



SMART
INDUSTRY
LABORATORY

Scheduling Algorithms (4)

- Production Planning Algorithm (1) -

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What is Production Planning?

Production Planning :

Evaluate the values of productivity, total profits and/or the other measurements, subject to the constraints of material quantities, facility abilities, man power, budget and/or the other factors.

Production Planning is to make a total balanced plan for given time period

Example of Production Planning

- Make two products called ProA and ProB
- Resources to make products, quantities of those for each product and the upper limits for those are given in the following table.

	ProA	ProB	Upper Limit
Material (kg)	4	11	440
Man Power (man-hour)	5	7	350
processing time (hour)	7	6	420

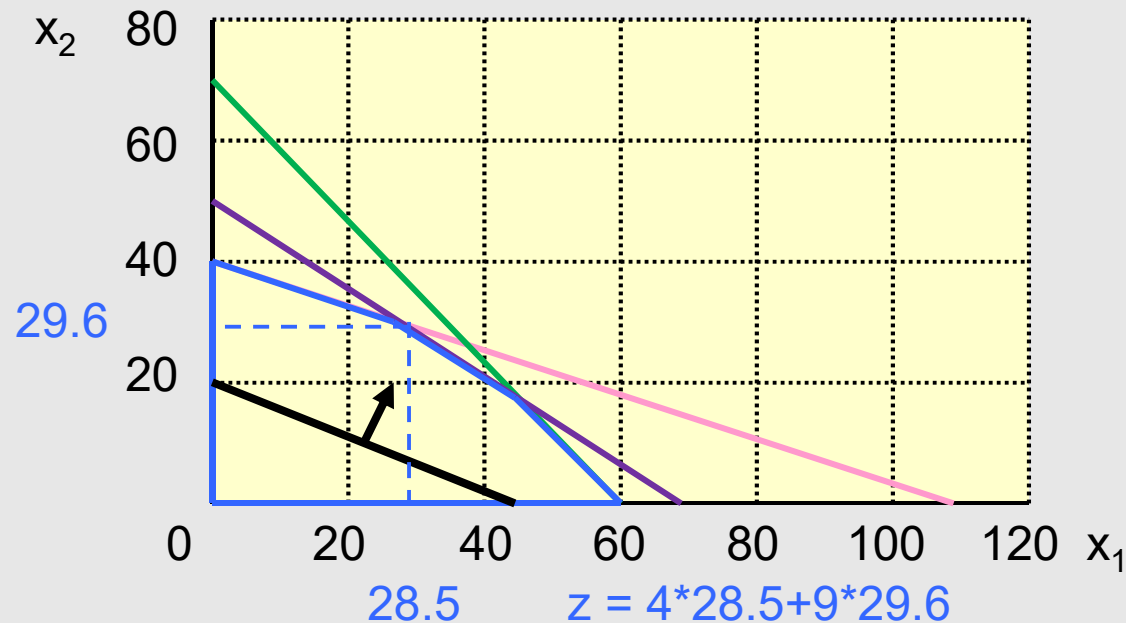
Profit for unit quantity of ProA is 40k¥.

Profit for unit quantity of ProB is 90k¥.

Decide the optimal quantities to get the largest total profit.

Solving with figure (2 dimensions)

Objective total benefit : $\max z = 4x_1 + 9x_2$
subject to material constraint : $4x_1 + 11x_2 \leq 440$ (a)
man power constraint : $5x_1 + 7x_2 \leq 350$ (b)
facility constraint : $7x_1 + 6x_2 \leq 420$ (c)
non-negative constraint : $x_1, x_2 \geq 0$



In the case of more than 3 variables ??

Terms for Linear Programming

Mathematical Programming:

Method to give the optimal solution that maximize or minimize the objective function, subject to constraints

Linear Programming:

One of Mathematical Programming

Objective function and constraints are given as linear real functions.

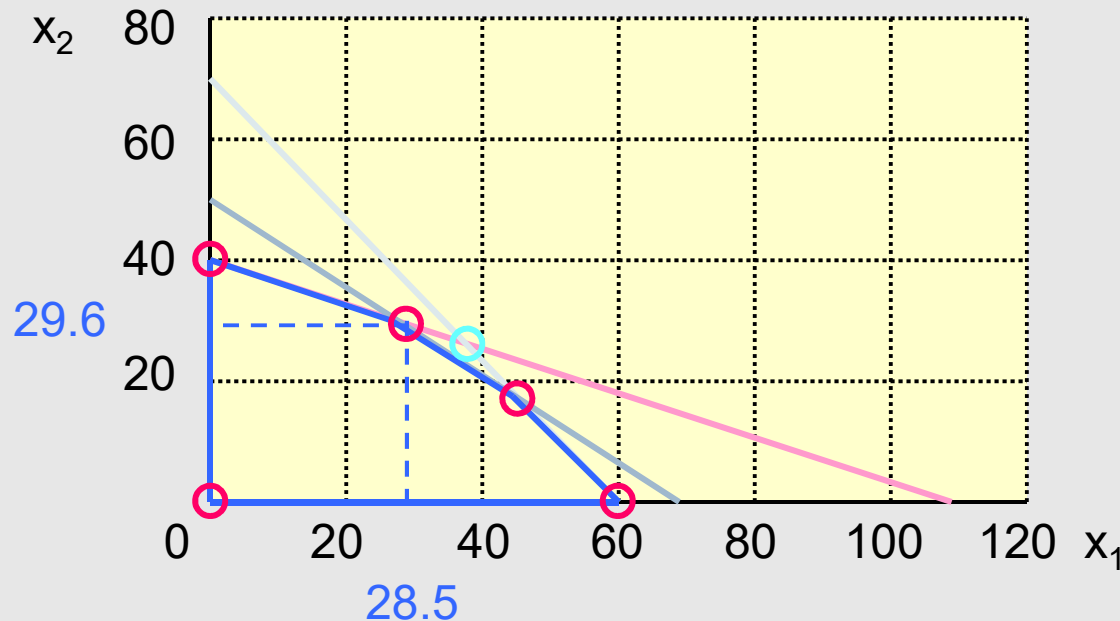
(In the graph, line for 2 dimensions
plane for 3 dimensions, ...)

$$f(x_1, x_2, \dots) = ax_1 + bx_2 + \dots \quad (a, b, \dots \text{ is fixed})$$

LP: Linear Programming

- ❑ Feasible Solution:
Solution that satisfies constraints
- ❑ Feasible Region:
Variable region including feasible solutions
- ❑ Optimal Solution :
the most optimized (maximized or minimized) solution of feasible solutions.
- ❑ Optimal Value:
the value of objective function at the optimal solution

Where is the optimal solution in the feasible region?



extreme point is:
Intersection point of the
boundaries of constraints
2 variables:

intersection point of
2 boundary lines

More than 3 variables:

intersection point of
boundary surfaces

(Extreme points are located at
the convex surface
of boundaries)

Extreme points of feasible region

= part of all the intersection points of constraint equations

Search extreme points and confirm if it is optimal solution.

Optimal Solution in the feasible region

- One optimal solution \Rightarrow at the extreme point.

- Many optimal solutions
 \Rightarrow on the boundary surface

Coefficients of constraint and objective function are same.

- No optimal solution \Rightarrow No feasible region

- No optimal solution \Rightarrow Feasible region is too wide.

The value of the objective function
is infinitely large.

Standard Form

$$\text{maximize } z = \sum_{j=1}^n c_j x_j$$

$$\text{subject to } \sum_{j=1}^n a_{ij} x_j = b_i \quad (i = 1, 2, \dots, m)$$

$$x_j \geq 0 \quad (j = 1, 2, \dots, n)$$

a_{ij}, c_j : fixed values,

b_i : fixed values (non-negative) ,

x_j : variables

m : num of equations, n : num of variables

Transformation to Standard form

Objective total benefit : $\max z = 4x_1 + 9x_2$
subject to material constraint : $4x_1 + 11x_2 \leq 440$
man power constraint : $5x_1 + 7x_2 \leq 350$
facility constraint : $7x_1 + 6x_2 \leq 420$
non-negative constraint : $x_1, x_2 \geq 0$



maximize $z = 4x_1 + 9x_2$
subject to $4x_1 + 11x_2 + s_1 = 440$
 $5x_1 + 7x_2 + s_2 = 350$
 $7x_1 + 6x_2 + s_3 = 420$
 $x_1, x_2, s_1, s_2, s_3 \geq 0$

Transformation to Standard form

- Constraint equation is expressed by inequality.
 - $=<$: use slack variable \Rightarrow use equal
 - $=>$: use surplus variable \Rightarrow use equal
- Right side fixed value of constraint equation is negative.
 - Times -1 for both sides
- Free variable that is not defined as non-negative exists.
 - Divide the variable to two non-negative variables (positive part and negative part).

Constraint equation is expressed by inequality.

Change to equal equation using not negative additional variable.

□ (the left side) $= < b$

(the left side) $+ s = b$

$s \geq 0$ slack variable

ex : $4x_1 - 7x_2 = < 12 \Rightarrow 4x_1 - 7x_2 + s = 12$

□ (the left side) $= > b$

(the left side) $- t = b$

$t \geq 0$ surplus variable

ex : $4x_1 - 7x_2 = > 12 \Rightarrow 4x_1 - 7x_2 - t = 12$

Free variable that is not defined as non-negative exists.

Free variable is replaced with the difference between two non-negative variables x^+ and x^- .

Free variable $x \Rightarrow x = x^+ - x^-$, $x^+ \geq 0$, $x^- \geq 0$

$$4x_1 - 7x_2 \leq 6, \quad x_1 \geq 0$$

$$4x_1 - 7(x_2^+ - x_2^-) \leq 6,$$

$$x_1 \geq 0, \quad x_2^+ \geq 0, \quad x_2^- \geq 0$$

Solution of simultaneous equations

$$\begin{aligned} \text{maximize } & z = 4x_1 + 9x_2 \\ \text{subject to } & 4x_1 + 11x_2 + s_1 = 440 \\ & 5x_1 + 7x_2 + s_2 = 350 \\ & 7x_1 + 6x_2 + s_3 = 420 \\ & x_1, x_2, s_1, s_2, s_3 \geq 0 \end{aligned}$$

Select x_1, x_2 as independent variables,
substitute 0 for these variables.
Values of s_1, s_2, s_3 are decided.

Relationship between the number of variables and the number of equations

m: number of variables,

n: number of equations

- $m \leq n$:

- Solution is decided uniquely or
 - is not decided uniquely
 - is not decided

- $m > n$:

- $m - n$ variables are not decided uniquely

- (these are called independent variables)

- \Leftrightarrow If the values of $m - n$ independent variables are given, the values of the remained variables are decided.

To find the extreme points

- 1 Transform constraints to Standard form
- 2 Select the independent variables and set 0 to decide the values of the other variables.

The solution of the simultaneous equations
= basic solution

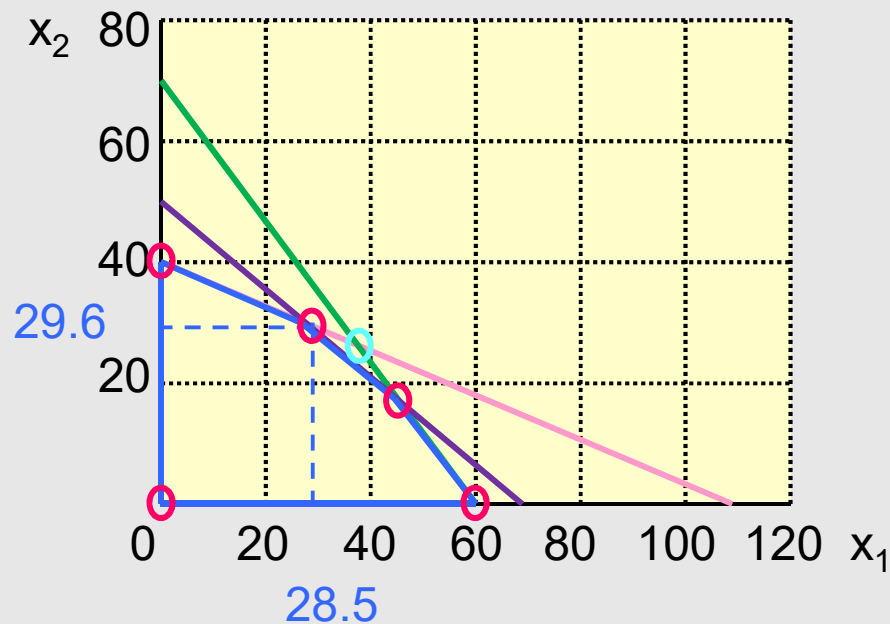
variables to find values = **basic variable**

⇔ variables to be set 0

(independent variables) = **non-basic variable**

- 3 When all of values of the basic solution are non-negative, this is one of the extreme points.

Extreme Points in the all of the intersection point of constraint equations



x_1	x_2	s_1	s_2	s_3
0	0	440	350	420
0	40	0	70	180
0	50	-110	0	120
0	70	-330	-140	0
110	0	0	-200	-350
70	0	160	0	-70
60	0	200	50	0
28.5	29.6	0	0	42.9
37.4	26.4	0	-21.8	0
44.2	18.4	60.8	0	0



In the next lecture,

Simplex method

Basic Method to derive the optimal solution for
Linear Programming problem,
finding only the extreme points increasing
the value of the objective function.

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Thank
you