

FUNDAMENTALS

BASIC PROGRAMMING MODEL

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JAVA SYNTAX REVIEW

LINEAR SEARCH

- Problem: Given an array of values, determine if a value is contained in that array.
- Mechanism: Check each element against the value.
- Basic Java Program:
 - Contained in a class.
 - Has a *main()* method.
 - Uses the Java API to read target value.
 - Uses a user defined method to implement search.

```

/**
 * This program provides an implementation of the linear search algorithm
 * and demonstrates it.
 *
 * @author Acuna
 * @version 1.0
 */
import java.util.Scanner;

public class LinearSearchExample
{
    public static boolean find(int target, int[] pool)
    {
        for(int i = 0; i < pool.length; i++)
            if(pool[i] == target)
                return true;

        return false;
    }

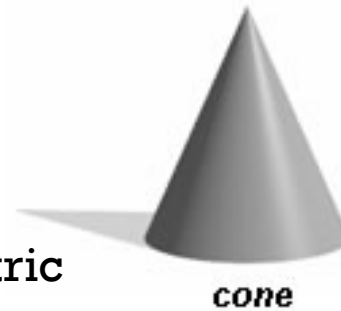
    public static void main(String args[])
    {
        Scanner scanner = new Scanner(System.in);
        int[] data = {4, 45, 8, 1, 3, 3, 22, 9};
        int target;

        System.out.println("What is the target number?");
        target = scanner.nextInt();
        if(find(target, data))
            System.out.println("Found.");
        else
            System.out.println("Missing.");
    }
}

```

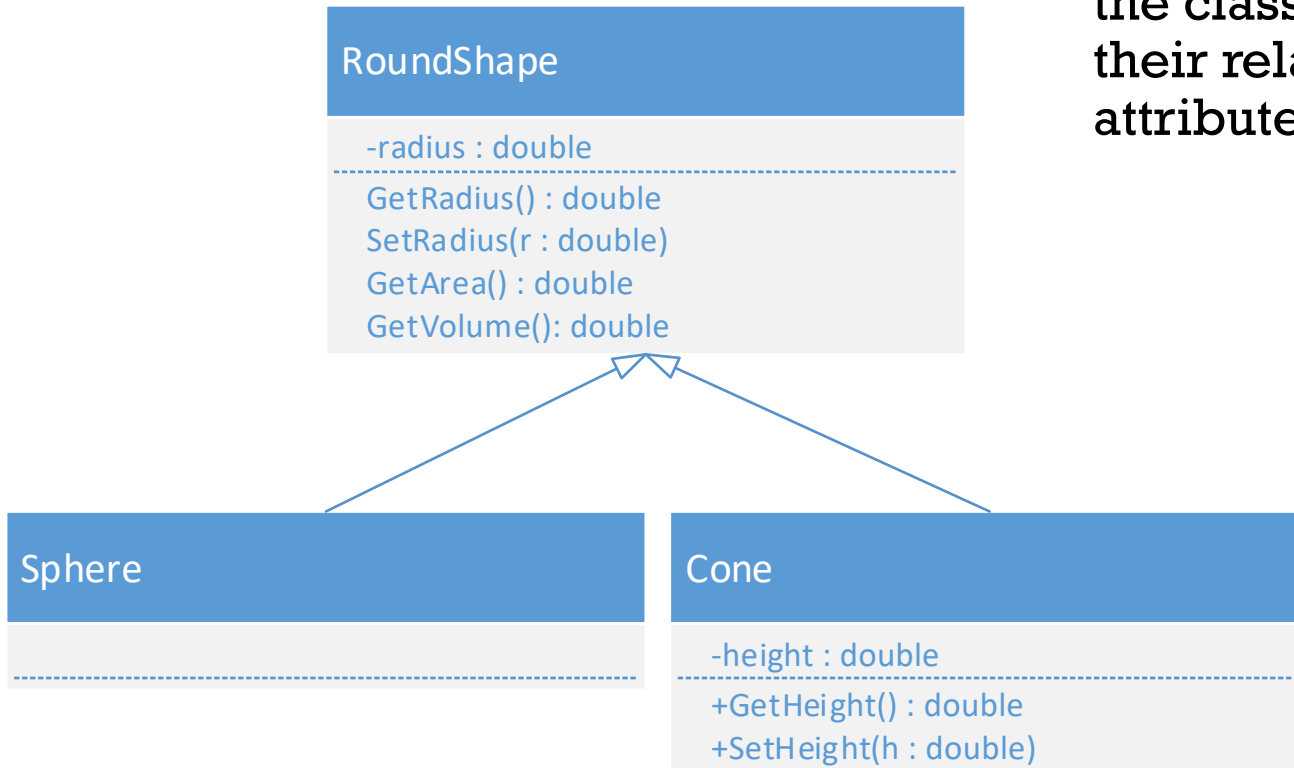
SHAPE LIBRARY

- Problem: Need a way to define geometric shapes in a program.
- In particular, shapes with a radius (*state*).
- Shapes should also be able to compute area and volume (exhibit *behavior*).
- One way to accomplish this aim is to implement shapes as **classes**.
- Organization of the classes should follow their concept(s) in terms of OOP principles.



SHAPE UML

- One standard way to model a system is UML.
- A class diagram describes the classes in a program, their relationships, and attributes/operations.



```
public abstract class RoundShape
{
    private double radius;

    public RoundShape(double r)
    {
        radius = r;
    }

    public double GetRadius()
    {
        return radius;
    }

    public void SetRadius(double r)
    {
        radius = r;
    }

    public abstract double GetArea();
    public abstract double GetVolume();
}
```

```
public class Cone extends RoundShape
{
    private double height;

    public Cone(double r, double h)
    {
        super(r);
        height = h;
    }

    public double GetHeight()
    {
        return height;
    }

    public void SetHeight(double r)
    {
        height = r;
    }

    public double GetArea()
    {
        return Math.PI * GetRadius() * (GetRadius())
    }

    public double GetVolume()
    {
        return Math.PI * Math.pow(GetRadius(), 2) *
    }

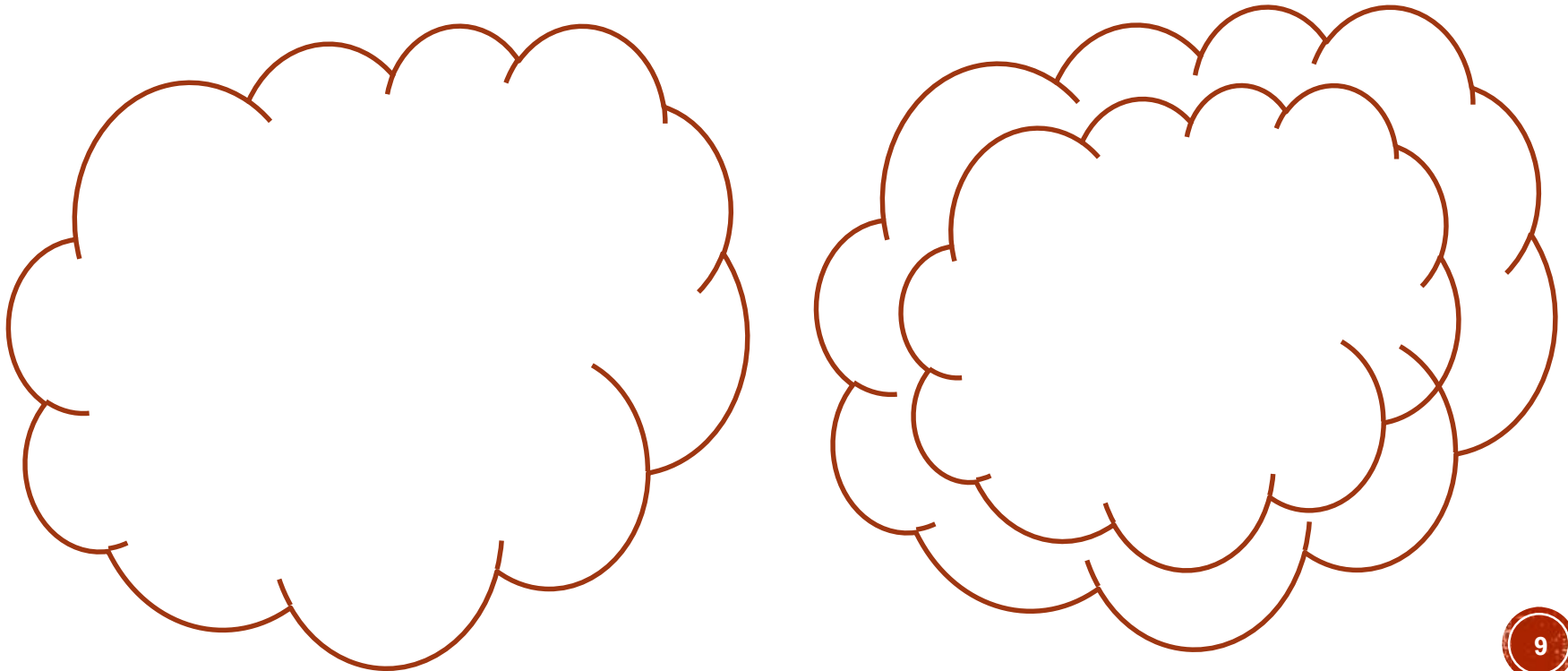
    String ToString()
    {
        return "A cone of radius " + GetRadius() +
    }
}
```



RECURSION

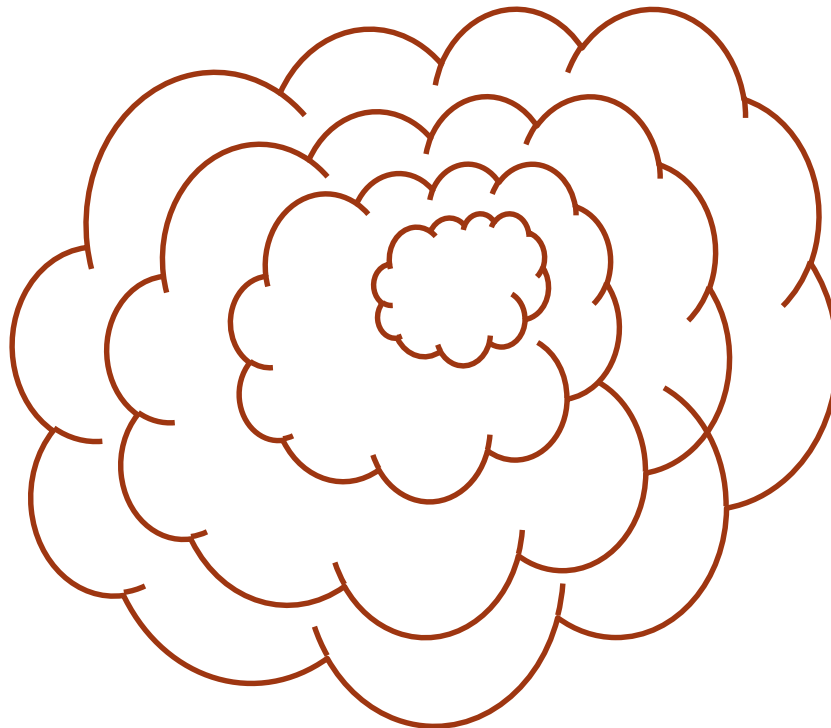
THE RECURSION CONCEPT

Suppose I want to draw a puffy cloud. Well, I can use PowerPoint's handy cloud tool. Only, it's not really puffy. So, I copy the first cloud, shrink it, and move it into the original cloud. That's okay, but that inner cloud doesn't look puffy. Maybe I can repeat the copy...



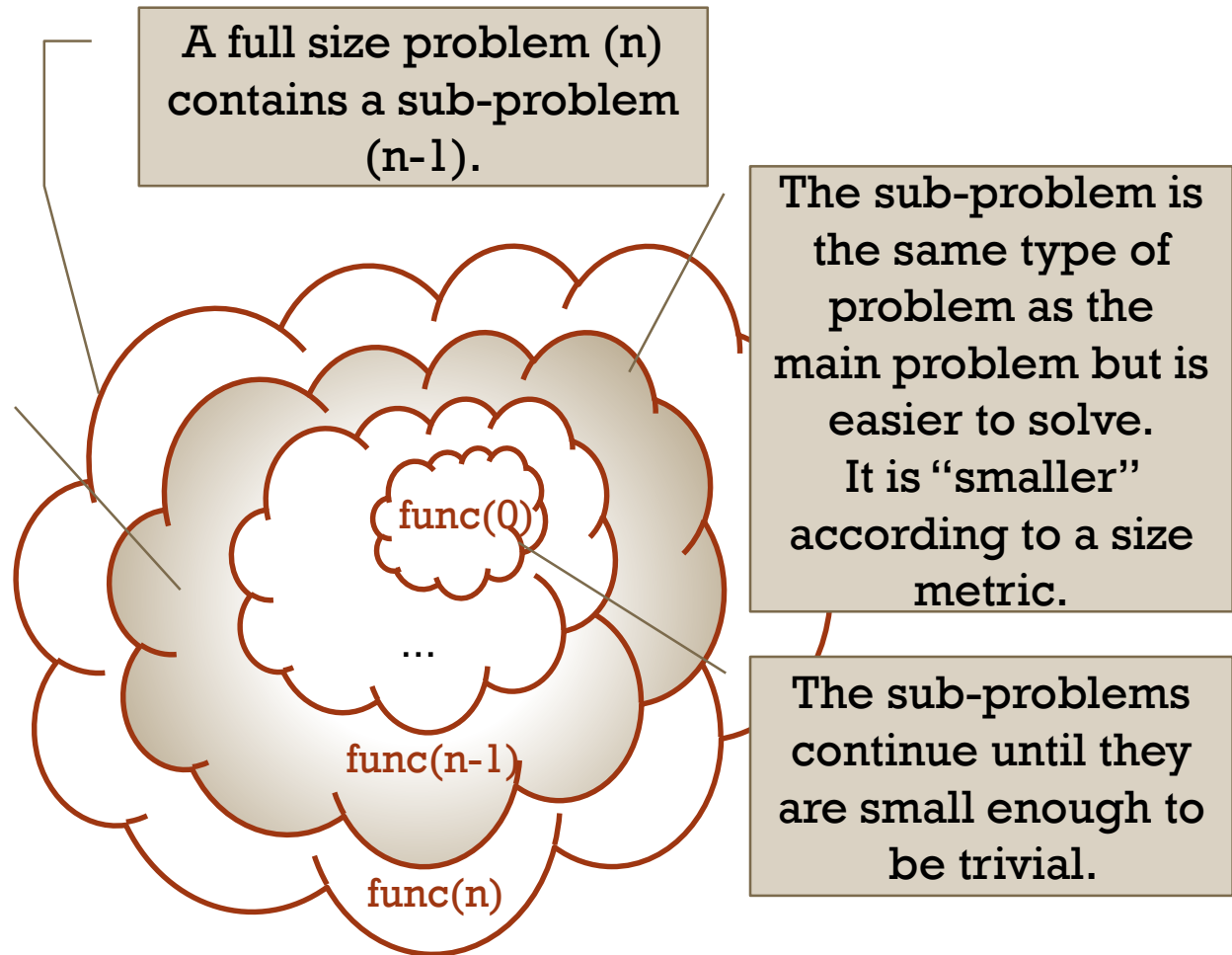
THE RECURSION CONCEPT

... and so I do. And again. And again. Until I've ended up shrinking the smaller cloud so much, I really don't need another one. So, I built a cloud out of a cloud! This is recursion.



RECURSION OVERVIEW

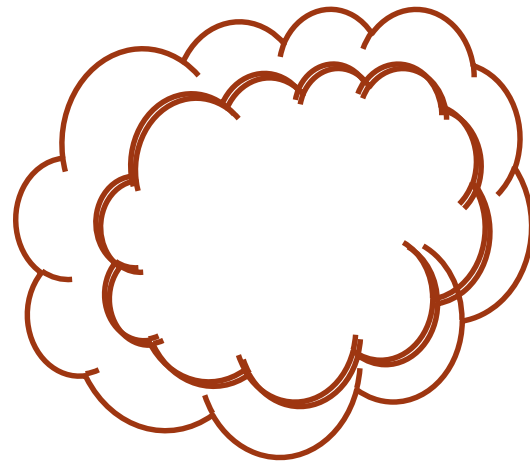
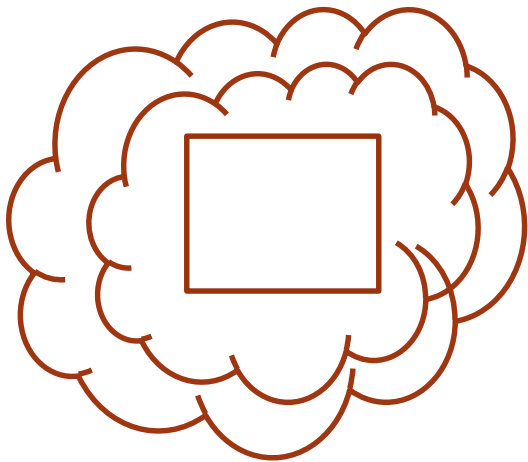
Each time the problem is analyzed, some part of the solution is computed and is combined with the solution for the sub-problem, to form the solution for the whole.



THE KEYS TO RECURSION

For all valid inputs, a recursive method must make the problem smaller during each call, and make **progress** towards some base case (**termination**).

The type of problem should not change, nor the sub-problem be the same as the original problem.



RECURSIVE METHODS

- Recursive structure is an important concept in computer science. **By definition, a method is recursive if it calls itself within the method.**
- There is a special type of recursive structure, which calls itself only once and in the last statement. We refer to this as the **tail-recursion**. **(Tail recursion is very similar to a do-while loop.)**
- Why do we need recursion?

THE FANTASTIC FOUR APPROACH

SCAFFOLDING THE DESIGN OF RECURSIVE METHODS

1. Formulate the size- n problem.
2. Find the stopping condition and the corresponding return value.
3. Formulate the size- m problem and find m . In many cases, $m = n - 1$;
4. Construct the solution of size- n problem from size- m problem.

Let's think about implementing **factorial** in terms of these steps.

(Chen and Tsai. *Introduction to Programming Languages*, Kendall Hunt Publishing.)

1. FORMULATE THE SIZE-N PROBLEM

THE FANTASTIC FOUR ABSTRACT OF WRITING RECURSIVE FUNCTIONS

- Like a loop, recursion is necessary only if you want to solve a problem that needs to repeat the same operations.
- We assume the number of iterations is **n**. In most cases, **n** is obvious. For example, if we want to compute factorial $n!$, the size **n** is already given.
- Formulating the size-**n** problem, in some cases, is merely choosing a function name and using **n** as the parameter of the function. Thus, the size-**n** problem for factorial problem is

int factorial(int n)

- The return value of the size-**n** is what the function is supposed to compute, or the value we are looking for. Obviously, in this step, we do not need to design the solution for size-**n** problem.

2. FIND THE STOPPING CONDITION AND RETURN VALUE

THE FANTASTIC FOUR ABSTRACT OF WRITING RECURSIVE FUNCTIONS

- Like a while-loop, every recursive function starts with checking the stopping condition.
- If the stopping condition is true, we return the corresponding value and exit the function.
- Otherwise, we enter the body of the recursive function.
- In some cases, identifying the stopping condition and corresponding value is trivial. For example, the stopping condition of **factorial(n)** is **n = 0** and the corresponding value is **1**

3. FORMULATE THE SIZE-M PROBLEM

THE FANTASTIC FOUR APPROACH OF WRITING RECURSIVE FUNCTIONS

- The size-m problem is simply the size-n problem with n replaced by m, where $m < n$ is determined by how much we can reduce the size in one step.
- If we can only reduce the problem size by one, m is n-1. For example: factorial(n-1).
- Do not try to define a solution, or the return value in this step! All we need to do here is **assume** the size-m problem will return a value and use the value, e.g., $n * \text{factorial}(n-1)$, to get us to the stopping condition.
- The problem may be resized in many different ways:
 - For the mergesort, $m = \frac{1}{2} n$.
 - For the maze homework, $m = \frac{1}{4} n$.

4. CONSTRUCT THE SOLUTION OF SIZE-N PROBLEM

THE FANTASTIC FOUR APPROACH OF WRITING RECURSIVE FUNCTIONS

- In this step, we will use the **assumed solution** or the return value for size-m or size-(n-1) problem to construct the solution of the size-n problem.
- This step is application-specific. In the case of factorial, the solution of the size-n problem is

`n*factorial(n-1);`

- Sometimes, we need to use the return values of multiple size-m problems, where $0 \leq m < n$ to construct the solution of size-n problem.

PUTTING FACTORIAL TOGETHER

```
int factorial(int n) {  
    if(n == 0)  
        return 1;  
    else  
        return n*factorial(n-1);  
}
```

```
// Mathematical Definition:  
// 1! = 1  
// N! = N * (N-1)!
```

MESSING WITH FACTORIAL

COMPUTING THE NTH FIBONNACI NUMBER

1. Formulate the size-n problem.

`long fib(int n)`

2. Find the stopping condition and the corresponding return value.

`if(n == 0) return 0;`

`if(n == 1) return 1;`

3. Formulate the size-m problem and find m. In many cases, $m = n - 1$;

`fib(n-1)`

`fib(n-2)`

4. Construct the solution of size-n problem from size-m problem.

`fib(n-1) + fib(n-2)`

THE NTH FIBONNACI NUMBER

```
public static long fib(int n) {  
  
    if(n == 0)  
        return 0;  
  
    if(n == 1)  
        return 1;  
  
    return fib(n-1)+ fib(n-2);  
}
```

At this point, recursion probably looks easy. And not that useful. Ha!

As a general rule, recursion is very useful whenever you are faced with a problem you do not know how to solve.

Due to the nature of our world, many problems have a naturally recursive structure, which a recursion algorithm is able to leverage.

PRINTING A LIST

1. Formulate the size-n problem.
`void displayList(LinearNode node)`
2. Find the stopping condition and the corresponding return value.
`if(node == null) and return;`
3. Formulate the size-m problem and find m. In many cases, $m = n - 1$;
`displayList(node.getNext());`
4. Construct the solution of size-n problem from size-m problem.
`System.out.println(node); displayList(node.getNext());`

PRINTING A LIST

```
//iterative
public static void
displayList(ListNode node) {
    ListNode iter = node;
    while(iter != null) {
        System.out.println(iter);
        iter = iter.getNext();
    }
}

//recursive
public static void displayList(ListNode
node) {
    if(node != null) {
        System.out.println(node);
        displayList(node.getNext());
    }
}
```


RECURSIVE DATA STRUCTURES

Here's a thought: we've made programs recursive, so might we make data recursive too?

```
public class Thing {  
    public Thing thing2;  
    private int alongfortheride;  
  
    public Thing(Thing t2, int i) {  
        thing2 = t2;  
        alongfortheride = i;  
    }  
}
```

This is the fundamental idea behind the “linked list” data structure.

- This class is self-referential – another form of recursion.
- This makes a list a recursive data structure.
- Made possible by references.

