

## Today's Content

T.C  $\rightarrow$  2  $\left\{ \begin{array}{l} \rightarrow \text{Big O} - O() \\ \rightarrow \text{TLE} - (\text{Time limit Exceeded error}) \\ \rightarrow \text{why TLE?} \end{array} \right.$

T.C  $\rightarrow$  1  $\{$  How to calculate iterations 3  $\rightarrow$  15/12  
Question

$\leftarrow$  Quote  $\rightarrow$

Don't tell your god how big your storm is.

Tell your storm how big your god is.

## Basic maths Revision

Ques 1: Sum of Natural numbers:

$$S_N = 1 + 2 + 3 + \dots + N = \frac{n(n+1)}{2}$$

$$\rightarrow : [3 \ 10] \rightarrow \{3, 4, 5, 6, 7, 8, 9, 10\} \rightarrow 8$$

$$[4 \ 8] = \{4, 5, 6, 7, 8\} \rightarrow 5$$

# no. of elements from  $[a \ b]$  included

$$: [a \ b] = b - a + 1 \quad \# \text{imp.}$$

GP

(when we divide any 2 cons. elements, ratio have to be same)  
 $\rightarrow$  Geometric Progression

$$S_1 = 3 \quad 6 \quad 12 \quad 24 \quad 48 \quad 96$$
$$\text{ratio} = \frac{6}{3} = \frac{12}{6} = \frac{24}{12} = \frac{48}{24} = \frac{96}{48} = 2$$

$$S_2 = 2 \quad 6 \quad 18 \quad 54 \quad 162$$
$$\text{ratio} = \frac{6}{2} = \frac{18}{6} = \frac{54}{18} = \frac{162}{54} = 3$$

// given gp

first term =  $a$

Common ratio =  $r$

1st 2nd 3rd ....  $n^{\text{th}}$   
 $a$   $ar$   $ar^2$   $ar^{n-1}$

$$S_n = a + ar + ar^2 + \dots + ar^{n-2} + ar^{n-1}$$

multiply both sides with  $r$ :-

$$S_n r = ar + ar^2 + ar^3 + \dots + ar^{n-1} + ar^n$$

$$-(S_n = a + ar + ar^2 + \dots + ar^{n-2} + ar^{n-1})$$

---

$$rS_n - S_n = ar^n - a$$

---

$$\Rightarrow S_n(r-1) = a(r^n-1)$$

$$\Rightarrow S_n = \frac{a(r^n-1)}{r-1}$$

Ex:-  $a=5, r=2, n=5$

$$\rightarrow 5, 10, 20, 40, 80 \Rightarrow 155$$

Let's use formula

$$\frac{a(r^n-1)}{r-1} \Rightarrow \frac{5(2^5-1)}{2-1} \Rightarrow 5 \times 31 \Rightarrow 155$$

with

Q1 *Ques* void func (int n) {

int S = 0;

(i = 0; i <= 100; i++) { *i [0, 100] → 100 - 0 + 1*

S = S + i

*⇒ 101*

→ 04)

3

Q2)

void func (int n) {

int S = 0;

(i = 35; i <= 87; i++) { *i [35 : 87] → 87 - 35 + 1*

S = S + i

*⇒ 59*

04)

3

Q3: *Ques*

void func (int n) {

int S = 0;

(i = 1; i <= n; i++) { *i [1 : n] → n - 1 + 1*

S = S + i

*⇒ n*

3

3

Q4:

```
void func (int n) {
```

```
    int S = 0;
```

```
    (i=1; i<=n; i++) {  $\rightarrow i[1:n] \rightarrow n-x+x \rightarrow n$ 
```

```
        if (i%2==0) {
```

```
            S = S+i
```

```
        }
```

```
    }
```

```
    (i=1; i<=m; i++) {  $\rightarrow i[1:m] \rightarrow m-x+x \rightarrow m$ 
```

```
        if (i%2==1) {
```

```
            S = S+i
```

```
        }
```

```
    }
```

```
}
```

$\rightarrow$

$n+m$

$a-b \rightarrow \underline{b-a+1}$

06) Quiz

```
void func (int n) {
```

$$i^2 \leq n \Rightarrow i \leq \sqrt{n}$$

```
    int S = 0;
```

```
    (i = 1; i * i <= n; i++) {
```

$i[1:\sqrt{n}]$   $\sqrt{n} - x + x$

```
        S = S + i
```

$\approx \sqrt{n}$  iter

```
    }
```

```
}
```

Q7)

$i = n$

void func (int n) {

iteration

i value after iteration

int i = n

1

$\frac{n}{2} : \frac{n}{2^1}$

while (i > 1) {

2

$\frac{n}{4} : \frac{n}{2^2}$

    i = i/2

3

$\frac{n}{8} : \frac{n}{2^3}$

3

4

$\frac{n}{16} : \frac{n}{2^4}$

3

k iterations

$\frac{n}{2^k}$

n

i

10 10 → 5 → 2 → 1 {break}

19 19 → 9 → 4 → 2 → 1 {break}

$$\frac{n}{2^k} = 1 \Rightarrow n = 2^k$$

$$\log_2 n = \log_2 2^k$$

$\log a^b = b \log a$

$$\Rightarrow \log_2 n = k \log_2 2$$

$$\Rightarrow \log_2 n = k$$

// After  $\log_2 n$  iterations code breaks.

ques)

```
void func (int n) {
```

```
    int i = n
```

```
    while (i > 0) {
```

```
        | i = i/2
```

```
    }
```

```
}
```

→ Todo & Try it out.

→ No. of iterations.



Q8)

Quiz

```
void func (int n) {
```

```
    int S = 0;
```

```
    (i = 0; i <= n; i = i * 2) {
```

```
        S = S + i
```

```
    }
```

```
}
```

i = 0	
iterations	value of i
1	0
2	0
3	0
	⋮
	∞

infinite

Quiz

i = 1

90) → { Near approximations }

iterations	value after each iteration
1	$i = 2 = 2^1$
2	$i = 4 = 2^2$
3	$i = 8 = 2^3$
4	$i = 16 = 2^4$
⋮	
after $k$ iterations	$i = 2^k$

$i > N$  it'll break.

$2^k > N$  after  $k$  iterations.

$$\log_2 2^k > \log_2 N \Rightarrow k \log_2 2 > \log_2 N$$

$$\Rightarrow k > \log_2 N$$

Break

100)

```
void func (int N) {
```

```
    (i = 1 ; i <= 4 ; i = i + 1) {
```

```
        J = 1 ; J <= 3 ; J++ ) {
```

```
            Print (" Hello World ");
```

```
        }
```

```
    }
```

```
}
```

i	J	Total iter.
1	[1-3]	3
2	[1-3]	3
3	[1-3]	3
4	[1-3]	3
		<u>12</u>

iterations.

110) Quiz

```
void func (int N) {
```

```
    (i = 1 ; i <= 10 ; i = i + 1) {
```

```
        (J = 1 ; J <= N ; J++ ) {
```

```
            Print (" Hello World ");
```

```
        }
```

```
    }
```

```
}
```

i	J	Total iter
1	[1 N]	2
2	[1 N]	2
3	[1 N]	2
...	...	...
10	[1 N]	2
		<u>20</u>
		20

10:08:10:18 pm

12)

*Quiz*

```
void func (int N) {
```

```
    (i = 1 ; i <= N ; i = i + 1) {
```

```
        (j = 1 ; j <= N ; j++) {
```

```
            Print ("Hello World");
```

```
        }
```

```
    }
```

```
}
```

i	J	Total iter
1	[1 N]	2
2	[1 N]	2
3	[1 N]	2
⋮		
N	[1 N]	2
		<u>2</u>
		<b><math>N * 2</math></b>

iterations.

13)

*Quiz*

```
void func (int N) {
```

```
    (i = 1 ; i <= N ; i = i + 1) {
```

```
        (j = 1 ; j <= i ; j++) {
```

```
            Print ("Hello World");
```

```
        }
```

```
    }
```

```
}
```

i	J	Total iter
1	[1 - 1]	1
2	[1 - 2]	2
3	[1 - 3]	3
⋮		+
N	[1 - N]	N
		<u>2</u>

**$\frac{n * (n + 1)}{2}$**



iterations  $\rightarrow \frac{n^2}{2} + \frac{n}{2} \rightarrow O(n^2)$

14)

Ques 3

```
void func (int N) {
```

```
    (i = 1 ; i <= N ; i = i + 1) {
```

```
        (j = 1 ; j <= N ; j = j + 2) {
```

```
            Print (" Hello World ");
```

```
        }
```

```
    }
```

```
}
```

i	J	Total iter
1	[1 to N]	$\log N$
2		$\log N$
⋮		
N		$\log N$
		<u><math>N * \log N</math></u>

Q15)

```
void func (int N) {
```

```
    (i = 1 ; i <= 2n ; i = i + 1) {
```

```
        Print ( )
```

```
    }
```

```
}
```

i [1 : 2<sup>n</sup>] → 2<sup>n</sup> - 1 + 1

→ 2<sup>n</sup> iterations

Q 16)

```
void func (int N) {
```

```
    (i=1; i<=N; i=i+1) {
```

```
        (j=1; j<=2i; j=j+1) {
```

```
            Print("Hello World");
```

```
        }
```

```
    }
```

```
}
```

i	J	Total iter
1	[1 to 2 <sup>1</sup> ]	2 <sup>1</sup>
2	[1 to 2 <sup>2</sup> ]	2 <sup>2</sup>
3	[1 to 2 <sup>3</sup> ]	2 <sup>3</sup>
⋮		
N	[1 to 2 <sup>N</sup> ]	2 <sup>N</sup>

$$2 + 2^2 + 2^3 + \dots + 2^N$$

$$\curvearrowright a = 2, a = 2, n = n$$

$$\frac{a(n^n - 1)}{a - 1} \approx \frac{2(2^n - 1)}{2 - 1} \approx 2(2^n - 1)$$

iterations

Comparison functions : { large  $N$  values }

$$\log N < \sqrt{N} < N < N \log N < N \sqrt{N} < N^2 < N^3 < 9^N$$

$\downarrow$   $N^{1/2}$   $\downarrow$   $N \times N^{1/2}$

Big O  $\rightarrow$  what ?

$\rightarrow$  why ?

$\rightarrow$  where do we use it ?

$\rightarrow$  How do we calculate Big O for any, code

$\rightarrow$  Calculate iterations

$\rightarrow$  Only take higher order term

$\rightarrow$  Neglect constant coefficients.

En1:- iterations  $\rightarrow 3n^2 + 5n + 10^4 \rightarrow O(n^2)$

En2: iterations  $\rightarrow 5n^2 + 10n^3 + 6n \log n + 100 \rightarrow O(n^3)$

En3: iterations  $\rightarrow 4n^2 + 3n + 10^6 \rightarrow O(n^2)$

En4: iterations  $\rightarrow 4n + 3n \log n + 10^6 \rightarrow O(n \log n)$

En5: iterations  $\rightarrow 10^3 \rightarrow \underline{O(1)} \rightarrow \text{constant.}$