

## General LPP - Simplex method

Simplex method is standard method for maximizing the linear function of several variables under the several constraints on other linear functions.

The simplex method was first used by George Dantzig in 1947. To understand the simplex method see Dantzig's simplex algorithm. The simplex algorithm is an iterative procedure carried systematically to determine the optimal solution of an LPP from a set of feasible solutions. The concept and algorithm is

introduced by T.S. Motzkin.

### Definition :

A set of values  $x_1, x_2, \dots, x_n$  which satisfies constraints of an LPP is called as solution.

### Feasible solution :

Any solution to LPP which satisfies non-negativity restriction of LPP is called its feasible solution.

### Optimal Feasible Solution :

Any feasible solution which optimizes the objective of the LPP is called optimal solution or optimum feasible solution.

### Definition :

If the constraints of general LPP be  $\sum a_{ij} x_{ij} \leq b_i$  for  $i=1, 2, \dots, K$ .

then the non-negativity variable  $s_i$  which are introduced to convert the inequalities to equalities called slack variables.

allowing for a margin of error.



Definition: If the constraints of LPP is

$$\sum_{j=1}^n a_{ij} b_{ij} \geq b_i \text{ for } i=1, 2, \dots, k$$

then the non-negative variable  $s_i$  which are introduced to convert the inequalities to equalities  $\sum_{j=1}^n a_{ij} b_{ij} - s_i = b_i$  is

called surplus variables.

## Two Different Forms of LPP

\* canonical form

\* standard form

canonical form is given by:

Is expressed as  $\text{Max } Z = c \times x$  (objective function)

$\text{s.t. } Ax \leq b$  (constraint)

&  $x \geq 0$  (not-negative restriction)

where ( $c \in C_1, C_2, \dots, C_m$ ) is given at II

$$\text{where } A = \begin{pmatrix} a_{11} & a_{12} & \dots & a_{1n} \\ a_{21} & a_{22} & \dots & a_{2n} \\ \vdots & \vdots & \ddots & \vdots \\ a_{m1} & a_{m2} & \dots & a_{mn} \end{pmatrix}, x = \begin{pmatrix} x_1 \\ x_2 \\ \vdots \\ x_n \end{pmatrix}, b = \begin{pmatrix} b_1 \\ b_2 \\ \vdots \\ b_m \end{pmatrix}$$

Characteristic of canonical form:

1. Objective function is of maximization type.
2. All the constraints are of less than or equal to type.

3. All variable  $x_i$  are non-negative

### Standard Form of LPP

Max  $Z = cx$  (objective function)

s.t  $Ax = b$  (constraints)

$$x \geq 0$$

### Characteristic of standard form

1. The objective function is of maximization type.

2. All the constraints are expressed as equation.

3. Right-hand side of each constraint are non-negative.

4. All the variables are non-negative.

(note: note: either  $\leq$  or  $\geq$  can be converted into  $=$ )

If the problem is minimization problem then it is equivalent to maximization of negative expression of that function.

$$\text{Min } Z = -\text{Max}(Z)$$

(or)

$$\text{Min } Z = -\text{Max } Z^*$$



Note 2: An inequality in one direction can be converted into an inequality to opposite direction by multiplying both sides by  $-1$ .

Note 3:  $f \geq g(x) + h(x)$

If the variable "is unconstraint or unrestricted it can always be expressed as the difference of two non-negative variables. (ie)  $x_2 = x_2' - x_2''$ , where

$x_2', x_2'' \geq 0$

Note 4:

when the slack & surplus variables are introduced in the constraints they should also appear in objective functions with zero co-efficient.

Express the foll./. LPP in canonical form

$$\text{Max } Z = 2x_1 + 3x_2 + x_3 = S \text{. W.M}$$

$$\text{s.t. } 4x_1 - 3x_2 + x_3 \leq 6$$

$$x_1 + 5x_2 - 7x_3 \geq -4$$

&  $x_1, x_2, x_3 \geq 0$ ,  $x_2$  is unrestricted.

$$0 = 2x_1 + 12x_2 + 6x_3 + 5x_4 - 7x_5$$



Sol: G.T.  $x_2$  is unrestricted S. T.  
 Matrix area is unbounded  
 $\therefore x_2 = x_2' - x_2''$ . Between 0 &  $\infty$   
 $4x_1 + 3(x_2' - x_2'') + x_3 \leq 6$  (i) Max Z  
 $-x_1 - 5x_2 + 7x_3 \leq 4$  (ii) Min Z

S.T:

$$4x_1 + 3x_2' + 3x_2'' + x_3 \leq 6 \text{ (i) } \min Z$$

$$-x_1 + 5x_2' + 5x_2'' + 7x_3 \leq 4 \text{ (ii) } \max Z$$

$$\text{and } x_1, x_2', x_2'', x_3 \geq 0.$$

2. Express the foll. LPP in standard form

$$\text{Max } Z = 4x_1 + 2x_2 + 6x_3$$

S.T: ~~Supply & demand~~ ~~Max Z~~  
~~Max Z~~  $2x_1 + 3x_2 + 2x_3 \geq 6$  Max Z  
~~Max Z~~  $3x_1 + 4x_2 = 8$  Max Z  
 $6x_1 - 4x_2 + x_3 \leq 10$  Max Z  
 $\& x_1, x_2, x_3 \geq 0.$

Sol: ~~Max Z~~ Max Z

$$\text{Max } Z = 4x_1 + 2x_2 + 6x_3$$

S.T: ~~Supply & demand~~ Max Z  
 $2x_1 + 3x_2 + 2x_3 - S_1 + OS_2 = 6$  Max Z

$$3x_1 + 4x_2 + OS_1 + OS_2 = 8$$

$$6x_1 - 4x_2 + x_3 + OS_1 + S_2 = 10$$

$$\& x_1, x_2, x_3, \geq 0.$$

$$S_1, S_2$$



Express the foll. LPP as standard form

$$\text{Min } Z = 4x_1 + 3x_2$$

subject to multiple of non-neg.

$$\text{s.t. } x_1 - x_2 \leq 10$$

$$3x_1 + x_2 \geq 18$$

$$5x_1 + 2x_2 \leq 20$$

$$x_1, x_2 \geq 0$$

Max Z is solution in given case

def:

$$\text{Min } Z = -\text{Max}(-Z)$$

$$\text{Max } Z^* = -4x_1 - 3x_2$$

with some slack variables also

s.t:

$$x_1 - x_2 + S_1 + OS_2 + OS_3 = 10$$

$$3x_1 + x_2 - S_2 + OS_1 = 18$$

$$5x_1 + 2x_2 + S_3 + OS_1 = 20$$

$$x_1, x_2, S_1, S_2, S_3 \geq 0$$

Matrix form of above problem

$$\text{Max } Z^* = Cx$$

where it is at basis in vertices also

$$\text{s.t. } Ax = b \quad \& \quad x \geq 0$$

where it is at basis in vertices also

where it is at basis in vertices also

$$A \begin{pmatrix} 1 & -1 & 1 & 0 & 0 \\ 3 & 1 & 0 & 1 & 0 \\ 5 & 2 & 0 & 0 & 1 \end{pmatrix} \begin{pmatrix} x_1 \\ x_2 \\ S_1 \\ S_2 \\ S_3 \end{pmatrix} = \begin{pmatrix} 10 \\ 18 \\ 20 \end{pmatrix}$$

using 3rd. basis in terms of vertices

$$C = (-4, -3, 0, 0, 0)$$



## Basic Solution:

Given a system of  $m$  linear equations with  $n$  variables ( $m < n$ ) solution obtained by setting  $n-m$  variables is equal to zero and solving for the remaining  $m$  variables is called basic solution. The  $m$  variables are called basic variables and they formed basic variables and their solution.  $m-n$  variables are called non-basic variables.

## Non-Degenerate Basic Solution:

A basic solution is said to be non-degenerate basic solution if the none of the basic variable is 0. A basic solution is said to be degenerate basic solution if one or more of basic variable are 0.

Feasible solution is also called basic and it is called basic feasible solution.