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Problem Definition of FSSP

Minimize C_{max}

We aim to minimize the makespan, which is the total time required to complete all jobs.

Constraint

Subject to

$$egin{aligned} C_{max} & \geq s_{j,m} + p_{j,m} \ \ (j=1,2,\ldots,J) \ s_{j,k+1} & \geq s_{j,k} + p_{j,k} \ \ (j=1,2,\ldots,J; k=1,2,\ldots,m-1) \ s_{j,k} & \geq s_{l,k} + p_{l,k} - M(1-x_{j,l}^{(k)}) \ s_{l,k} & \geq s_{j,k} + p_{j,k} - Mx_{j,l}^{(k)} \ \ \ (j
eq l; j, l=1,2,\ldots,J; k=1,2,\ldots,m) \end{aligned}$$

Step 1: Makespan Constraint

$$C_{max} \geq s_{j,m} + p_{j,m} ~~(j=1,2,\ldots,J)$$

The makespan must be greater than or equal to the finish time of the last operation of every job

Step 2: Job Precedence Constraint (Intra-job ordering)

$$s_{j,k+1} \geq s_{j,k} + p_{j,k}$$

A job's next operation on machine M_{k+1} can only start after it finishes on machine M_k

Step 3: Machine Conflict Constraints (Inter-job sequencing)

$$egin{aligned} s_{j,k} &\geq s_{l,k} + p_{l,k} - M(1-x_{j,l}^{(k)}) \ s_{l,k} &\geq s_{j,k} + p_{j,k} - Mx_{j,l}^{(k)} \ \ (j
eq l; j,l = 1,2,\ldots,J; k = 1,2,\ldots,m) \end{aligned}$$

Jobs cannot overlap on the same machine. Either job j goes before job l , or vice versa.

Binary variable $x_{j,l}^{(k)}$ determines order. Big-M logic ensures at least one ordering is enforced.

Fixed values

J: number of jobs

m: number of machines

 $p_{j,k}$: process time of job j on machine M_k

Decision Variables

 C_{max} : makespan

 $s_{j,k}$: start time of job j on machine M_k

 $x_{j,l}^{(k)} \in \{0,1\}$: 1 if job j is before job l on machine M_k