

# 字符串

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- ◆
- ◆ KMP算法

## 1. KMP模式匹配

- 1) 为了方便判断边界条件, 输入字符串时**跳过下标0**
- 2) 预处理 $next[i] = \max\{j\}$ ,  $j-1$ 为前缀串的最后一个下标, 且该前缀串与 $i-1$ 为结尾的前缀串的后缀串能匹配, 若不存在这种匹配, 则令 $next[i]=0$
- 3) 引: 设 $j_0$ 为 $next[i]$ 的候选项, 则 $next[j_0]$ 之前的都不是最大候选项
- 4) 引:  $j$ 是 $next[i]$ 的候选项 iff  $j-1$ 是 $next[i-1]$ 的候选项 (不越界时)
- 5) 应用:  $next[i-1]$ 及之前都打表成功后,  $next[i]$ 的可能值是 $next[i-1]+1$ ,  $next[next[i-1]]+1$ ,  $next[next[next[i-1]]]+1$ .....逐一判断即可
- 6) 求next数组
  - 1) 初值:  $next[1]=0$  (此处的下标为字符串中字符的序号, 没有0)
  - 2) 

```
for(int i=2, j=0; i<=n; ++i){
    while(j>0 && a[i]!=a[j+1])
        j=next[j];
    if(a[i]==a[j+1])
        ++j;
    next[i]=j;}
```
- 7) 求f数组 (长串中各字符为结尾时能匹配到的最长原串前缀长度)
  - 1) 

```
for(int i=1, j=0; i<=m; ++i){
    while(j>0 && (j==n || b[i]!=a[j+1]))
        j= next[j];
    if(b[i]==a[j+1])
        ++j;
    f[i]= j;
    //    if(f[i]==n)    cout<<"找到模式串了";
}
```

## 2. 例: 串串香匹配: 判断t个s中能否找到w或反向的w, s中可能有?代表任意字符

- 1) 将匹配从简单的相等改成相等或左边为问号
- 2) 将w反过来, 存在另一个数组, 分别求next数组
- 3) 

```
const int MN = 100005;
char w[MN], rw[MN], s[MN];
int nxt[MN], rnxt[MN]; //正向和反向的next数组
int t;
inline bool pp(char c, char c2){
    return c=='?' || c==c2;
} //重新定义匹配
```
- 4) 

```
scanf("%s%d", w+1, &t);
```

```

s[0]= w[0]= rw[0]= '!'; //空开下标0的字符
nxt[1]= 0;
int j= 0;
int len= strlen(w)-1;
for__(i,2,len){
    while(j>0 && w[i]!=w[j+1])
        j= nxt[i];
    if(w[i]==w[j+1])
        ++j;
    nxt[i]= j;}
for__(i,1,len)
    rw[i]= w[len+1-i];
rnxt[1]= 0;
j= 0;
for__(i,2,len){
    while(j>0 && rw[i]!=rw[j+1])
        j= rnxt[i];
    if(rw[i]==rw[j+1])
        ++j;
    rnxt[i]= j;}
while(t--){
    scanf("%s",s+1);
    int lens= strlen(s)-1;
    j= 0;
    bool ys= 0;
    for__(i,1,lens){
        while(j>0 && !pp(s[i], w[j+1]))
            j= nxt[j];
        if(pp(s[i], w[j+1]))
            ++j;
        if(j==len){
            ys= 1;
            break;}}
    if(ys){
        printf("YES\n");
        continue;}
    j= 0;
    for__(i,1,lens){
        while(j>0 && !pp(s[i], rw[j+1]))
            j= rnxt[j];
        if(pp(s[i], rw[j+1]))
            ++j;
        if(j==len){
            ys= 1;
            break;}}
    if(ys)
        printf("YES\n");
    else
        printf("NO\n");}

```

◆

◆ 常用函数

### 3. 求子串

```

std::string::substr
string substr (size_t pos = 0, size_t len = npos) const;

```

## Generate substring

- 1) Returns a newly constructed [string](#) object with its value initialized to a copy of a substring of this object.
- 2) The substring is the portion of the object that starts at character position pos and spans len characters (or until the end of the string, whichever comes first).
- 3) Parameters
  - 1) pos
    - 1) Position of the first character to be copied as a substring.
    - 2) If this is equal to the [string length](#), the function returns an empty string.
    - 3) If this is greater than the [string length](#), it throws [out\\_of\\_range](#).
    - 4) Note: The first character is denoted by a value of 0 (not 1).
  - 2) len
    - 1) Number of characters to include in the substring (if the string is shorter, as many characters as possible are used).
    - 2) A value of [string::npos](#) indicates all characters until the end of the string.
- 4) [size\\_t](#) is an unsigned integral type (the same as member type [string::size\\_type](#)).  
来自 <<http://www.cplusplus.com/reference/string/string/substr/>>
- 5) 顺带一提java除了substr外还有substring(l,r)自动给负数下标转换成0, l>r时自动交换lr, 再返回[l,r)

## 4. 找子串

[std::string::find](#)

### • C++11

```
string (1) size_t find (const string& str, size_t pos = 0) const noexcept;  
c-string (2) size_t find (const char* s, size_t pos = 0) const;  
buffer (3) size_t find (const char* s, size_t pos, size_type n) const;  
character (4) size_t find (char c, size_t pos = 0) const noexcept;
```

- 1) Find content in string
- 2) Searches the [string](#) for the first occurrence of the sequence specified by its arguments.
- 3) When pos is specified, the search only includes characters at or after position pos, ignoring any possible occurrences that include characters before pos.
- 4) Notice that unlike member [find\\_first\\_of](#), whenever more than one character is being searched for, it is not enough that just one of these characters match, but the entire sequence must match.
- 5) Parameters
  - 1) str
    - 1) Another [string](#) with the subject to search for.
  - 2) pos
    - 1) Position of the first character in the string to be considered in the search.
    - 2) If this is greater than the [string length](#), the function never finds

matches.

- 3) Note: The first character is denoted by a value of 0 (not 1): A value of 0 means that the entire string is searched.

3) s

- 1) Pointer to an array of characters.
- 2) If argument n is specified (3), the sequence to match are the first n characters in the array.
- 3) Otherwise (2), a null-terminated sequence is expected: the length of the sequence to match is determined by the first occurrence of a null character.

4) n

- 1) Length of sequence of characters to match.

5) c

- 1) Individual character to be searched for.

6) Return Value

- 1) The position of the first character of the first match.
- 2) If no matches were found, the function returns [string::npos](#).
- 3) [size\\_t](#) is an unsigned integral type (the same as member type [string::size\\_type](#)).

来自 <http://www.cplusplus.com/reference/string/string/find/>



◆ 例题

1. 求长度n的随机字符串中，长度不小于m的不同子串数量

- 1) 因为是随机的，所以长度大于8的其实没有重复子串.....
- 2) 本来想用(t=s.find(sb,t+1))!=s.npos来遍历每个相同字符串的，超时
- 3) 果然还是set大法好

```
4) ull ans=0;
   for__(l,m,min(n,10)){
       for__(i,0,n-l)
           ss.insert(s.substr(i,l));
       ans+=ss.size();
       ss.clear();}
   for__(l,max(m,11),n)
       ans+=n-l+1;
   cout<<ans;
```

i.

ii.

iii.

iv.

v.

vi.

vii.

viii.

ix.

x. -----我是底线-----