Let TTRT = T (average token interval time)

Let  $\alpha_0, \alpha_1, ..., \alpha_{m-1}$  be the THT for each of the m stations

$$\alpha_0 + \alpha_1 + \ldots + \alpha_{m-1} \le T$$

Let  $t_0, t_1, ..., t_{m-1}$  be the time of arrival of token at stations 0, 1, ..., m-1

 $t_i$ , i > 0 is the time at which token reaches station i

= i mod m in cycle i/m

 $t_{-m},...,t_{-1}$ , be the times at which token arrives

at m, ..., 1 in the previous cycle

If  $t_i - t_{i-m} < T$ , low priority frames transmitted If  $t_i - t_{i-m} > T$ , no low priority frames transmitted Both case high priority traffic transmitted Time at which token reaches next node is  $t_{i+1} = t_{i-m} + T + \alpha_i$ , for  $t_i - t_{i-m} < T$ ,  $i \ge 0$  $t_{i+1} = t_i + \alpha_i$ , for  $t_i - t_{i-m} > T, i \ge 0$ where  $\alpha_i = \alpha_{i \mod m}$  is the allocated transmission plus propagation time for node (i mod m)

Special case : 
$$\alpha_i = 0$$
, for all i  $t_{i+1} \leq \max (t_i, t_{i-m} + T)$ ,  $i \geq 0$   
Since  $t_{i-m} \leq t_i$   $t_{i+1} \leq t_i + T$   
Similarly for  $1 \leq j \leq m+1$   $t_{i+j} \leq t_i + T$   
Hence  $t_{i+m+1} \leq t_i + T$ , for all  $i \geq 0$ 

```
Iterate over multiples of m + 1
t_i \le t_{i \mod (m+1)} + i/(m+1)T all i > 0
The m + 1 occurs to ensure that when stations
are heavily loaded every cycle a different transmits
First cycle station 0 transmits
Next cycle station 1 transmits ,...
t<sub>m</sub> - station 0 transmits T
t_{2m} = T \Rightarrow station 0 cannot transmit - token late
station 1 transmits \Rightarrow fair share to all stations
```

### Utilisation

$$U = \frac{1}{1 + a/N}$$

N - number of stations

a - propagation delay

1 - time take to transmit a packet

$$N \to \infty U \to 1$$