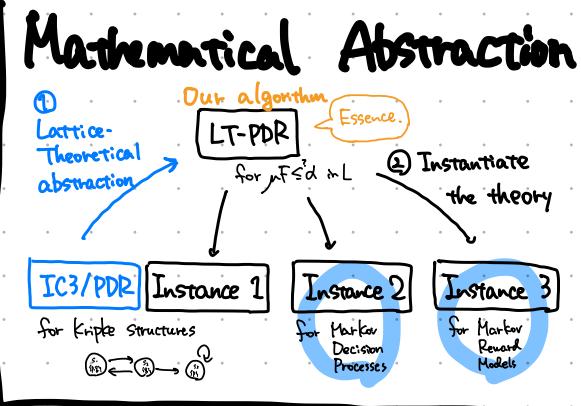
$$L := [0,1]^S$$

$$F(d: \delta \rightarrow [0,1])$$

$$= \left(S \mapsto \begin{cases} \text{rew}(s) + \sum_{s' \in S} d(s') \cdot \delta(s,s') & \text{if } s \in a \\ 0 & \text{otherwise} \end{cases} \right)$$



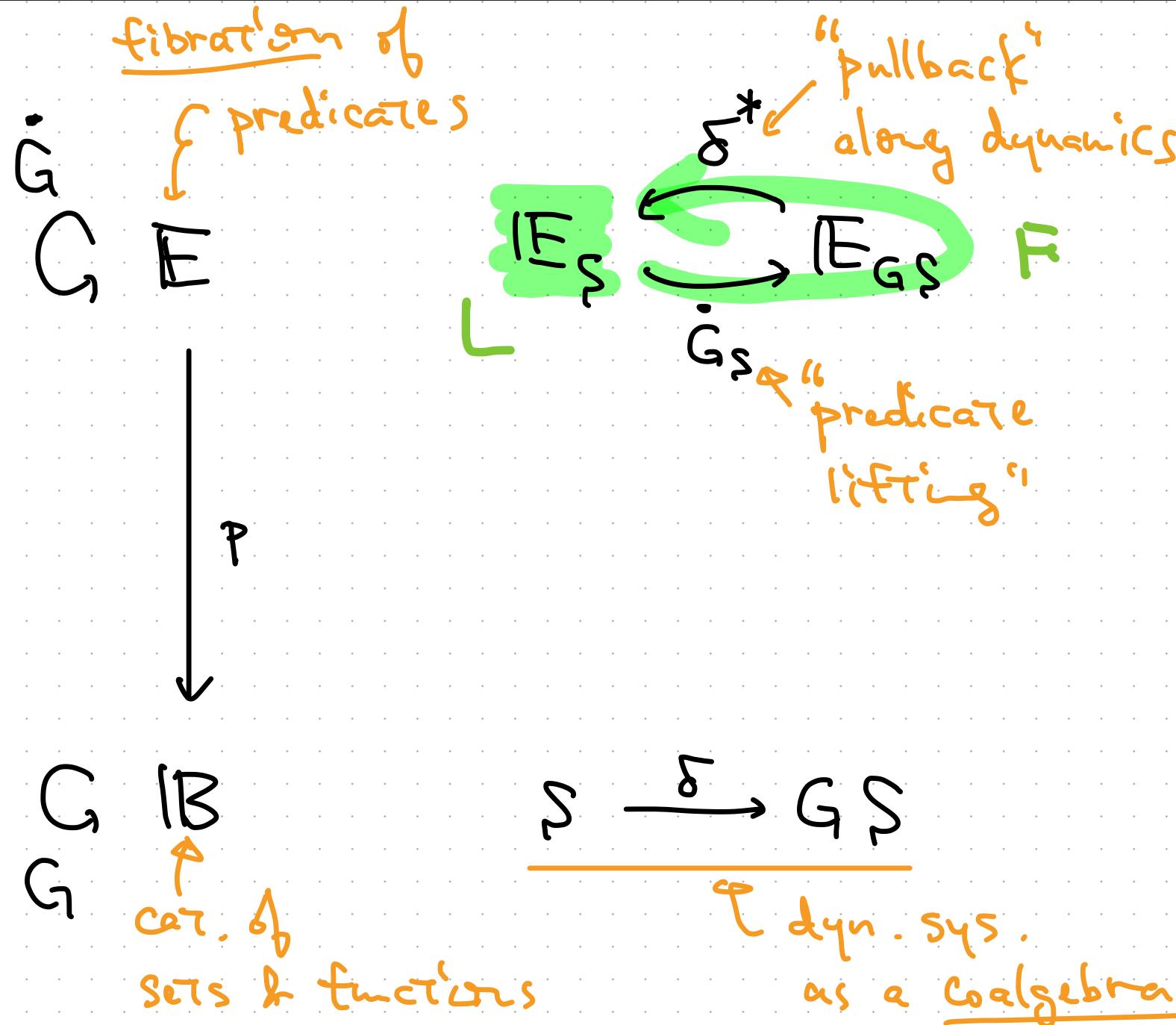
New PDR variant for MRM.



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LT-PDR for Dynamical Systems



A Categorical framework

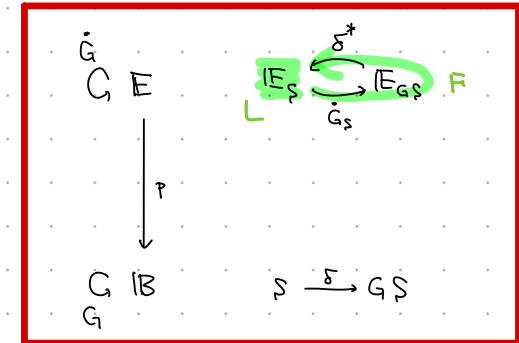
- Coalgebras as dyn. sys.
[Rutten, Jacobs, ...]
- Fibrations for predicates
[Bénabou, Jacobs, ...]

[Bénabou, Jacobs, ...]

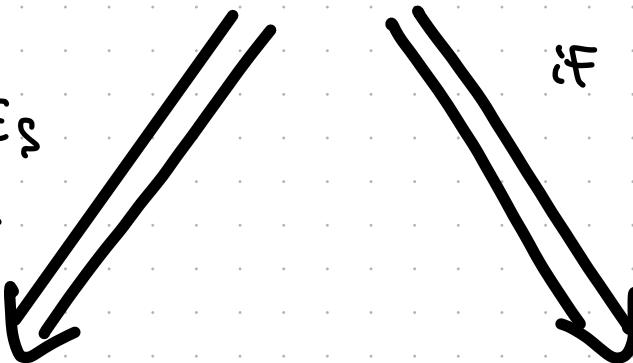
Unifying View of PDR Variants

Bwd PDR

$$C \leq \cdot \cup \alpha \wedge \delta^* G \models ?$$



if $E_S \xrightarrow{P} E_S$
involution



$$E_S \xrightarrow{\begin{array}{l} GS \\ T \\ \exists H_S \end{array}} E_{GS}$$

$$\mu \lambda. (\lambda \alpha)^V$$

$$(\lambda \delta^* G_S \circ \lambda \alpha)$$

$$\leq ? \quad \lambda L$$

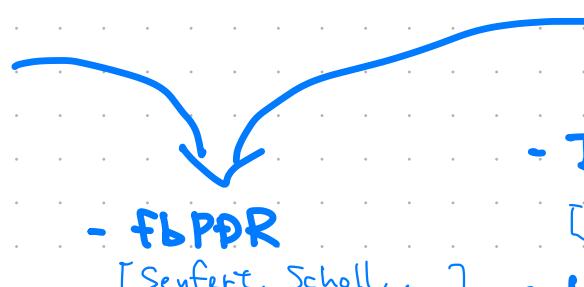
Inv.-Bwd PDR

- Reverse PDR
[Seufert, Scholl, ...]
- P+IC3
[Betz, Junges, ...]

$$\mu \lambda. L^V$$

$$H_S \delta^* \chi \leq ? \alpha$$

Fwd PDR



- IC3/PDR
[Bradley, Ein, ...]
- fBPPDR
[Seufert, Scholl, ...]
- Unlikely for MDPs

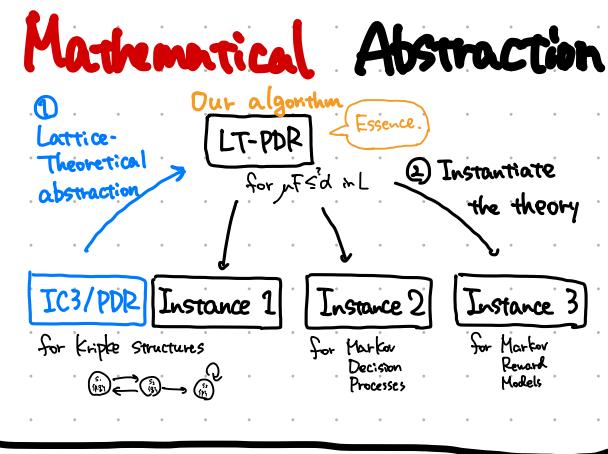
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Programming Abstraction

Mathematical abstraction yields it naturally.

We can get instances of LT-PDR by small efforts.



Lattice-Theoretical abstraction

type classes allows generic code

Generic Haskell code of LT-PDR

≡ Pseudo code of LT-PDR

Instantiate the theory

~900 lines
(C++)

~130 lines

~50 lines

~80 lines

~80 lines

Instance 2
PDR^{MDP}

Instance 3
PDR^{MRM}

IC3/PDR

Instance 1
PDR^{KF}

for Kripke structures

for Markov Decision Processes

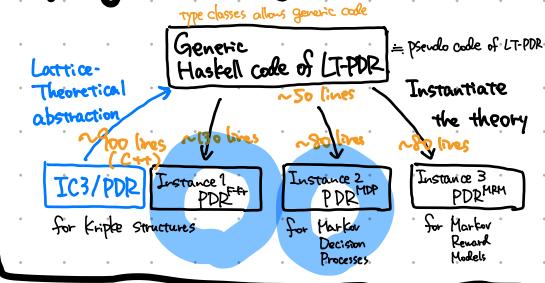
for Markov Reward Models

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Experiments

Programming Abstraction



target system	ours	To compare w/	machine
Kripke structure	PDR ^{F-Kr}	IC3 ref [Bradley, VMCAI'11]	Apple M1 chip with 16GB memory
MDPs	PDR ^{MDP}	Pr IC3 [Barzil, CAV'20]	1.2GHz Quad-Core Intel Core i7 with 10GB memory using Docker.

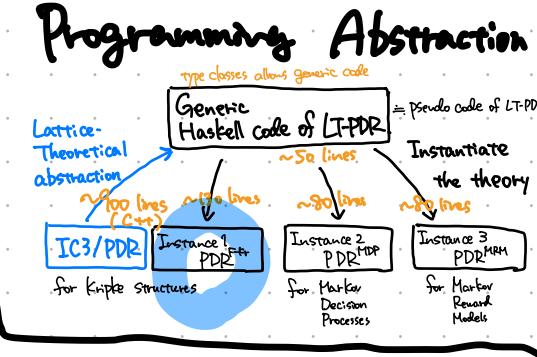
Observations:

Performance is not too bad,
especially given that the coding effort is minimal

Results

Comparison to IC3ref

(<https://github.com/arbrad/IC3ref>)



Benchmark	$ S $	Result	PDR ^{F-Kr}	IC3ref
latch0.smv	2^3	True	317 μ s	270 μ s
counter.smv	2^5	False	1.620 s	3.27 ms
power2bit8.smv	2^{15}	True	1.516 s	4.13 ms
ndista128.smv	2^{17}	True	TO	73.1 ms
shift1add256.smv	2^{21}	True	TO	174 ms

Benchmarks:

HwMC'15 competition

+ our own (latch0, counter)

TO: 600 sec.

IC3ref uses 23.

PDR^{Fkr} uses toysolver

(<https://github.com/msakai/toysolver>)

for Haskell compatibility

IC3ref outperforms our PDR^{F-k_r}.

Potential improvement: use generalization techniques used in IC3ref.

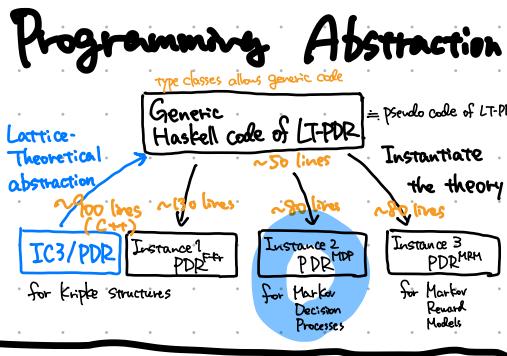
use another SAT-Solver.

Results

Comparison to PrIC3

[Batz+, CAV'20]

Benchmark	$ S $	GT pr.	λ	PDR ^{IB-MDP}	PrIC3			
					none	lin.	pol.	hyb.
Grid	10^2	$1.2E^{-3}$	0.3	0.31	1.31	19.34	—	—
			0.2	0.48	1.75	24.62	—	—
Grid	10^3	$4.4E^{-10}$	0.3	122.29	—	—	—	—
			0.2	136.46	—	—	—	—
BRP	10^3	0.035	0.1	—	—	—	—	—
			0.01	18.52	56.55	594.89	—	722.38
			0.005	1.36	11.68	238.09	—	—
ZeroConf	10^4	0.5	0.9	—	—	—	0.58	0.51
			0.75	—	—	—	0.55	0.46
			0.52	—	—	—	0.48	0.46
			0.45	<0.1	<0.1	<0.1	<0.1	<0.1
Chain	10^3	0.394	0.9	—	72.37	—	0.91	0.70
			0.4	—	80.83	—	0.93	—
			0.35	177.12	115.98	—	—	—
			0.3	88.27	66.89	557.68	—	—
DoubleChain	10^3	0.215	0.9	—	—	—	1.83	1.99
			0.3	—	—	—	1.88	1.96
			0.216	—	—	—	139.76	—
			0.15	7.46	—	—	—	—



Benchmarks are from [Batz+, CAV'20].

T0: 600 sec.

PrIC3 uses \mathbb{Z}^3

PDR^{MDP} uses GLPK

No clear comparison. Ours are often better with smaller benchmarks.

Potential improvement:

- express MDPs symbolically.
- use other (so-called) generalization (linear \Rightarrow polynomial / hybrid)

Conclusions

- LT-PDR.: lattice-theoretic abstraction of PDR.
 - Essence of PDR.: ingenious combination of FT & Kleene theorem.
 - Gives many instances for qualitative & quantitative systems.
- Generic implementation of LT-PDR in Haskell.
 - Easily-obtained instances have at least reasonable performance.

Future work

- Relation to Abstract Interpretation.
- Get instances for other systems (e.g. hybrid systems [Suenaga I, VMCAI'20])