## ExamPaper3-W16

Kārlis Kreilis

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**Example 11.11** Let us compare, on the basis of average response time, the performance of two identical servers each with its own separate queue, to the case when there is only a single queue to hold customers for both servers. The systems to compare are illustrated in Figure 11.17. We shall also check to see how these two possibilities compare to a single processor working twice as fast.

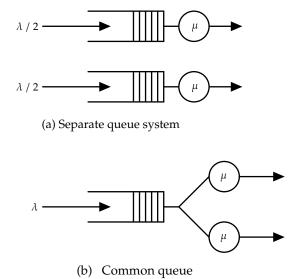


Figure 11.17. Two Configurations.

In the first case, we have two independent M/M/1 queues, each with arrival rate  $\lambda/2$  and service rate  $\mu$ . It follows that  $\rho = (\lambda/2)/\mu = \lambda/(2\mu)$ . The mean number in each M/M/1 queue is given by  $\rho/(1-\rho)$  so that the mean number of customers in this first scenario is given as

$$L_1 = E[N_1] = 2 \times \frac{\rho}{1 - \rho} = \frac{2\rho}{1 - \rho}.$$

The average response time can now be found using Little's law. We have

$$E[R_1] = \frac{1}{\lambda} E[N_1] = \frac{1}{\lambda} \frac{2\rho}{(1-\rho)} = \frac{2}{2\mu - \lambda}.$$

Now consider the second scenario in which the system may be represented as an M/M/2 queue. To find  $E[R_2]$ , we first must find  $E[N_2]$  (=  $L_2$ ). The mean number of customers in an M/M/c queue with arrival rate  $\lambda$  and service rate  $\mu$  per server is given by

$$E[N_2] = \frac{\lambda}{\mu} + \frac{(\lambda/\mu)^c \lambda \mu}{(c-1)!(c\mu - \lambda)^2} p_0 \quad \text{with} \quad \frac{\lambda}{c\mu} = \rho \quad \text{or} \quad \lambda/\mu = c\rho$$

With c = 2, we obtain

$$L_{2} = E[N_{2}] = \frac{\lambda}{\mu} + \frac{(\lambda/\mu)^{2}\lambda\mu}{(2\mu - \lambda)^{2}}p_{0} = \frac{\lambda}{\mu} + \frac{(\lambda/\mu)^{2}(\lambda/\mu)}{(1/\mu^{2})(2\mu - \lambda)^{2}}p_{0}$$
$$= \frac{\lambda}{\mu} + \frac{(\lambda/\mu)^{3}}{(2 - \lambda/\mu)^{2}}p_{0}$$
$$= 2\rho + \frac{(2\rho)^{3}}{(2 - 2\rho)^{2}}p_{0}.$$

**Example 11.11** Let us compare, on the basis of average response time, the performance of two identical servers each with its own separate queue, to the case when there is only a single queue to hold customers for both servers. The systems to compare are illustrated in Figure 11.17. We shall also check to see how these two possibilities compare to a single processor working twice as fast.

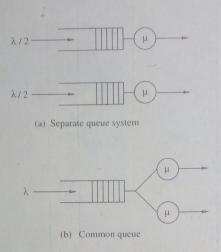


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With c = 2, we obtain

$$L_2 = E[N_2] = \frac{\lambda}{\mu} + \frac{(\lambda/\mu)^2 \lambda \mu}{(2\mu - \lambda)^2} p_0 = \frac{\lambda}{\mu} + \frac{(\lambda/\mu)^2 (\lambda/\mu)}{(1/\mu^2)(2\mu - \lambda)^2} p_0$$
$$= \frac{\lambda}{\mu} + \frac{(\lambda/\mu)^3}{(2 - \lambda/\mu)^2} p_0$$
$$= 2\rho + \frac{(2\rho)^3}{(2 - 2\rho)^2} p_0.$$

```
\documentclass{report}
\usepackage[utf8]{inputenc}
\usepackage{geometry}
\usepackage{amsmath}
\geometry{paper=a4paper,
left=32mm,
top=40mm,
right=30mm,
bottom=38mm,
\usepackage{fancyhdr,amssymb,txfonts,pxfonts}
\usepackage{pgf,tikz,pgfplots}
\pgfplotsset{compat=1.15}
\usepackage{mathrsfs}
\usepackage{caption}
\captionsetup[figure]{labelformat=empty}
\usetikzlibrary{arrows}
\fancyhf{}
\setcounter{page}{423}
\rhead{\small{\textbf{11.4 Multiserver Systems}\hspace{8mm}\thepage}}
\usepackage{graphicx}
\usepackage{listings}
\usepackage{microtype}
\hyphenpenalty=100
\pagestyle{fancy}
\title{ExamPaper3-W16}
\author{Kārlis Kreilis}
\date{May 2019}
\begin{document}
\maketitle
\setcounter{page}{423}
\noindent
\textbf{Example 11.11} Let us compare, on the basis of average response time, the performance
of two identical servers each with its own separate queue, to the case when there is only a
single queue to hold customers for both servers. The systems to compare are illustrated in
Figure 11.17. We shall also check to see how these two possibilities compare to a single
processor working twice as fast.
\begin{figure}[!h]
\begin{tikzpicture}[line cap=round,line join=round,>=triangle 45,x=1cm,y=1cm,
scale=0.8, transform shape]
clip(-12, -0.55) rectangle (5.8,3.2);
\frac{-5,2}{-75,2};
\frac{1}{100} [line width=1pt] (-3.5,1.5)-- (-1,1.5);
\draw [line width=1pt] (-1,1.5)-- (-1,2.5);
\frac{1}{100} [line width=1pt] (-1,2.5)-- (-3.5,2.5);
\draw [line width=1pt] (-1.8,2.5)-- (-1.8,1.5);
\draw [line width=1pt] (-1.2,2.5)-- (-1.2,1.5);
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```
\draw [line width=1pt] (-1.6,2.5)-- (-1.6,1.5);
 draw [line width=1pt] (-1.4,2.5)-- (-1.4,1.5);
\draw [line width=1pt] (0,2) circle (0.5cm);
\draw [line width=1pt] (-1,2)-- (-0.5,2.);
\frac{-}{1.6,2};
\frac{1}{2} \cdot \frac{1}{5} \cdot \frac{1}
\draw (-6,2.25) node[anchor=north west] {\lambda / 2\};
\draw (-0.2,2.22) node[anchor=north west] {<math>\mwedge mu$};
\frac{-5.0,-0}{-5.0,-0} -- (-2.75,-0);
\frac{1}{100} [line width=1pt] (-1,-0.5)-- (-1,0.5);
\frac{1}{100} \draw [line width=1pt] (-1,0.5)-- (-3.5,0.5);
\draw [line width=1pt] (-0,-0) circle (0.5cm);
\draw [line width=1pt] (-1,-0)-- (-0.5,-0);
\frac{-}{1.6,-0};
\draw [line width=1pt] (-1,-0.5)-- (-3.5,-0.5);
\frac{1}{100} [line width=1pt] (-1.2,0.5)-- (-1.2,-0.5);
\frac{1}{100} [line width=1pt] (-1.4,0.5)-- (-1.4,-0.5);
\frac{1}{100} [line width=1pt] (-1.6,0.5)-- (-1.6,-0.5);
\frac{1}{100} [line width=1pt] (-1.8,0.5)-- (-1.8,-0.5);
\frac{1}{2} \cdot \frac{1}
\draw (-6,0.25) node[anchor=north west] {$ \lambda / 2$};
\draw (-0.2,0.22) node[anchor=north west] {<math>\mws};
\end{tikzpicture}
\caption{\small{(a) Separate queue system}}
\end{figure}
\begin{figure}[h!]
\begin{tikzpicture}[line cap=round,line join=round,>=triangle 45,x=1cm,y=1cm,
 scale=0.8, transform shape]
clip(-12,-0.55) rectangle (9,3.55);
\frac{-5,2}{-75,2};
\frac{1}{100} [line width=1pt] (-3.5,1.5)-- (-1,1.5);
draw [line width=1pt] (-1,1.5)-- (-1,2.5);
\det[1] (-1,2.5) -- (-3.5,2.5);
\draw [line width=1pt] (-1.8,2.5)-- (-1.8,1.5);
\frac{1}{100} [line width=1pt] (-1.2,2.5)-- (-1.2,1.5);
\frac{1.6,2.5}{-1.6,1.5};
\draw [line width=1pt] (-1.4,2.5)-- (-1.4,1.5);
\draw [line width=1pt] (1,3) circle (0.5cm);
\frac{1}{2} \cdot \frac{1}
\draw (-5.5,2.25) node[anchor=north west] {\shape \lambda\shapers};
\draw (0.75,3.22) node[anchor=north west] {<math>\mbox{mu}};
\draw (0.75, 1.22) node[anchor=north west] {<math>\mbox{mu}};
\draw [line width=1pt] (1,1) circle (0.5cm);
\frac{1}{100} [line width=1pt] (-1,2)-- (-0.5,2);
\draw [line width=1pt] (-0.5,2)-- (0.5,1);
\frac{1}{2} \operatorname{draw} [\lim \operatorname{dish} -1 \operatorname{pt}] (-0.5, 2) -- (0.5, 3);
\frac{-}{1.5,3} -- (3,3);
\frac{-}{1} (1.5,1) -- (3,1);
\draw (-4,0.2) node[anchor=north west] {\large{(b)}
                                                                                                                                                                                                                                                                                                                    Common queue}};
\end{tikzpicture}
\caption{Figure 11.17. Two Configurations.}
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```
\end{figure}
 \hfill
In the first case, we have two independent $M/M/1$ queues, each with arrival rate $\lambda/2$
 and service rate \mu. It follows that \rho = (\lambda/2)/\mu = \lambda/2 /(2\mu).
The mean number in each M/M/1 queue is given by \rho(1-\rho) so that the mean number of
 customers in this first scenario is given as
\begin{equation*}
                    L_{1}= E[N_1]= 2 \times \frac{rho}{1-rho}=\frac{2\rho}{1-rho}.
\end{equation*}
\noindent
The average response time can now be found using Little's law. We have
\begin{equation*}
                     E[R_1] = \frac{1}{\lambda}E[N_1] = \frac{1}{\lambda}
                      \frac{2\rho}{(1-\rho)}=\frac{2}{2\mu-\lambda}.
\end{equation*}
\noindent
Now consider the second scenario in which the system may be represented as an $M/M/2$ queue.
To find E[R_2], we first must find E[N_2]: (= L_2)$. The mean number of customers in
 an $M/M/c$ queue with arrival rate $\lambda$ and service rate $\mu$ per server is given by
\begin{equation*}
                      E[N_2] = \frac{\lambda}{\mu} + \frac{(\lambda)^2}{\mu} + \frac{(\lambda)^2}{\mu}^2 + \frac{(\lambda)^2}{\mu
                     \hspace{3mm} \label{lambda/mu=c} \hspace{3mm} \hspace{mu=c\rho}
\end{equation*}
\noindent
With c = 2, we obtain
\begin{align*}
 L_2 = E[N_2] = \frac{\lambda}{\lambda}^2 \left[ \frac{\lambda}{\mu} + \frac{(\lambda/\mu)^2 \ln da/\mu}{(2\mu/\mu)^2} \right] 
 $$ = \frac{\lambda}{\mu} + \frac{(\lambda)^2(\lambda)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2(1-\mu)^2
= \frac{\sum_{u \in \mathbb{Z}} mu}{+}
\frac{(\lambda/\mu)^3}{(2-\lambda/\mu)^2}p_0 \
= 2\rho+ \frac{(2\rho)^3}{(2-2\rho)^2}p_0.
\end{align*}
```