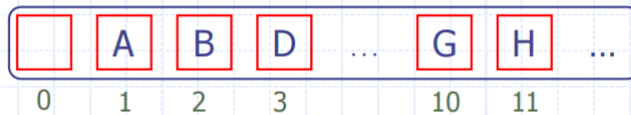


# Exercise on Trees

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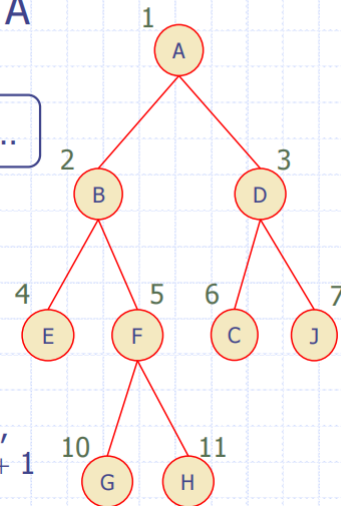
## Array-Based Representation of Binary Trees

- Nodes are stored in an array  $A$



- Node  $v$  is stored at  $A[\text{rank}(v)]$

- $\text{rank}(\text{root}) = 1$
- if node is the left child of  $\text{parent}(\text{node})$ ,  
 $\text{rank}(\text{node}) = 2 \cdot \text{rank}(\text{parent}(\text{node}))$
- if node is the right child of  $\text{parent}(\text{node})$ ,  
 $\text{rank}(\text{node}) = 2 \cdot \text{rank}(\text{parent}(\text{node})) + 1$



A binary tree is stored in an array  $A$  of  $n$  cells, using the array-based representation Ch.7-1, p.18. Now

(1) Show the maximum number of nodes  $A$  can store.

In array-based representation, nodes are stored at index  $i$ , with:

- left child at  $2i$
- right child at  $2i + 1$

This structure assumes a complete binary tree format. Therefore:

The maximum number of nodes is  $n$ , when the tree is completely filled and all array positions from index 1 to  $n$  are used.

(2) Show the minimum number of nodes  $A$  can store.

The minimum number of nodes occurs when the tree is highly unbalanced.

This forms a single chain using only right children, and the index for the next right child is always:

$$\text{index}_{k+1} = 2 \cdot \text{index}_k + 1$$

Define the number of such nodes as  $h$ . The last index used would be:

$$\text{index}_h = 2^h - 1$$

To ensure index does not exceed  $n$ , we must find the largest  $h$  such that:

$$2^h - 1 \leq n \Rightarrow h \leq \log_2(n + 1)$$

Thus, the minimum number of nodes is:

$$\lfloor \log_2(n + 1) \rfloor$$