	Linked Lists
1.	Describe an algorithm that reverses a singly linked list using only constant extra space. Note: you cannot use recursion to solve this problem (the function call stack involved in recursion has a space cost).
	Hint: use a new linked list.
2.	Suppose that you have a pointer to a node in a singly linked list that is guaranteed not to be the last node in the list. You do not have pointers to any other nodes (except by following links). Describe ar $O(1)$ algorithm that logically removes the $value$ stored in such a node from the linked list, maintaining the integrity of the linked list.
	Hint: involve the next node.
3.	Lazy deletion is a style of deletion whereby data isn't actually deleted from the data structure. Rather it is simply marked as deleted. The rest of the data structure is then updated to ignore the marked elements as if they had been deleted. Would it be worthwhile to add lazy deletion to our linked list implementation? Why or why not?
4.	Linked lists can be kept sorted by finding the correct position when inserting. For a sorted list, how does the time complexity of the insert(), contains(), remove() operations change?
5.	Design an $O(n)$ algorithm to determine whether a linked list of unknown size contains a cycle. You may only use $O(1)$ extra space.
	Note: A linked list contains a cycle if, starting from some node p, following a sufficient number of next links brings us back to node p. Node p does not have to be the first node in the list.
	Hint: use two pointers that are initially at the start of the list, but advance at different speeds.

Homework 2

CompSci 404.1

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	Stacks an	nd Queues		
6.	Write the contents of each data structure after the following operations are performed:			
	add(1) add(2) remove() add(3) add(4) remove() remove() add(5)			
	(a) Stack			
	(b) Queue			
7.		queues backed by linked lists perform bette	er than queues backed by arrays.	
8.	The time comp through every	plexity for size() in an inefficient impl	ementation of a stack or queue is $O(n)$ (iterate cribe an $O(1)$ algorithm to return the size of a	
	Hint: keep trac	ck of the size independently, do not calcula	ate it every time.	
9.		thod to store two separate stacks within a from both stacks independently. Be as spa	a single array. Your method should be able to ce efficient as possible.	