

Today's Content

→ Pair Sum = k

→ Distinct elements in every
window of len = k.

Problem Solving → Sunday Morning → 10 am.

Today's Quote :-

'All of man's troubles arise because he can't sit in a room quietly by himself.' If you could just sit of 30 mins & be happy you are very successful.

Ques) Given n array elements, check if there exists a pair (i, j) such that $arr[i] + arr[j] = k$ & $i \neq j$

$arr[] =$ 0 1 2 3 4 5 6 7 8 9
 8 9 1 -2 4 5 11 -6 7 5

$k = 11$ yes $(i, j) = (4, 8)$

$k = 6$ yes $(i, j) = (2, 5)$

$k = 22$ No.

$a + b == k$

idea 1:- check for every pair $sum = k$.

T.C $\rightarrow O(n^2)$ S.C $\rightarrow O(1)$.

```

for (i = 0; i < N; i++) {
    a = arr[i], b = k - a
    for (j = i + 1; j < N; j++) {
        if (arr[j] == b) {
            return True
        }
    }
}
return false

```

arr[] =

0	1	2	3	4	5	6	7	8	9
8	9	1	-2	4	5	11	-6	7	5

→ This will fail.

idea 2 :- Optimization using HashSet.

→ Insert all the elements in HashSet.

HS { 8, 9, 1, -2, 4, 5, 11, -6, 7, 5 }

k=11

$$a + b = 11$$

a	b(k-a)	
8	3	X
9	2	X
1	10	X
-2	13	X

4 7 ✓

{ return true }

k=5

$$a + b = 5$$

a	b(k-a)	
8	-3	X
9	-4	X

1 4 ✓

k=22

$$a + b = 22$$

a	b(k-a)	
8	14	X
9	13	*
1	21	*
-2	24	*

$$a+b=10$$

a b(x-a)

8 2

9 1

4 6

-2 12

4 6

5 5 $a==b, \text{freq}(a) \geq 1, \text{freq}(s) \geq 1$

↓
we can get reqd sum.

T.C $\rightarrow O(n)$, S.C $\rightarrow O(n)$

```
bool pairSum(int arr[], int k) {
    Hmap <int, int> hm;  $\rightarrow$  T.C  $\rightarrow O(n)$ 
    Insert all arr[]  $\rightarrow$  hm // Todo.
    for (i=0; i<n; i++) {  $\rightarrow$  T.C  $\rightarrow O(n)$ 
        a = arr[i], b = k-a
        if (a != b && hm.search(b) == T)
            return True
        3 else if (a == b && hm[a] > 1) {
            return True
            3
        2
    }
    return false
    3
```

idea 4:- optimization using hashset:-

arr[] = ^{0 1 2 3 4 5 6 7 8 9}
8 9 5 -2 11 5 4 -6 7 4

↳ idea:- if we are at ith idx,
insert only [0, i-1]
elements in hash.

a + b = 10

a	b	hash
8	2	{ 2 }
9	1	{ 8, 2 }
5	5	{ 8, 9 }
-2	12	{ 8, 9, 5 }
11	-1	{ 8, 9, 5, -2 }
5	5	{ 8, 9, 5, -2, 11 }

↳ return True.

$$a + b = 11$$

a	b	H ₃
8	3	{ 3 }
9	2	{ 8, 3 }
4	7	{ 8, 9, 3 }
-2	13	{ 8, 9, 4, 3 }
4	7	{ 8, 9, 4, -2, 3 }
5	6	{ 8, 9, 4, -2, 3 }
11	0	{ 8, 9, 4, -2, 5, 3 }
-6	17	{ 8, 9, 4, -2, 5, 11, 3 }
7	4	{ 8, 9, 4, -2, 5, 11, -6, 3 }

$$a + b = 22$$

a	b	H ₃
8	14	{ 3 }
9	13	{ 8, 3 }
4	18	{ 8, 9, 3 }
-2	24	{ 8, 9, 4, 3 }
11	11	{ 8, 9, 4, -2, 3 }

↳ not failing here.

Pseudo Code :- $T.C \rightarrow O(n^2)$, $S.C \rightarrow O(n)$.

```
bool TargetSum( int arr[], int k) {  
    int n = arr.length;  
    HashSet<int> hs;  
    for (i = 0; i < n; i++) {  
        a = arr[i], b = k - a ;  
        if (hs.search(b) == True) {  
            return True  
        }  
        hs.insert(a)  
    }  
    return false  
}
```

Break 9:53 pm to 10:05 pm

Q0) Given n elements, calculate no. of distinct elements in every subarray of size $= k$.

Ex:- arr[10] = ^{0 1 2 3 4 5 6 7 8 9}
2 4 3 8 3 9 7 9 4 10

$k=4$

Subarray

Print

idea | :

[0-3]

4

for every subarray of len = k ,

[1-4]

3

insert all elements into

[2-5]

3

hashset, and get no. of

[3-6]

4

distinct elements.

[4-7]

3

$\Rightarrow \boxed{(n-k+1) \cdot k}$ when $k = \frac{n}{2}$.

[5-8]

2

$\Rightarrow (n - \frac{n}{2} + 1) (\frac{n}{2}) \Rightarrow \frac{(n^2)}{2}$
T.C \rightarrow

[6-9]

3

S.C $\rightarrow O(k)$.

first subarray

Last subarray.

0

x $\Rightarrow n-k$

a b

\Downarrow

[x n-1] = k

Total

$n-x - x + 1 = k$

Subarrays

$n-x = k$

= $n-k+1$

$n-k = x$

idea 2 :- Optimization using hashset.

Ex:- arr[10] = ^{0 1 2 3 4 5 6 7 8 9}
2 4 3 8 3 9 4 9 4 10
k=4

Subarray

(0-3) \longrightarrow HS = {2, 4, 3, 8} : 4
(1-4) del arr[0] insert arr[4] HS = {4, 3, 8} : 3
(2-5) del arr[1] ins arr[5] HS = {3, 8, 9} : 3
(3-6) del arr[2] ins arr[6] HS = {8, 9, 4} : 3

Note :- In hs, deleting an element will
indirectly delete all occurrences.

idea 3 :- Optimization using hashmap.

Ex:- arr[10] = ^{0 1 2 3 4 5 6 7 8 9}
2 4 3 8 3 9 4 9 4 10
k=4

Subarray

Hmap

(0-3) { <2, 1> <4, 1> <3, 1> <8, 1> } : 4

(1-4) del (0) add (4)

{ <4, 1> <3, 2> <8, 1> } : 3

(2-5) del (1) add (5) <3, 2> <8, 1> <9, 1> : 3

(3-6) del (2) add (6) <3, 1> <8, 1> <9, 1> <4, 1> : 4

(4 - 7) del(3) add(7) <3,1><9,2><4,1> : 3

(5 - 8) del(4) add(8) <9,2><4,2> : 2

(6 - 9) del(5) add(9) <9,1><4,2><10,1> : 3

(s e) → del(s-1), add(e)

→ T.C → $O(n)$ S.C → $O(k)$.

```
void distinctMap ( int arr[], int n ) {
```

```
    int m = arr.length;
```

```
    Hash Map <int, int> hm;
```

```
    // Insert first k elements.
```

```
    for ( i = 0; i < k; i++ ) {      →  $O(k)$ 
```

```
        if ( hm.search ( arr[i] ) == true ) {
```

```
            hm [ arr[i] ] + = 1
```

```
        } else {
```

```
            hm.insert ( arr[i], 1 )
```

```
        }
```

```
    }
```

```
    print ( hm.size() );
```

```
    s = 1, e = k
```

```
    while ( e < n ) →  $(n - k)$ .
```

```
    // sub [s e]
```

```
    hm [ arr[s-1] ] - = 1
```

```
    if ( !hm [ arr[s-1] ] == 0 ) {
```

```
        // remove key arr[s-1]
```

```
        hm.remove ( arr[s-1] )
```

```
    }
```

```

if (hm.containsKey(x[i])) {
    hm.put(x[i], hm.get(x[i]) + 1);
} else {
    hm.put(x[i], 1);
}
i++;
}
return hm.size();

```

3

3

Ques) N 2d points, calculate no. of distinct points.

x[5] = { 2, 3, 2, 5, 3 } y[5] = { 3, 6, 3, 7, 6 } Ans = 9

2, 13
2, 13

hm < >

has <String>

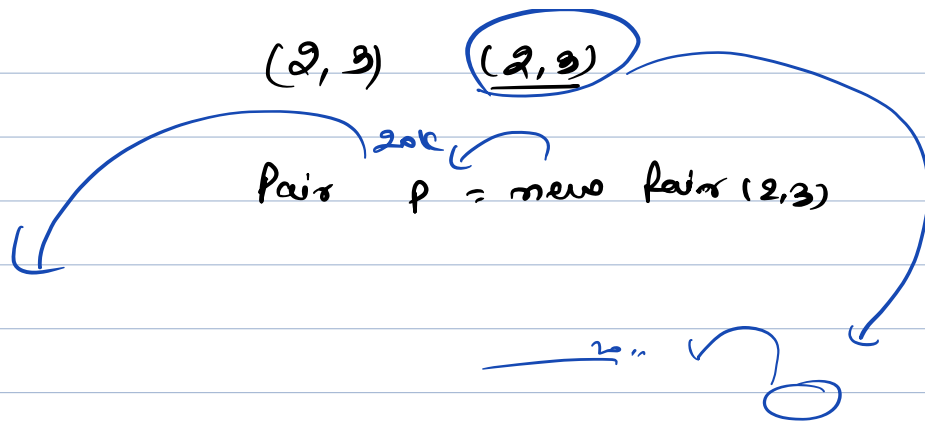
x + @ + y

(3, 6)

(2, 3)

203 306
507 307
302

23 36
57
213 213



$$\begin{array}{r} 123 \rightarrow \text{true} \\ \hline 6 \end{array}$$

$$234 \rightarrow 9$$

$$33 \rightarrow \underline{6}$$