

CS450 Computer Networks

The slides used in class are derived from the slides available on our text book companion website:

http://wps.pearsoned.com/ecs_kurose_compnetw_6/

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CS450 Computer Networks

Lesson 2

Application Layer Overview

Consciousness is the field of all
possibilities

Lecture 2

Application Layer Overview

Our goal:

Understand key Application Layer concepts:

- ❑ application architectures
 - client-server paradigm
 - peer-to-peer paradigm
 - hybrid client-server/P2P
- ❑ application process communication
- ❑ what services do applications need?

Some network apps

- ❖ e-mail
- ❖ web
- ❖ instant messaging
- ❖ remote login
- ❖ P2P file sharing
- ❖ multi-user network games
- ❖ streaming stored video (YouTube)
- ❖ voice over IP
- ❖ real-time video conferencing
- ❖ cloud computing
- ❖ IPTV
- ❖ ...
- ❖

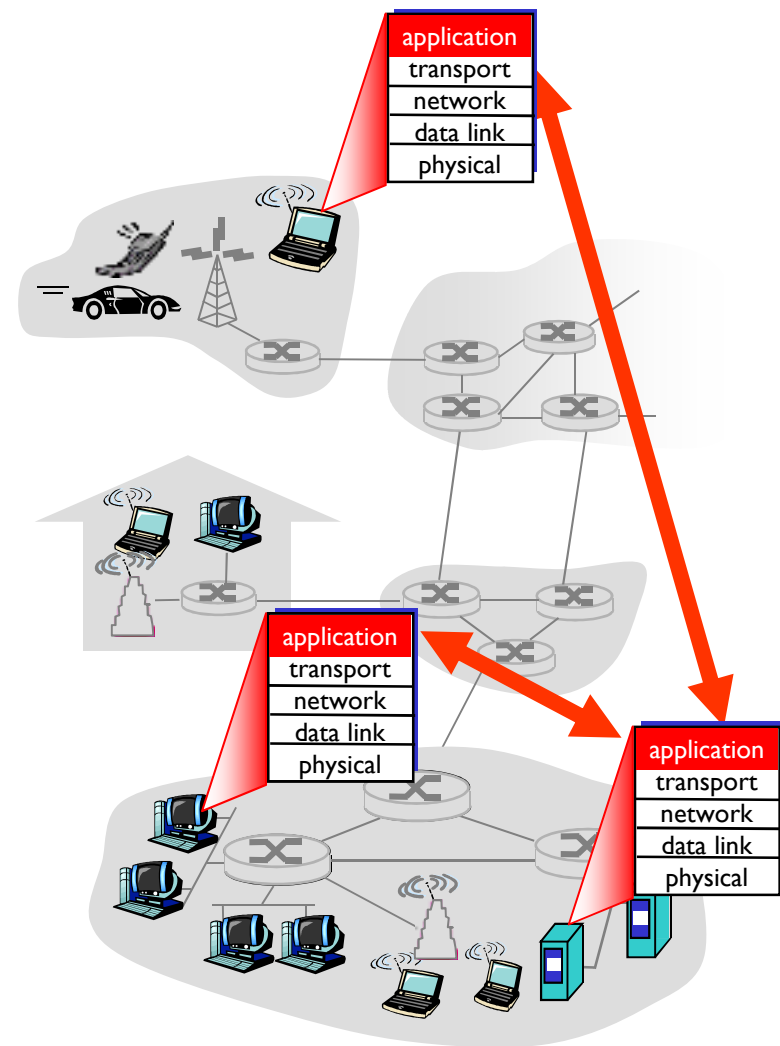
Creating a network app

write programs that

- run on (different) *end systems*
- communicate over network
- e.g., web server software communicates with browser software

No need to write software for network-core devices

- network-core devices do not run user applications
- applications on end systems allows for rapid app development, propagation

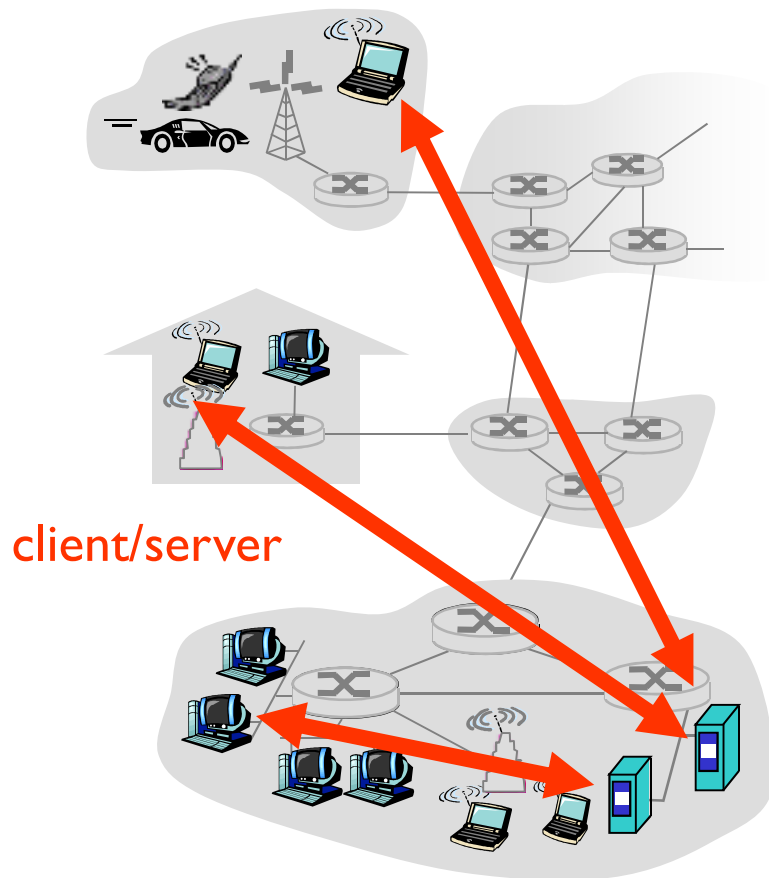


Application architectures

possible structure of applications:

- ❖ client-server
- ❖ peer-to-peer (P2P)

Client-server architecture



server:

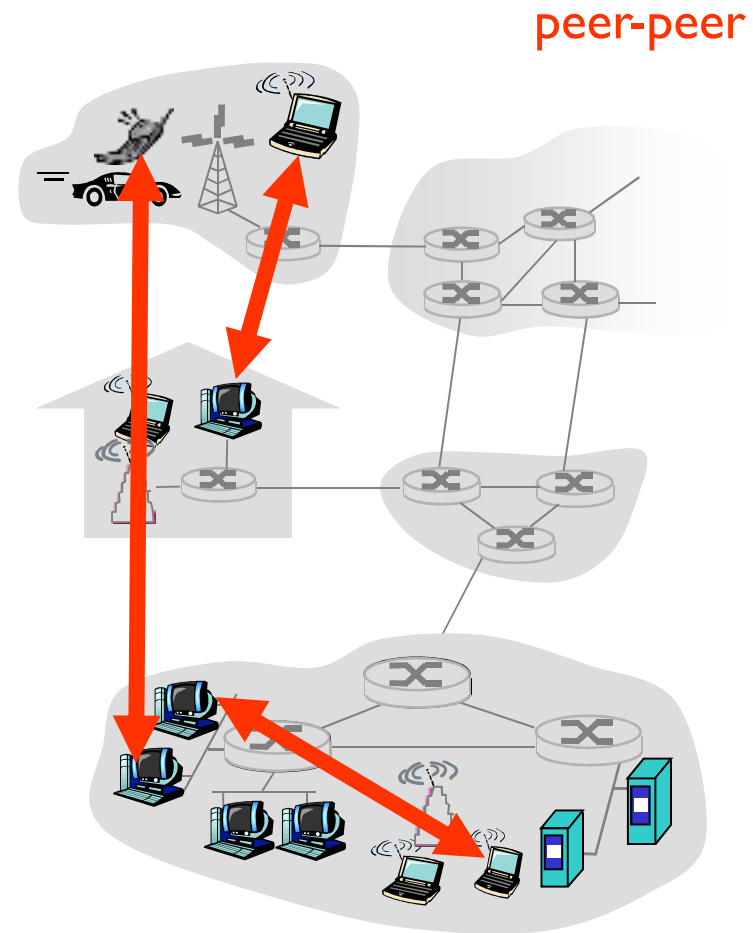
- always-on host
- permanent IP address
- server farms for scaling

clients:

- communicate with server
- may be intermittently connected
- may have dynamic IP addresses
- do not communicate directly with each other

P2P architecture

- ❖ no always-on server
- ❖ arbitrary end systems directly communicate
- ❖ peers request service from other peers, provide service in return to other peers
 - *self scalability* – new peers bring new service capacity, as well as new service demands
- ❖ peers are intermittently connected and change IP addresses
 - complex management



Hybrid of client-server and P2P

Skype

- voice-over-IP P2P application
- centralized server: finding address of remote party:
- client-client connection: direct (not through server)

Instant messaging

- chatting between two users is P2P
- centralized service: client presence detection/location
 - user registers its IP address with central server when it comes online
 - user contacts central server to find IP addresses of buddies

Processes communicating

process: program running within a host.

- ❖ within same host, two processes communicate using **inter-process communication** (defined by OS).
- ❖ processes in different hosts communicate by exchanging **messages**

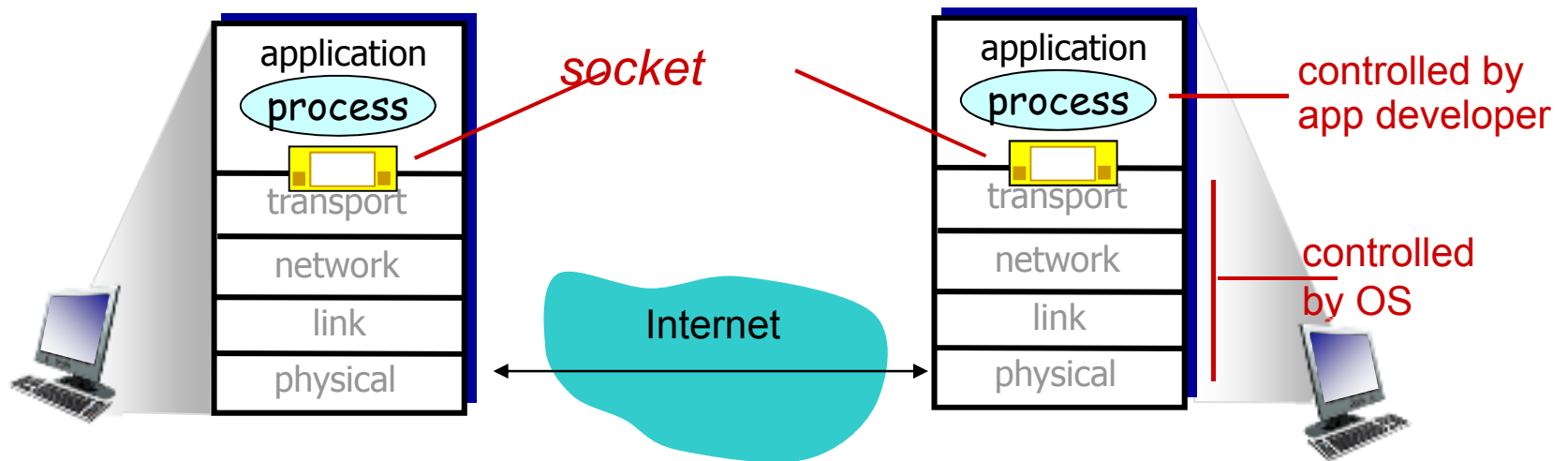
client process: process that initiates communication

server process: process that waits to be contacted

- ❖ aside: applications with P2P architectures have client processes & server processes

Sockets

- ❖ process sends/receives messages to/from its **socket**
- ❖ socket analogous to door
 - sending process shoves message out door
 - sending process relies on transport infrastructure on other side of door to deliver message to socket at receiving process



Addressing processes

- ❖ to receive messages, process must have *identifier*
- ❖ host device has unique 32-bit IP address
- ❖ Q: does IP address of host on which process runs suffice for identifying the process?

Addressing processes

- ❖ to receive messages, process must have *identifier*
- ❖ host device has unique 32-bit IP address
- ❖ Q: does IP address of host on which process runs suffice for identifying the process?
 - A: No, *many* processes can be running on same host
- ❖ *identifier* includes both **IP address** and **port numbers** associated with process on host.
- ❖ example port numbers:
 - HTTP server: 80
 - Mail server: 25
- ❖ to send HTTP message to gaia.cs.umass.edu web server:
 - **IP address**: 128.119.245.12
 - **Port number**: 80
- ❖ more shortly...

App-layer protocol defines

- ❖ types of messages exchanged,
 - e.g., request, response
- ❖ message syntax:
 - what fields in messages & how fields are delineated
- ❖ message semantics
 - meaning of information in fields
- ❖ rules for when and how processes send & respond to messages

public-domain protocols:

- ❖ defined in RFCs
- ❖ allows for interoperability
- ❖ e.g., HTTP, SMTP

proprietary protocols:

- ❖ e.g., Skype

What transport service does an app need?

Data Integrity

- ❖ some apps (e.g., audio) can tolerate some loss
- ❖ other apps (e.g., file transfer, telnet) require 100% reliable data transfer

Timing

- ❖ some apps (e.g., Internet telephony, interactive games) require low delay to be “effective”

Throughput

- ❖ some apps (e.g., multimedia) require minimum amount of throughput to be “effective”
- ❖ other apps (“elastic apps”) make use of whatever throughput they get

Security

- ❖ encryption, data integrity, ...

Transport service requirements of common apps

Application	Data loss	Throughput	Time Sensitive
file transfer			
e-mail			
Web documents			
real-time audio/video			
stored audio/video			
interactive games			
instant messaging			

Transport service requirements of common apps

Application	Data loss	Throughput	Time Sensitive
file transfer	no loss	elastic	no
e-mail	no loss	elastic	no
Web documents	no loss	elastic	no
real-time audio/video	loss-tolerant	audio: 5kbps-1Mbps video: 10kbps-5Mbps	yes, 100's msec
stored audio/video	loss-tolerant	same as above	yes, few secs
interactive games	loss-tolerant	few kbps up	yes, 100's msec
instant messaging	no loss	elastic	yes and no

Internet transport protocols services

TCP service:

- ❖ *connection-oriented*: setup required between client and server processes
- ❖ *reliable transport* between sending and receiving process
- ❖ *flow control*: sender won't overwhelm receiver
- ❖ *congestion control*: throttle sender when network overloaded
- ❖ *does not provide*: timing, minimum throughput guarantees, security

UDP service:

- ❖ unreliable data transfer between sending and receiving process
- ❖ does not provide: connection setup, reliability, flow control, congestion control, timing, throughput guarantee, or security

Q: why bother? Why is there a UDP?

Securing TCP

TCP & UDP

- ❖ no encryption
- ❖ cleartext passwds sent into socket traverse Internet in cleartext

SSL

- ❖ provides encrypted TCP connection
- ❖ data integrity
- ❖ end-point authentication

SSL is at app layer

- ❖ Apps use SSL libraries, which “talk” to TCP

SSL socket API

- ❖ cleartext passwds sent into socket traverse Internet encrypted
- ❖ See Chapter 7

Internet apps: application, transport protocols

Application	Application layer protocol	Underlying transport protocol
e-mail	SMTP [RFC 2821]	TCP
remote terminal access	Telnet [RFC 854]	TCP
Web	HTTP [RFC 2616]	TCP
file transfer	FTP [RFC 959]	TCP
streaming multimedia	HTTP (e.g., YouTube), RTP [RFC 1889]	TCP or UDP
Internet telephony	SIP, RTP, proprietary (e.g., Skype)	typically UDP