

# **Dijkstra's Algorithms**

**Textbook: Weiss Chapter 9.3**

**Byoungyoung Lee**

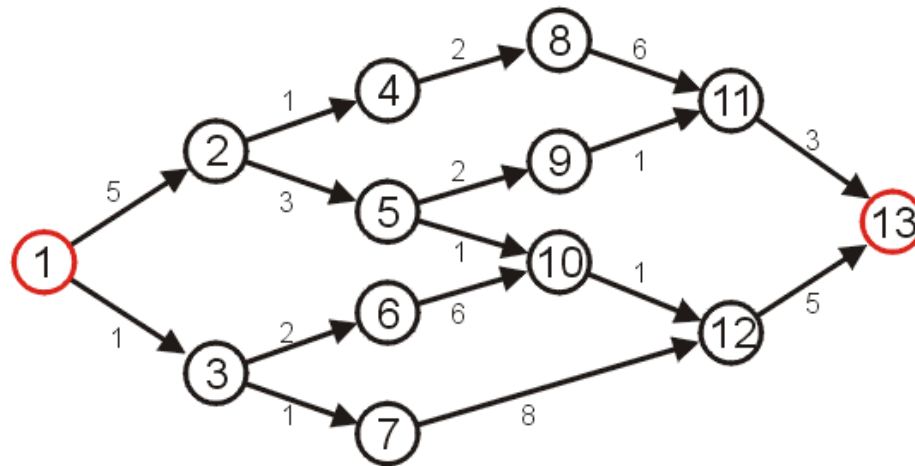
**<https://compsec.snu.ac.kr>**

**byoungyoung@snu.ac.kr**

# Shortest Path

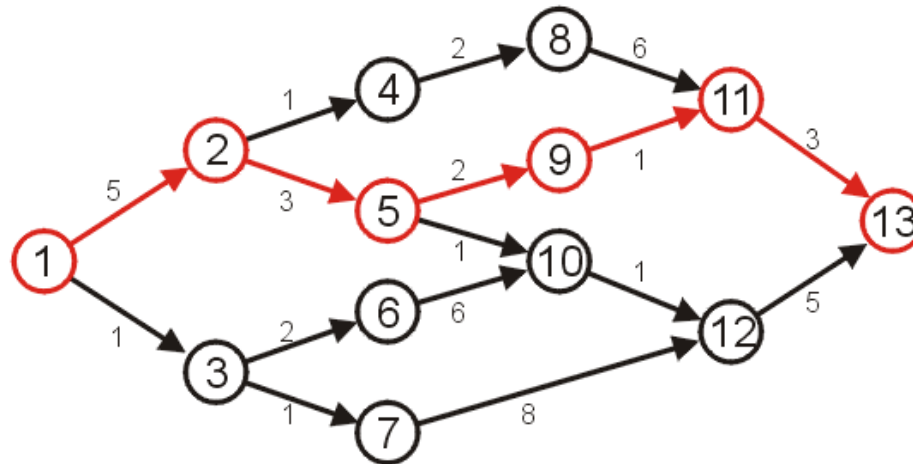
Given the graph, suppose you wish to **find the shortest path** from vertex 1 to vertex 13

- Length of the path is the sum of edge weights
- The graph is directed acyclic graph



# Shortest Path

The shortest path has length 14



Other paths exists, but they are longer

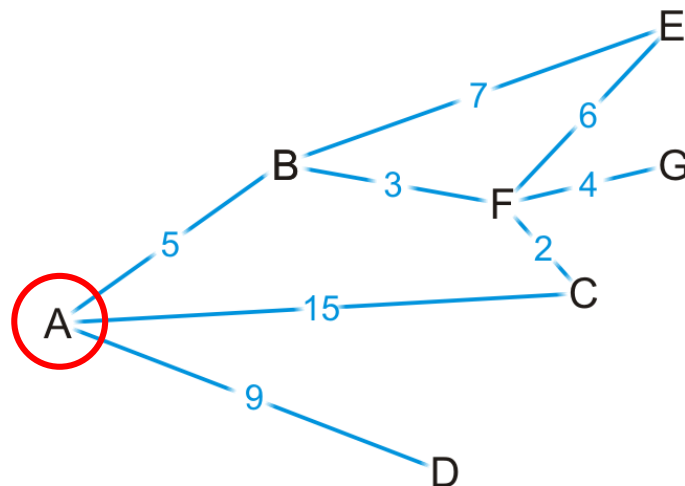
# Dijkstra's algorithm

**Dijkstra's algorithm** solves the **single-source shortest path** problem

- It is very similar to Prim's algorithm
- Assumption: All the weights are positive

Suppose you are at vertex A

- You are aware of all adjacent vertices to A
- This information is either in an adjacency list or adjacency matrix



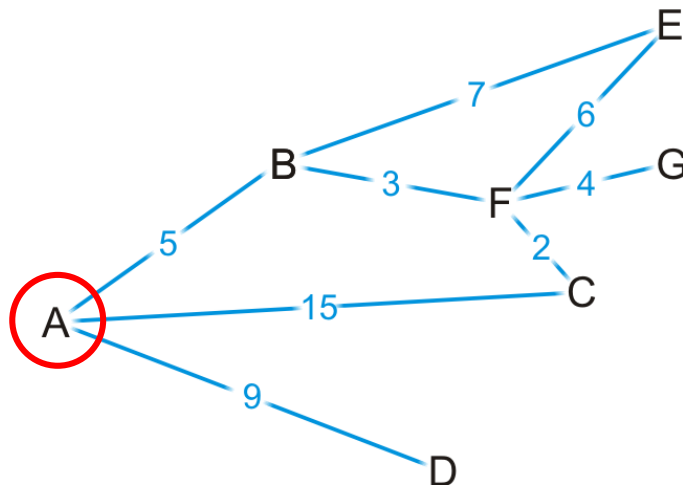
# Strategy of Dijkstra's Algorithm

By looking at all the adjacent edges of A, what can you be sure?

Is 5 the shortest distance from A to B?

Is 15 the shortest distance from A to C?

Is 9 the shortest distance from A to D?



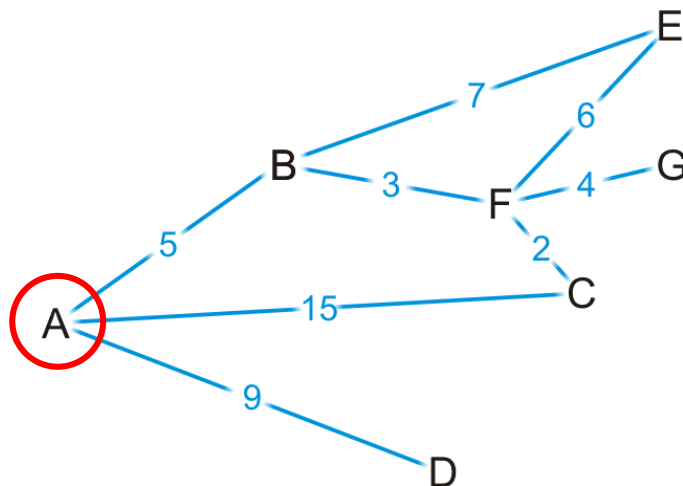
# Strategy of Dijkstra's Algorithm

By looking at all the adjacent edges of A, what can you be sure?

**Is 5 the shortest distance from A to B? (Yes)**

**Is 15 the shortest distance from A to C? (No)**

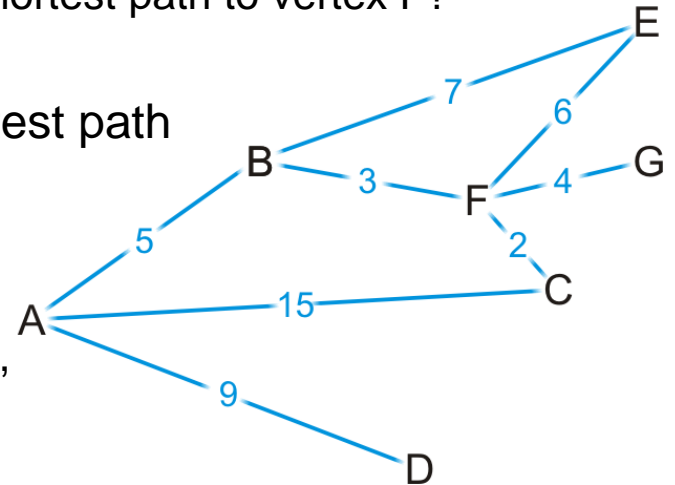
**Is 9 the shortest distance from A to D? (No)**



# Dijkstra's algorithm

Like Prim's algorithm, we initially don't know the distance to any vertex except vertices adjacent to the initial vertex

- We require an array of distances, all initialized to infinity except for the source vertex, which is initialized to 0
- Each time we visit a vertex, we will examine all adjacent vertices
  - We need to track visited vertices—a Boolean table of size  $|V|$
- Do we need to track the shortest path to each vertex?
  - That is, do I have to store (A, B, F) as the shortest path to vertex F?
- We really only have to record that the shortest path to vertex F came from vertex B
  - We would then determine that the shortest path to vertex B came from vertex A
  - Thus, we need an array of previous vertices, all initialized to null



# Dijkstra's algorithm

Thus, we will iterate  $|V|$  times:

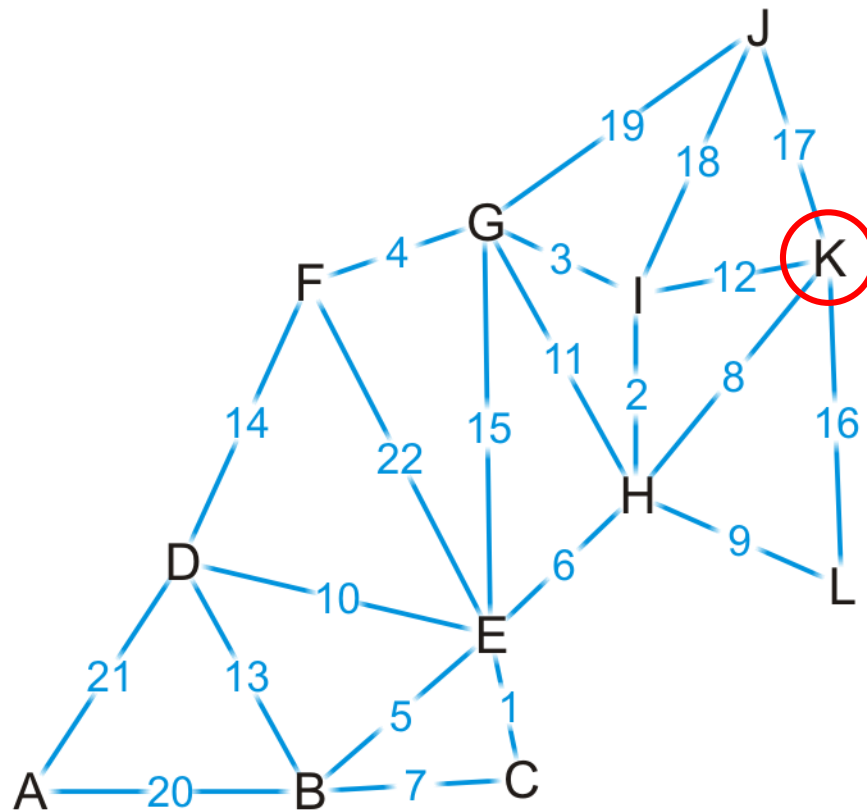
- Find that unvisited vertex  $v$  that has a minimum distance to it
- Mark it as having been visited
- Consider every adjacent vertex  $w$  that is unvisited:
  - Is the distance to  $v$  plus the weight of the edge  $(v, w)$  less than our currently known shortest distance to  $w$
  - If so, update the shortest distance to  $w$  and record  $v$  as the previous pointer
- Continue iterating until all vertices are visited or all remaining vertices have a distance to them of infinity



# Example

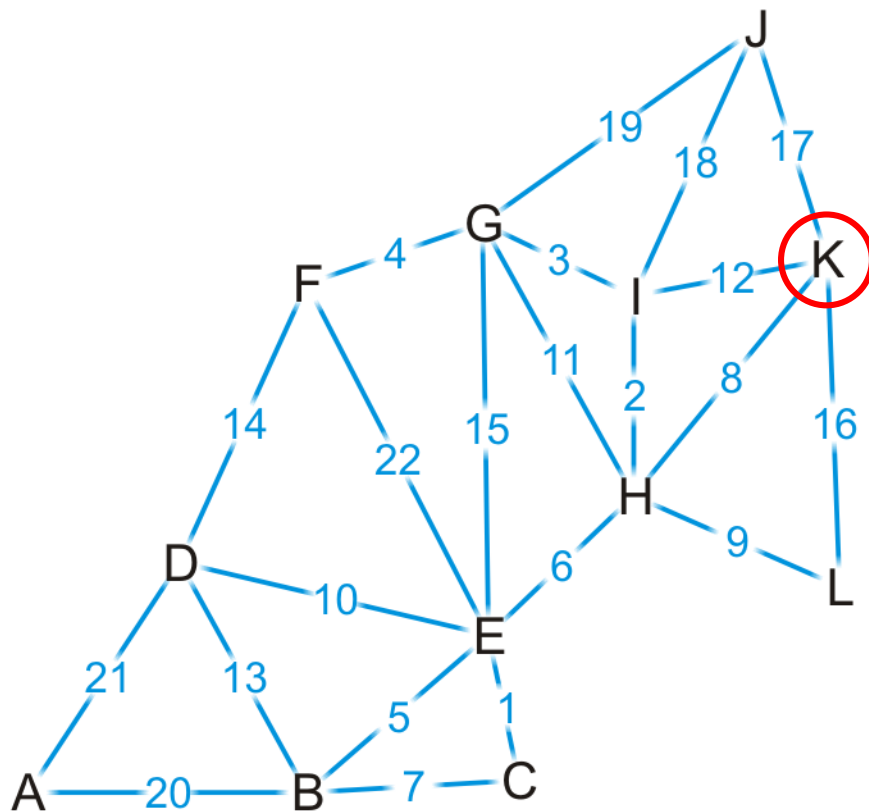
Let's take an example with this graph

We will find the shortest distance from **the vertex K** to every other vertices



# Example

We first set up the table.

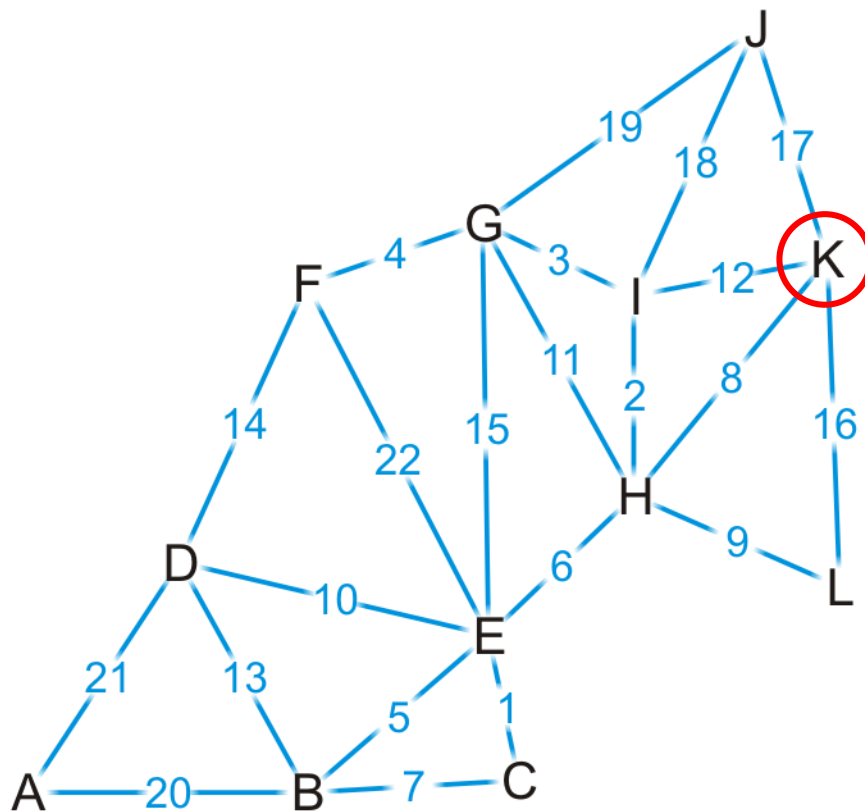


Vertex	Visited	Distance	Previous
A	F	$\infty$	$\emptyset$
B	F	$\infty$	$\emptyset$
C	F	$\infty$	$\emptyset$
D	F	$\infty$	$\emptyset$
E	F	$\infty$	$\emptyset$
F	F	$\infty$	$\emptyset$
G	F	$\infty$	$\emptyset$
H	F	$\infty$	$\emptyset$
I	F	$\infty$	$\emptyset$
J	F	$\infty$	$\emptyset$
K	F	0	$\emptyset$
L	F	$\infty$	$\emptyset$

# Example

We visit vertex K.

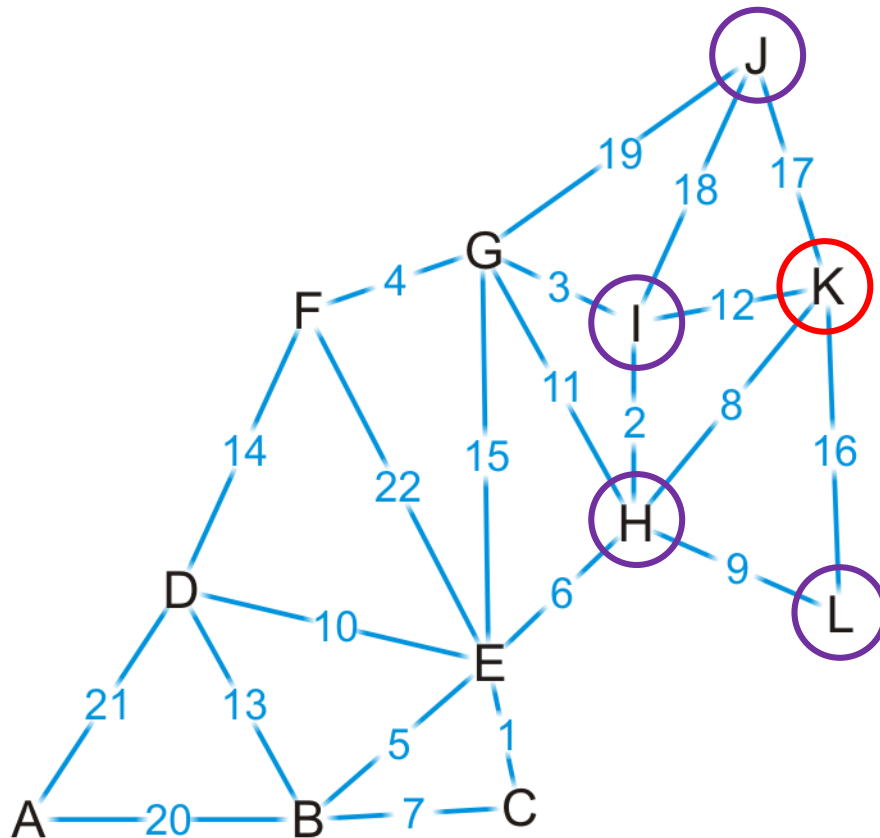
- Which unvisited vertex has the minimum distance to the vertex k?



Vertex	Visited	Distance	Previous
A	F	$\infty$	$\emptyset$
B	F	$\infty$	$\emptyset$
C	F	$\infty$	$\emptyset$
D	F	$\infty$	$\emptyset$
E	F	$\infty$	$\emptyset$
F	F	$\infty$	$\emptyset$
G	F	$\infty$	$\emptyset$
H	F	$\infty$	$\emptyset$
I	F	$\infty$	$\emptyset$
J	F	$\infty$	$\emptyset$
<b>K</b>	<b>T</b>	<b>0</b>	<b><math>\emptyset</math></b>
L	F	$\infty$	$\emptyset$

# Example

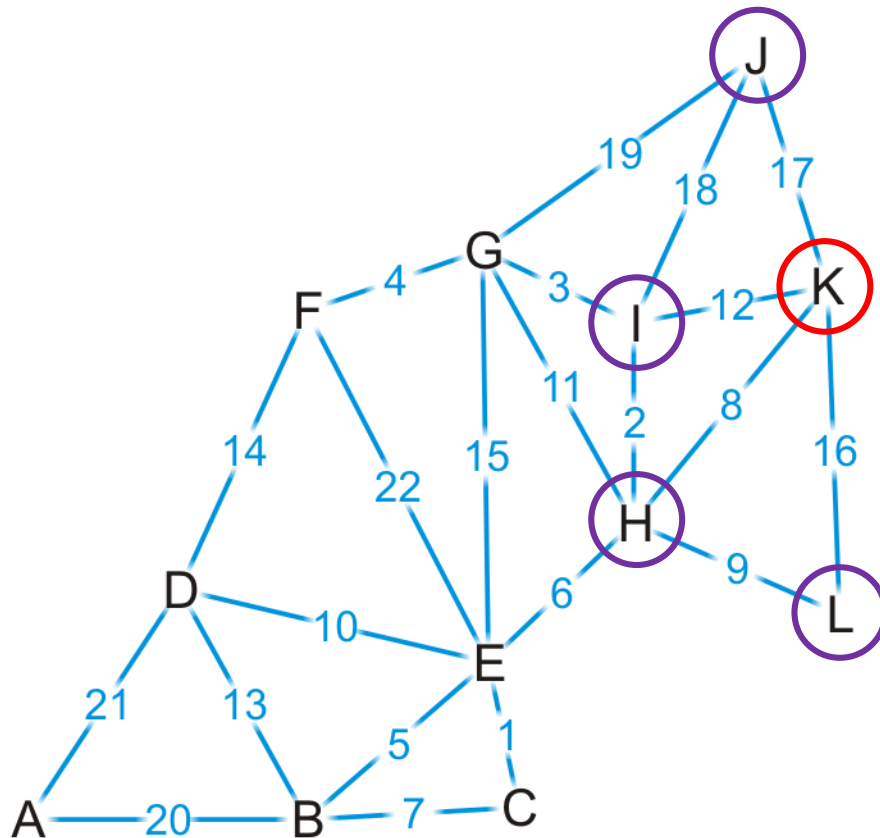
Vertex K has four neighbors: H, I, J and L



Vertex	Visited	Distance	Previous
A	F	$\infty$	$\emptyset$
B	F	$\infty$	$\emptyset$
C	F	$\infty$	$\emptyset$
D	F	$\infty$	$\emptyset$
E	F	$\infty$	$\emptyset$
F	F	$\infty$	$\emptyset$
G	F	$\infty$	$\emptyset$
H	F	$\infty$	$\emptyset$
I	F	$\infty$	$\emptyset$
J	F	$\infty$	$\emptyset$
<b>K</b>	<b>T</b>	<b>0</b>	<b><math>\emptyset</math></b>
L	F	$\infty$	$\emptyset$

# Example

We have now found at least one path to each of these vertices

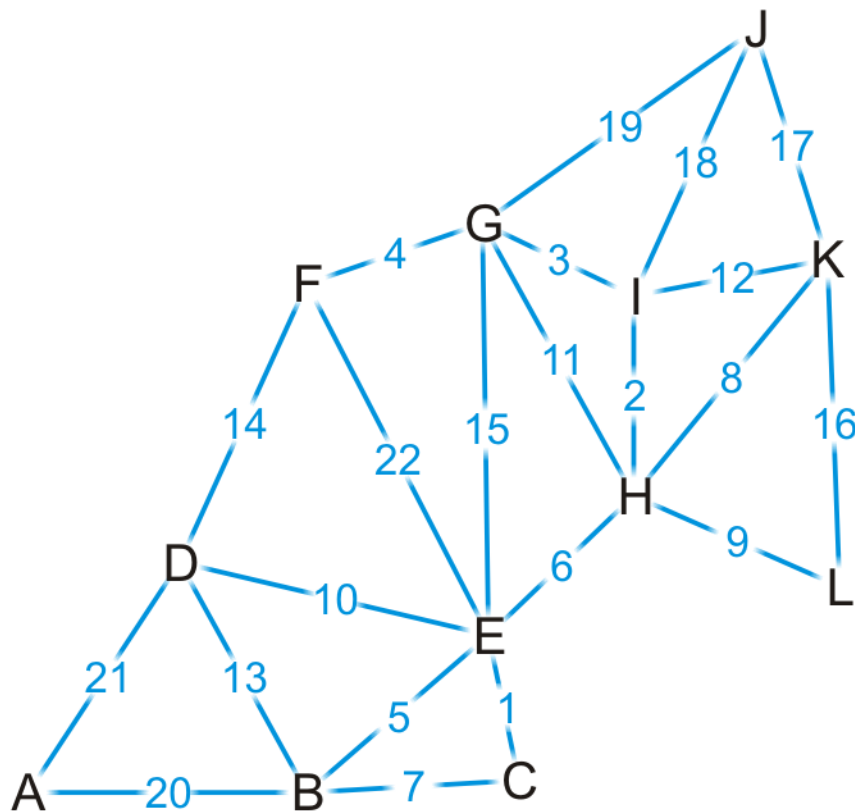


Vertex	Visited	Distance	Previous
A	F	$\infty$	$\emptyset$
B	F	$\infty$	$\emptyset$
C	F	$\infty$	$\emptyset$
D	F	$\infty$	$\emptyset$
E	F	$\infty$	$\emptyset$
F	F	$\infty$	$\emptyset$
G	F	$\infty$	$\emptyset$
H	F	8	K
I	F	12	K
J	F	17	K
K	T	0	$\emptyset$
L	F	16	K

# Example

We're finished with vertex K

- To which vertex are we now guaranteed we have the shortest path?

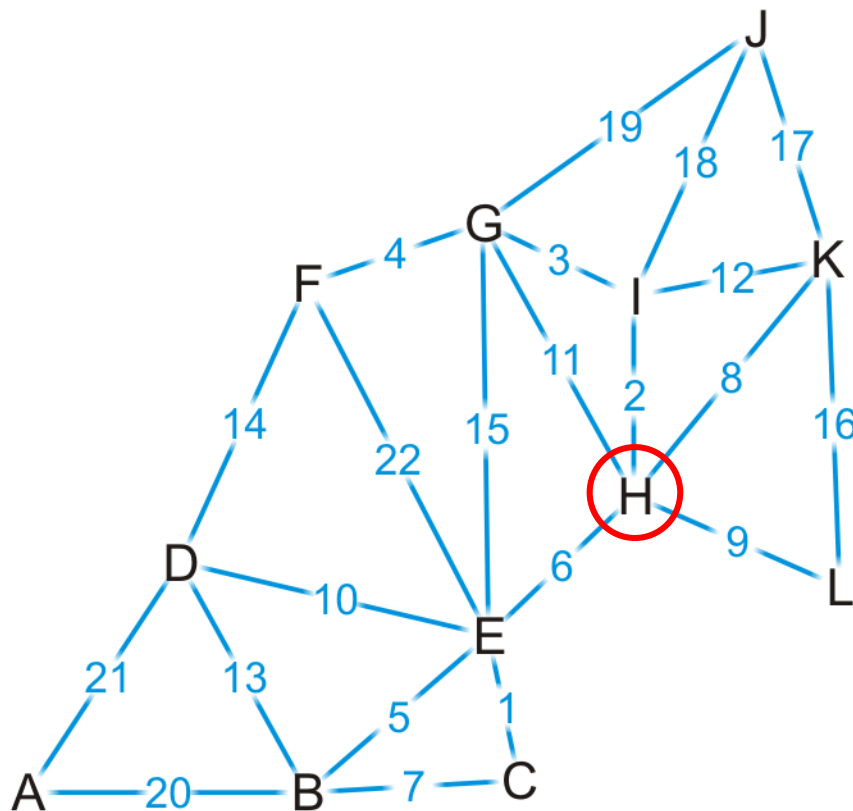


Vertex	Visited	Distance	Previous
A	F	$\infty$	$\emptyset$
B	F	$\infty$	$\emptyset$
C	F	$\infty$	$\emptyset$
D	F	$\infty$	$\emptyset$
E	F	$\infty$	$\emptyset$
F	F	$\infty$	$\emptyset$
G	F	$\infty$	$\emptyset$
H	F	8	K
I	F	12	K
J	F	17	K
K	T	0	$\emptyset$
L	F	16	K

# Example

We visit vertex H: the shortest path is (K, H) of length 8

- Vertex H has four unvisited neighbors: E, G, I, L



Vertex	Visited	Distance	Previous
A	F	$\infty$	$\emptyset$
B	F	$\infty$	$\emptyset$
C	F	$\infty$	$\emptyset$
D	F	$\infty$	$\emptyset$
E	F	$\infty$	$\emptyset$
F	F	$\infty$	$\emptyset$
G	F	$\infty$	$\emptyset$
<b>H</b>	<b>T</b>	<b>8</b>	<b>K</b>
I	F	12	K
J	F	17	K
K	T	0	$\emptyset$
L	F	16	K

# Example

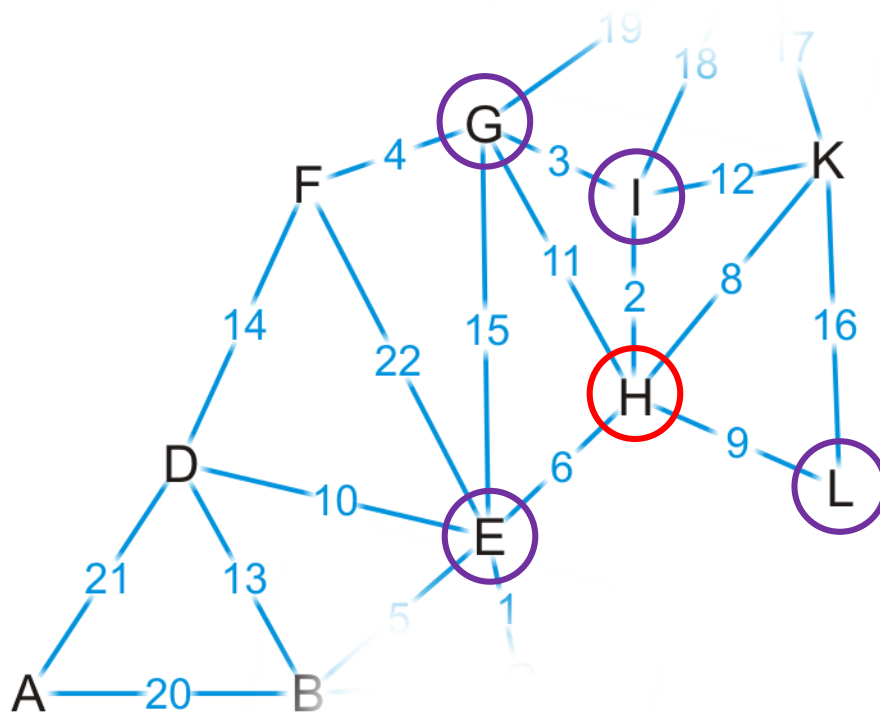
Consider these paths:

(K, H, E) of length  $8 + 6 = 14$

(K, H, G) of length  $8 + 11 = 19$

(K, H, I) of length  $8 + 2 = 10$

(K, H, L) of length  $8 + 9 = 17$

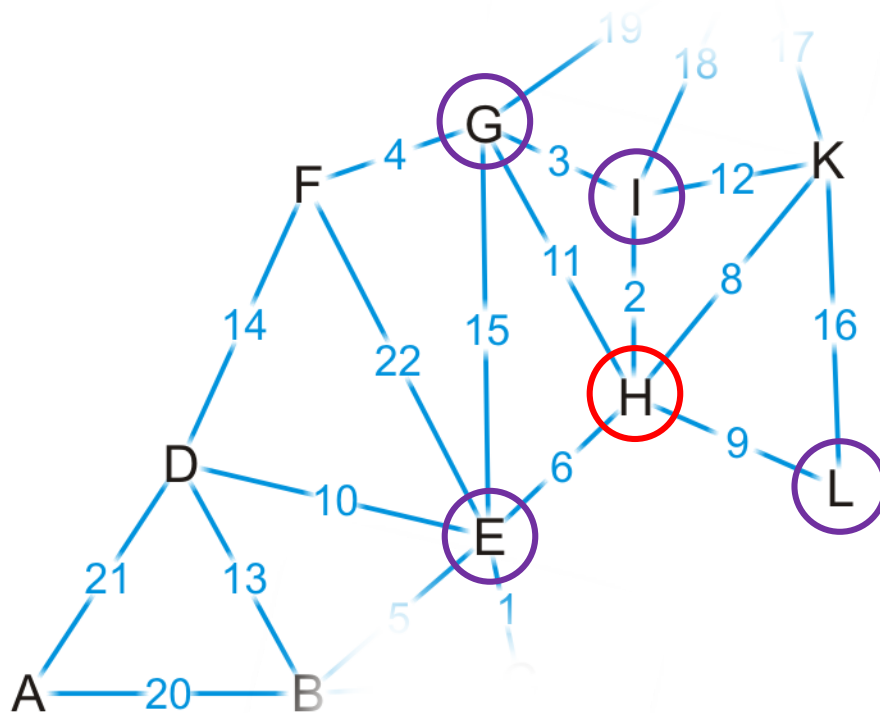


Vertex	Visited	Distance	Previous
A	F	$\infty$	$\emptyset$
B	F	$\infty$	$\emptyset$
C	F	$\infty$	$\emptyset$
D	F	$\infty$	$\emptyset$
E	F	$\infty$	$\emptyset$
F	F	$\infty$	$\emptyset$
G	F	$\infty$	$\emptyset$
H	T	8	K
I	F	12	K
J	F	17	K
K	T	0	$\emptyset$
L	F	16	K



# Example

We already have a shorter path (K, L), so L is not updated.  
But the other three (e.g., E, G, and I) are updated.

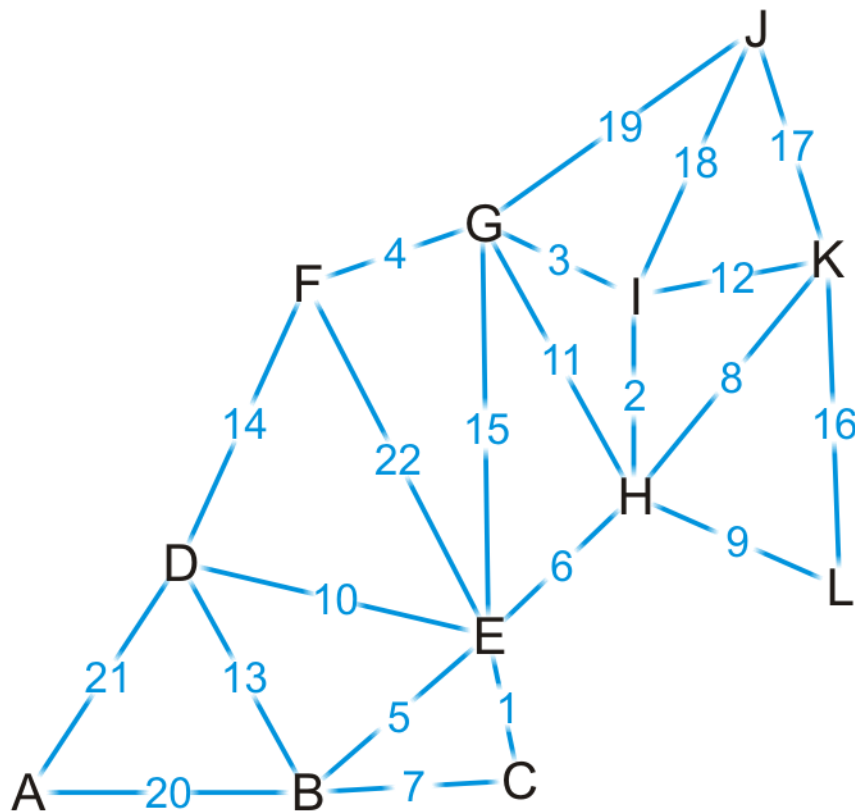


Vertex	Visited	Distance	Previous
A	F	$\infty$	$\emptyset$
B	F	$\infty$	$\emptyset$
C	F	$\infty$	$\emptyset$
D	F	$\infty$	$\emptyset$
E	F	14	H
F	F	$\infty$	$\emptyset$
G	F	19	H
H	T	8	K
I	F	10	H
J	F	17	K
K	T	0	$\emptyset$
L	F	16	K

# Example

We are finished with vertex H

- Which vertex do we visit next?

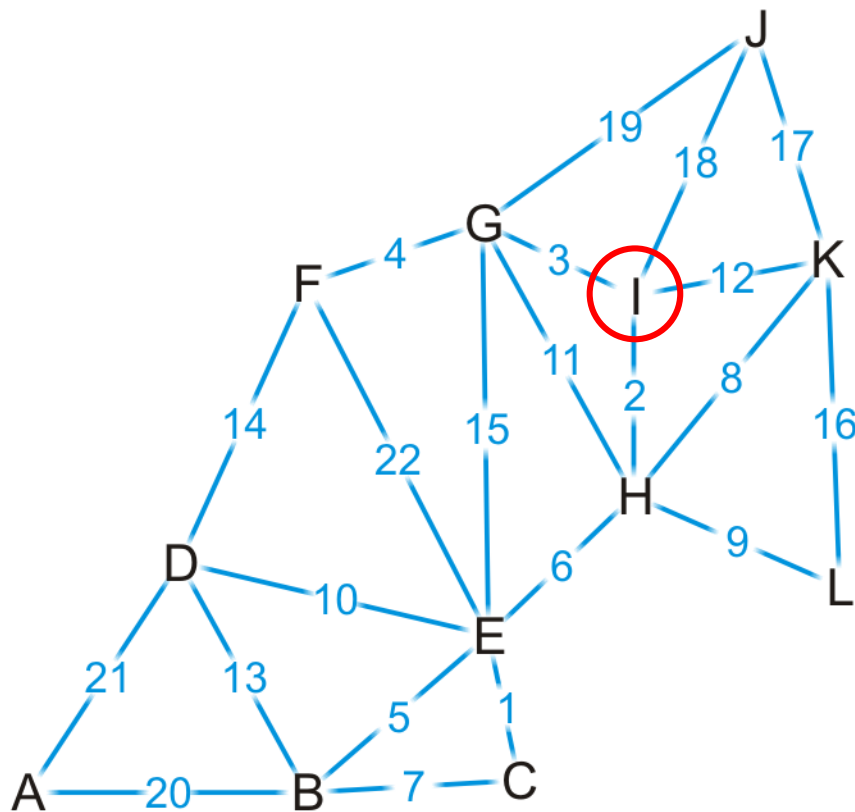


Vertex	Visited	Distance	Previous
A	F	$\infty$	$\emptyset$
B	F	$\infty$	$\emptyset$
C	F	$\infty$	$\emptyset$
D	F	$\infty$	$\emptyset$
E	F	14	H
F	F	$\infty$	$\emptyset$
G	F	19	H
H	T	8	K
I	F	10	H
J	F	17	K
K	T	0	$\emptyset$
L	F	16	K

# Example

The path (K, H, I) is the shortest path from K to I of length 10

- Vertex I has two unvisited neighbors: G and J



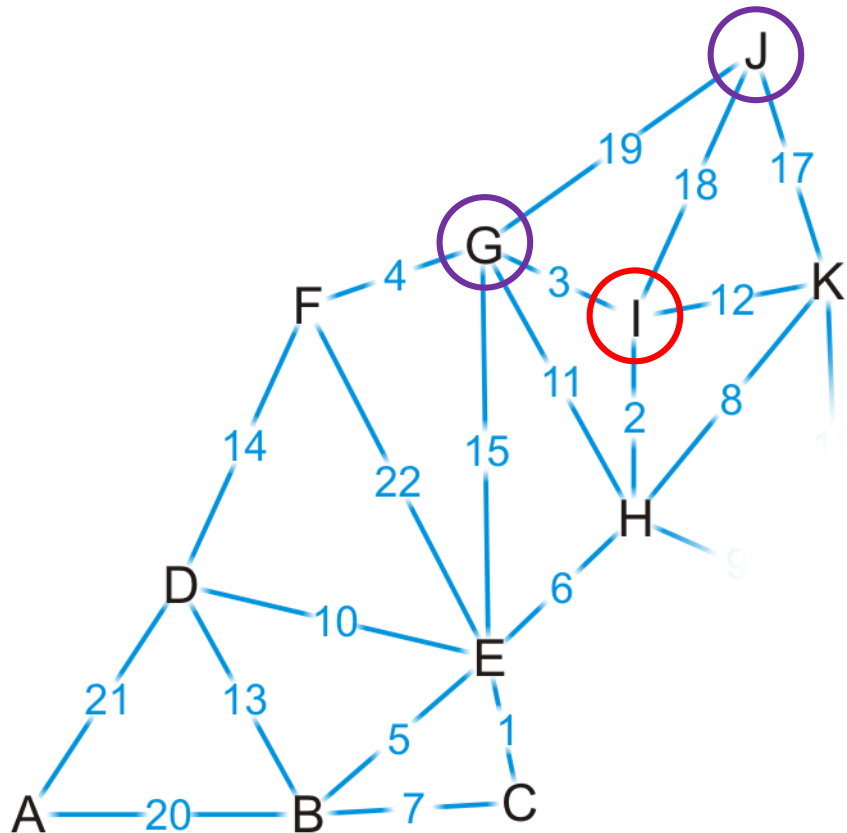
Vertex	Visited	Distance	Previous
A	F	$\infty$	$\emptyset$
B	F	$\infty$	$\emptyset$
C	F	$\infty$	$\emptyset$
D	F	$\infty$	$\emptyset$
E	F	14	H
F	F	$\infty$	$\emptyset$
G	F	19	H
H	T	8	K
<b>I</b>	<b>T</b>	<b>10</b>	<b>H</b>
J	F	17	K
K	T	0	$\emptyset$
L	F	16	K

# Example

Consider these paths:

(K, H, I, G) of length  $10 + 3 = 13$

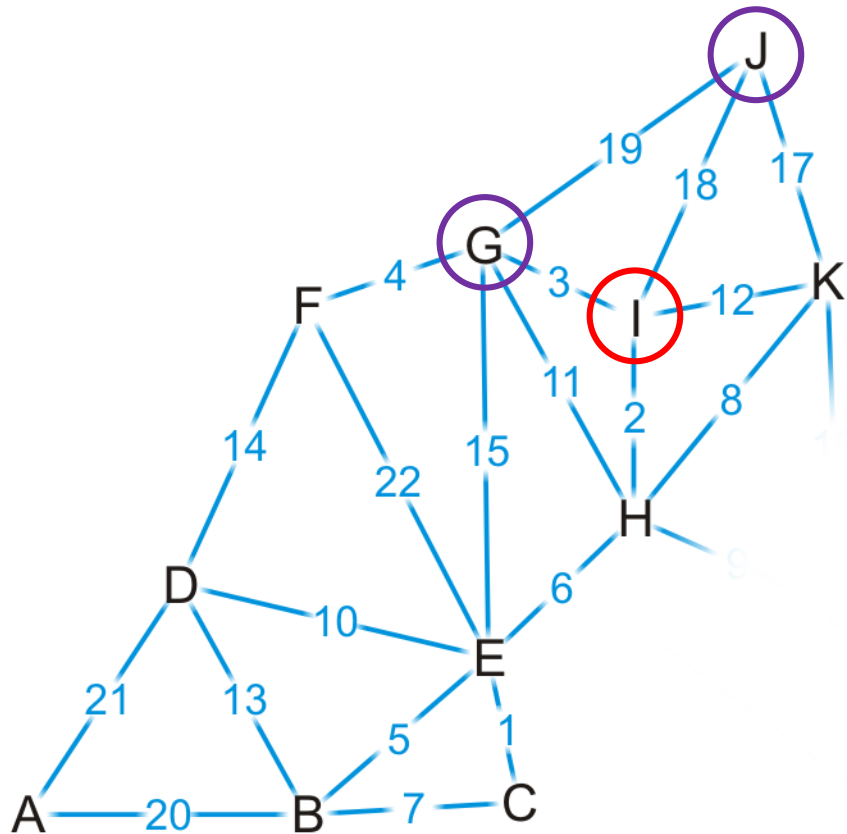
(K, H, I, J) of length  $10 + 18 = 28$



Vertex	Visited	Distance	Previous
A	F	$\infty$	$\emptyset$
B	F	$\infty$	$\emptyset$
C	F	$\infty$	$\emptyset$
D	F	$\infty$	$\emptyset$
E	F	14	H
F	F	$\infty$	$\emptyset$
G	F	19	H
H	T	8	K
I	T	10	H
J	F	17	K
K	T	0	$\emptyset$
L	F	16	K

# Example

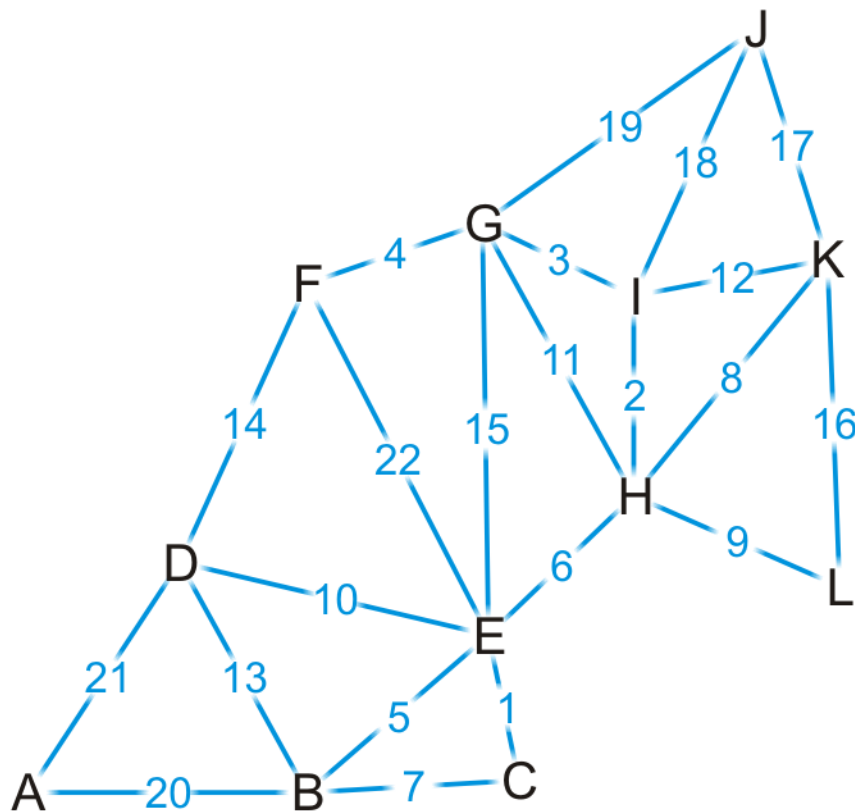
We have discovered a shorter path to vertex G, but (K, J) is still the shortest known path to vertex J



Vertex	Visited	Distance	Previous
A	F	$\infty$	$\emptyset$
B	F	$\infty$	$\emptyset$
C	F	$\infty$	$\emptyset$
D	F	$\infty$	$\emptyset$
E	F	14	H
F	F	$\infty$	$\emptyset$
<b>G</b>	<b>F</b>	<b>13</b>	<b>I</b>
H	T	8	K
<b>I</b>	<b>T</b>	<b>10</b>	<b>H</b>
<b>J</b>	<b>F</b>	<b>17</b>	<b>K</b>
K	T	0	$\emptyset$
L	F	16	K

# Example

Which vertex can we visit next?

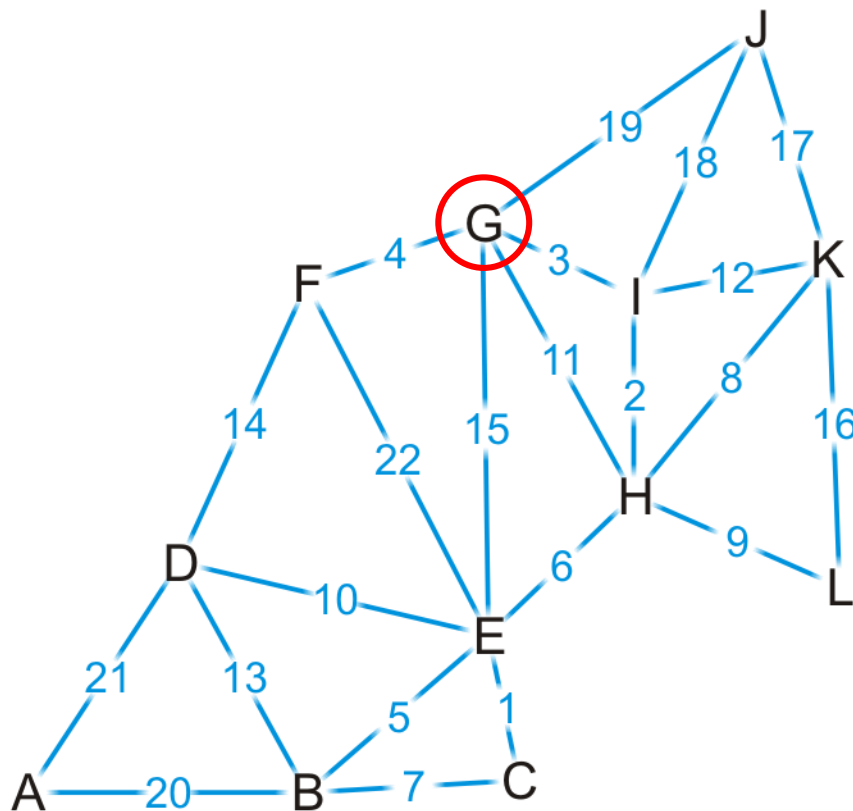


Vertex	Visited	Distance	Previous
A	F	$\infty$	$\emptyset$
B	F	$\infty$	$\emptyset$
C	F	$\infty$	$\emptyset$
D	F	$\infty$	$\emptyset$
E	F	14	H
F	F	$\infty$	$\emptyset$
G	F	13	I
H	T	8	K
I	T	10	H
J	F	17	K
K	T	0	$\emptyset$
L	F	16	K

# Example

The path (K, H, I, G) is the shortest path from K to G of length 13

- Vertex G has three unvisited neighbors: E, F and J



Vertex	Visited	Distance	Previous
A	F	$\infty$	$\emptyset$
B	F	$\infty$	$\emptyset$
C	F	$\infty$	$\emptyset$
D	F	$\infty$	$\emptyset$
E	F	14	H
F	F	$\infty$	$\emptyset$
<b>G</b>	<b>T</b>	<b>13</b>	<b>I</b>
H	T	8	K
I	T	10	H
J	F	17	K
K	T	0	$\emptyset$
L	F	16	K

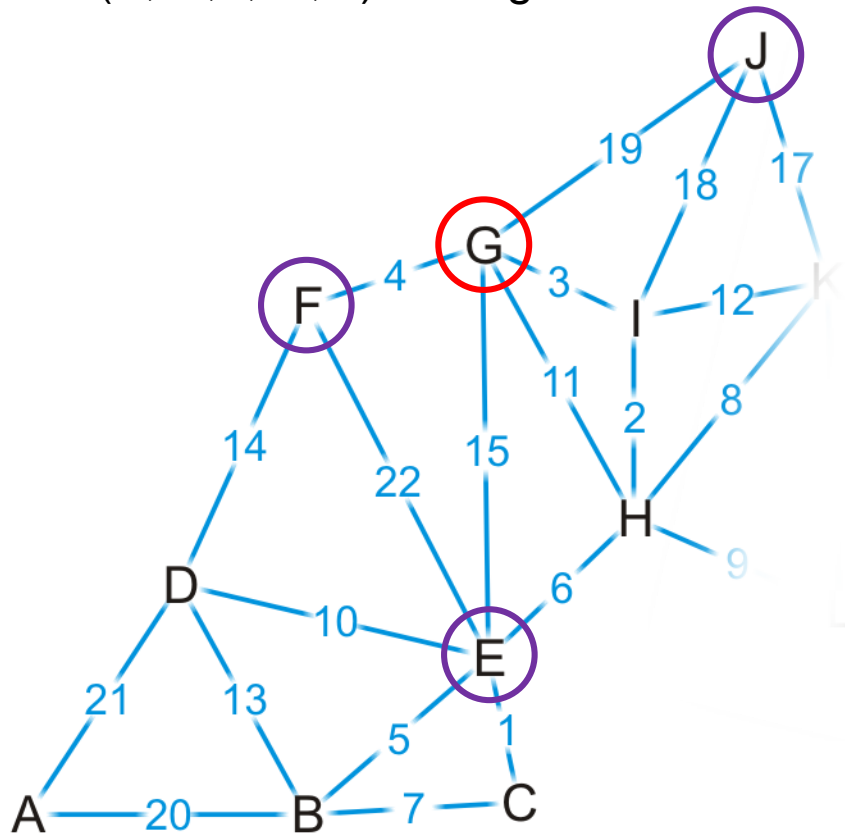
# Example

Consider these paths:

(K, H, I, G, E) of length **13** + 15 = 28

(K, H, I, G, F) of length **13** + 4 = 17

(K, H, I, G, J) of length **13** + 19 = 32

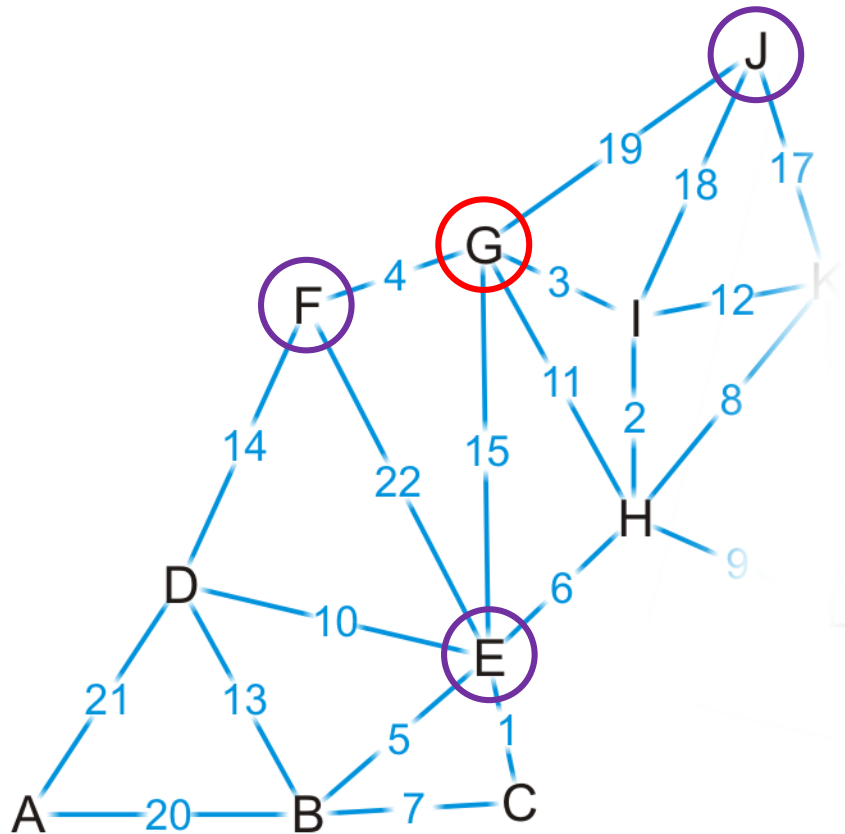


Vertex	Visited	Distance	Previous
A	F	$\infty$	$\emptyset$
B	F	$\infty$	$\emptyset$
C	F	$\infty$	$\emptyset$
D	F	$\infty$	$\emptyset$
E	F	14	H
F	F	$\infty$	$\emptyset$
G	T	13	I
H	T	8	K
I	T	10	H
J	F	17	K
K	T	0	$\emptyset$
L	F	16	K



# Example

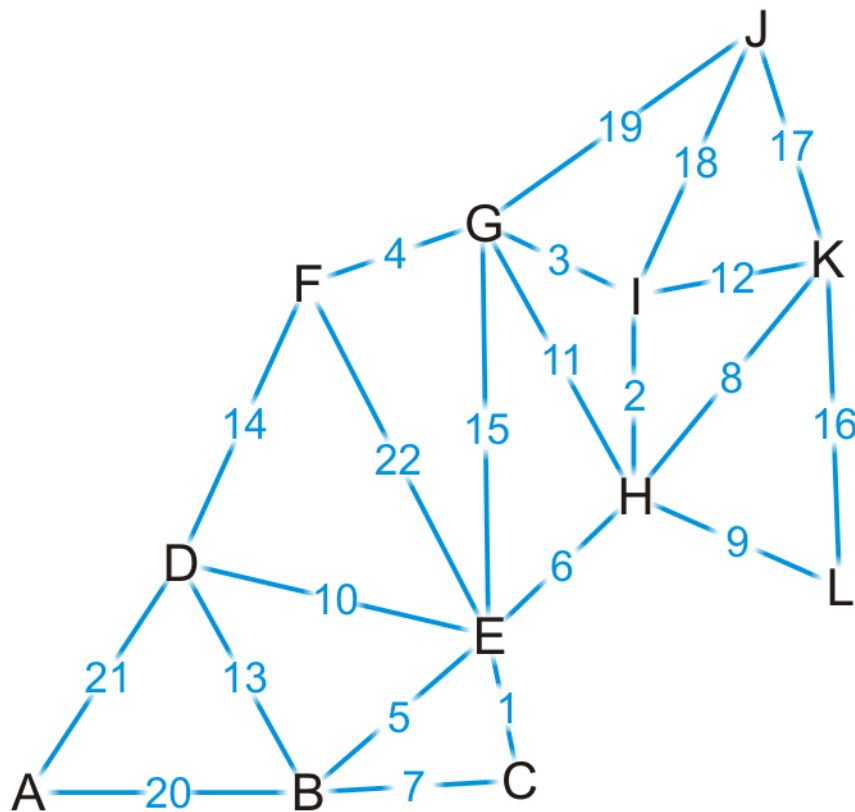
We have now found a path to vertex F



Vertex	Visited	Distance	Previous
A	F	$\infty$	$\emptyset$
B	F	$\infty$	$\emptyset$
C	F	$\infty$	$\emptyset$
D	F	$\infty$	$\emptyset$
E	F	14	H
F	F	17	G
G	T	13	I
H	T	8	K
I	T	10	H
J	F	17	K
K	T	0	$\emptyset$
L	F	16	K

# Example

Where do we visit next?

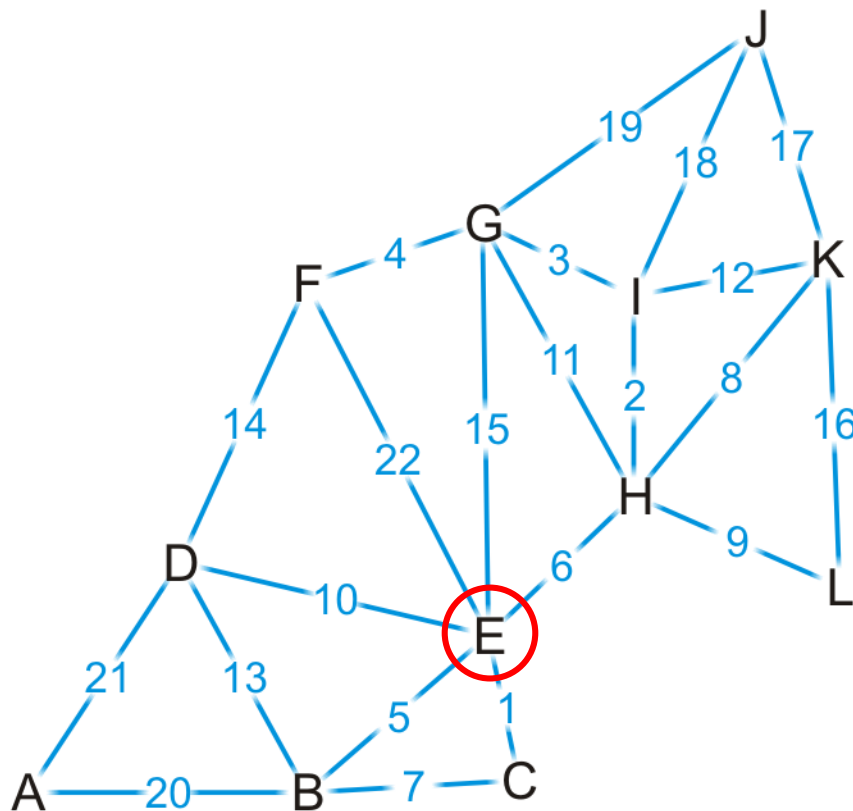


Vertex	Visited	Distance	Previous
A	F	$\infty$	$\emptyset$
B	F	$\infty$	$\emptyset$
C	F	$\infty$	$\emptyset$
D	F	$\infty$	$\emptyset$
E	F	14	H
F	F	17	G
G	T	13	I
H	T	8	K
I	T	10	H
J	F	17	K
K	T	0	$\emptyset$
L	F	16	K

# Example

The path (K, H, E) is the shortest path from K to E of length 14

- Vertex G has four unvisited neighbors: B, C, D and F

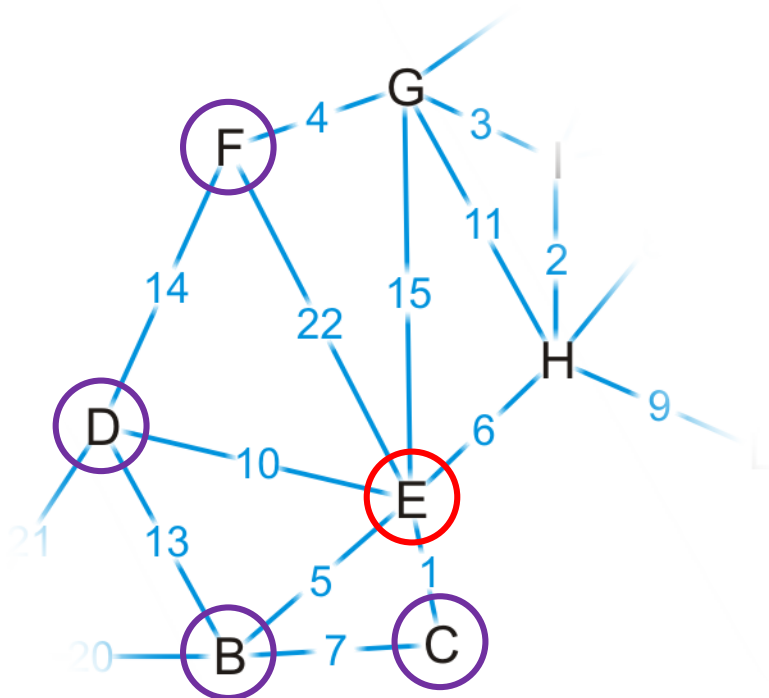


Vertex	Visited	Distance	Previous
A	F	$\infty$	$\emptyset$
B	F	$\infty$	$\emptyset$
C	F	$\infty$	$\emptyset$
D	F	$\infty$	$\emptyset$
<b>E</b>	<b>T</b>	<b>14</b>	<b>H</b>
F	F	17	G
G	T	13	I
H	T	8	K
I	T	10	H
J	F	17	K
K	T	0	$\emptyset$
L	F	16	K

# Example

The path (K, H, E) is the shortest path from K to E of length 14

- Vertex G has four unvisited neighbors: B, C, D and F

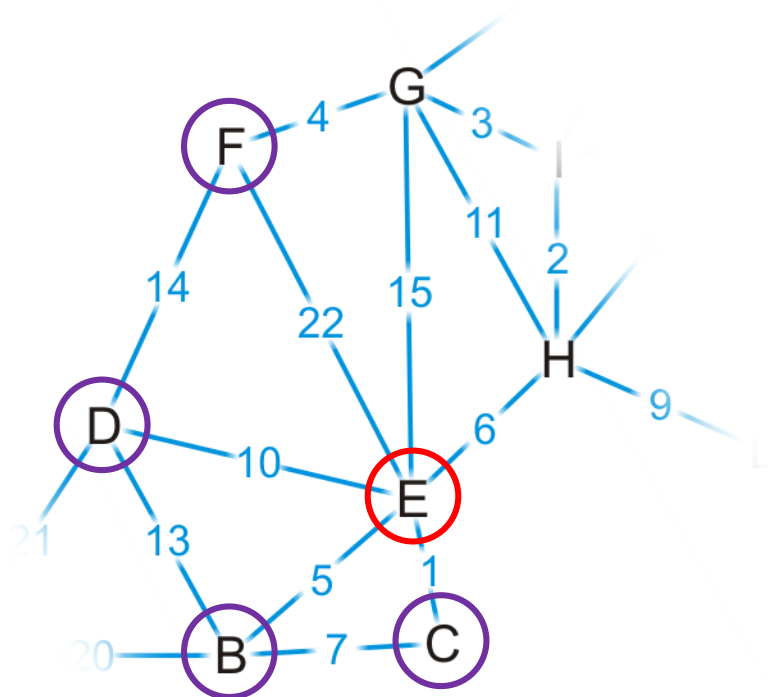


Vertex	Visited	Distance	Previous
A	F	$\infty$	$\emptyset$
B	F	$\infty$	$\emptyset$
C	F	$\infty$	$\emptyset$
D	F	$\infty$	$\emptyset$
E	T	14	H
F	F	17	G
G	T	13	I
H	T	8	K
I	T	10	H
J	F	17	K
K	T	0	$\emptyset$
L	F	16	K

# Example

Consider these paths:

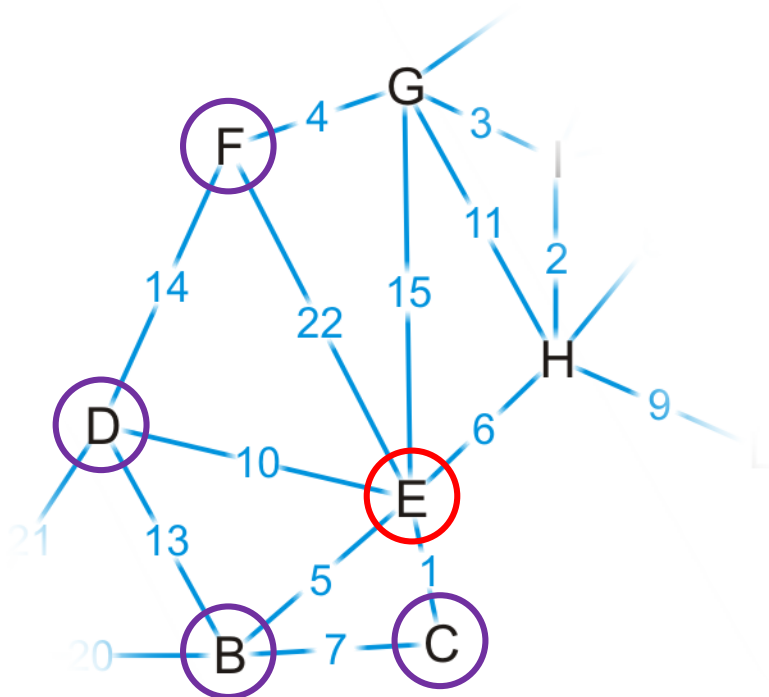
- (K, H, E, B) of length  $14 + 5 = 19$
- (K, H, E, C) of length  $14 + 1 = 15$
- (K, H, E, D) of length  $14 + 10 = 24$
- (K, H, E, F) of length  $14 + 22 = 36$



Vertex	Visited	Distance	Previous
A	F	$\infty$	$\emptyset$
B	F	$\infty$	$\emptyset$
C	F	$\infty$	$\emptyset$
D	F	$\infty$	$\emptyset$
E	T	14	H
F	F	17	G
G	T	13	I
H	T	8	K
I	T	10	H
J	F	17	K
K	T	0	$\emptyset$
L	F	16	K

# Example

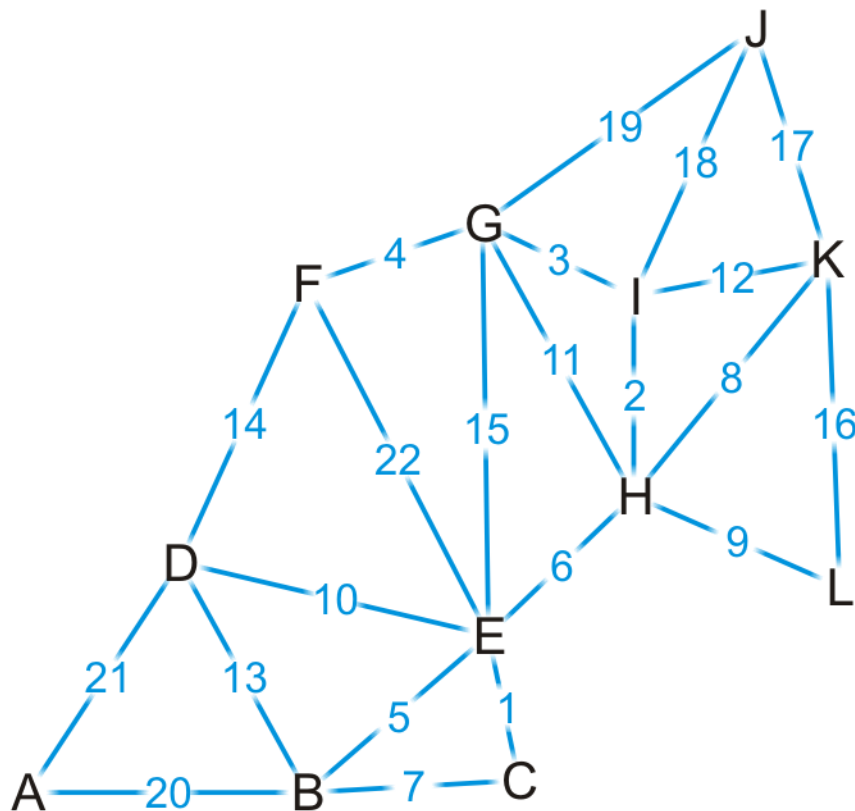
We've discovered paths to vertices B, C, D



Vertex	Visited	Distance	Previous
A	F	$\infty$	$\emptyset$
B	F	19	E
C	F	15	E
D	F	24	E
E	T	14	H
F	F	17	G
G	T	13	I
H	T	8	K
I	T	10	H
J	F	17	K
K	T	0	$\emptyset$
L	F	16	K

# Example

Which vertex is next?

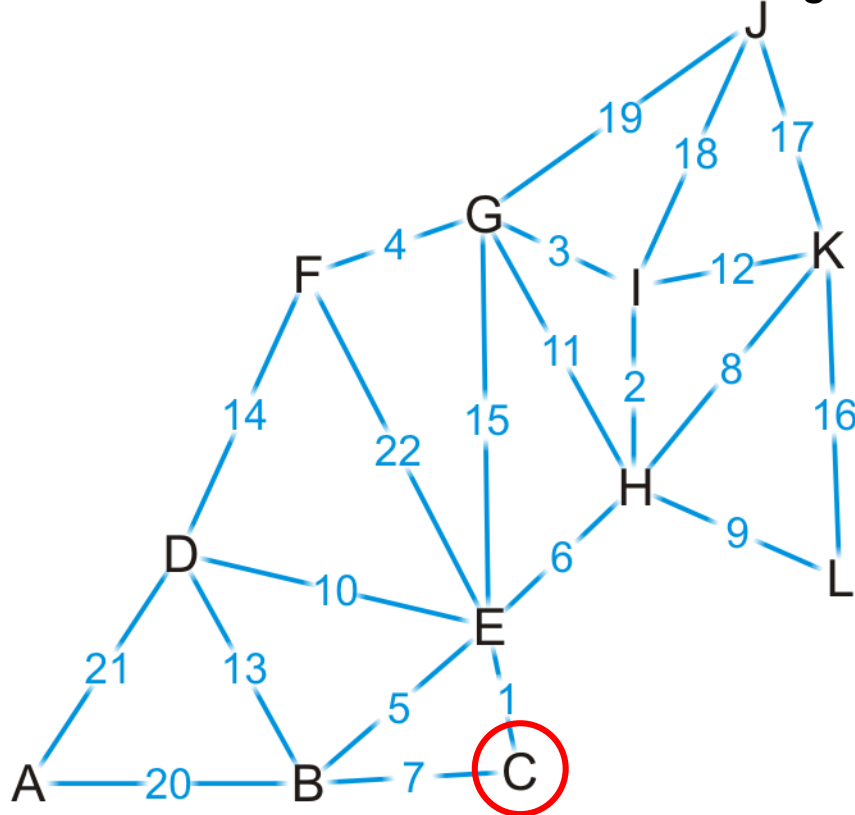


Vertex	Visited	Distance	Previous
A	F	$\infty$	$\emptyset$
B	F	19	E
C	F	15	E
D	F	24	E
E	T	14	H
F	F	17	G
G	T	13	I
H	T	8	K
I	T	10	H
J	F	17	K
K	T	0	$\emptyset$
L	F	16	K

# Example

We've found that the path (K, H, E, C) of length 15 is the shortest path from K to C

- Vertex C has one unvisited neighbor, B



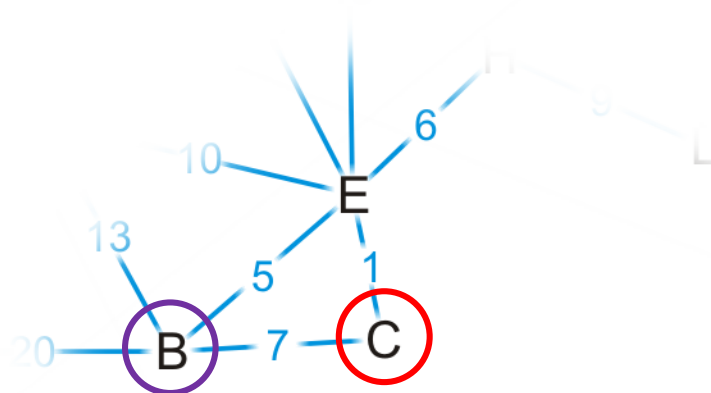
Vertex	Visited	Distance	Previous
A	F	$\infty$	$\emptyset$
B	F	19	E
<b>C</b>	<b>T</b>	<b>15</b>	<b>E</b>
D	F	24	E
E	T	14	H
F	F	17	G
G	T	13	I
H	T	8	K
I	T	10	H
J	F	17	K
K	T	0	$\emptyset$
L	F	16	K



# Example

The path (K, H, E, C, B) is of length  $15 + 7 = 22$

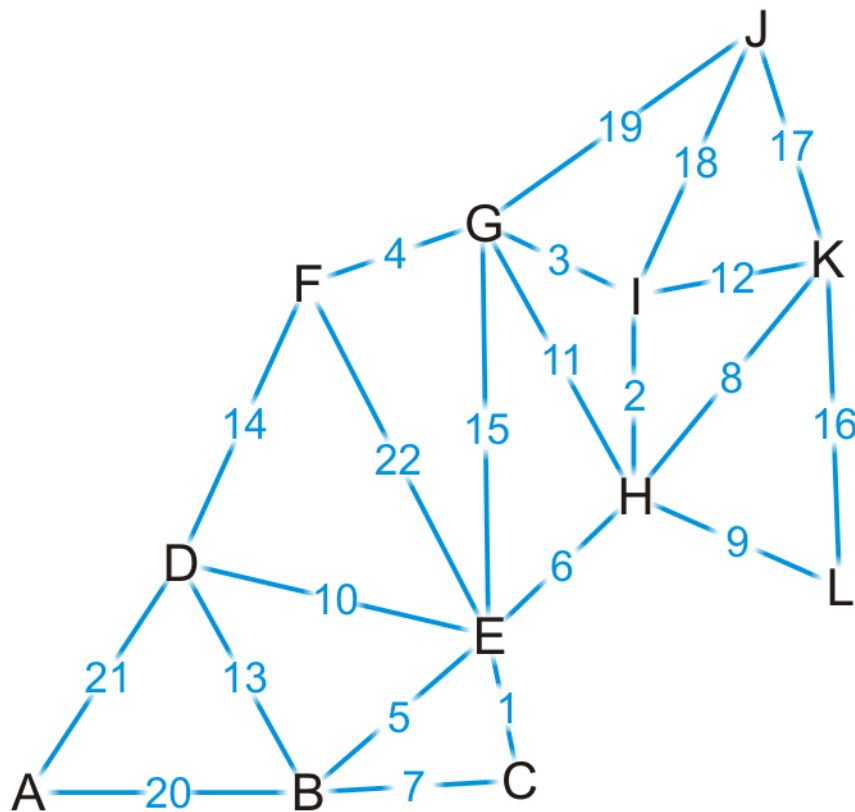
- We have already discovered a shorter path through vertex E



Vertex	Visited	Distance	Previous
A	F	$\infty$	$\emptyset$
B	F	19	E
C	T	15	E
D	F	24	E
E	T	14	H
F	F	17	G
G	T	13	I
H	T	8	K
I	T	10	H
J	F	17	K
K	T	0	$\emptyset$
L	F	16	K

# Example

Where to next?

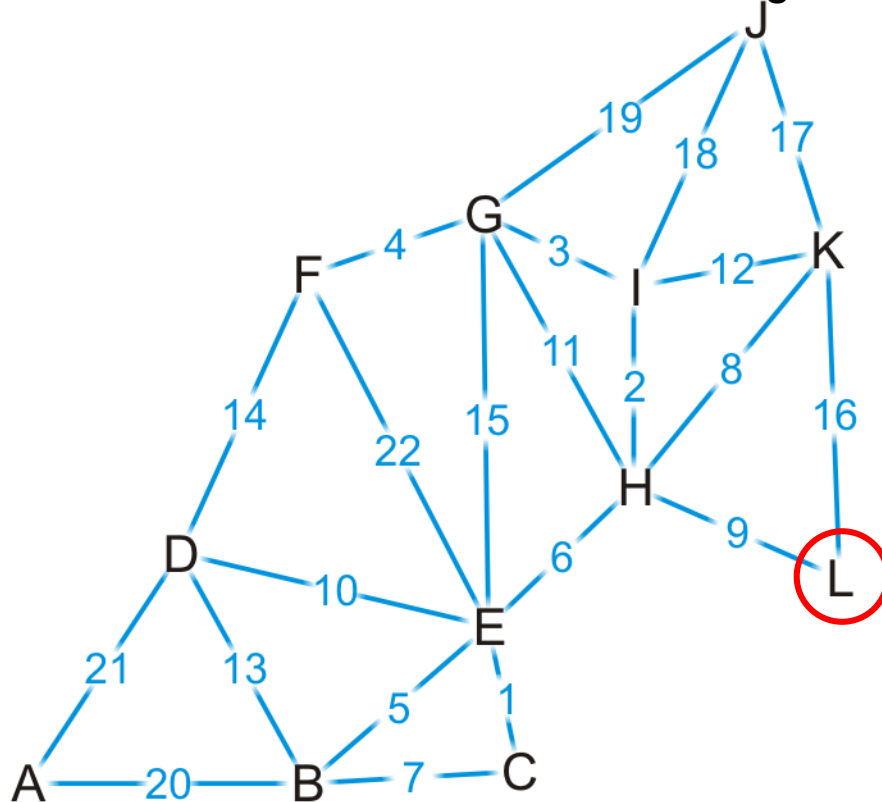


Vertex	Visited	Distance	Previous
A	F	$\infty$	$\emptyset$
B	F	19	E
C	T	15	E
D	F	24	E
E	T	14	H
F	F	17	G
G	T	13	I
H	T	8	K
I	T	10	H
J	F	17	K
K	T	0	$\emptyset$
L	F	16	K

# Example

We now know that (K, L) is the shortest path between these two points

- Vertex L has no unvisited neighbors

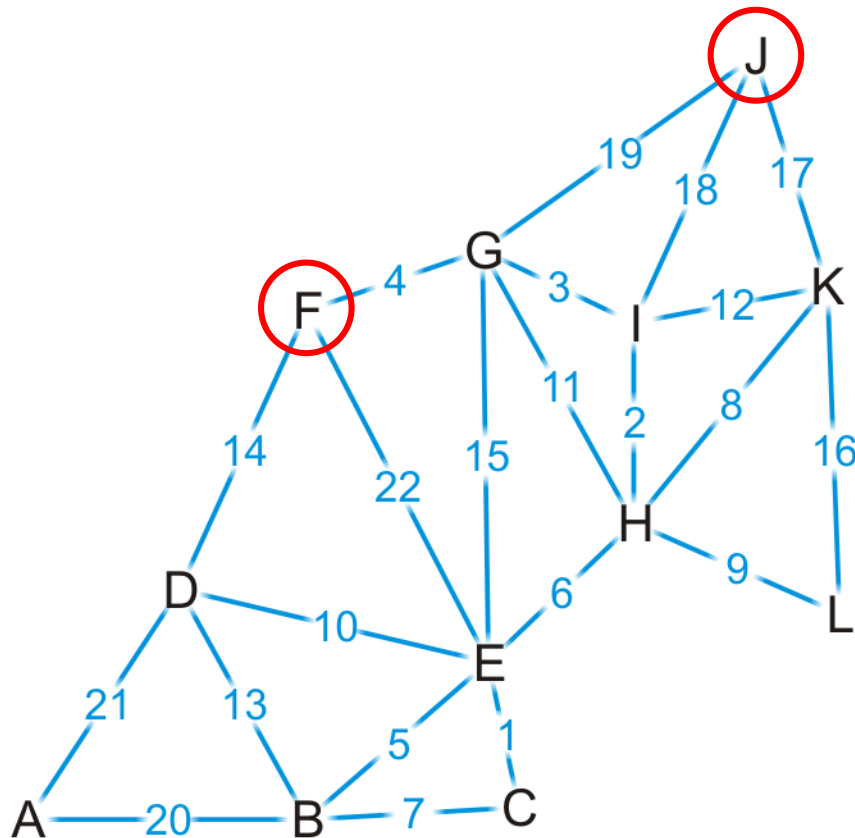


Vertex	Visited	Distance	Previous
A	F	$\infty$	$\emptyset$
B	F	19	E
C	T	15	E
D	F	24	E
E	T	14	H
F	F	17	G
G	T	13	I
H	T	8	K
I	T	10	H
J	F	17	K
K	T	0	$\emptyset$
<b>L</b>	<b>T</b>	<b>16</b>	<b>K</b>

# Example

Where to next?

- Does it matter if we visit vertex F first or vertex J first?

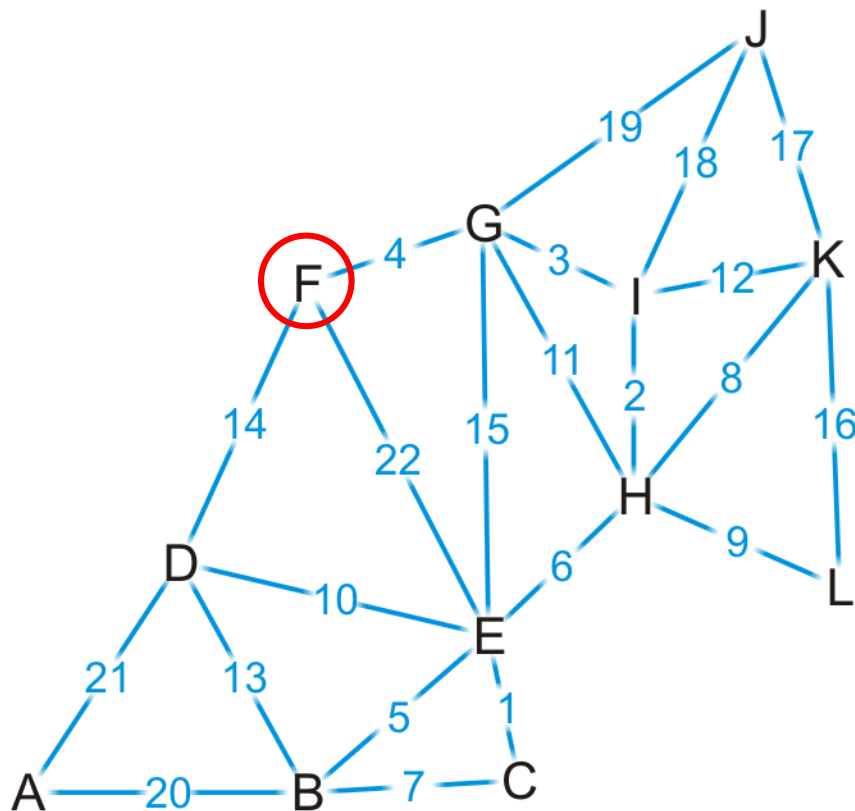


Vertex	Visited	Distance	Previous
A	F	$\infty$	$\emptyset$
B	F	19	E
C	T	15	E
D	F	24	E
E	T	14	H
F	F	17	G
G	T	13	I
H	T	8	K
I	T	10	H
J	F	17	K
K	T	0	$\emptyset$
L	T	16	K

# Example

Let's visit vertex F first

- It has one unvisited neighbor, vertex D

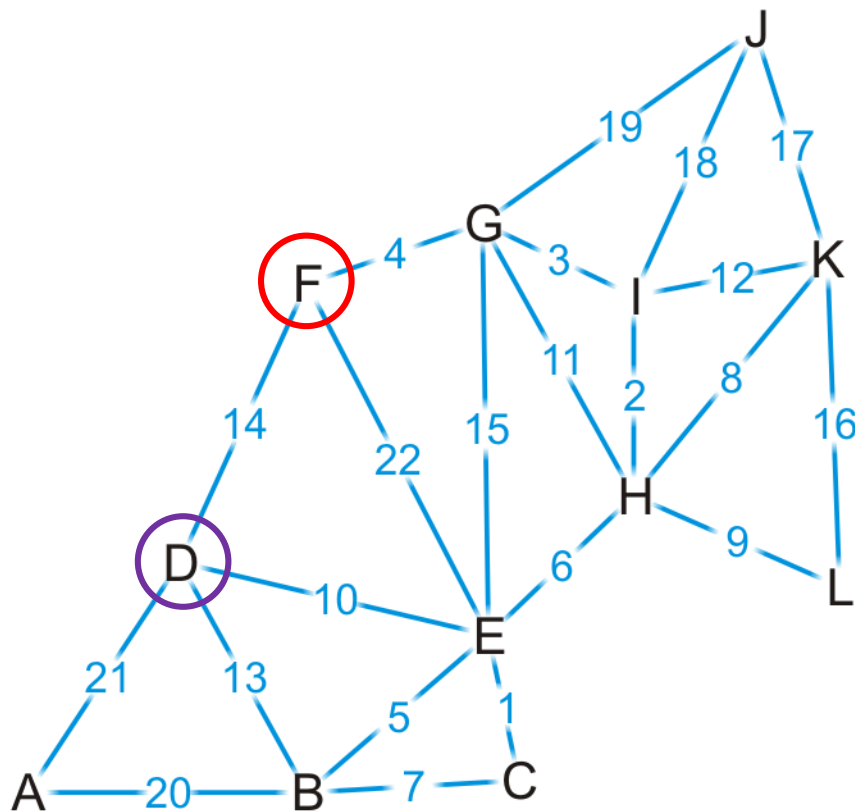


Vertex	Visited	Distance	Previous
A	F	$\infty$	$\emptyset$
B	F	19	E
C	T	15	E
D	F	24	E
E	T	14	H
<b>F</b>	<b>T</b>	<b>17</b>	<b>G</b>
G	T	13	I
H	T	8	K
I	T	10	H
J	F	17	K
K	T	0	$\emptyset$
L	T	16	K

# Example

The path (K, H, I, G, F, D) is of length  $17 + 14 = 31$

- This is longer than the path we've already discovered

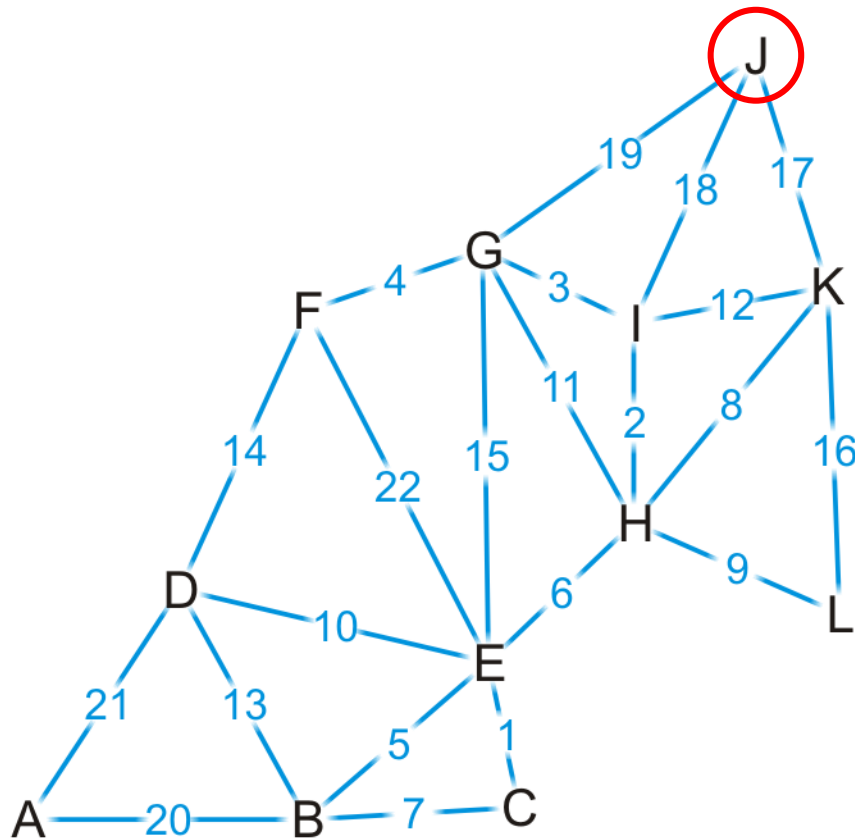


Vertex	Visited	Distance	Previous
A	F	$\infty$	$\emptyset$
B	F	19	E
C	T	15	E
D	F	24	E
E	T	14	H
F	T	17	G
G	T	13	I
H	T	8	K
I	T	10	H
J	F	17	K
K	T	0	$\emptyset$
L	T	16	K

# Example

Now we visit vertex J

- It has no unvisited neighbors



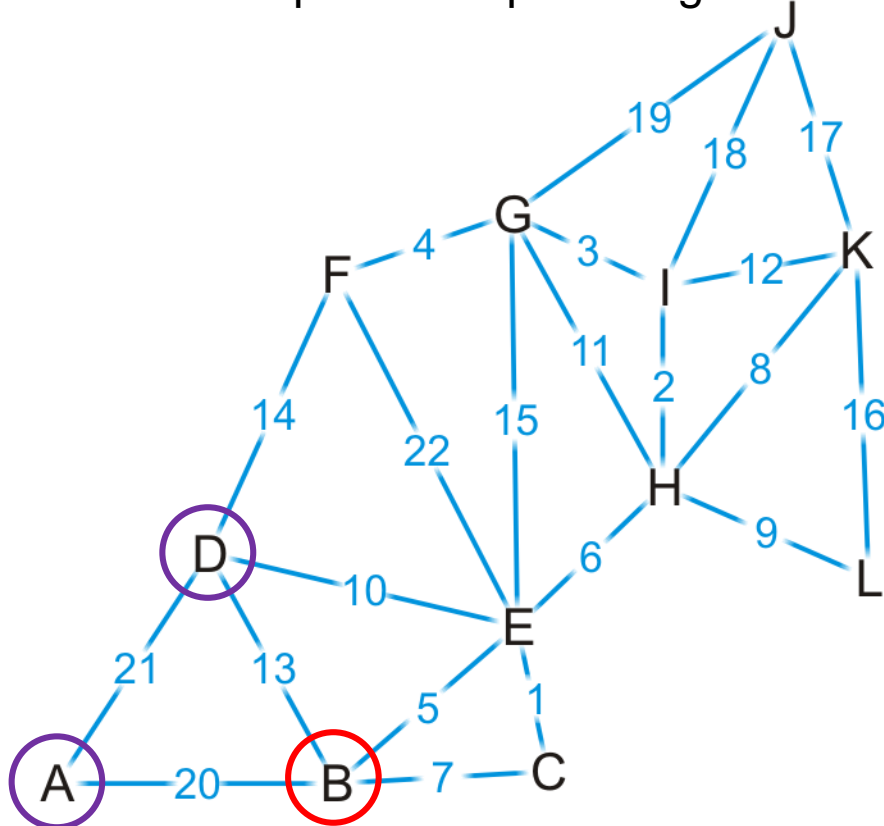
Vertex	Visited	Distance	Previous
A	F	$\infty$	$\emptyset$
B	F	19	E
C	T	15	E
D	F	24	E
E	T	14	H
F	T	17	G
G	T	13	I
H	T	8	K
I	T	10	H
<b>J</b>	<b>T</b>	<b>17</b>	<b>K</b>
K	T	0	$\emptyset$
L	T	16	K

# Example

Next we visit vertex B, which has two unvisited neighbors:

(K, H, E, B, A) of length **19** + 20 = 39      (K, H, E, B, D) of length **19** + 13 = 32

– We update the path length to A



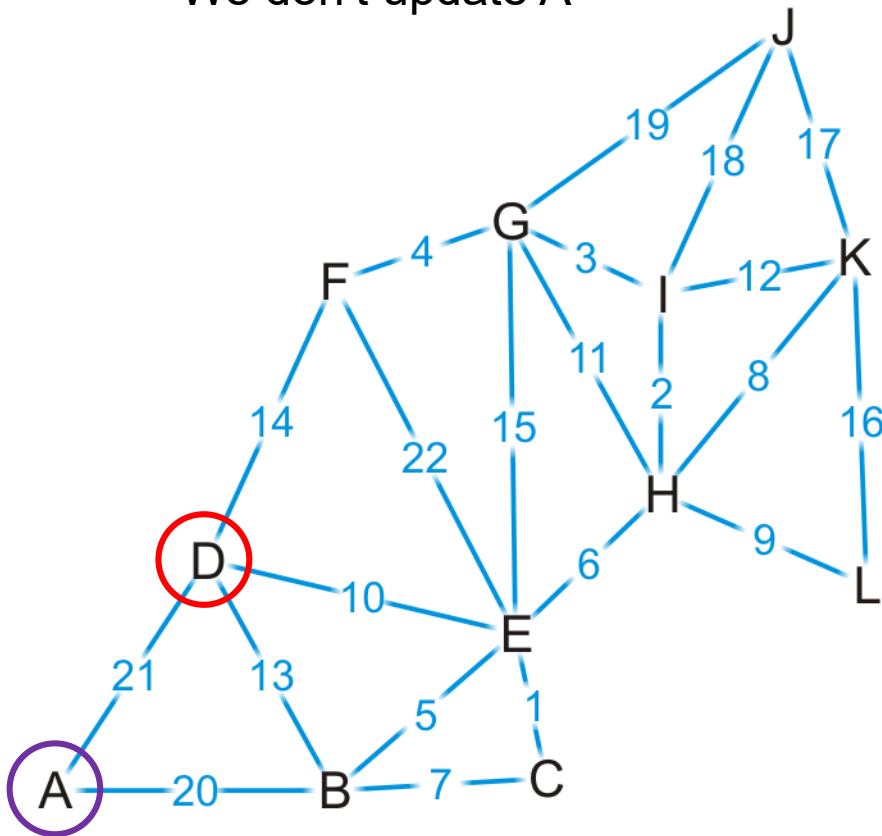
Vertex	Visited	Distance	Previous
A	F	<b>39</b>	<b>B</b>
<b>B</b>	<b>T</b>	<b>19</b>	<b>E</b>
C	T	15	E
D	F	<b>24</b>	<b>E</b>
E	T	14	H
F	T	17	G
G	T	13	I
H	T	8	K
I	T	10	H
J	T	17	K
K	T	0	∅
L	T	16	K



# Example

Next we visit vertex D

- The path (K, H, E, D, A) is of length  $24 + 21 = 45$
- We don't update A

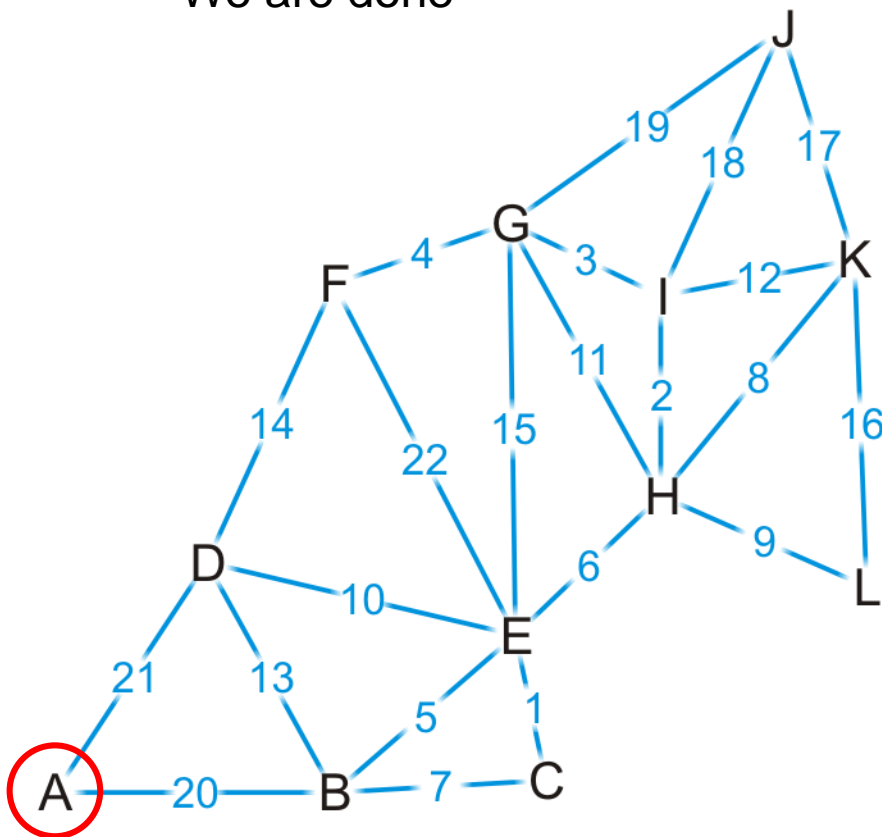


Vertex	Visited	Distance	Previous
A	F	39	B
B	T	19	E
C	T	15	E
D	T	24	E
E	T	14	H
F	T	17	G
G	T	13	I
H	T	8	K
I	T	10	H
J	T	17	K
K	T	0	∅
L	T	16	K

# Example

Finally, we visit vertex A

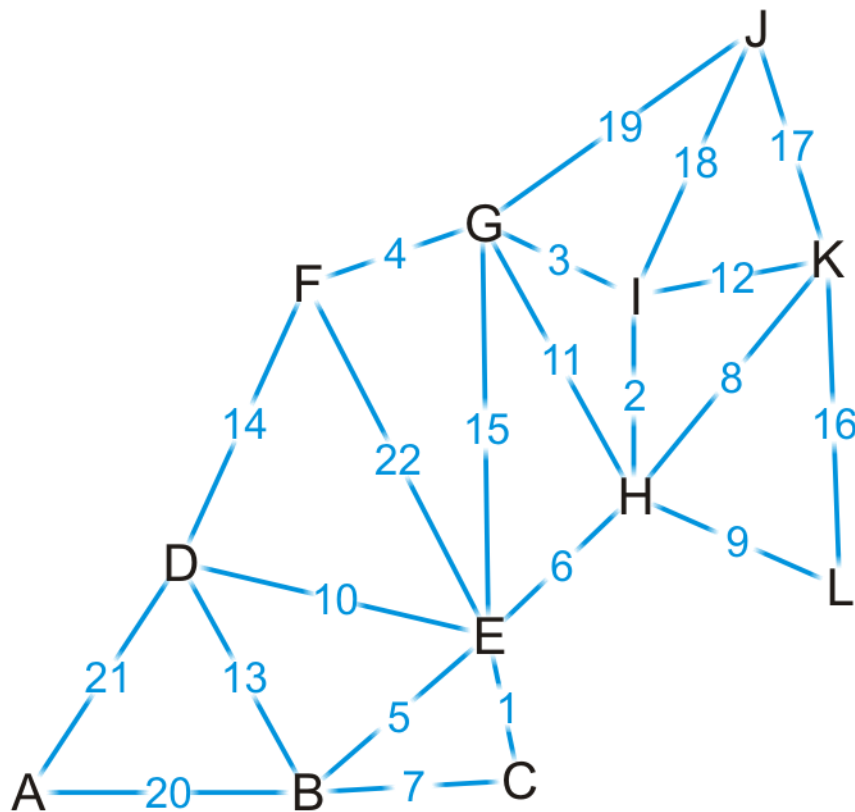
- It has no unvisited neighbors and there are no unvisited vertices left
- We are done



Vertex	Visited	Distance	Previous
<b>A</b>	<b>T</b>	<b>39</b>	<b>B</b>
B	T	19	E
C	T	15	E
D	T	24	E
E	T	14	H
F	T	17	G
G	T	13	I
H	T	8	K
I	T	10	H
J	T	17	K
K	T	0	∅
L	T	16	K

# Example

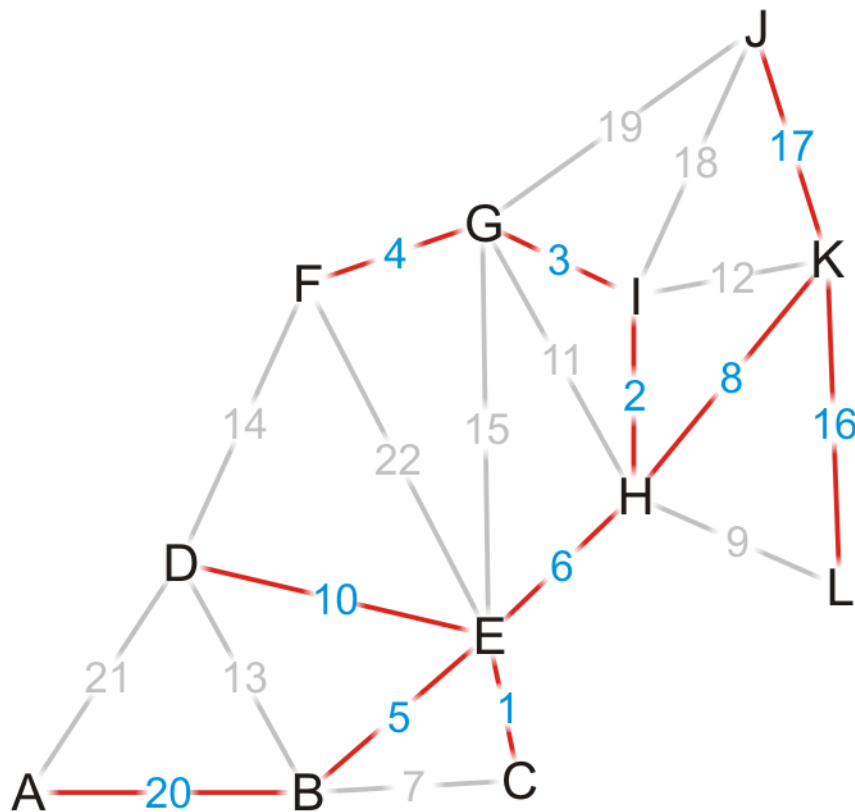
Thus, we have found the shortest path from vertex K to each of the other vertices



Vertex	Visited	Distance	Previous
A	T	39	B
B	T	19	E
C	T	15	E
D	T	24	E
E	T	14	H
F	T	17	G
G	T	13	I
H	T	8	K
I	T	10	H
J	T	17	K
K	T	0	Ø
L	T	16	K

# Example

Using the *previous* pointers, we can reconstruct the paths



Vertex	Visited	Distance	Previous
A	T	39	B
B	T	19	E
C	T	15	E
D	T	24	E
E	T	14	H
F	T	17	G
G	T	13	I
H	T	8	K
I	T	10	H
J	T	17	K
K	T	0	∅
L	T	16	K

# Comments on Dijkstra's algorithm

## Questions:

- Q1. What if at some point, all unvisited vertices have a distance  $\infty$ ?
- Q2. What if we just want to find the shortest path from  $v_j$  to  $v_k$ ?
- Q3. Does the algorithm change if we have a directed graph?

# Runtime Analysis

The runtime of Dijkstra's algorithm is the same as Prim's algorithm

- With an adjacency matrix, the run time is  $\Theta(|V|(|V| + |V|)) = \Theta(|V|^2)$
- With an adjacency list, the run time is  $\Theta(|V|^2 + |E|) = \Theta(|V|^2)$  as  $|E| = O(|V|^2)$

```
for _ in range(|V|-1): // until visiting all vertices
    // visit the closest vertex
    v = table.find_min_dist_vertex()
    table.mark_visit(v)

    // update the shortest path if needed
    for j in graph.get_adj_vertices(v):
        if (table.get_dist(v) + graph.get_dist(v, j) < table.get_dist(j)):
            new_dist = table.get_dist(v) + graph.get_dist(v, j)
            table.set_dist(j, new_dist)
```

Again, using the binary heap may show the better runtime

- $O(|V| \ln(|V|) + |E| \ln(|V|)) = O(|E| \ln(|V|))$

# Summary

We have seen an algorithm for finding single-source shortest paths

- Start with the initial vertex
- Continue finding the next vertex that is closest

Dijkstra's algorithm always finds the next closest vertex

- It solves the problem in  $O(|E| \ln(|V|))$  time

# References

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