

# Assignment 1

## COMP 206

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### Instructions

Submit a single pdf file named yourFirstName.pdf

There are 9 questions. Q1-Q8 are for 10 marks each, and Q9 is for 20 marks. All the programs must be written in  $C_0$ .

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1. Provide a one line answer for the following:

- (a) What is the type of an array of boolean values? **bool [ ]**
- (b) Declare an int array of size 100 and store it in a variable named **x**.
- (c) What is the use of **-d** flag in the **coin -d** command? **enables contracts.**
- (d) What is the type of a function with no return value? **void**
- (e) Enumerate all the constants of type **bool**. **true/false**

2. Write a function named **abs** which takes an int and returns its absolute value.

$$\text{abs}(x) = \begin{cases} x & \text{if } x \geq 0 \\ -x & \text{if } x < 0 \end{cases}$$

3. Write a function named **log**, which is defined as follows for int values:

$$\text{log}(x) = \begin{cases} 0 & \text{if } x = 0 \\ 1 + \text{log}(x/2) & \text{if } x \geq 1 \\ ??? & \text{undefined otherwise} \end{cases}$$

The function should have a precondition contract to handle the undefined case. Note that the division operator used above is the same as  $C_0$  division.

4. Write a function named **max2**, which takes two positive ints and returns their maximum. Write a precondition to make sure that the caller supplied positive numbers. Write a postcondition for the fact that when  $x$  is the maximum of two numbers  $a$  and  $b$ , then we have  $x \geq a$  and  $x \geq b$ .
5. Write a postcondition for the following function which adds two numbers. The postcondition should check that the returned value is **a+b**.

2. int abs (int x){  
    if ( $x \geq 0$ ) {  
        return x;  
    }  
    else {  
        return -x;  
    }  
}

3. int log (int x)  
// requires  $x \geq 0$ ;  
{ if ( $x == 0$ ) {  
    return 0;  
}  
else {  
    return 1 + log( $x/2$ );  
}  
}

4)

```
int max2( int a, int b )
```

//@ requires  $a > 0 \wedge b > 0$ ;

{ //@ ensures  $\text{result} \geq a \wedge \text{result} \geq b$ ;

//@ ensures  $\text{result} = a$

$\text{result} = b$ ;

{ if ( $a > b$ ) {

    return a;

}

else {

    return b;

}

{

logical  
OR.



5) int add (int a, int b)  
// @ensures result == a+b;  
{  
 return a+b;  
}

6) put extra param in  
func :-

bool binarySearch( int [ ] a,  
 int item,  
 int len )

       and use the foll :-

False  $\leftrightarrow$  false

not found  $\leftrightarrow$  !found

```
int add(int a, int b) {  
    a = a+b;  
    return a;  
}
```

6. Convert the following python function (line by line) to search an integer array in  $C_0$ . As there is no len function in  $C_0$  to find the length of an array, write a function which takes length of the array also as an extra argument.

```
def binarySearch(alist, item):  
    """  
    alist: list to search in  
    item: element to search for  
    """  
  
    first = 0  
    last = len(alist)-1  
    found = False  
  
    while first<=last and not found:  
        mid = (first + last)//2  
        if alist[mid] == item:  
            found = True  
        else:  
            if item < alist[mid]:  
                last = mid-1  
            else:  
                first = mid+1  
  
    return found
```

7. Write a proof for the loop invariant  $b^e r = x^y$  when  $e$  is an even number as discussed in the class on 27th August 2020. You can collaborate with classmates for this question only, but write your own proof after discussion (copying is not allowed). Also mention the names of your classmates you discussed with.
8. (a) Rewrite the following code without using the `while` keyword by converting it into a `for` loop.

```
int x = 1;  
int t = 0;  
while (x<10) {  
    t = t + x;  
    x = x+1;  
}
```

- (b) Rewrite the following code without using the `for` keyword by converting it into a `while` loop.

7) Modify the proof  
in notes, about  
odd nos., saved  
on canvas.

8) The question is about  
conversion for  $\leftrightarrow$  while  
loops.

```
[for (s1; s2; s3){}  
    body;  
}
```

```
    }  
    s1;  
    while (s2){  
        body;  
        s3;
```

use this idea  
to convert  
while  $\rightarrow$  for.

```

int x;
for (x=1; x<10 && x*x<5; x=x+1) {
    x = x+1;
}

```

9. The function `sumInt` takes an int  $n$  as input and returns the sum of ints from 1 to  $n - 1$ . You are already provided with the pre and postconditions along with the loop invariants. By using the code provided to you answer the following.

- (a) Explain informally why the loop terminates.
- (b) **INIT Step:** Prove that each loop invariant is true immediately before the loop condition is tested for the first time.
- (c) **PRES Step:** Prove that if each loop invariant is true at the start of a loop iteration, then the loop invariants are also all true at the end of that iteration.
- (d) Show that if the loop terminates, the postcondition must hold.

```

int sumInt(int n)
//@requires n > 0;
//@ensures 2 * \result == n * (n - 1);
{
    int sum = 0;
    int i = 0;
    while (i < n)
        //@loop_invariant 0 <= i && i <= n;
        //@loop_invariant 2 * sum == i * (i - 1);
    {
        sum = sum + i;
        i = i + 1;
    }
    return sum;
}

```

- g) a)
- ①  $i = 0$  by L4.
  - ②  $i$  is strictly increasing by L9.
  - ③ ZG fails when  $i \geq n$

Combining ①, ②, ③ and noticing  
that a strictly increasing  
int sequence is finite  
when it is upper bounded.

9b) INIT :-

L I 1 :-

T.P.  $0 \leq i \leq n$  is true before  
LG is reached.

Proof:

①  $i = 0$  by L4

②  $0 \leq 0$  fact

①+②  $\Rightarrow 0 \leq i$ , first half of ineq.  
proved.

③  $n > 0$  by C1.

①+③  $\Rightarrow n > i$

$\Rightarrow n >= i$ , LI 1 proved.

L12:-

$$2\text{sum} = i(i-1).$$

Proof :-

①  $i=0$  by L4

②  $\text{sum}=0$  by L3

③  $2\text{sum}=0$  by ②

④  $i * (i-1)=0$  by ①

by ③ + ④ we have

$$2\text{sum} = i(i-1)$$

Hence proved.

### 9c) PRES step :-

L I 1 ,

- Assumptions**
- (1) Assume  $0 \leq i' \leq n$  before the loop is executed.
  - (2) assume Value of  $i^o$  is  $i'$  at the end of the loop .
  - (3) notice  $n$  doesn't change in the loop Body .

(4)  $i' = i^o + 1$  by LG .

(5)  $i' > i$  using (4)

(1)+(5) gives us  $0 < i'$   
 $0 \leq i'$  using (6)

(7)  $i < n$  by LG as we are going to execute the loop :

(8)  $i + 1 \leq n$  using (8)

(9)  $i' \leq n$  using (4)

(7)+(10) proves L I 1 .

## LI2 : PRES step :

Assumptions

(1) Assume  $2\sum = i(i-1)$   
 (2) assume  $i, \sum$  get changed  
 to  $i', \sum'$  by the  
 end of the loop.

$$(3) i' = i+1 \text{ by L9}$$

$$(4) \sum' = \sum + i \text{ by L8}$$

$$(5) 2\sum' = 2\sum + 2i \text{ using (4)}$$

$$(6) 2\sum' = i(i-1) + 2i \text{ using (1) and (5)}$$

$$\begin{aligned} (7) 2\sum' &= i(i-1+2) \text{ using (6)} \\ &= i(i+1) \\ &= (i'-1)i' \text{ using (3)} \end{aligned}$$

$$\therefore 2\sum' = i'(i'-1)$$

Hence proved.

9d)

The loop terminates when  $LG.$  is false.

① This gives us that  $i < n$  is false i.e.  $i \geq n$ .

② LI1 gives us that  $i \leq n$ .  
 $\therefore n \leq i \leq n$  by ①+②  
 $\Rightarrow i = n$

③ Outside the loop at L10  
we have  $i = n$ .

④  $\backslash result = sum$  by L10

⑤ By LI2 & ③ we have  
 $sum = i(i-1) = n(n-1)$

⑥  $2 \backslash result = n(n-1)$  by ④  
 $\therefore C2$  holds. Hence proved.

# Attention!!

Are we done?

Or does the picture  
change because of the  
fact that int is  
Co one upper

bounded by  $2^{31} - 1$ ?

The proofs above  
work well when  
ints are unbounded.

Remember that overflow  
will occur when nos.  
become larger than  $2^{31} - 1$ .

Exercise :- assignment 2 ?

Can you find out  
the places where our  
proofs fail for 32 bit  
ints? Which of these  
places are fixable?

1. we may have a bug  
at each place the proofs  
fail.

Q. Debug the program!  
(Hint: contracts??).