Shamir's Secret Sharing Algorithm - Complete Implementation

Overview

Shamir's Secret Sharing (SSS) is a cryptographic algorithm that splits a secret into multiple shares, where a minimum threshold of shares is required to reconstruct the original secret. This implementation handles high-precision numbers, share verification, and wrong share detection.

Core Algorithm Components

1. Polynomial Generation

The secret is embedded as the constant term (a₀) of a polynomial:

$$f(x) = a_0 + a_1x + a_2x^2 + ... + a_{-1}x^{\wedge}(k-1)$$

Where:

- a₀ = secret
- a₁, a₂, ..., a_{□-1} = random coefficients
- k = minimum shares required (threshold)

2. Share Generation Process

```
For i = 1 to n:
share_i = (i, f(i))
hash_i = SHA256(i + "," + f(i))
```

3. Secret Reconstruction using Lagrange Interpolation

```
secret = f(0) = \Sigma(i=1 \text{ to } k) y_i \times L_i(0)
```

Where Lagrange basis polynomial:

$$L_i(0) = \Pi(j=1 \text{ to } k, j\neq i) (-x\square) / (x_i - x\square)$$

Implementation Steps

Step 1: High-Precision Arithmetic Functions

```
javascript
// Custom precision operations
function preciseAdd(a, b) {
```

```
return parseFloat((parseFloat(a) + parseFloat(b)).toFixed(15));
}
function preciseMultiply(a, b) {
  return parseFloat((parseFloat(a) * parseFloat(b)).toFixed(15));
}
function preciseDivide(a, b) {
  if (Math.abs(b) < 1e-15) throw new Error("Division by zero");
  return parseFloat((parseFloat(a) / parseFloat(b)).toFixed(15));
}
Step 2: Polynomial Evaluation
javascript
function evaluatePolynomial(coefficients, x) {
  let result = 0;
  for (let i = 0; i < coefficients.length; i++) {
     result += coefficients[i] * Math.pow(x, i);
  }
  return parseFloat(result.toFixed(15));
}
Step 3: Share Generation
javascript
function generateShares(secret, n, k) {
  // Generate k-1 random coefficients
  const coefficients = [parseFloat(secret)];
  for (let i = 1; i < k; i++) {
     coefficients.push(Math.random() * 1000);
  }
  // Generate n shares
  const shares = [];
  for (let x = 1; x \le n; x++) {
     const y = evaluatePolynomial(coefficients, x);
     const hash = createHash(x, y);
     shares.push({x: x, y: y, hash: hash});
  }
  return shares;
}
```

Step 4: Lagrange Interpolation for Reconstruction

```
javascript function lagrangeInterpolation(shares) {
```

```
let secret = 0;
  const n = shares.length;
  for (let i = 0; i < n; i++) {
     let xi = shares[i].x;
     let yi = shares[i].y;
     // Calculate Lagrange basis polynomial Li(0)
     let li = 1;
     for (let j = 0; j < n; j++) {
        if (i !== j) {
           let xj = shares[j].x;
           li = preciseMultiply(li, preciseDivide(-xj, xi - xj));
        }
     }
     secret = preciseAdd(secret, preciseMultiply(yi, li));
  return secret;
}
```

Share Verification Methods

1. Vandermonde Matrix Method

```
javascript
function createVandermondeMatrix(xValues, k) {
  const matrix = [];
  for (let i = 0; i < xValues.length; i++) {
     const row = [];
     for (let j = 0; j < k; j++) {
        row.push(Math.pow(xValues[i], j));
     }
     matrix.push(row);
  return matrix;
}
function calculateDeterminant(matrix) {
  // Recursive determinant calculation
  const n = matrix.length;
  if (n === 1) return matrix[0][0];
  if (n === 2) return matrix[0][0] * matrix[1][1] - matrix[0][1] * matrix[1][0];
  let det = 0;
  for (let i = 0; i < n; i++) {
     const subMatrix = getSubMatrix(matrix, 0, i);
     const cofactor = matrix[0][i] * calculateDeterminant(subMatrix);
```

```
det += (i % 2 === 0) ? cofactor : -cofactor;
}
return det;
}
```

2. Wrong Share Detection Algorithm

```
javascript
function findWrongShares(shares, k) {
  const combinations = getCombinations(shares, k);
  const reconstructionResults = [];
  // Test all k-combinations
  for (let combo of combinations) {
     try {
       const secret = lagrangeInterpolation(combo);
       reconstructionResults.push({
          shares: combo,
          secret: secret.toFixed(10)
       });
     } catch (error) {
       // Invalid combination
     }
  }
  // Find consensus
  const secretCounts = {};
  for (let result of reconstructionResults) {
     if (!secretCounts[result.secret]) {
       secretCounts[result.secret] = [];
     }
     secretCounts[result.secret].push(result);
  }
  // Identify most common result
  let mostCommon = null;
  let maxCount = 0;
  for (let [secret, results] of Object.entries(secretCounts)) {
     if (results.length > maxCount) {
       maxCount = results.length;
       mostCommon = secret:
     }
  }
  return {
     correctSecret: mostCommon,
     validCombinations: maxCount,
```

```
wrongShares: identifyWrongShares(shares, secretCounts[mostCommon])
};
}
```

Configuration Parameters

Standard Configuration (n=4, k=3):

- Total Shares (n): 4
- Minimum Required (k): 3
- Secret Precision: Up to 20 digits with decimals
- Hash Algorithm: SHA-256

JSON Share Format:

Error Detection Logic

Combination Testing Process:

- 1. Generate all possible k-combinations from available shares
- 2. Reconstruct secret for each combination using Lagrange interpolation
- 3. Group results by reconstructed secret value (rounded to 10 decimal places)
- 4. Find the most frequent result (consensus)
- 5. Mark shares not participating in consensus combinations as potentially wrong

Validation Criteria:

- Matrix Determinant ≠ 0: Ensures shares are linearly independent
- Consensus Threshold: Majority of combinations produce same result
- Hash Verification: SHA-256 integrity check for each share
- **Precision Tolerance**: Allow small floating-point errors (< 1e-10)

Usage Process

1. Generate Shares:

Input: secret = "123.456789012345"

n = 4, k = 3

Process: Generate polynomial, evaluate at x=1,2,3,4

Output: 4 shares with x,y coordinates and hashes

2. Verify and Reconstruct:

Input: JSON array of shares

Process: Apply Lagrange interpolation

Output: Reconstructed secret with error analysis

3. Detect Wrong Shares:

Input: Potentially corrupted shares

Process: Test all k-combinations, find consensus

Output: Identification of valid/invalid shares

Mathematical Foundation

Polynomial Mathematics:

• **Degree:** k-1 (one less than threshold)

• Uniqueness: Any k points uniquely determine a polynomial of degree k-1

• Security: Fewer than k shares reveal no information about the secret

Linear Algebra:

• Vandermonde Matrix: Ensures share uniqueness

• **Determinant Test:** Non-zero determinant confirms linear independence

• Matrix Inversion: Required for coefficient recovery

Precision Handling:

• Floating Point: 15-decimal precision to handle 20-digit secrets

• **Error Tolerance:** 1e-10 threshold for equality comparisons

• Rounding Strategy: Consistent rounding to avoid accumulation errors