Supplement 01.04 - Big-Oh

This supplement is intended to assist in your understanding of Big-Oh and how to figure out the Big-Oh of a formula, and code. Figure 1 is the hierarchal order of Big-Oh. This shows how fast an algorithm will preform. Such as O(1) will be faster than O(n).

$$1 < logn < n < nlog < n^c < c^n < c^{c^n}$$

Figure 1: Hierarchical order of Big-Oh

Finding the Big-Oh from a growth function can be done in two steps. 1) Find the largest factor and dump the other factors. 2) Remove any coefficients. What is left will be the Big-Oh. Figure 2 below shows an example.

Big-Oh method:

$$f(n) = 3n^2 + 5n + 18$$

 $g(n) = 3n^2$
 $g(n) = O(n^2)$

Figure 2: Big-Oh growth function example

Why can we just drop the other factors and the coefficients to find the Big-Oh? Using the proof shown below in Algorithm 1, let k equal the coefficients of f(n). Following the algorithm, we to see that $3n^2$ will quickly overcome the other factors with any large input.

Algorithm 1 The proof for the Big-Oh method showed in Figure 2

Proof:
$$k|n| \ge f(n)$$

 $f(n) = 3n^2 + 5n + 18$
 $k|n^2| \ge f(n) = 3n^2 + 5n + 18$
let $k = 3 + 5 + 18$
 $k|n^2| = 3n^2 + 5n^2 + 18n^2$
 $3n^2 + 5n^2 + 18n^2 \ge 3n^2 + 5n + 18$
 $3n^2 \ge 3n^2$
 $5n^2 \ge 5n$
 $18n^2 \ge 18$

Looking at the last three lines in Algorithm 1, you see that 5n and 18 will quickly be overcome by the largest factor: $3n^2$. So at n goes to infinity, the only factor that will matter is the largest.

How do we get a Big-Oh expression from code? First, a line of code must be chosen to create a benchmark. Below is a simple example of this. Since the method only has code that will be executed once (no loops) this code would be O(1).

```
public static int foobar(int n) {
    System.out.println(n); //Benchmark
}
```

Now, lets add a loop to the code. Where is a good place to set the benchmark? If Benchmark 1 is selected, the Big-Oh would be O(1) which is not an effective way to see how long the code would take to execute. If Benchmark 2 is selected, the Big-Oh would be O(n) since the loop is being included.

So how would we find code that would give $O(\log n)$ if a loop that increments is always going to be O(n)? Well, it depends on *how* is is being incrementing. The code below has a multiple of 4 for the loop. If a large input was given, it would traverse it at a much higher rate than if it was just a normal increment of 1. This creates the $\log(n)$ curve since it is much faster than just n.

Nested loops will give a Big-Oh of $O(n^2)$ assuming you are incrementing at a constant rate for both loops. However, applying what is shown above, a O(nlogn) can be shown in the code below. The inner for loop will be O(logn) due to the multiplier m being the incrementing factor. The outer while loop will continue to execute as a standard incrementing loop which gives O(n). With the loops combined, the O(nlogn) is achieved. The order does not matter in ascertaining the Big-Oh. The while loop could increment by a multiple, and the for loop increment by 1, and it will still be O(nlogn).

The final segment of code shown below has a O(n). The 1500 in the for loop does not matter to the size of the Big-Oh. Although n must be a large input before the print statement will be executed, the code segment below is still O(n). Remember that the Big-Oh is the upper-bound and constants do not have any influence over it.