template

```
#pragma GCC optimize("Ofast")
#include <bits/stdc++.h>
using namespace std;
typedef long long 11;
#define repx(i, a, b) for (int i = a; i < b; i++)
#define rep(i, n) repx(i, 0, n)
#define invrepx(i, a, b) for (int i = b - 1; i \ge a; i - -)
#define invrep(i, n) invrepx(i, 0, n)
int main() {
 ios::sync_with_stdio(false);
 cin.tie(NULL);
```

dsu

```
struct Dsu {
 vector<int>p.r:
  // initialize the disjoint-set-union to all unitary sets
 void reset(int N) {
   p.resize(N), r.assign(N, 0);
 rep(i, N) p[i] = i;
  // find the leader node corresponding to node `i`
 int find(int i) {
    if (p[i] != i) p[i] = find(p[i]);
    return p[i];
 // perform union on the two sets with leaders `i` and `j` // note: `i` and `j` must be GROUP LEADERS
 void unite(int i, int j) {
    if (i == j) return;
    if (r[i] > r[j]) swap(i, j);
    if (r[i] == r[j]) r[j] += 1;
   p[i] = j;
```

segtree

```
template <class T
struct StSum {
  vector<T> node;
  void reset(int N) { node.assign(4 * N, 0); }
  void build(const vectorT>\& a, int v = 1, int vl = 0, int vr = -1) {
    node.resize(4 * a.size());
    if (vr == -1) vr = node.size() / 4;
    if (vr - vl == 1) {
      node[v] = a[v1]; // construction
      return;
    int vm = (vl + vr) / 2;
    build(a, 2 * v, vl, vm);
    build(a, 2 * v + 1, vm, vr);
node[v] = node[2 * v] + node[2 * v + 1]; // query op
  // query for range [l, r)
  T query(int 1, int r, int v = 1, int vl = 0, int vr = -1) {
    if (vr == -1) vr = node.size() / 4;
if (1 == vl && r == vr) return node[v];
    int vm = (vl + vr) / 2:
```

```
T val = 0; // neutral element
    if (1 >= vr | | r <= vl) return val;
    val += query(1, min(r, vm), 2 * v, v1, vm);  // query op
val += query(max(1, vm), r, 2 * v + 1, vm, vr);  // query op
    return val;
  // set element i to val
  void update(int i, T val, int v = 1, int vl = 0, int vr = -1) {
    if (vr == -1) vr = node.size() / 4;
    if (vr - vl == 1) {
      node[v] = val;
       return;
    int vm = (vl + vr) / 2;
    if (i < vm) {</pre>
       update(i, val, 2 * v, vl, vm);
    }else {
      update(i, val, 2 * v + 1, vm, vr);
    node[v] = node[2 * v] + node[2 * v + 1]; // query op
  }
};
const 11 INF = 1e18:
template <class T>
struct StMax {
  vector<T> node;
  void reset(int N) { node.assign(4 * N, -INF); }
  void build(const vectorT>\& a, int v = 1, int vl = 0, int vr = -1) {
    node.resize(4 * a.size());
    if (vr == -1) vr = node.size() / 4;
    if (vr - vl == 1) {
       node[v] = a[v1]; // construction
    int vm = (vl + vr) / 2;
    build(a, 2 * v, vl, vm);
    build(a, 2 * v + 1, vm, vr);
    node[v] = max(node[2 * v], node[2 * v + 1]); // query op
   // query for range [l, r)
  T query (int 1, int r, int v = 1, int vl = 0, int vr = -1) {
    if (vr == -1) vr = node.size() / 4;
    if (1 == vl && r == vr) return node[v];
    int vm = (vl + vr) / 2;
    T val = -INF; // neutral element
    if (1 >= vr | | r <= vl) return val;
    val = max(val, query(1, min(r, vm), 2 * v, v1, vm));
                                                             // query op
    val = max(val, query(max(1, vm), r, 2 * v + 1, vm, vr)); // query op
    return val:
  // set element i to val
  void update(int i, T val, int v = 1, int vl = 0, int vr = -1) {
    if (vr == -1) vr = node.size() / 4;
    if (vr - vl == 1) {
       node[v] = val:
      return:
    }
    int vm = (vl + vr) / 2;
    if (i < vm) {
      update(i, val, 2 * v, vl, vm);
    lelse (
       update(i, val, 2 * v + 1, vm, vr);
    node[v] = max(node[2 * v], node[2 * v + 1]); // query op
  }
}:
template <class T>
struct StMin {
  vector<T> node;
  void reset(int N) { node.assign(4 * N, INF); }
  void build(const vectorT>\& a, int v = 1, int vl = 0, int vr = -1) {
    node.resize(4 * a.size());
     if (vr == -1) vr = node.size() / 4;
    if (vr - vl == 1) {
       node[v] = a[v1]; // construction
       return;
    int vm = (vl + vr) / 2;
    build(a, 2 * v, v1, vm);
```

```
build(a, 2 * v + 1, vm, vr);
    node[v] = min(node[2 * v], node[2 * v + 1]); // query op
  // query for range [l, r)
  T query(int 1, int r, int v = 1, int vl = 0, int vr = -1) {
    if (vr == -1) vr = node.size() / 4;
if (1 == vl && r == vr) return node[v];
    int vm = (vl + vr) / 2;
    T val = INF; // neutral element
    if (1 >= vr | | r <= vl) return val;</pre>
    val = min(val, query(1, min(r, vm), 2 * v, vl, vm));
                                                              // query op
    val = min(val, query(max(1, vm), r, 2 * v + 1, vm, vr)); // query op
    return val:
  7
  // set element i to val
  void update(int i, T val, int v = 1, int vl = 0, int vr = -1) {
    if (vr == -1) vr = node.size() / 4;
    if (vr - vl == 1) {
      node[v] = val;
      return;
    7
    int vm = (vl + vr) / 2;
    if (i < vm) {
      update(i, val, 2 * v, vl, vm);
    } else {
      update(i, val, 2 * v + 1, vm, vr);
    node[v] = min(node[2 * v], node[2 * v + 1]); // query op
};
```

segtreelazy

```
template <class T>
struct StlSumSum {
  // immediate (result of querying in the segment)
  // lazy (value that has not been pushed to the children)
 vector<pair<T, T>> node;
  void reset(int N) { node.assign(4 * N, {0, 0}); }
  void build(const vectorT>\& a, int v = 1, int vl = 0, int vr = -1) {
    node.resize(4 * a.size());
    if (vr == -1) vr = node.size() / 4;
    if (vr - vl == 1) {
      node[v].first = a[v1]; // construction
      return;
    int vm = (vl + vr) / 2;
    build(a, 2 * v, vl, vm);
    build(a, 2 * v + 1, vm, vr);
    \verb|node[v].first = \verb|node[2 * v].first + \verb|node[2 * v + 1].first; // query op|
    node[v].second = 0; // update-zero
  // helper: propagate lazy values in vertex `v` to both of its children
  void push(int v, int vl, int vr) {
    int vm = (vl + vr) / 2;
    T& lazy = node[v].second;
    node[2 * v].first += lazy * (vm - vl); // update-op & query-op mix
                                        // update-op
    node[2 * v].second += lazy;
    node[2 * v + 1].first += lazy * (vr - vm); // update-op & query-op mix
node[2 * v + 1].second += lazy; // update-op
    lazy = 0;
                                 // update-zero
  // update range [l, r) using val
  void update(int l, int r, T val, int v = 1, int vl = 0, int vr = -1) {
    if (vr == -1) vr = node.size() / 4;
    if (1 >= vr || r <= vl || r <= 1) return;
    if (1 == v1 && r == vr) {
      node[v].first += val * (vr - vl); // update-op & query-op mix
node[v].second += val; // update-op
      return:
    push(v, v1, vr);
    int vm = (vl + vr) / 2;
    update(1, min(r, vm), val, 2 * v, vl, vm);
    update(max(1, vm), r, val, 2 * v + 1, vm, vr);
    node[v].first = node[2 * v].first + node[2 * v + 1].first; // query-op
  // query range [l, r)
 T query(int 1, int r, int v = 1, int vl = 0, int vr = -1) {
```

```
if (vr == -1) vr = node.size() / 4;
    if (1 <= vl && r >= vr) return node[v].first;
    int vm = (vl + vr) / 2:
    T val = 0; // query-zero
    if (1 >= vr | | r <= vl | | r <= 1) return val;
    push(v. vl. vr):
    val += query(1, min(r, vm), 2 * v, v1, vm);  // query-op
    val += query(max(1, vm), r, 2 * v + 1, vm, vr); // query-op
    return val:
  }
};
const 11 INF = 1e18:
template <class T>
struct StlSetSum {
  // immediate (result of querying in the segment)
  // lazy (value that has not been pushed to the children)
  vector<pair<T, T>> node;
  void reset(int N) { node.assign(4 * N, {0, INF}); }
  void build(const vectorT>\& a, int v = 1, int vl = 0, int vr = -1) {
    node.resize(4 * a.size());
    if (vr == -1) vr = node.size() / 4;
    if (vr - vl == 1) {
      node[v] = {a[v1], INF}; // construction
    int vm = (vl + vr) / 2;
    build(a, 2 * v, vl, vm);
    build(a, 2 * v + 1, vm, vr);
    node[v].first = node[2 * v].first + node[2 * v + 1].first; // query op
    node[v].second = INF; // update-zero
  // helper: propagate lazy values in vertex `v` to both of its children
  void push(int v, int vl, int vr) {
    int vm = (vl + vr) / 2;
    T& lazy = node[v].second;
    if (lazy != INF) {
      node[2 * v].first = lazy * (vm - vl); // update-op & query-op mix
      node[2 * v] .second = lazy;
                                    // update-op
      node[2 * v + 1].first =
        lazy * (vr - vm);
                                // update-op & query-op mix
      node[2*v+1].second = lazy; // update-op
    lazy = INF; // update-zero
  // update range [l, r) using val
  void update(int 1, int r, T val, int v = 1, int vl = 0, int vr = -1) {
    if (vr == -1) vr = node.size() / 4;
    if (1 >= vr | | r <= vl | | r <= 1) return;
    if (1 == v1 && r == vr) {
      node[v].first = val * (vr - vl); // update-op & query-op mix
node[v].second = val; // update-op
      return;
    push(v, vl, vr);
    int vm = (vl + vr) / 2;
    update(1, min(r, vm), val, 2 * v, vl, vm);
    update(max(1, vm), r, val, 2 * v + 1, vm, vr);
    node[v].first = node[2 * v].first + node[2 * v + 1].first; // query-op
  // query range [l, r)
  T query(int 1, int r, int v = 1, int vl = 0, int vr = -1) {
    if (vr == -1) vr = node.size() / 4;
    if (1 <= vl && r >= vr) return node[v].first;
    int vm = (vl + vr) / 2;
    T val = 0; // query-zero
    if (1 >= vr || r <= vl || r <= 1) return val;
    push(v, v1, vr);
    val += query(1, min(r, vm), 2 * v, v1, vm); // query-op
    val += query(max(1, vm), r, 2 * v + 1, vm, vr); // query-op
    return val;
  }
};
template <class T>
struct StlSumMin {
  // immediate (result of querying in the segment)
  // lazy (value that has not been pushed to the children)
  vector<pair<T, T>> node;
  void reset(int N) { node.assign(4 * N, {0, 0}); }
  void build(const vectorT>\& a, int v = 1, int vl = 0, int vr = -1) {
```

```
node.resize(4 * a.size());
     if (vr == -1) vr = node.size() / 4;
    if (vr - vl == 1) {
      node[v] = {a[vl], 0}; // construction
      return;
    int vm = (vl + vr) / 2:
    build(a, 2 * v, vl, vm);
    build(a, 2 * v + 1, vm, vr);
    node[v].first =
      min(node[2 * v].first, node[2 * v + 1].first); // query op
    node[v].second = 0;
                                          // update-zero
  // helper: propagate lazy values in vertex `v` to both of its children
  void push(int v, int vl, int vr) {
    int vm = (vl + vr) / 2;
    T& lazy = node[v].second;
    node[2 * v].first += lazy;  // update-op & query-op mix
node[2 * v].second += lazy;  // update-op
    node[2 * v + 1].first += lazy; // update-op & query-op mix
node[2 * v + 1].second += lazy; // update-op
    lazy = 0;
                          // update-zero
  // update range [l, r) using val
  void update(int 1, int r, T val, int v = 1, int vl = 0, int vr = -1) {
     if (vr == -1) vr = node.size() / 4;
     if (1 >= vr || r <= vl || r <= 1) return;
     if (1 == vl && r == vr) {
       node[v].first += val; // update-op & query-op mix
       node[v].second += val; // update-op
    push(v, vl, vr);
     int vm = (vl + vr) / 2;
     update(1, min(r, vm), val, 2 * v, vl, vm);
     update(max(1, vm), r, val, 2 * v + 1, vm, vr);
    node[v].first =
      min(node[2 * v].first, node[2 * v + 1].first); // query-op
  // query range [l, r)
  T query(int 1, int r, int v = 1, int vl = 0, int vr = -1) {
    if (vr == -1) vr = node.size() / 4;
     if (1 <= vl && r >= vr) return node[v].first;
    int vm = (vl + vr) / 2;
    T val = INF; // query-zero
if (1 >= vr || r <= vl || r <= 1) return val;
    push(v, vl, vr);
    val = min(val, query(1, min(r, vm), 2 * v, v1, vm)); // query-op
    val = min(val, query(max(1, vm), r, 2 * v + 1, vm, vr)); // query-op
    return val;
};
```

sparse

```
// handle immutable range maximum queries (or any idempotent query) in O(1)
template <class T>
struct Sparse {
  vector<vector<T>>> st;
  Sparse() {}
  void reset(int N) { st = {vector<T>(N)}; }
  void set(int i, T val) { st[0][i] = val; }
  // O(N log N) time
  // O(N log N) memory
  void init() {
    int N = st[0].size();
    int npot = N <= 1 ? 1 : 32 - __builtin_clz(N);</pre>
    st.resize(npot);
    repx(i, 1, npot) rep(j, N + 1 - (1 << i)) st[i].push_back(</pre>
      \max(st[i-1][j], st[i-1][j+(1 << (i-1))])); // query op
  // query maximum in the range [l, r) in O(1) time
 T query(int 1, int r) {
  int i = 31 - __builtin_clz(r - 1);
    return max(st[i][1], st[i][r - (1 << i)]); // query op</pre>
}:
```

umap

```
// hackproof rng
static mt19937 rng(chrono::steady_clock::now().time_since_epoch().count());
// deterministic rng
uint64_t splitmix64(uint64_t* x) {
  uint64_t z = (*x += 0x9e3779b97f4a7c15);
  z = (z^{(z)} (z)) * 0xbf58476d1ce4e5b9;
  z = (z^{(z)} (z)) * 0x94d049bb133111eb;
 return z^(z >> 31);
}
// hackproof unordered map hash
struct Hash {
  size_t operator()(const ll& x) const {
    static const uint64_t RAND =
      chrono::steady_clock::now().time_since_epoch().count();
    uint64_t z = x + RAND + 0x9e3779b97f4a7c15;
    z = (z^{(z)}) * 0xbf58476d1ce4e5b9;
    z = (z^{(z)} (z) * 0x94d049bb133111eb;
    return z^(z >> 31);
 }
};
// hackproof unordered_map
template <class T, class U>
using umap = unordered_map<T, U, Hash>;
// hackproof unordered_set
template <class T>
using uset = unordered set<T, Hash>;
// an unordered map with small integer keys that avoids hashing, but allows
\hookrightarrow \Omega(N)
\ensuremath{/\!/} iteration and clearing, with N being the amount of items (not the maximum
// key).
template <class T>
struct Map {
  int N:
  vector<bool> used:
  vector<int> kevs:
  vector<T> vals;
  Map() {}
  II \cap (C)
  void resize(int C) {
    C += 1, used.resize(C), keys.resize(C), vals.resize(C);
  // 0(1)
  T& operator [] (int k) {
    if (!used[k]) used[k] = true, keys[N++] = k, vals[k] = T();
    return vals[k];
  // O(N)
  void clear() {
    while (N) used[keys[--N]] = false;
  // O(N)
  template <class OP>
  void iterate(OP op) {
    rep(i, N) op(keys[i], vals[keys[i]]);
  }
};
mo
struct Query {
 int 1, r, idx;
// answer segment queries using only `add(i)`, `remove(i)` and `get()`
// functions.
//
// complexity: O((N + Q) * sqrt(N) * F)
//
    N = length of the full segment
   Q = amount of queries
F = complexity of the `add`, `remove` functions
template <class A, class R, class G, class T>
void mo(vector<Query>& queries, vector<T>& ans, A add, R remove, G get) {
  int Q = queries.size(), B = (int)sqrt(Q);
```

```
sort(queries.begin(), queries.end(), [&] (Query& a, Query& b) {
    return make_pair(a.l / B, a.r) < make_pair(b.l / B, b.r);
});
ans.resize(Q);

int l = 0, r = 0;
for (auto& q : queries) {
    while (r < q.r) add(r), r++;
    while (l > q.l) l--, add(l);
    while (r > q.r) r--, remove(r);
    while (l < q.l) remove(l), l++;
    ans[q.idx] = get();
}
</pre>
```

math

<u>arith</u>

// floor(log2(n)) without precision loss

```
inline int floor_log2(int n) { return n <= 1 ? 0 : 31 - __builtin_clz(n); }</pre>
// ceil(log2(n)) without precision loss
inline int ceil_log2(int n) { return n <= 1 ? 0 : 32 - __builtin_clz(n - 1); }</pre>
inline 11 floordiv(11 a, 11 b) {
  11 d = a / b;
  return d * b == a ? d : d - ((a < 0) ^ (b < 0));
inline ll ceildiv(ll a, ll b) {
  11 d = a / b;
  return d * b == a ? d : d - ((a < 0) ^ (b < 0)) + 1;
// binary exponentiation
11 binexp(ll a, ll m) {
  11 res = 1; // neutral element
  while (m) {
    if (m & 1) res = res * a; // multiplication
    a = a * a;
                      // multiplication
    m >>= 1;
  return res;
// counts the divisors of a positive integer in O(\operatorname{sqrt}(n))
11 count_divisors(11 x) {
  11 \, \text{divs} = 1;
  11 i = 2;
  while (x > 1) {
    if (i * i > x) {
      divs *= 2;
      break;
    while (x \% i == 0) {
     x /= i;
      n += 1;
    divs *= n;
    i += 1;
  return divs;
// gets the prime factorization of a number in O(sqrt(n))
void factorize(vector<pair<11, int>>& facts, 11 x) {
  11 k = 2:
  while (x > 1) {
    if (k * k > x) {
      facts.push_back({x, 1});
      break:
    int n = 0:
    while (x \% k == 0) x /= k, n++;
    if (n > 0) facts.push_back({k, n});
```

```
k += 1:
 }
}
// iterate over all divisors of a number.
//
// divisor count upper bound: n^(1.07 / ln ln n)
template <class OP>
void divisors(11 x, OP op) {
  vector<pair<11, int>> facts;
  factorize(facts, x);
  vector<int> f(facts.size()):
  while (true) {
    11 y = 1;
    rep(i, f.size()) rep(j, f[i]) y *= facts[i].first;
    op(y);
    int i = 0;
    while (i < f.size()) {
      f[i] += 1;
      if (f[i] > facts[i].second)
        f[i++] = 0;
      else
        break;
    }
    if (i == f.size()) break;
 }
}
mod
// take the modulo for possibly negative numbers.
11 mod(11 a, 11 M) { return (a % M + M) % M; }
// binary exponentiation modulo M.
11 binexp(ll a, ll m, ll M) {
  assert(m >= 0);
  ll res = 1 \% M;
  while (m) {
   if (m & 1) res = (res * a) % M;
    a = (a * a) \% M;
    m >>= 1;
  7
  return res;
}
/\!/ compute the modular multiplicative inverse, assuming M is prime.
11 multinv(ll a, ll M) { return binexp(a, M - 2, M); }
// calculate gcd(a, b).
// also, calculate x and y such that:
// a * x + b * y == gcd(a, b)
//
// time: O(\log min(a, b))
// (ignoring complexity of arithmetic)
ll ext_gcd(ll a, ll b, ll& x, ll& y) {
 if (b == 0) {
    x = 1, y = 0;
    return a;
  11 d = ext_gcd(b, a % b, y, x);
  y = a / b * x;
  return d;
// compute the modular multiplicative inverse, assuming a and MOD are
   coprime.
// MOD may not be prime.
11 multinv_euc(ll a, ll M) {
 11 x, y;
  11 g = ext_gcd(a, M, x, y);
  assert(g == 1);
 return (x % M + M) % M;
// computes euler totative function phi(x), counting the amount of integers
// [1, x] that are relatively prime to x.
// time: 0(sqrt(x))
ll eulerphi(ll x) {
  11 \text{ phi} = 1;
  11 k = 2:
  while (x > 1) {
   if (k * k > x) {
```

phi *= x - 1;

return a:

```
break:
                                                                                         11 d = ext_gcd_small(b, a % b, y, x);
    11 k1 = 1 . k0 = 0:
                                                                                         v = a / b * x:
    while (x \% k == 0) x /= k, k0 = k1, k1 *= k;
                                                                                         return d:
    phi *= k1 - k0;
    k += 1:
  }
                                                                                       pair<11, 11> solve_crt(const vector<pair<11, 11>>& eqs) {
                                                                                         11 a0 = eqs[0].first, p0 = eqs[0].second;
  return phi;
                                                                                         repx(i, 1, eqs.size()) {
                                                                                           ll a1 = eqs[i].first, p1 = eqs[i].second;
11 k1, k0;
                                                                                           11 d = ext_gcd(p1, p0, k1, k0);
  return fact[n] * multinv(fact[k] * fact[n - k] % M, M) % M;
                                                                                           a0 = a1:
                                                                                           if (a0 % d != 0) return {-1, -1};
                                                                                            p0 = p0 / d * p1;
// multiply two big numbers (~10^18) under a large modulo, without resorting
                                                                                            a0 = a0 / d * k1 % p0 * p1 % p0 + a1;
                                                                                           a0 = (a0 \% p0 + p0) \% p0;
// bigints.
11 bigmul(11 x, 11 y, 11 M) {
                                                                                         return {a0, p0};
  11z = 0;
  while (y) {
   if (y & 1) z = (z + x) % M;
    x = (x << 1) % M, y >>= 1;
                                                                                       gauss
  return z:
                                                                                       const double EPS = 1e-9;
// discrete logarithm.
// solve a \hat{x} = b \pmod{M} for the smallest x.
                                                                                        // complexity: O(min(N, M) * N * M)
// returns -1 if no solution is found.
                                                                                        // `a` is a list of rows
                                                                                        // the last value in each row is the result of the equation
// time: O(sqrt(M))
                                                                                       // return values:
11 dlog(11 a, 11 b, 11 M) {
                                                                                        // 0 -> no solutions
  11 k = 1 \% M, s = 0;
                                                                                           1 -> unique solution, stored in `ans`
  while (true) {
                                                                                        // -1 -> infinitely many solutions, one of which is stored in `ans`
    11 g = __gcd(a, M);
                                                                                       int gauss (vector < double >> a, vector < double > & ans) {
    if (g <= 1) break;
                                                                                         int N = a.size(), M = a[0].size() - 1;
    if (b == k) return s;
    if (b \% g != 0) return -1;
                                                                                         vector<int> where(M, -1);
    b = g, M = g, s = 1, k = a / g * k % M;
                                                                                         for (int col = 0, row = 0; col < M && row < N; col++, row++) {
                                                                                           int sel = row:
  ll N = sqrt(M) + 1;
                                                                                           repx(i, row, N) if (abs(a[i][col]) > abs(a[sel][col])) sel = i;
                                                                                            if (abs(a[sel][col]) < EPS) continue;</pre>
  static umap<11, 11> r;
                                                                                           repx(i, col, M + 1) swap(a[sel][i], a[row][i]);
  r.clear();
                                                                                            where[col] = row:
  11 baq = b;
  rep(q, N+1) {
                                                                                           rep(i, N) if (i != row) {
    r[baq] = q;
                                                                                              double c = a[i][col] / a[row][col];
    baq = baq * a % M;
                                                                                              repx(j, col, M + 1) a[i][j] -= a[row][j] * c;
  11 aN = binexp(a, N, M), aNp = k;
  repx(p, 1, N + 1) {
                                                                                         ans.assign(M, 0);
    aNp = aNp * aN % M;
                                                                                         rep(i, M) if (where[i] != -1) ans[i] = a[where[i]][M] / a[where[i]][i];
    if (r.count(aNp)) return N * p - r[aNp] + s;
                                                                                         rep(i, N) {
                                                                                           double sum = 0;
                                                                                           rep(j, M) sum += ans[j] * a[i][j];
  return -1:
                                                                                           if (abs(sum - a[i][M]) > EPS) return 0;
                                                                                         rep(i, M) if (where[i] == -1) return -1;
                                                                                         return 1;
crt
// given a set of modular equations, each of the form x = Ai \pmod{Pi},
   compute
// `x mod LCM(PO, P1, ..., Pn) `.
                                                                                       poly
// returns a pair of `x` and `LCM(PO, P1, ..., Pn)`.
// if the equations cannot be satisfied, {-1, -1} is returned.
                                                                                        using cd = complex<double>;
pair<bigint, bigint> solve_crt_big(const vector<pair<11, 11>>& eqs) {
                                                                                       const double PI = acos(-1);
  bigint a0 = eqs[0].first, p0 = eqs[0].second;
  repx(i, 1, eqs.size()) {
                                                                                        // compute the DFT of a power-of-two-length sequence.
    bigint a1 = eqs[i].first, p1 = eqs[i].second;
                                                                                       // if `inv` is true, computes the inverse DFT.
    bigint k1, k0;
                                                                                       //
                                                                                       // the DFT of a polynomial A(x) = A0 + A1*x + A2*x^2 + ... + An*x^n is the
    bigint d = ext_gcd(p1, p0, k1, k0);
    a0 -= a1;
                                                                                       \hookrightarrow array
    if (a0.divmod(d) != 0) return {-1, -1};
                                                                                       // of the polynomial A evaluated in all nths roots of unity: [A(w0), A(w1),
                                                                                       // A(w2), ..., A(wn-1)], where w0 = 1 and w1 is the nth principal root of
    p0 /= d, p0 *= p1;
    a0 *= k1, a0 *= p1, a0 += a1;
                                                                                       \hookrightarrow unity.
    a0 %= p0, a0 += p0, a0 %= p0;
                                                                                       void fft(vector<cd>& a. bool inv) {
                                                                                         int N = a.size();
                                                                                         assert(N == 1 << __builtin_ctz(N));</pre>
  return {a0, p0};
                                                                                         int k = 0:
ll ext_gcd_small(ll a, ll b, ll& x, ll& y) {
                                                                                         rep(i, N) {
  if (b == 0) {
                                                                                           int bit = N >> 1:
    x = 1, y = 0;
                                                                                            while (k & bit) k ^{=} bit, bit >>= 1;
```

k ^= bit:

```
if (i < k) swap(a[i], a[k]);</pre>
  for (int len = 2; len <= N; len <<= 1) {
    double ang = 2 * PI / len * (inv ? -1 : 1);
    cd wlen(cos(ang), sin(ang));
    for (int i = 0; i < N; i += len) {
      cdw(1);
      repx(j, 0, len / 2) {
        cd u = a[i + j], v = a[i + j + len / 2] * w;
        a[i + j] = u + v;
a[i + j + len / 2] = u - v;
        w *= wlen;
      i += len;
    }
    len <<= 1;
  if (inv)
    for (cd&x : a) x /= N;
const 11 MOD = 7340033, ROOT = 5, ROOTPOW = 1 << 20;
// compute the DFT of a power-of-two-length sequence, modulo a special prime
// number with principal root.
// the modulus _must_ be a prime number with an Nth root of unity, where N
// power of two. the FFT can only be performed on arrays of size \leftarrow N.
void modfft(vector<11>& a, bool inv) {
  int N = a.size();
  assert(N == 1 << __builtin_ctz(N) && N <= ROOTPOW);</pre>
  rep(i, N) a[i] = (a[i] % MOD + MOD) % MOD;
  int k = 0;
  repx(i, 1, N) {
    int bit = N >> 1:
    while (k \& bit) k = bit, bit >>= 1;
    if (i < k) swap(a[i], a[k]);</pre>
  for (int len = 2; len <= N; len <<= 1) {
    11 wlen = inv ? multinv(ROOT, MOD) : ROOT;
    for (ll i = ROOTPOW; i > len; i >>= 1) wlen = wlen * wlen % MOD;
    for (int i = 0; i < N; i += len) {</pre>
      11 w = 1;
      repx(j, 0, len / 2) {
        11 u = a[i + j], v = a[i + j + len / 2] * w % MOD;
         a[i + j] = (u + v) \% MOD;
         a[i + j + len / 2] = (u - v + MOD) % MOD;
        w = w * wlen % MOD;
      }
   }
  if (inv) {
    11 ninv = multinv(N, MOD):
    for (11\& x : a) x = x * ninv % MOD;
using T = 11;
T pmul(T a, T b) { return a * b % MOD; }
T padd(T a, T b) { return (a + b) % MOD; }
T psub(T a, T b) { return (a - b + MOD) % MOD; }
T pinv(T a) { return multinv(a, MOD); }
struct Poly {
  vector<T> a;
  Poly() {}
  Poly(T c) : a(c) { trim(); }
  Poly(vector<T>c): a(c) { trim(); }
  void trim() {
    while (!a.empty() && a.back() == 0) a.pop_back();
  int deg() const { return a.empty() ? -1000000 : a.size() - 1; }
  Poly sub(int 1, int r) const {
    r = min(r, (int)a.size()), l = min(l, r);
    return vector<T>(a.begin() + 1, a.begin() + r);
  Poly trunc(int n) const { return sub(0, n); }
  Poly shl(int n) const {
    Poly out = *this;
    out.a.insert(out.a.begin(), n, 0);
```

```
return out:
Poly rev(int n, bool r = false) const {
  Polvout(*this):
  if (r) out.a.resize(max(n, (int)a.size()));
  reverse(out.a.begin(), out.a.end());
  return out.trunc(n):
Poly& operator+=(const Poly& rhs) {
  auto& b = rhs.a;
  a.resize(max(a.size(), b.size()));
  rep(i, b.size()) a[i] = padd(a[i], b[i]); // add
  trim():
  return *this;
Poly& operator == (const Poly& rhs) {
  auto& b = rhs.a;
  a.resize(max(a.size(), b.size()));
  rep(i, b.size()) a[i] = psub(a[i], b[i]); // sub
  trim();
  return *this;
Poly& operator*=(const Poly& rhs) {
  int n = deg() + rhs.deg() + 1;
  if (n <= 0) return *this = Poly();</pre>
  n = 1 \ll (n \ll 1?0:32 - \_builtin_clz(n-1));
  vector<T> b = rhs.a;
  a.resize(n), b.resize(n);
  modfft(a, false), modfft(b, false);
  rep(i, a.size()) a[i] = pmul(a[i], b[i]); // mul
  modfft(a, true), trim();
  return *this;
Poly inv(int n) const {
  assert(deg() >= 0);
  Poly ans = pinv(a[0]); // inverse
  int b = 1;
  while (b < n) {
    Poly C = (ans * trunc(2 * b)).sub(b, 2 * b);
    ans -= (ans * C).trunc(b).shl(b);
  return ans.trunc(n);
Poly operator+(const Poly& rhs) const { return Poly(*this) += rhs; }
Poly operator-(const Poly& rhs) const { return Poly(*this) -= rhs; }
Poly operator*(const Poly& rhs) const { return Poly(*this) *= rhs; }
pair<Poly, Poly> divmod(const Poly& b) const {
  if (deg() < b.deg()) return {Poly(), *this};</pre>
  int d = deg() - b.deg() + 1;
  Poly D = (rev(d) * b.rev(d).inv(d)).trunc(d).rev(d, true);
  return \{D, *this - D * b\}:
Poly operator/(const Poly& b) const { return divmod(b).first: }
Poly operator (const Poly b) const { return divmod(b).second; }
Poly& operator/=(const Poly& b) { return *this = divmod(b).first; }
Poly& operator%=(const Poly& b) { return *this = divmod(b).second; }
Teval(Tx) {
  T v = 0:
  invrep(i, a.size()) y = padd(pmul(y, x), a[i]); // add, mul
  return y;
Poly& build(vector<Poly>& tree, vector<T>& x, int v, int l, int r) {
  if (1 == r) return tree[v] = vector<T>{-x[1], 1};
  int m = (1 + r) / 2;
  return tree[v] = build(tree, x, 2 * v, 1, m) *
           build(tree, x, 2 * v + 1, m + 1, r);
void subeval(vector<Poly>& tree, vector<T>& x, vector<T>& y, int v, int l,
      int r) {
  if (1 == r) {
    y[1] = eval(x[1]);
    return;
  }
  int m = (1 + r) / 2;
  (*this \frac{1}{4} tree [2 * v]).subeval(tree, x, y, 2 * v, 1, m);
  (*this \frac{1}{2} tree [2 * v + 1]).subeval(tree, x, y, 2 * v + 1, m + 1, r);
// evaluate m points in O(k (log k)^2) with k = max(n, m).
vector<T> multieval(vector<T>& x) {
  int N = x.size();
  if (deg() < 0) return vector < T > (N, 0);
  vector<Poly> tree(4 * N);
  build(tree, x, 1, 0, N - 1);
```

```
vector<T> y(N);
  subeval(tree, x, y, 1, 0, \mathbb{N} - 1);
  return v;
friend ostream& operator<<(ostream&s, const Poly&p) {
  bool first = true;
  rep(i, p.a.size()) {
    if (p.a[i] == 0) continue;
     if (!first) s << " + ";
    s \ll p.a[i];
    if (i > 0) s << " x";
if (i > 1) s << "^" << i;
    first = false;
  }
  s << ")";
  return s;
}
```

biaint

```
using u32 = uint32 t:
using u64 = uint64 t;
// signed bigint
struct bigint {
  vector<u32> digits;
  u32 neg;
  bigint() : neg(0) {}
  \label{local_bigint}  \mbox{bigint(ll x) : digits{lo(x), hi(x)}, neg(x < 0 ? ~0 : 0) { this->trim(); } } 
  bigint(vector<u32> d) : digits(d), neg(0) {}
  static u32 lo(u64 dw) { return (u32)dw; }
  static u32 hi(u64 dw) { return (u32)(dw >> 32); }
  // remove leading zeros from representation
  void trim() {
    while (digits.size() && digits.back() == neg) digits.pop_back();
  void add(const bigint& rhs, u32 c = 0) {
    int ls = digits.size();
    int rs = rhs.digits.size();
    rep(i, max(ls, rs)) {
      if (i >= ls) digits.push_back(neg);
      u64 r = (u64) digits[i] + (i < rs ? rhs.digits[i] : rhs.neg) + c;
      digits[i] = lo(r), c = hi(r);
    u64 ec = (u64)c + neg + rhs.neg;
neg = ((hi(ec) ^ neg ^ rhs.neg) & 1 ? ~0:0);
    if (lo(ec) != neg) digits.push_back(lo(ec));
  bigint& operator+=(const bigint& rhs) {
    this->add(rhs);
    return *this;
  bigint& operator+=(u32 rhs) {
    this->add({}, rhs);
    return *this;
  void negate() {
    rep(i, digits.size()) digits[i] = ~digits[i];
    neg = ~neg;
    this->add({}, 1);
  bigint negated() const {
    bigint out = *this;
    out.negate();
    return out;
  bigint& operator == (const bigint& rhs) {
    this->negate();
    *this += rhs;
    this->negate():
   return *this;
  bigint& operator*=(bigint& rhs) {
    static bigint lhs;
    swap(*this, lhs), digits.clear(), neg = 0;
```

```
u32 r = rhs.neg, s = 0;
if (lhs.neg) s ^= lhs.neg, lhs.negate();
  if (rhs.neg) s ^= rhs.neg, rhs.negate();
  rep(j, rhs.digits.size()) {
    u64 c = 0;
    int ls = digits.size();
    int rs = lhs.digits.size();
    repx(i, j, max(ls, rs + j)) {
      if (i >= ls) digits.push_back(0);
      u64 r =
         (u64)digits[i] +
         (u64)(i-j < rs? lhs.digits[i-j]:0) * rhs.digits[j] +
      digits[i] = lo(r), c = hi(r);
    if (c != 0) digits.push_back(c);
  if (r) rhs.negate();
  if (s) negate();
  return *this;
bigint& operator/=(bigint& rhs) {
  divmod(rhs);
  return *this:
bigint& operator%=(bigint& rhs) {
  *this = divmod(rhs);
  return *this;
int divmod_trunc(int rhs) {
  u32 s = (rhs < 0 ? ~0 : 0) ^ this->neg, q = abs(rhs);
  u64 r = 0;
  if (this->neg) this->negate();
  invrep(i, digits.size()) {
    r = (r << 32) | digits[i];
    digits[i] = r/q, r \% = q;
  if (s) {
    this->negate();
    return -(int)r;
 return (int)r;
// compares `this` with `rhs`
    `this < rhs`: -1
   `this == rhs`: 0
// `this > rhs`: 1
int cmp(const bigint& rhs) const {
  if (neg && !rhs.neg) return -1;
  if (!neg && rhs.neg) return 1;
  int ls = digits.size(), rs = rhs.digits.size();
  invrep(i, max(ls, rs)) {
    u321 = i < ls ? digits[i] : neg;
    u32 r = i < rs ? rhs.digits[i] : rhs.neg;
    if (1 < r) return -1;
    if (1 > r) return 1;
 return 0;
bool operator == (const bigint this) const { return cmp(rhs) == 0; }
bool operator!=(const bigint& rhs) const { return cmp(rhs) != 0; }
bool operator<(const bigint& rhs) const { return cmp(rhs) == -1; }</pre>
bool operator>=(const bigint& rhs) const { return cmp(rhs) != -1; }
bool operator>(const bigint& rhs) const { return cmp(rhs) == 1; }
bool operator <= (const bigint ths) const { return cmp(rhs) != 1; }
friend ostream& operator<<(ostream& s, const bigint& self) {</pre>
  if (self == bigint()) return s << "0";</pre>
  bigint x = self;
  if (x.neg) {
    x.negate();
    s << "-";
  vector<int> digs;
  while (x != bigint()) digs.push_back(x.divmod_trunc(10));
  invrep(i, digs.size()) s << digs[i];</pre>
// truncating division and modulo
bigint divmod(bigint& rhs) {
  assert(rhs != bigint());
  u32 sr = rhs.neg, s = neg ^ rhs.neg;
  if (neg) negate();
  if (sr) rhs.negate();
```

```
bigint l = 0, r = *this, x:
    r += 1u;
    while (1 != r) {
      bigint m = 1;
      m += r;
      rep(i, m.digits.size()) m.digits[i] =
         (m.digits[i] >> 1)
         (i + 1 \le m.digits.size() ? m.digits[i + 1] \le 31 : 0);
      x = m, x *= rhs:
      if (x \le *this) {
        1 = (m += 1);
      } else {
        r = m;
    1 -= 1, swap(1, *this);
    r = *this, r *= rhs, 1 -= r;
    trim(), 1.trim();
    if (sr) rhs.negate();
    if (s) negate(), l.negate();
    return 1;
}:
// calculate gcd(a, b).
// also, calculate x and y such that:
// a * x + b * y == gcd(a, b)
bigint ext_gcd(bigint a, bigint b, bigint&x, bigint&y) {
  if (b == bigint()) {
    x = 1, y = 0;
    return a;
  bigint c = a.divmod(b), d = ext_gcd(b, c, y, x);
  a *= x, y -= a;
  return d;
```

segment

```
// in-place segment intersection.
void intersect(pair<int, int>& a, pair<int, int> b) {
 a = {max(a.first, b.first), min(a.second, b.second)};
// in-place segment "union".
// finds the shortest segment that contains both `a` and `b`.
// for [a, b) segments: change > to >= and <= to <
void unite(pair<int, int>& a, pair<int, int> b) {
 if (a.first > a.second)
    a = b:
  else if (b.first <= b.second)
    a = {min(a.first, b.first), max(a.second, b.second)};
// segment containment.
// [a, b] in [c, d]
// subset or equal: a >= c && b <= d // a > b
// proper subset: a > c && b < d || a > b && c <= d
// [a, b) in [c, d)
// subset or equal: a >= c && b <= d // a >= b
// proper subset: a > c \otimes b < d \mid a >= b \otimes c < d
bool is_subset(pair<int, int> sub, pair<int, int> sup) {
 return sub.first >= sup.first && sub.second <= sup.second ||
     sub.second < sub.first;</pre>
bool is_subset_proper(pair<int, int> sub, pair<int, int> sup) {
 return sub.first > sup.first && sub.second < sup.second | |
     sub.second < sub.first && sup.first <= sup.second;</pre>
```

theorems

```
Burnside lemma

For a set X, with members x in X, and a group G, with operations g in G, where g(x) 

→ : X -> X.

F_g is the set of x which are fixed points of g (ie. { x in X / g(x) = x }).

The number of orbits (connected components in the graph formed by assigning eac 

→ h x a node and a directed edge between x and g(x) for every g) is called M.
```

```
M = the average of the fixed points of all g = (|F_g1| + |F_g2| + ... + |F_gn|) / |F_g2|
GI
  If x are images and g are simmetries, then M corresponds to the amount of objects |
  corresponds to the amount of simmetries, and F_g corresponds to the amount of si \parallel

→ mmetrical

  images under the simmetry g.
Rational root theorem
  All rational roots of the polynomials with integer coefficients:
  a0 * x^0 + a1 * x^1 + a2 * x^2 + ... + an * x^n = 0
  If these roots are represented as p / q, with p and q coprime,
  - p is an integer factor of a0
  - q is an integer factor of an
  Note that if a0 = 0, then x = 0 is a root, the polynomial can be divided by x and th

→ e theorem

  applies once again.
Legendres formula
  Considering a prime p, the largest power p^k that divides n! is given by:
  k = floor(n/p) + floor(n/p^2) + floor(n/p^3) + ...
  O(\log n / \log p)
```

graph

dist

```
const 11 INF = 1e18:
// calculate distances between every pair of nodes in O(n^3) time and O(n^2)
// memory.
// requires an NxN array to store results.
void floyd(const vector<vector<pair<int, ll>>>& G, vector<vector<ll>>& dists) {
  int N = G.size();
  rep(i, N) rep(j, N) dists[i][j] = i == j ? 0 : INF;
  rep(i, N) for (auto edge : G[i]) dists[i] [edge.first] = edge.second;
  rep(k, N) rep(i, N) rep(j, N) {
    dists[i][j] = min(dists[i][j], dists[i][k] + dists[k][j]);
}
// calculate shortest distances from a source node to every other node in
// O(m log n). requires an array of size N to store results.
void dijkstra(const vector<vector<pair<int, ll>>>& G, vector<ll>& dists,
       int src) {
  dists.assign(G.size(), INF);
  dists[src] = 0;
  priority_queue<pair<11, int>> q;
  q.push({0, src});
  while (!q.empty()) {
    11 d = q.top().first;
    int v = q.top().second;
    q.pop();
     if (d > dists[v]) continue;
    for (auto edge : G[v]) {
      int to = edge.first;
      11 w = edge.second;
      if (d + w < dists[to]) {</pre>
        dists[to] = d + w;
        q.push({dists[to], to});
      }
 }
```

lca

flow

const ll INF = 1e18;

struct Flow {

```
// calculates the lowest common ancestor for any two nodes in O(\log N) time,
// with O(N log N) preprocessing
struct Lca {
 int L:
 vector<vector<int>> up;
 vector<pair<int, int>> time;
 Lca() {}
  void init(const vector<vector<int>>& G) {
    int N = G.size();
    L = N \le 1 ? 0 : 32 - \_builtin_clz(N - 1);
    up.resize(L + 1);
   rep(1, L + 1) up[1].resize(N);
    time.resize(N);
    int t = 0;
    visit(G, 0, 0, t);
    rep(1, L) rep(i, N) up[1+1][i] = up[1][up[1][i]];
  void visit(const vector<vector<int>>& G, int i, int p, int& t) {
    up[0][i] = p;
    time[i].first = t++
    for (int edge : G[i]) {
      if (edge == p) continue;
      visit(G, edge, i, t);
    time[i].second = t++;
 bool is_anc(int up, int dn) {
   return time[up].first <= time[dn].first &&
       time[dn].second <= time[up].second;
  int get(int i, int j) {
    if (is_anc(i, j)) return i;
    if (is_anc(j, i)) return j;
    int l = L;
    while (1 \ge 0) {
      if (is_anc(up[1][i], j))
       1--:
      else
        i = up[1][i];
    }
    return up[0][i];
matching
vector<bool> seen;
vector<int> mt;
bool subkuhn(vector<vector<int>>& adj, int i) {
 if (seen[i]) return false:
 seen[i] = true;
 for (int to : adj[i])
    if (mt[to] == -1 || subkuhn(adj, mt[to])) {
     mt[to] = i;
      return true:
 return false;
// get a maximum matching out of a **bipartite** graph.
// returns the size of the matching (ie. the covered vertices).
// runs in O(n * m) time.
int kuhn(vector<vector<int>>& adi) {
 int N = adj.size(), total = 0;
 mt.assign(N, -1);
 rep(i, N) {
    seen.assign(N, -1);
    total += subkuhn(adj, i);
 return total;
```

```
vector<11> excess;
  vector<int> height:
  Flow() {}
  void resize(int N) { cap.assign(N, vector<11>(N)); }
  // push as much excess flow as possible from u to v.
  void push(int u, int v) {
    11 f = min(excess[u], cap[u][v] - flow[u][v]);
    flow[u][v] += f;
    flow[v][u] -= f;
    excess[v] += f;
    excess[u] -= f;
  7
  // relabel the height of a vertex so that excess flow may be pushed.
  void relabel(int u) {
    int d = INT32_MAX;
    rep(v, cap.size()) if (cap[u][v] - flow[u][v] > 0) d =
      min(d, height[v]);
    if (d < INF) height [u] = d + 1;
  7
  // get the maximum flow on the network specified by `cap` with source `s`
  // and sink `t'
  // node-to-node flows are output to the `flow` member.
  // time: O(V^2 \ sqrt(E)) \iff O(V^3)
  // memory: 0(V^2)
  11 maxflow(int s, int t) {
    int N = cap.size(), M;
    flow.assign(N, vector<11>(N));
    height.assign(N, 0), height[s] = N;
    excess.assign(N, 0), excess[s] = INF;
    rep(i, N) if (i != s) push(s, i);
    vector<int> q;
    while (true) {
      // find the highest vertices with excess
      q.clear(), M = 0;
      rep(i, N) {
        if (excess[i] <= 0 || i == s || i == t) continue;
         if (height[i] > M) q.clear(), M = height[i];
        if (height[i] >= M) q.push_back(i);
      if (q.empty()) break;
       // process vertices
      for (int u : q) {
        bool relab = true;
        rep(v, N) {
           if (excess[u] <= 0) break;</pre>
           if (cap[u][v] - flow[u][v] > 0 && height[u] > height[v])
             push(u, v), relab = false:
        if (relab) {
          relabel(u):
           break;
        }
      }
    11 f = 0:
    rep(i, N) f += flow[i][t];
    return f;
 }
};
hld
struct Hld {
  vector<int> parent, heavy, depth, pos, top;
  Hld() {}
  void init(vector<vector<int>>& G) {
    int N = G.size():
    parent.resize(N), heavy.resize(N), depth.resize(N), pos.resize(N),
      top.resize(N):
    depth[0] = -1, dfs(G, 0);
    int t = 0:
    rep(i, N) if (heavy[parent[i]] != i) {
      int j = i;
      while (j != -1) {
        top[j] = i, pos[j] = t++;
```

vector<vector<ll>>> cap. flow:

```
j = heavy[j];
 }
int dfs(vector<vector<int>>& G. int i) {
  int w = 1, mw = 0;
  depth[i] = depth[parent[i]] + 1, heavy[i] = -1;
  for (int c : G[i]) {
   if (c == parent[i]) continue;
    parent[c] = i;
    int sw = dfs(G, c):
    if (sw > mw) heavy[i] = c, mw = sw;
    w += sw:
  7
  return w;
template <class OP>
void path(int u, int v, OP op) {
  while (top[u] != top[v]) {
    if (depth[top[u]] > depth[top[v]]) swap(u, v);
    op(pos[top[v]], pos[v]);
    v = parent[top[v]];
  }
  if (depth[u] > depth[v]) swap(u, v);
  op(pos[u], pos[v]); // value on vertex
     op(pos[u]+1, pos[v]); // value on path
// segment tree
template <class T, class S>
void update(S& seg, int i, T val) {
  seg.update(pos[i], val);
// segment tree lazy
template <class T, class S>
void update(S& seg, int u, int v, T val) {
 path(u, v, [&](int 1, int r) { seg.update(1, r, val); });
template <class T, class S>
T query(S& seg, int u, int v) {
  path(u, v, [k](int l, int r) \{ ans += seg.query(l, r); \}); // query op
 return ans;
```

strings

sufarr

```
// build the suffix array
// suffixes are sorted, with each suffix represented by its starting position
vector<int> suffixarray(const string& s) {
 int N = s.size() + 1; // optional: include terminating NUL
 vector<int> p(N), p2(N), c(N), c2(N), cnt(256);
 rep(i, N) cnt[s[i]] += 1;
 repx(b, 1, 256) cnt[b] += cnt[b-1];
 rep(i, N) p[--cnt[s[i]]] = i;
 repx(i, 1, N) c[p[i]] = c[p[i-1]] + (s[p[i]]! = s[p[i-1]]);
 for (int k = 1; k < N; k <<= 1) {
   int C = c[p[N-1]] + 1;
    cnt.assign(C + 1, 0);
   for (int& pi : p) pi = (pi - k + N) % N;
   for (int cl : c) cnt[cl + 1] += 1;
   rep(i, C) cnt[i + 1] += cnt[i];
   rep(i, N) p2[cnt[c[p[i]]]++] = p[i];
   c2[p2[0]] = 0;
   repx(i, 1, N) c2[p2[i]] =
     c2[p2[i-1]] + (c[p2[i]] != c[p2[i-1]] ||
               c[(p2[i]+k) \% N] != c[(p2[i-1]+k) \% N]);
    swap(c, c2), swap(p, p2);
```

```
p.erase(p.begin()); // optional: erase terminating NUL
  return p;
// build the lcp
// `lcp[i]` represents the length of the longest common prefix between
\hookrightarrow suffix i
// and suffix i+1 in the suffix array \ \ p \ \  the last element of \ \ \ \  is zero
\hookrightarrow by
// convention
vector<int> makelcp(const string& s, const vector<int>& p) {
  int N = p.size(), k = 0;
  vector<int> r(N), lcp(N);
  rep(i, N) r[p[i]] = i;
  rep(i, N) {
    if(r[i] + 1 >= N) {
       k = 0:
       continue;
    7
    int j = p[r[i] + 1];
     while (i + k < N & j + k < N & s[i + k] == s[j + k]) k += 1;
    lcp[r[i]] = k;
    if (k) k = 1;
  return lcp;
kmp
// compute the prefix function for string `s`:
// for every character substring [0 : i], compute the longest proper prefix
// is also a suffix.
// O(N)
11
/// computing `prefunc` on a string of the type `wwww#ttttttt` will give for // every `t` the amount of characters from `w` that match (ie. search for the
// string `w` inside the string `t`).
void prefunc(const string&s, vector<int>&p) {
  int N = s.size(), j;
  p.resize(N), p[0] = 0;
  repx(i, 1, N) {
    for (j = p[i - 1]; j > 0 \&\& s[j] != s[i];) j = p[j - 1];
    p[i] = j + (s[j] == s[i]);
 }
```

hash

}

```
// compute substring hashes in O(1).
// hashes are compatible between different strings.
struct Hash {
  11 HMOD;
  int N;
  vector<int>h;
  vector<int>p;
  Hash() {}
  Hash(const string& s, 11 HMOD_ = 1000003931)
    : N(s.size() + 1), HMOD(HMOD_), p(N), h(N) {
    static const 11 P =
      chrono::steady_clock::now().time_since_epoch().count() % (1 << 29);</pre>
    rep(i, N - 1) p[i + 1] = p[i] * P \% HMOD;
    rep(i, N - 1) h[i + 1] = (h[i] + (ll)s[i] * p[i]) % HMOD;
  pair<11, int> get(int i, int j) { return {(h[j] - h[i] + HMOD) % HMOD, i}; }
  bool cmp(pair<11, int> x0, pair<11, int> x1) {
    int d = x0.second - x1.second;
    ll& lo = d < 0 ? x0.first : x1.first;
    lo = lo * p[abs(d)] % HMOD;
    return x0.first == x1.first:
 }
};
// compute hashes in multiple prime modulos simultaneously, to reduce the
// of collisions.
```

```
struct HashM {
 int N;
 vector < Hash > sub:
 HashM() {}
  // O(K N)
 HashM(const string&s, const vector<11>& mods): N(mods.size()), sub(N) {
   rep(i, N) sub[i] = Hash(s, mods[i]);
  // \Omega(K)
 vector<pair<11, int>> get(int i, int j) {
    vector<pair<11, int>> hs(N);
   rep(k, N) hs[k] = sub[k].get(i, j);
 bool cmp(const vector<pair<11, int>>& x0, const vector<pair<11, int>>& x1) {
    rep(i, N) if (!sub[i].cmp(x0[i], x1[i])) return false;
    return true;
 bool cmp(int i0, int j0, int i1, int j1) {
    rep(i, N) if (!sub[i].cmp(sub[i].get(i0, j0),
                  sub[i].get(i1, j1))) return false;
    return true;
 }
};
```

hash2d

```
using Hash = pair<11, int>;
struct Block {
    int x0, y0, x1, y1;
struct Hash2d {
     11 HMOD;
      int W. H:
     vector<int>h;
     vector<int> p;
     Hash2d() {}
     Hash2d(const string&s, int W_, int H_, 11 HMOD_ = 1000003931)
            : W(W_ + 1), H(H_ + 1), HMOD(HMOD_) {
            static const 11 P =
                 chrono::steady_clock::now().time_since_epoch().count() % (1 << 29);</pre>
           p.resize(W * H);
           p[0] = 1;
           rep(i, W * H - 1) p[i + 1] = p[i] * P % HMOD;
           h.assign(W * H, 0);
           repx(y, 1, H) repx(x, 1, W) {
                 11c = (11)s[(y-1)*(W-1)+x-1]*p[y*W+x]%HMOD;
                 h[y * W + x] = (HMOD + h[y * W + x - 1] + h[(y - 1) * W + x] -
                                        h[(y-1)*W+x-1]+c)%
                                        HMOD;
      bool isout(Block s) {
           return s.x0 < 0 || s.x0 >= W || s.x1 < 0 || s.x1 >= W || s.y0 < 0 ||
                      s.y0 >= H | | s.y1 < 0 | | s.y1 >= H;
     Hash get(Block s) {
          return \{(2 * HMOD + h[s.y1 * W + s.x1] - h[s.y1 * W + s.x0] - h[s.y1 *
                        h[s.y0 * W + s.x1] + h[s.y0 * W + s.x0]) %
                       s.y0 * W + s.x0;
     bool cmp(Hash x0, Hash x1) {
           int d = x0.second - x1.second;
           11& lo = d < 0 ? x0.first : x1.first;</pre>
           lo = lo * p[abs(d)] % HMOD;
           return x0.first == x1.first;
    }
struct Hash2dM {
     int N:
     vector < Hash2d > sub:
     Hash2dM() {}
     Hash2dM(const string&s, int W, int H, const vector<11>& mods)
```

```
: N(mods.size()), sub(N) {
   rep(i, N) sub[i] = Hash2d(s, W, H, mods[i]);
}

bool isout(Block s) { return sub[0].isout(s); }

vector<Hash> get(Block s) {
   vector<Hash> hs(N);
   rep(i, N) hs[i] = sub[i].get(s);
   return hs;
}

bool cmp(const vector<Hash>& x0, const vector<Hash>& x1) {
   rep(i, N) if (!sub[i].cmp(x0[i], x1[i])) return false;
   return true;
}

bool cmp(Block s0, Block s1) {
   rep(i, N) if (!sub[i].cmp(sub[i].get(s0), sub[i].get(s1))) return false;
   return true;
}
};
```

palin

```
// find maximal palindromes (and therefore all palindromes) in \mathcal{O}(n).
// returns a vector of positions, with one position for every character and
// between characters
// a b c
// 0 1 2 3 4 5 6 7 8
// 101012321
void manacher(const string&s, vector<int>&p) {
  int N = s.size(), P = 2 * N - 1;
  p.assign(P, 0);
  int 1 = 0, r = -1;
 rep(i, P) {
    int d = (r >= i ? min(p[1+r-i], r-i+2) : i % 2);
    while (i - d \ge 0 & i + d \le P & s[(i - d) / 2] == s[(i + d) / 2])
    p[i] = d;
    if (i+d-2>r) 1 = i-d+2, r=i+d-2;
 rep(i, P) p[i] -= 1;
```

geo

2d

```
typedef double T;
struct P {
  Tx;
  Тy;
  P(T x_{-}, T y_{-}) : x\{x_{-}\}, y\{y_{-}\} \{\}
  P() : x{0}, y{0} {}
  friend ostream& operator << (ostream& s, const P& self) {
    s << self.x << "
                     " << self.y;
    return s;
  friend istream& operator>>(istream& s, P& self) {
    s >> self.x;
    s >> self.y;
    return s;
  P& operator+=(const P&r) {
    this->x += r.x;
    this->y += r.y;
```

```
return *this:
  friend P operator+(P1, const P&r) { return {1.x+r.x, 1.y+r.y}; }
  P& operator == (const P&r) {
    this->x -= r.x:
    this->y -= r.y;
    return *this;
  friend P operator-(P1, const P&r) { return {1.x-r.x, 1.y-r.y}; }
  P operator-() { return {-this->x, -this->y}; }
  P& operator*=(const T&r) {
    this->x *= r;
    this->y *= r;
    return *this;
  friend P operator*(P1, const T&r) { return {1.x*r, 1.y*r}; }
  friend P operator*(const T&1, Pr) { return {1 * r.x, 1 * r.y}; }
  P& operator/=(const T&r) {
    this->x /= r;
    this->y /= r;
    return *this;
  friend P operator/(P1, const T&r) { return {1.x/r, 1.y/r}; }
  // Dot product
  friend T operator*(P1, const P&r) { return 1.x * r.x + 1.y * r.y; }
  // Cross product (equiv to l.rotated() * r in 2D)
  friend Toperator (P1, const P&r) { return 1.x * r.y - 1.y * r.x; }
  T \text{ magsq() } \{ \text{ return this->x * this->x + this->y * this->y; } \}
  T mag() { return sqrt(this->magsq()); }
  P unit() { return (1. / this->mag()) * (*this); }
  Protated() { return {-this->y, this->x}; }
  double angle() { return atan2((double)this->y, (double)this->x); }
  float angle_float() { return atan2((float)this->y, (float)this->x); }
  static P from_angle(T angle) { return {(T)cos(angle), (T)sin(angle)}; }
// receives two segments in origin/distance format.
// the second segment is extended to infinity
// returns a weight, representing how much along the first line segment do
    the
// lines intersect.
// if it is in the [0, 1] range, they intersect.
T seg_line(P o0, P d0, P o1, P d1) {
  T d = d0 ^d1;
  if (d == 0) return 0;
 return ((o1 - o0) ^ d0) / d;
// returns true if two segments intersect.
// handles all corner cases correctly, including parallel and zero
// sized segments.
// in the non-parallel case, the intersection point can be retrieved // as 'o0 + n0 / d * d0' or 'o1 + n1 / d * d1'.
bool seg_seg(P o0, P d0, P o1, P d1) {
 P \circ = 01 - 00:
  T d = d0 ^d1;
  if (d == 0) {
    if ((o \hat{ } d0) != 0) return false;
    T = 0 = 0 * d0, e1 = (o + d1) * d0;
    return (0 <= e1 && e0 <= d0 * d0) || (0 <= e0 && e1 <= d0 * d0);
  T n0 = o ^d0, n1 = o ^d1;
  if (d < 0) n0 = -n0, n1 = -n1, d = -d;
  return 0 <= n0 && n0 <= d && 0 <= n1 && n1 <= d;
// iterate over all slopes, keeping points sorted with respect to the signed
// distance to the slope.
template <class OP>
void iter_slopes(vector<P>& points, OP op) {
  int N = points.size();
  vector<pair<int, int>> slopes;
  rep(i, N) rep(j, N) {
    if (i == j) continue;
    slopes.push_back({i, j});
  vector<int> perms(N);
  rep(i, N) perms[i] = i;
  sort(slopes.begin(), slopes.end(), [&](pair<int, int> i, pair<int, int> j) {
    P d1 = points[i.second] - points[i.first];
    Pd2 = points[j.second] - points[j.first];
```

```
return (d1 ^ d2) > 0;
  });
  for (auto&s : slopes) {
    int i = perms[s.first], j = perms[s.second];
    op(i, j);
    swap(points[perms[i]], points[perms[j]]);
    swap(perms[i], perms[j]);
}
3d
typedef double T;
struct Vec3 {
  Tx;
  Тy;
  Τz;
  Vec3(Tx_,Ty_,Tz_):x\{x_\},y\{y_\},z\{z_\}\{\}
  Vec3(): x{0}, y{0}, z{0}{}
  friend ostream& operator<<(ostream& out, const Vec3& self) {</pre>
    out << "(" << self.x << ", " << self.y << ", " << self.x << ")";
    return out;
  friend istream& operator>>(istream& in, Vec3& self) {
    in >> self.x;
    in >> self.v:
    in >> self.z:
    return in;
  Vec3& operator+=(const Vec3& r) {
    this->x += r.x:
    this->y += r.y;
    this->z += r.z:
    return *this;
  friend Vec3 operator+(Vec3 1, const Vec3& r) {
    return {1.x + r.x, 1.y + r.y, 1.z + r.z};
  Vec3& operator-=(const Vec3& r) {
    this->x -= r.x;
    this->y -= r.y;
    this->z -= r.z;
    return *this;
  1
  friend Vec3 operator-(Vec3 1, const Vec3&r) {
    return {1.x - r.x, 1.y - r.y, 1.z - r.z};
  Vec3 operator-() { return {-this->x, -this->y, -this->z}; }
  Vec3& operator*=(const T&r) {
    this->x *= r;
    this->y *= r;
    this->z *= r;
    return *this;
  friend Vec3 operator*(Vec3 1, const T&r) {
    return {1.x * r, 1.y * r, 1.z * r};
  friend Vec3 operator*(const T& 1, Vec3 r) {
    return {1 * r.x, 1 * r.y, 1 * r.z};
  Vec3& operator/=(const T&r) {
    this->x /= r;
    this->y /= r;
    this->z \neq r;
    return *this;
  friend Vec3 operator/(Vec3 1, const T&r) {
    return {1.x/r, 1.y/r, 1.z/r};
  // Dot product
  friend T operator*(Vec31, const Vec3&r) {
    return 1.x * r.x + 1.y * r.y + 1.z * r.z;
  // Cross product
  friend Vec3 operator^(Vec3 1, const Vec3& r) {
    return {1.y * r.z - r.y * 1.z, 1.x * r.z - r.x * 1.z,
        1.x * r.y - r.x * 1.y;
```

```
}
  T magsq() {
   return this->x * this->x + this->y * this->y + this->z * this->z;
  T mag() { return sqrt(this->magsq()); }
  Vec2 unit() { return (1. / this->mag()) * (*this); }
circle
struct Circle {
  Vec2o;
  Tr;
/// Find the pair of tangent points on circumference `c` to point `p`.
/// That is, find the tangent lines that cross point `p` and intersect the
/// circumference `c`, and return the points where these lines intersect `c`.
/// The first point returned is the counterclockwise tangent, followed by the
/// clockwise tangent.
///
/// If the point is inside the circle, NaN is returned.
pair<Vec2, Vec2> tangents(Circle c, Vec2 p) {
  Vec2 d = p - c.o;
  Tr2d2 = c.r * c.r / d.magsq();
  Vec2 mid = c.o + r2d2 * d;
  Vec2 dif = sqrt(r2d2 * (1. - r2d2)) * d.rotated();
  return {mid + dif, mid - dif};
search
// searches for a value in an [l, r] range (both inclusive).
/\!/ the `isleft(m)` function evaluates whether `m` is strictly to the left of
\hookrightarrow the
// target value.
int binsearch_left(int 1, int r, bool isleft(int)) {
  while (1 != r) {
    int m = (1 + r) / 2;
    if (isleft(m)) {
      1 = m + 1;
    } else {
      r = m;
    }
  return 1;
}
// searches for a value in an [l, r] range (both inclusive).
// the `isright(m)` function evaluates whether `m` is strictly to the right
    of
// the target value.
// note the `+1` when computing `m`, which avoids infinite loops.
// the only difference with `binsearch_left` is how the evaluation function
// specified. both are functionally identical.
int binsearch_right(int 1, int r, bool isright(int)) {
  while (1 != r) {
    int m = (1 + r + 1) / 2;
    if (isright(m)) {
     r = m - 1;
    } else {
      1 = m;
    }
 return 1;
// continuous ternary (golden section) search.
// searches for a minimum value of the given unimodal function (monotonic
// positive derivative).
template <typename T, typename U>
pair<T, U> ctersearch(int iter, Ta, Tb, Uf(T)) {
  const T INVG = 0.61803398874989484820:
  Uav = f(a):
  U bv = f(b):
```

```
T \text{ mid} = a + (b - a) * INVG;
  U midv = f(mid):
  for (int i = 0; i < iter; i++) {</pre>
    T \text{ new\_mid} = a + (mid - a) * INVG;
    U new_midv = f(new_mid);
    if (new_midv > midv) {
       // Search the right interval
       a = b;
       av = bv;
       b = new_mid;
       bv = new_midv;
    }else {
       // Search the left interval
       b = mid;
       bv = midv:
       mid = new_mid;
       midv = new_midv;
  1
  return {mid, midv};
}
```