

# Generic Pointers

# **Generic Data Structures**

# Stacks

- We defined stacks of **strings**
- But,
  - the code for stacks of **ints** would be identical except for **string** changed to **int**
  - the code for stacks of *any type* would be identical except for **string** changed to this type
  - ...

```
/****** Implementation *****/

typedef struct list_node list;
struct list_node {
    string data;
    list* next;
};

typedef struct stack_header stack;
...

/****** Interface *****/

// typedef _____* stack_t;

bool stack_empty(stack_t S)
/* @requires S != NULL; @*/ ;

stack_t stack_new()
/* @ensures \result != NULL; @*/
/* @ensures stack_empty(\result); @*/ ;

void push(stack_t S, string x)
/* @requires S != NULL; @*/
/* @ensures !stack_empty(S); @*/ ;

string pop(stack_t S)
/* @requires S != NULL; @*/
/* @requires !stack_empty(S); @*/ ;
```

# Stacks

- Each time we need a stack for a new element type, we need to make a **copy** of the stack library
- This is **bad**
  - It's easy to make a mistake
  - We need to come up with new names
    - `int_stack_t`, `int_stack_empty`, ...
    - `int_list`, `int_list_node`, `is_int_segment`, ...
  - If we discover a bug, we need to fix it in **every copy** of the library
    - same if we discover a better implementation
- For a large application, this quickly becomes unmanageable

```
/****** Implementation *****/

typedef struct list_node list;
struct list_node {
    string data;
    list* next;
};

typedef struct stack_header stack;
...

/****** Interface *****/

// typedef _____* stack_t;

bool stack_empty(stack_t S)
/* @requires S != NULL; @*/ ;

stack_t stack_new()
/* @ensures \result != NULL; @*/
/* @ensures stack_empty(\result); @*/ ;

void push(stack_t S, string x)
/* @requires S != NULL; @*/
/* @ensures !stack_empty(S); @*/ ;

string pop(stack_t S)
/* @requires S != NULL; @*/
/* @requires !stack_empty(S); @*/ ;
```

# Generic Data Structures

- Stacks are intrinsically **generic data structures**
  - They work the same way no matter the type of their elements
  - They do not modify elements
    - they only store them in the data structure and give them back
- We would like to implement them as a **generic library**
  - a **single** stack implementation that can be used for elements of any type
    - without copying it over and over
    - without a proliferation of function and type names
  - if we find a bug, there is **one place** where to fix it
    - if we are told of a better implementation, there is **one file** to change

# Generic Stacks -- Take 1

- Here's an idea:
  - use a generic type name **elem** in the library
  - let the client define what **elem** is
- We note the type **elem** is to be defined by the client in the **client interface**
- The client needs to define what **elem** actually is in the **client definition code**

```
/****** Client definitions *****/  
  
typedef string elem;
```

```
/****** Client Interface *****/  
// typedef _____ elem;  
  
/****** Implementation *****/  
typedef struct list_node list;  
struct list_node {  
    elem data;  
    list* next;  
};  
  
typedef struct stack_header stack;  
...  
  
/****** Interface *****/  
  
// typedef _____* stack_t;  
  
bool stack_empty(stack_t S)  
/* @requires S != NULL; @*/ ;  
  
stack_t stack_new()  
/* @ensures \result != NULL; @*/  
/* @ensures stack_empty(\result); @*/ ;  
  
void push(stack_t S, elem x)  
/* @requires S != NULL; @*/  
/* @ensures !stack_empty(S); @*/ ;  
  
elem pop(stack_t S)  
/* @requires S != NULL; @*/  
/* @requires !stack_empty(S); @*/ ;
```

# Generic Stacks -- Take 1

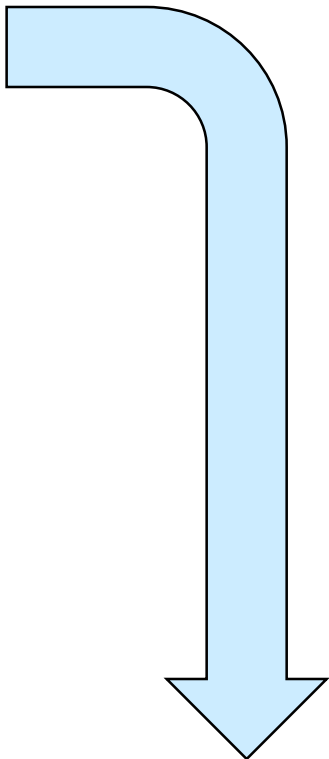
## Pros:

- A **single** library for any kind of stack

- If the client needs a stack of **ints** in a different application,

➤ simply define **elem** as **int**

- If another application requires a different stack type, just define **elem** appropriately



```
/****** Client definitions *****/  
  
typedef int elem;
```

```
/****** Client Interface *****/  
// typedef _____ elem;  
  
/****** Implementation *****/  
typedef struct list_node list;  
struct list_node {  
    elem data;  
    list* next;  
};  
  
typedef struct stack_header stack;  
...  
  
/****** Interface *****/  
  
// typedef _____* stack_t;  
  
bool stack_empty(stack_t S)  
/* @requires S != NULL; @*/ ;  
  
stack_t stack_new()  
/* @ensures \result != NULL; @*/  
/* @ensures stack_empty(\result); @*/ ;  
  
void push(stack_t S, elem x)  
/* @requires S != NULL; @*/  
/* @ensures !stack_empty(S); @*/ ;  
  
elem pop(stack_t S)  
/* @requires S != NULL; @*/  
/* @requires !stack_empty(S); @*/ ;
```

# Generic Stacks -- Take 1

## Cons:

- Client application has to be split into **two** files

➤ Client definition file

```
/* Client definitions */  
typedef string elem;
```

➤ Rest of the client application

```
/* Client application */  
  
int main() {  
    ... push ... pop ...  
}
```

## ○ because

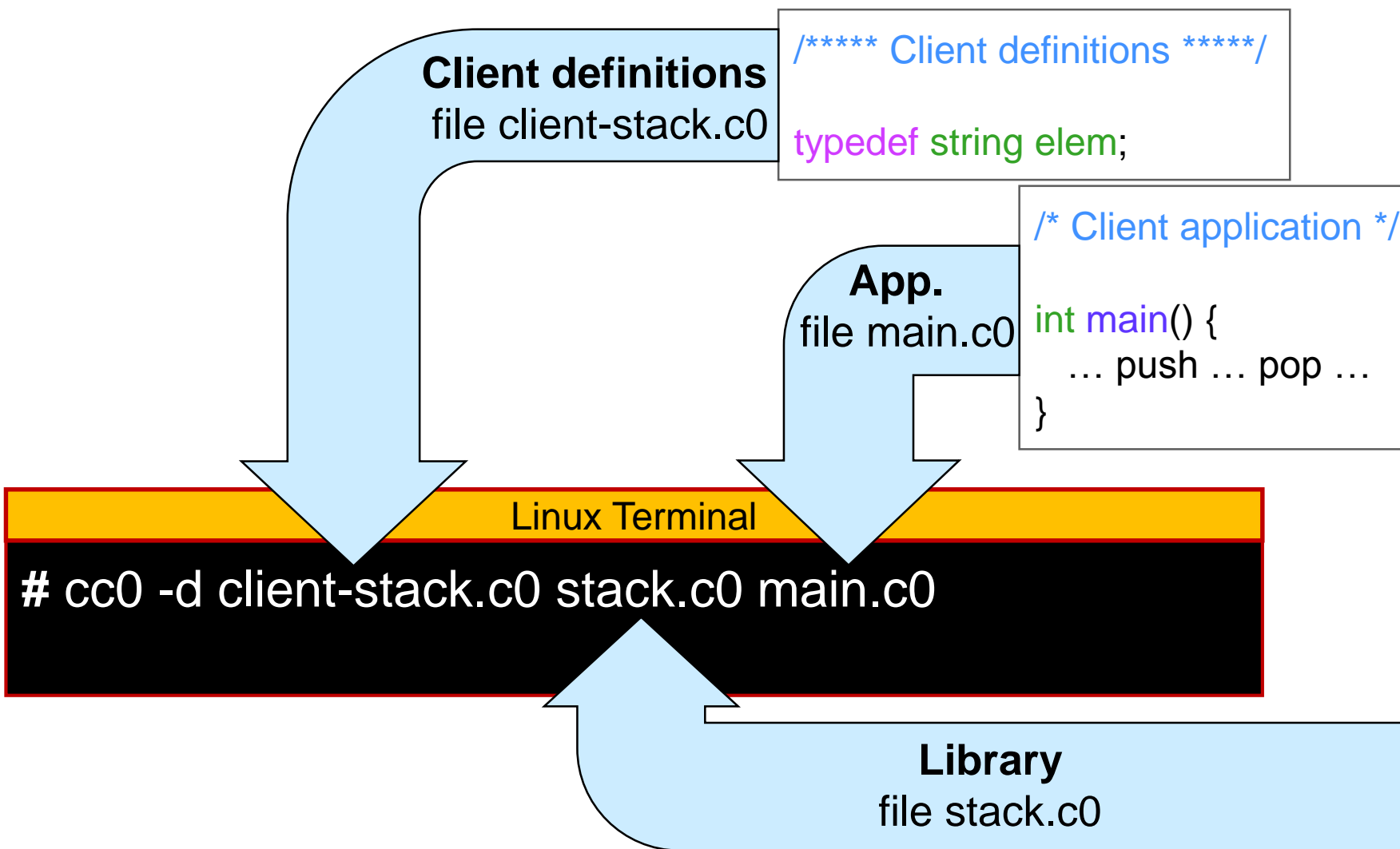
- the library needs **elem** to be defined
  - ❑ This must occur **before** the library
- the client application needs the types and functions provided by the library to be defined
  - ❑ This must occur **after** the library

```
/* Client Interface */  
// typedef _____ elem;  
  
/* Implementation */  
typedef struct list_node list;  
struct list_node {  
    elem data;  
    list* next;  
};  
  
typedef struct stack_header stack;  
...  
  
/* Interface */  
  
// typedef _____* stack_t;  
  
bool stack_empty(stack_t S)  
/* @requires S != NULL; @*/ ;  
  
stack_t stack_new()  
/* @ensures \result != NULL; @*/  
/* @ensures stack_empty(\result); @*/ ;  
  
void push(stack_t S, elem x)  
/* @requires S != NULL; @*/  
/* @ensures !stack_empty(S); @*/ ;  
  
elem pop(stack_t S)  
/* @requires S != NULL; @*/  
/* @requires !stack_empty(S); @*/ ;
```

This is mildly annoying



# Generic Stacks -- Take 1



```
/* Client Interface */
// typedef _____ elem;

/* Implementation */
typedef struct list_node list;
struct list_node {
    elem data;
    list* next;
};

typedef struct stack_header stack;
...

/* Interface */

// typedef _____* stack_t;

bool stack_empty(stack_t S)
/* @requires S != NULL; @*/ ;

stack_t stack_new()
/* @ensures \result != NULL; @*/
/* @ensures stack_empty(\result); @*/ ;

void push(stack_t S, elem x)
/* @requires S != NULL; @*/
/* @ensures !stack_empty(S); @*/ ;

elem pop(stack_t S)
/* @requires S != NULL; @*/
/* @requires !stack_empty(S); @*/ ;
```

- This forces an **unnatural compilation pattern**

This is mildly annoying

# Generic Stacks -- Take 1

## Cons:

- Client application can contain at most **one type** of stacks
  - no way to have both a stack of **strings** and a stack of **ints** in the **same** application
    - but we can have multiple stacks of **ints**
  - because there can be only **one definition** for **elem**

This is a big deal

```
/***** Client definitions *****/  
  
typedef string elem;  
typedef int elem;
```

?

The compiler won't know which **elem** to use when

Compilation error!

```
/****** Client Interface *****/  
// typedef _____ elem;  
  
/****** Implementation *****/  
typedef struct list_node list;  
struct list_node {  
    elem data;  
    list* next;  
};  
  
typedef struct stack_header stack;  
...  
  
/****** Interface *****/  
  
// typedef _____* stack_t;  
  
bool stack_empty(stack_t S)  
/* @requires S != NULL; @*/ ;  
  
stack_t stack_new()  
/* @ensures \result != NULL; @*/  
/* @ensures stack_empty(\result); @*/ ;  
  
void push(stack_t S, elem x)  
/* @requires S != NULL; @*/  
/* @ensures !stack_empty(S); @*/ ;  
  
elem pop(stack_t S)  
/* @requires S != NULL; @*/  
/* @requires !stack_empty(S); @*/ ;
```

# Generic Stacks -- Take 1

## Summary

### ● Pros:

- A single library for any kind of stacks

### ● Cons:

- Client application is split into **two** files
  - Unnatural compilation pattern
- Client application can contain at most **one type** of stacks

This is mildly annoying

This is a big deal

```
/****** Client Interface *****/
// typedef _____ elem;

/****** Implementation *****/
typedef struct list_node list;
struct list_node {
    elem data;
    list* next;
};

typedef struct stack_header stack;
...

/****** Interface *****/

// typedef _____* stack_t;

bool stack_empty(stack_t S)
/* @requires S != NULL; @*/ ;

stack_t stack_new()
/* @ensures \result != NULL; @*/
/* @ensures stack_empty(\result); @*/ ;

void push(stack_t S, elem x)
/* @requires S != NULL; @*/
/* @ensures !stack_empty(S); @*/ ;

elem pop(stack_t S)
/* @requires S != NULL; @*/
/* @requires !stack_empty(S); @*/ ;
```

# Can we do Better?

- Not in C0 ...



- ... but the language C1 extends C0 with a mechanism to address these issues



**C1**

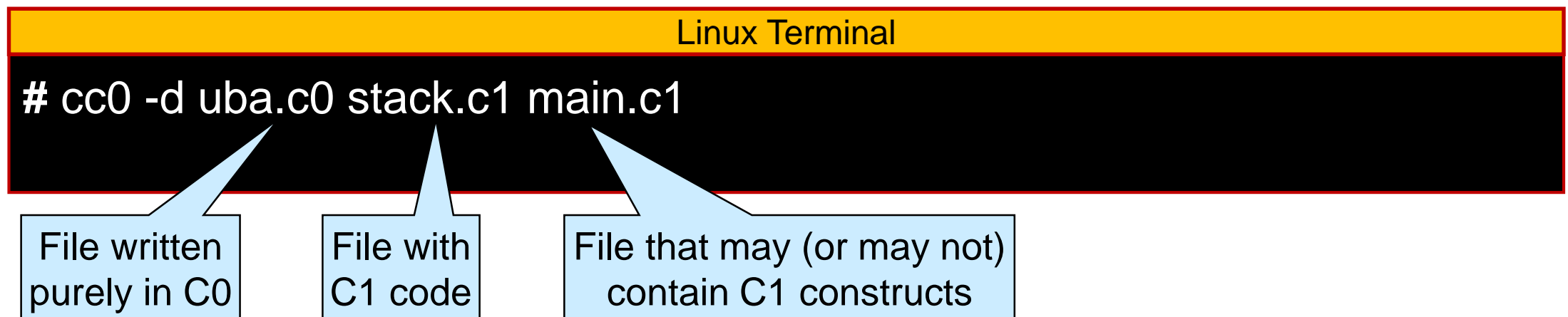
# The language C1

- C1 is an **extension** of C0
  - Every C0 program is a C1 program
- C1 provides two additional mechanisms
  - Generic pointers
  - Function pointers

Both help with genericity
- Right now, we will only examine generic pointers

# Running C1 Programs

- C1 programs are compiled with cc0
  - but C1-only constructs are only allowed in files with a **.c1** extension
    - C0-only code can appear in files with either a .c0 or a .c1 extension
- Example



- The coin interpreter does not currently support C1 constructs
  - no way to experiment with them in coin

# Generic Pointers



# void\*

This is **not** a pointer to **void**:  
**void** is not a type  
*Blame C for the confusing name!*

- C1 provides a new **pointer type**: **void\***

**void\*** **q**; ————— q is a variable of type **void\***

- a value of this type is a **generic pointer**

- Any pointer can be turned into a **void\*** using a **cast**

**int\*** **p** = alloc(**int**);

\*p = 7;

**q** = (**void\***)p;

Cast p to **void\***

q still has type **void\***,  
but contains the same address of p

- and later back to its original type

**int\*** **r** = (**int\***)q;

Cast q to **int\***

q still has type **void\***,  
but r contains the same address  
as p

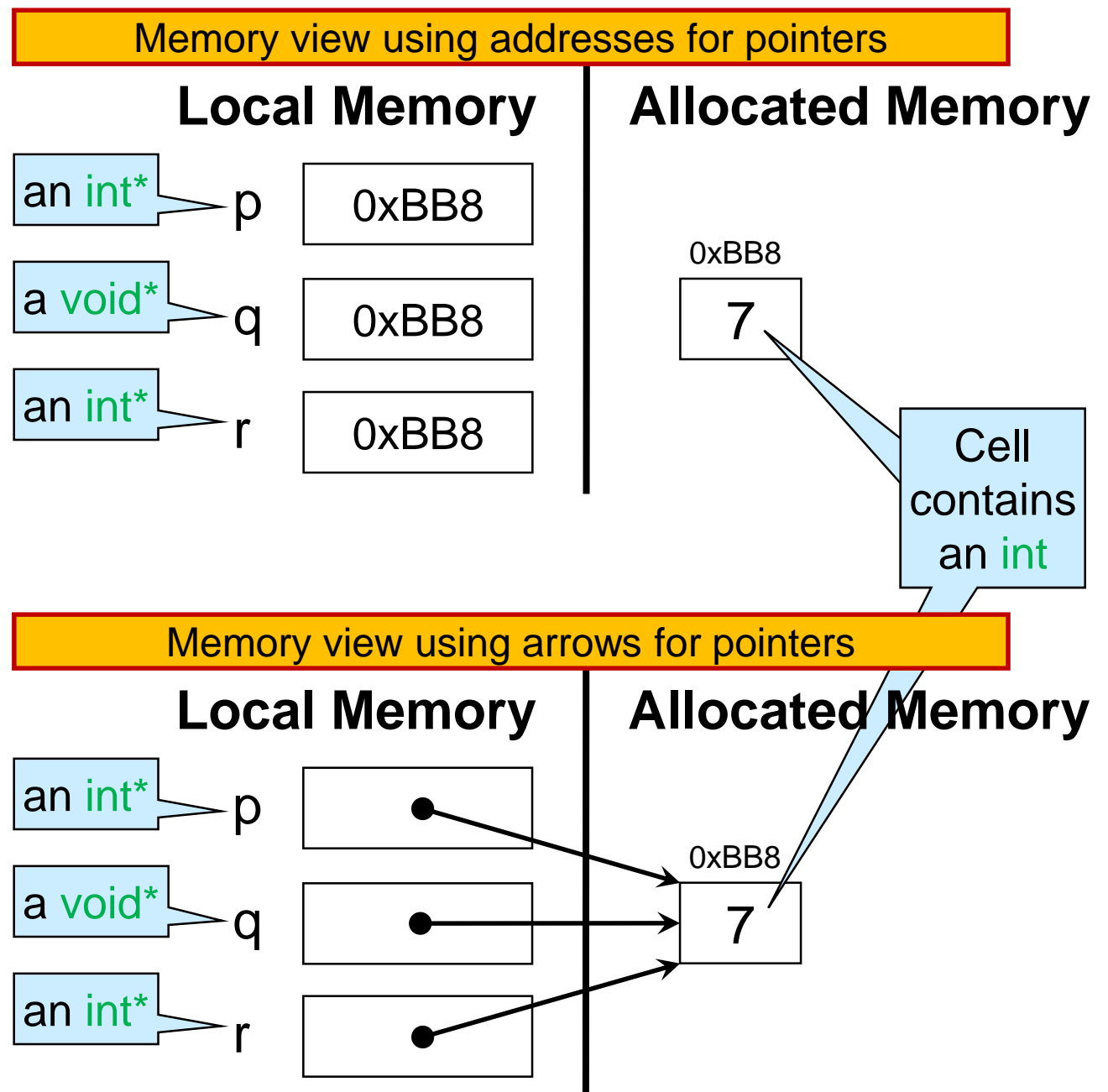
# void\*

```
void* q;  
int* p = alloc(int);  
*p = 7;  
q = (void*)p;  
int* r = (int*)q;
```

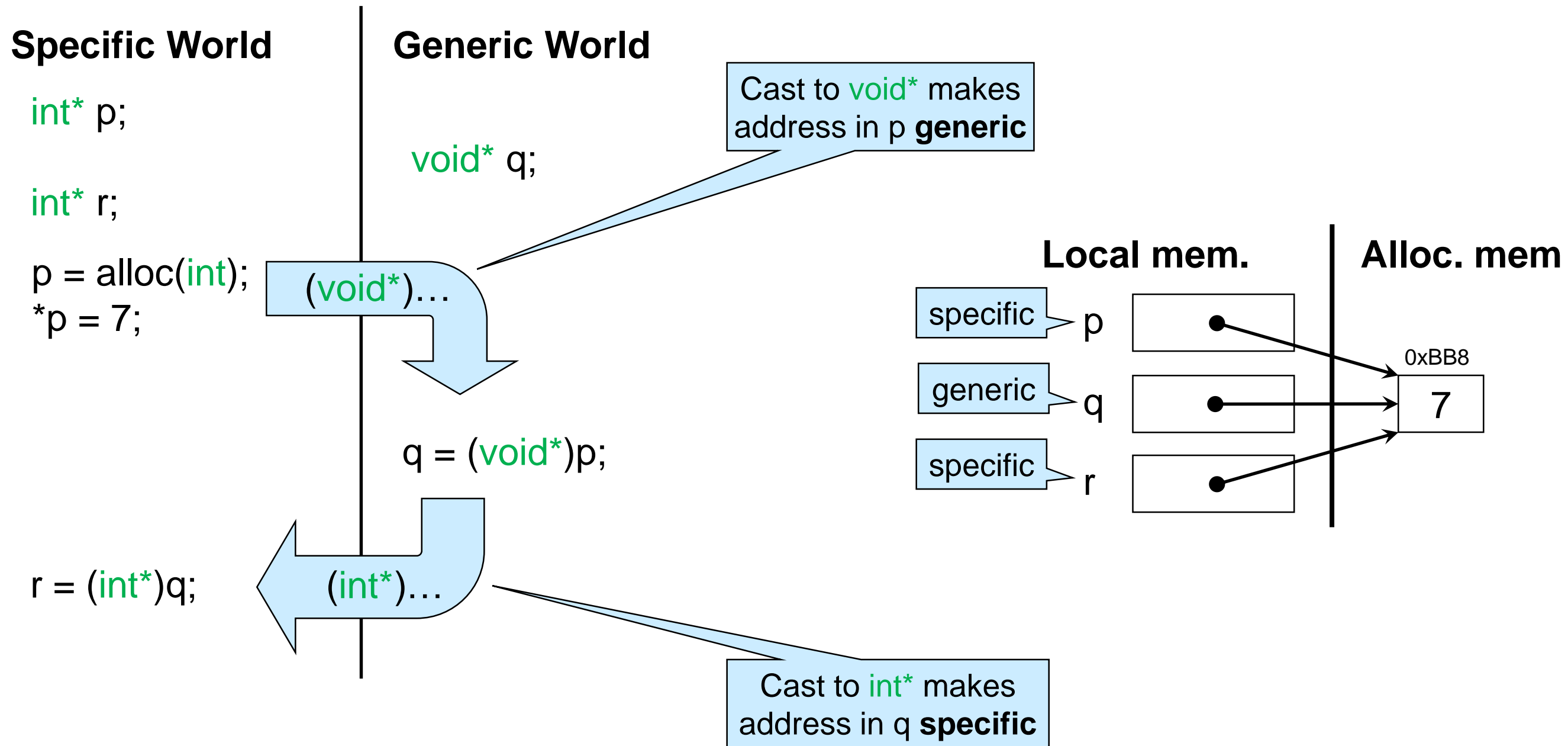
- p, q and r contain the same address
  - they are aliases
- but

- p and r have type `int*`
- q has type `void*`

- With casting, we can pretend that a **specific** pointer (e.g., an `int*`) is a **generic** pointer (`void*`)
  - a controlled way for a pointer to have two types
  - *only for pointers*



# The Specific/Generic Divide



# What can we do with `void*` Pointers?

## Allowed

- Cast to original type

```
int* p = alloc(int);
```

```
void* q1 = (void*)p;
```

```
int* r = (int*)q1; ✓
```

- Compare for equality

```
void* q2 = (void*)alloc(int);
```

```
if (q1 == q2) println("same"); ✓
```

- Assign to a `void*` variable

```
➤ void* q3 = q1; ✓
```

## Not Allowed

- Dereference

```
void x = *q1; ✗
```

○ `void` is **not** a type in C0/C1/C

- Allocate

```
void* q4 = alloc(void); ✗
```

○ `void` is not a type

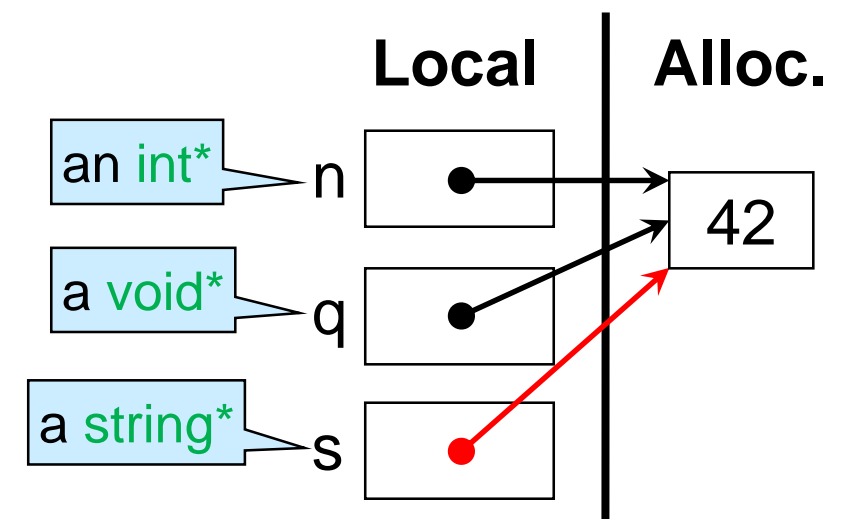
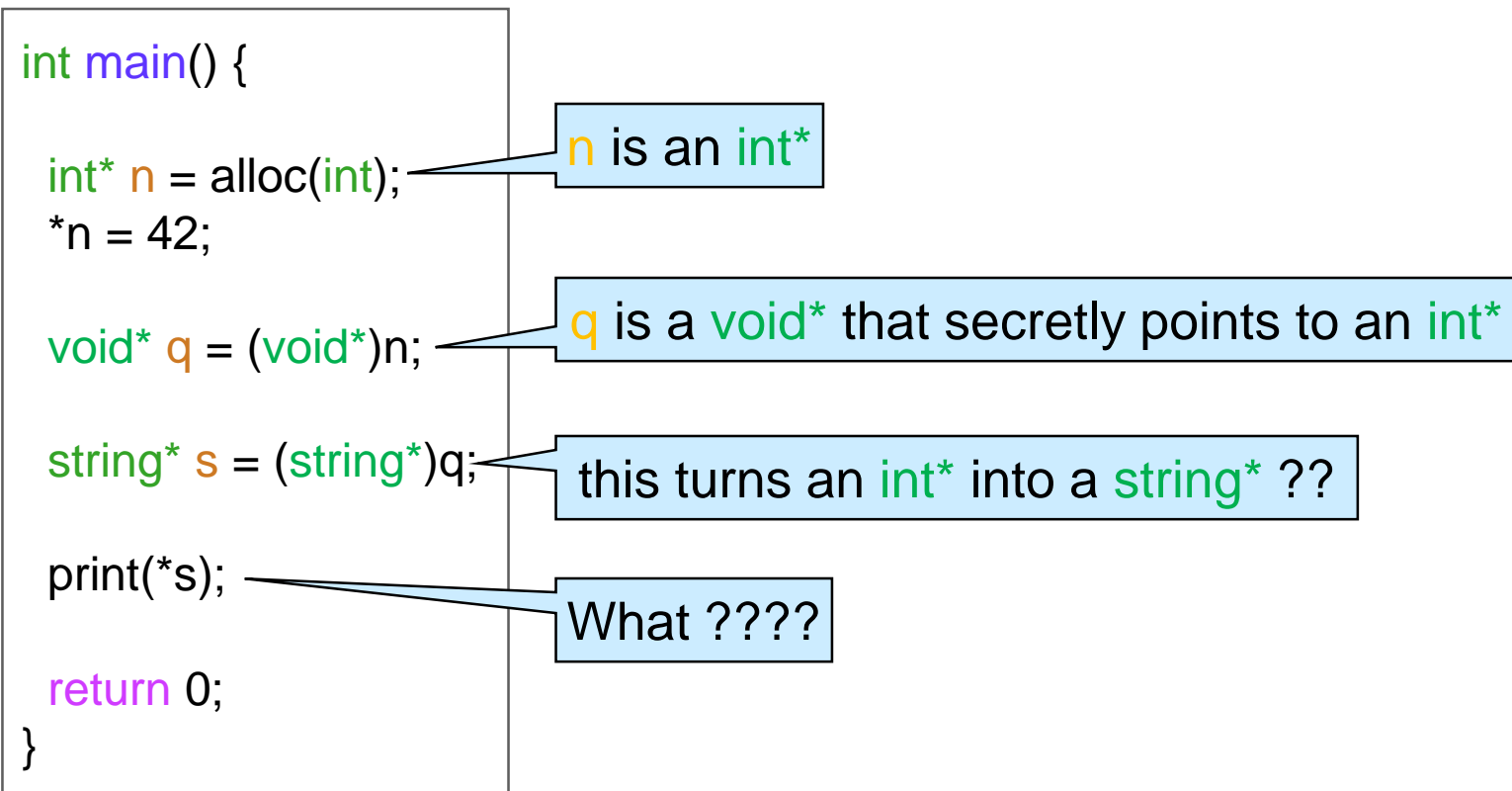
- Cast to type other than original

```
println((string*)q1; ✗
```

○ *(see next)*

# **Safety of Generic Pointers**

# Casting back to the Wrong Type



- This makes no sense!!!
    - dereferencing **s**, we get to an **int** (42)
    - but print expects a **string**
- This doesn't feel right
- a safety violation maybe?

# Casting back to the Wrong Type

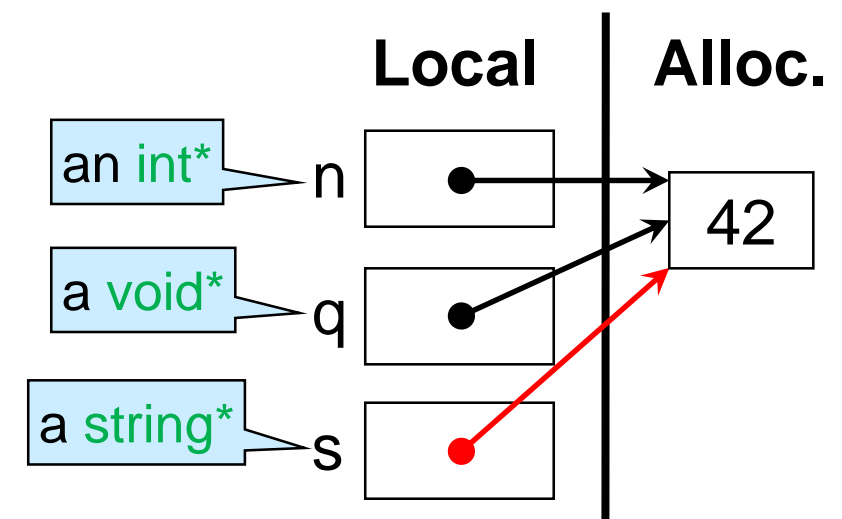
```
int main() {  
    int* n = alloc(int);  
    *n = 42;  
    void* q = (void*)n;  
    string* s = (string*)q;  
    print(*s);  
    return 0;  
}
```

`n` is an `int*`

`q` is a `void*` that secretly points to an `int*`

this turns an `int*` into a `string*` ??

What ????



- Let's run it

```
Linux Terminal  
# cc0 -d bad-casting.c1  
# ./a.out  
untagging pointer failed  
Segmentation fault (core dumped)
```

- We get a memory error
- This **is** a safety violation

# Tags

```

int main() {
    int* n = alloc(int);
    *n = 42;

    void* q = (void*)n;
    string* s = (string*)q;
    print(*s);

    return 0;
}

```

**n** is an **int\***

**q** is a **void\*** that secretly points to **n**

this turns an **int\*** into a **string\***

What ????

Linux Terminal

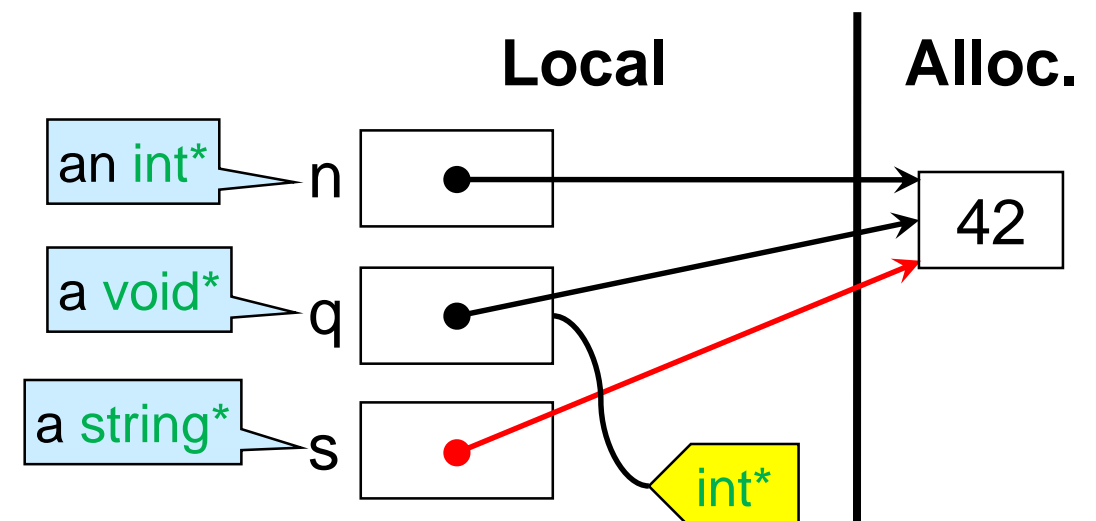
```

# cc0 -d bad-casting.c1
# ./a.out
untagging pointer failed
Segmentation fault (core dumped)

```

## ● *Untagging pointer failed?*

- At run time, values of type **void\*** carry a **tag** that records the original type of the pointer
- C1 checks that the tag is correct before casting back

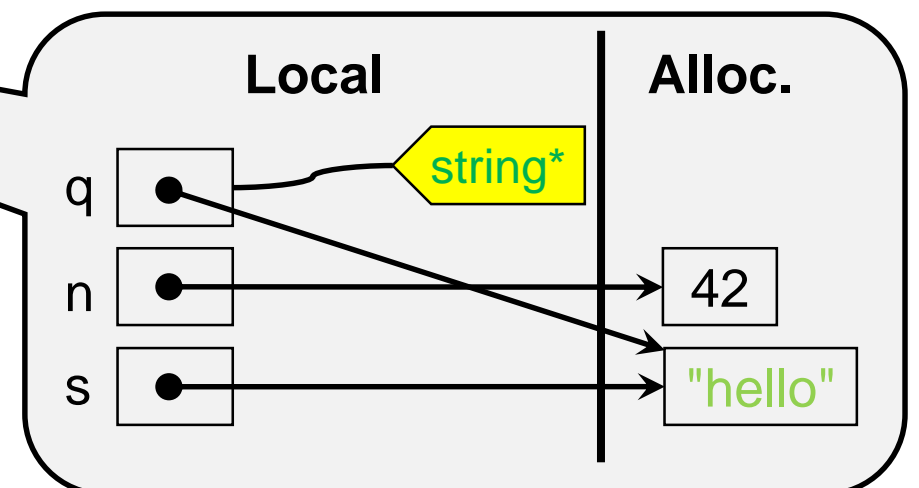
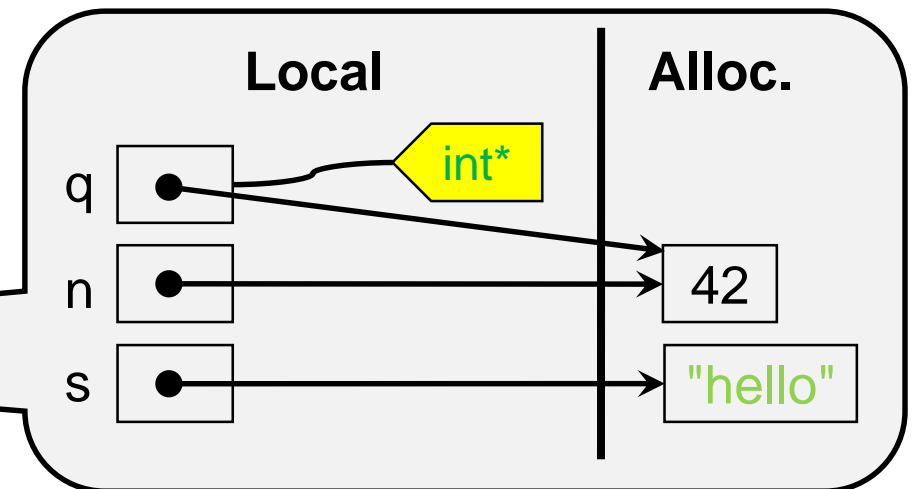
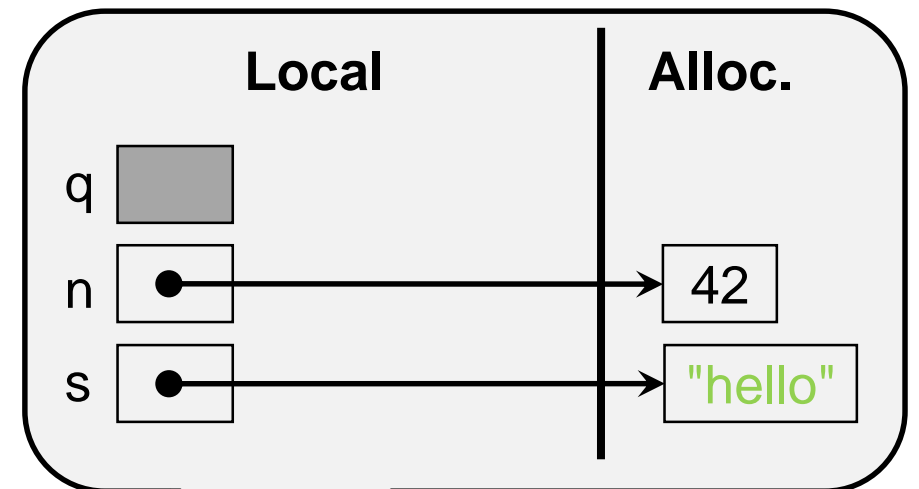




# Tags

- The tag of a **void\*** changes as execution proceeds

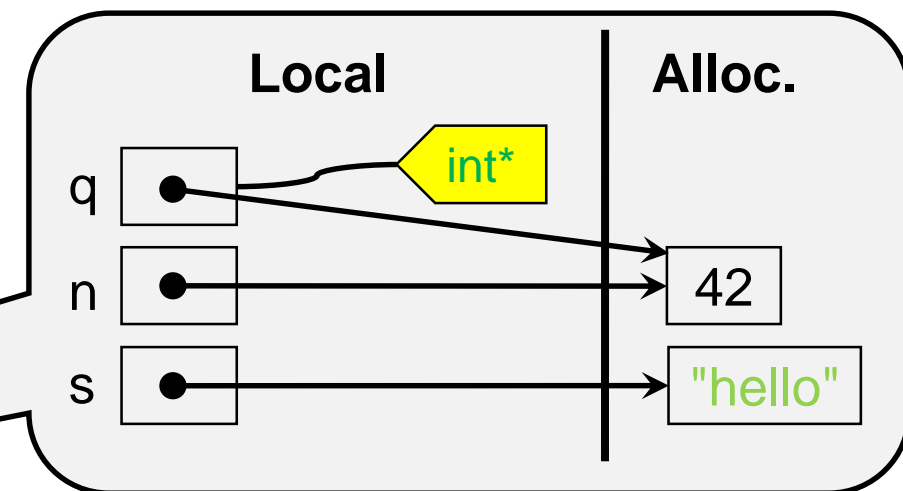
```
int main() {  
    void* q;  
  
    int* n = alloc(int);  
    *n = 42;  
  
    string* s = alloc(string);  
    *s = "hello";  
  
    q = (void*)n;  
    printint(*(int*)q);  
  
    q = (void*)s;  
    print(*(string*)q);  
  
    return 0;  
}
```



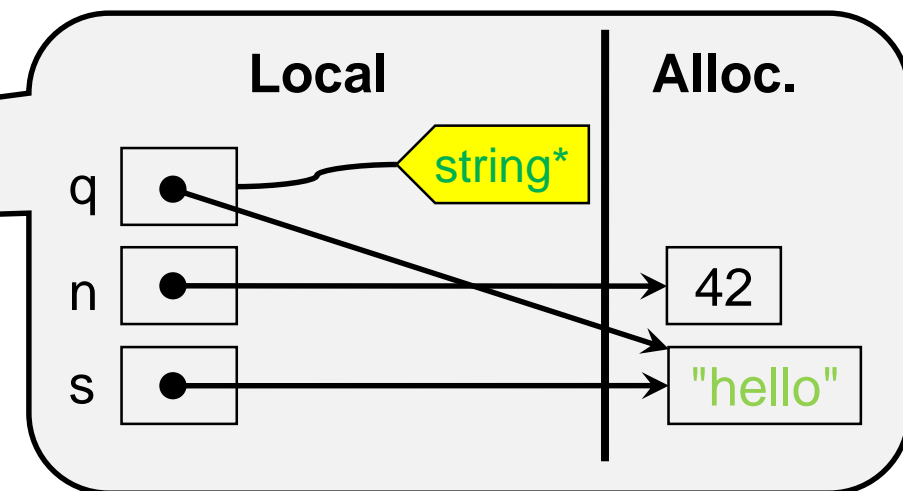
# \hastag

- Annotation-only function `\hastag(tp, ptr)` can be used to check that generic pointer `ptr` has type `tp` *in debugging mode*

```
int main() {  
    void* q;  
  
    int* n = alloc(int);  
    *n = 42;  
  
    string* s = alloc(string);  
    *s = "hello";  
  
    q = (void*)n;  
    //@assert \hastag(int*, q);  
    printint(*(int*)q);  
  
    q = (void*)s;  
    //@assert \hastag(string*, q);  
    //@assert !\hastag(int*, q);  
    print(*(string*)q);  
  
    return 0;  
}
```



Checks that q has tag `int*`



Checks that q now has tag `string*`

Checks that q does not have tag `int*`

# \hastag

- Annotation-only function `\hastag(tp, ptr)` can be used to check that generic pointer `ptr` has type `tp` *in debugging mode*

```
int main() {
```

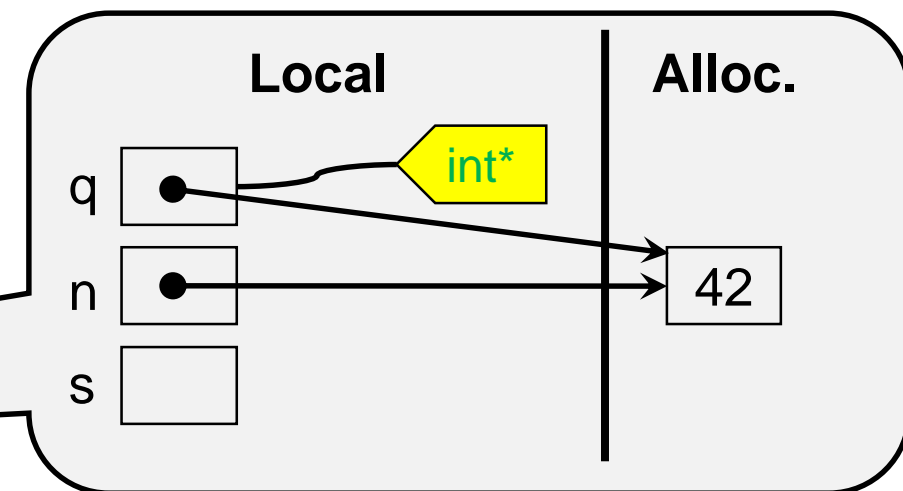
```
    int* n = alloc(int);  
    *n = 42;
```

```
    void* q = (void*)n;  
    //@assert \hastag(int*, q);
```

```
    //@assert \hastag(string*, q);  
    string* s = (string*)q;
```

```
    print(*s);
```

```
    return 0;  
}
```



Checks that q has tag `int*` ✓

Checks that q has tag `string*` ✗

- Use `\hastag` before casting a `void*` back to a specific type

# NULL

- NULL is a pointer of any type, including `void*`

- We can cast NULL back and forth as we please

```
int* p = NULL;
```

```
void* q = (void*)p;
```

```
string* r = (string*)q;
```

This is legal because q is NULL

- or do even wilder things

```
void* q = NULL;
```

```
void* r = (void*)(int*)(void*)(string*)q;
```

This is legal because q is NULL

- A NULL variable of type `void*` has *every* tag

```
void* v = NULL;
```

```
//@assert \hastag(int*, v);
```

```
//@assert \hastag(string*, v);
```

This is legal because v is NULL

- except `void*`

```
//@assert \hastag(void*, v);
```

This causes a **compilation error**

# Contracts of Cast Operations

- Casts are **potentially unsafe** operations over pointer expressions
  - With `\hastag`, we can write contracts for them

- Casting from specific to generic type
  - `(void*)x` where `x` was declared of type `tp*`

```
(void*)x  
//@ensures \hastag(tp*, \result);
```

- Casting from generic to specific type
  - `(tp*)q` where `q` was declared of type `void*`

```
(tp*)q  
//@requires \hastag(tp*, q);
```

# **Generic Stacks in C1**

# Generic Stacks

Use `void*` as the type of the elements

## ● Pros:

- Simple change to the library

`typedef void* elem;`

That's it!

- A single library for any kind of stacks

## ● Cons:

- Stack elements must be pointers

- We cannot have a stack of `ints`
- We need to turn them into `int*`

- This is the best we will be able to do

- genericity is limited to pointers
- not just in C1, but also in C

```
/* ***** Implementation ***** */
typedef void* elem; // Element type

typedef struct list_node list;
struct list_node {
    elem data;
    list* next;
};

typedef struct stack_header stack;
...

/* ***** Interface ***** */

// typedef _____* stack_t;

bool stack_empty(stack_t S)
/* @requires S != NULL; @ */ ;

stack_t stack_new()
/* @ensures \result != NULL; @ */
/* @ensures stack_empty(\result); @ */ ;

void push(stack_t S, elem x)
/* @requires S != NULL; @ */
/* @ensures !stack_empty(S); @ */ ;

elem pop(stack_t S)
/* @requires S != NULL; @ */
/* @requires !stack_empty(S); @ */ ;
```

# Converting an `int*` Stack to Generic

- Cast elements to `void*` when pushing
- Cast them back to `int*` when popping

main.c0

```
int main() {  
    stack_t l = stack_new();  
  
    int* x = alloc(int);  
    *x = 42;  
    push(l, x);  
  
    int* y = alloc(int);  
    *y = 15122;  
    push(l, y);  
  
    int* z = pop(l);  
    printint(*z); println("");  
  
    printint(*pop(l));  
  
    return 0;  
}
```

push now expects a `void*` ;  
cast x to a temp variable of type `void*`

Same thing, done inline

pop returns a `void*`, but z has type `int*`

Same thing, inline

main.c1

```
int main() {  
    stack_t l = stack_new();  
  
    int* x = alloc(int);  
    *x = 42;  
    void* p = (void*)x;  
    push(l, p);  
  
    int* y = alloc(int);  
    *y = 15122;  
    push(l, (void*)y);  
  
    int* z = (int*)pop(l);  
    printint(*z); println("");  
  
    printint(*(int*)pop(l));  
  
    return 0;  
}
```



# Compilation

**Application**  
file main.c1

```
int main() {  
    stack_t l = stack_new();  
  
    int* x = alloc(int);  
    *x = 42;  
    void* p = (void*)x;  
    push(l, p);  
  
    int* y = alloc(int);  
    *y = 15122;  
    push(l, (void*)y);  
  
    int* z = (int*)pop(l);  
    printint(*z); printf("\n");  
    printint(*(int*)pop(l));  
  
    return 0;  
}
```

Linux terminal

```
# cc0 -d stack.c1 main.c1
```

**Library**  
file stack.c1

```
/* ***** Implementation ***** */  
typedef void* elem; // Element type  
  
typedef struct list_node list;  
struct list_node {  
    elem data;  
    list* next;  
};  
  
typedef struct stack_header stack;  
...  
  
/* ***** Interface ***** */  
  
// typedef _____* stack_t;  
  
bool stack_empty(stack_t S)  
/* @requires S != NULL; @*/ ;  
  
stack_t stack_new()  
/* @ensures \result != NULL; @*/  
/* @ensures stack_empty(\result); @*/ ;  
  
void push(stack_t S, elem x)  
/* @requires S != NULL; @*/  
/* @ensures !stack_empty(S); @*/ ;  
  
elem pop(stack_t S)  
/* @requires S != NULL; @*/  
/* @requires !stack_empty(S); @*/ ;
```

*No need for a client-stack.c1 file!*

# Converting an `int` Stack to Generic

- No way to store an `int` into a generic stack

- We need to convert elements to `int*` first

➤ And cast them to `void*` to use the stack

This is annoying  
... but that's the best we can do

main.c0

```
int main() {  
    stack_t l = stack_new();
```

```
    push(l, 42);
```

```
    int y = 15122;  
    push(l, y);
```

```
    int z = pop(l);  
    printint(z); println("");
```

```
    printint(pop(l));
```

```
    return 0;  
}
```

We must store 42 in allocated memory, and  
cast its pointer to a temp variable of type `void*`

We can inline the cast, but not the allocation

pop returns a `void*`, but z has type `int*`

Same thing, inline

main.c1

```
int main() {  
    stack_t l = stack_new();
```

```
    int* x = alloc(int);  
    *x = 42;  
    void* p = (void*)x;  
    push(l, p);
```

```
    int* y = alloc(int);  
    *y = 15122;  
    push(l, (void*)y);
```

```
    int z = *(int*)pop(l);  
    printint(*z); println("");
```

```
    printint(*(int*)pop(l));
```

```
    return 0;  
}
```

# Using two Stacks of Different Type in C0

... in the same application

- We need to have two copies of the stack library
  - *int\_stack* for *ints* and *str\_stack* for *strings*

```
main.c0
int main() {
    int_stack_t I = int_stack_new(); // a stack of ints
    int_push(I, 42);
    int y = 15122;
    int_push(I, y);
    int z = int_pop(I);
    printint(z); println("");
    printint(int_pop(I));

    str_stack_t S = str_stack_new(); // a stack of strings
    str_push(S, "hello");
    string s = "world";
    str_push(S, s);
    string w = str_pop(S);
    println(w);
    println(str_pop(S));

    return 0;
}
```

# Using two Stacks of Different Type in C1

... in the same application

- The one generic stack library is enough
- but we need to convert elements to be pointers

main.c1

```
int main() {  
    stack_t I = stack_new(); // a stack for ints
```

```
    int* x = alloc(int);  
    *x = 42;  
    void* p = (void*)x;  
    push(I, p);
```

```
    int* y = alloc(int);  
    *y = 15122;  
    push(I, (void*)y);
```

```
    int z = *(int*)pop(I);  
    printint(z); println("");
```

```
    printint(*(int*)pop(I));
```

```
    // continued to the right
```

```
// continued from left
```

```
    stack_t S = stack_new(); // a stack for strings
```

```
    string* s1 = alloc(string);  
    *s1 = "hello";  
    push(S, (void*)s1);
```

```
    string* s = alloc(string);  
    *s = "world";  
    push(S, (void*)s);
```

```
    string w = *(string*) pop(S);  
    println(w);  
    println(*(string*)(pop(S)));
```

```
    return 0;
```

```
}
```

# Bad Uses of Generic Stacks

- Nothing prevents pushing elements of *different* type in the same generic stack
  - *but why would you want to do that???*



main.c1

```
int main() {  
    stack_t X = stack_new(); // one stack  
  
    int* i = alloc(int);  
    *i = 42;  
    push(X, (void*)i); // push an int onto X  
  
    string* s = alloc(string);  
    *s = "Ouch!";  
    push(X, (void*)s); // now push a string onto X!  
  
    string w = *(string*)pop(X);  
    println(w); // pop the string and print it  
  
    printint(*(int*)pop(X)); // pop the int and print it  
  
    return 0;  
}
```

- Extremely error-prone

In general, how do we remember  
this element will be a **string**?

... and this one an **int**?

- There is always a cleaner  
way to do this