

* Priority Queue

Priority Queue is a container adaptor that provides constant time lookup of the largest (by default) element, at the expense of logarithmic insertion and extraction.

A user provided compare can be supplied to change the ordering. Eg - using `std::greater<T>` would cause the smallest element to appear as the top.

STL priority queue is the implementation of heap Data structure.

PS: Container Adaptor - Container adaptors provide a different interface for sequential container. limited accessibility. Eg - stack, queue, priority queue.

Sequential containers - Sequence containers implement data structures which can be accessed sequentially. Eg - array, vectors, list.

- Creating max Heap -

```
priority_queue<int> maxHeap;
```

By default, priority queue max heap create ~~parti~~ hai, that's why container type and comparator like ~~ki~~ need nahi hai.

- Creating min Heap -

```
priority_queue<int, vector<int>, greater<int>> minHeap;
```

- = int = defines type of stored element.
- = vector<int> = container type.
- = greater<int> = comparator.

Accessibilities -

Push

Pop

top

size

empty

* **kth smallest element (gfg) -**

I/P - arr[] = 7 10 4 3 20 15

else - K = 3

l = 0, r = 5 (l and r denotes starting and ending index of array)

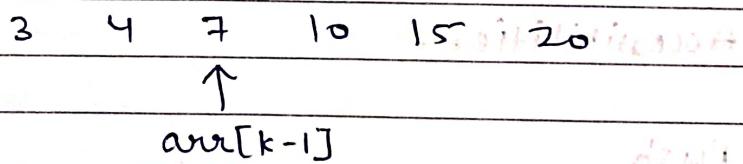
O/P - 7

Explanation -

3	4	7	10	15	20
0	1	2	3	4	5

- 3rd smallest value consists of arr[2]
- 7th smallest value consists of arr[6]

Approach 1 - Sort the array. Return arr[K-1].



$$T.C. = O(n \log n)$$

Approach 2 - Create min heap. Pop from min heap (K-1) times. Now, the top of min heap has kth smallest element.

$$T.C. = O(n) + O((K-1) \log n)$$

$$S.C. = O(n)$$

Approach 3 - Space optimised approach

Create max-heap of first k elements in array to maintain initial state of heap.

For remaining elements, compare the top of max heap with current element in array.

```
if (maxHeap.top() > arr[i]) {
```

```
    maxHeap.pop();
```

```
    maxHeap.push(arr[i]);
```

```
}
```

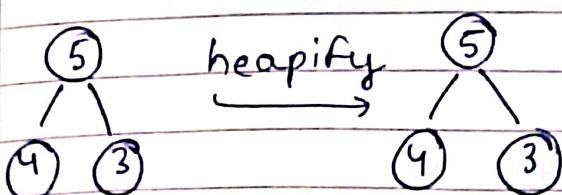
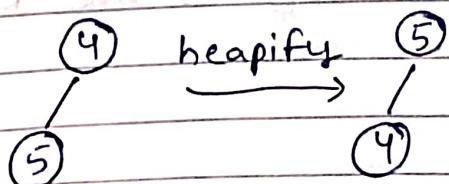
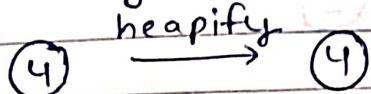
So, at last, the top of max Heap is the k th smallest element.

$$T.C. = O(k) + O((n-k)\log k) \approx O(n \log k)$$

$$S.C. = O(k)$$

= let array, $arr[] = \{4, 5, 3, 1, 2\}$, $K = 3$

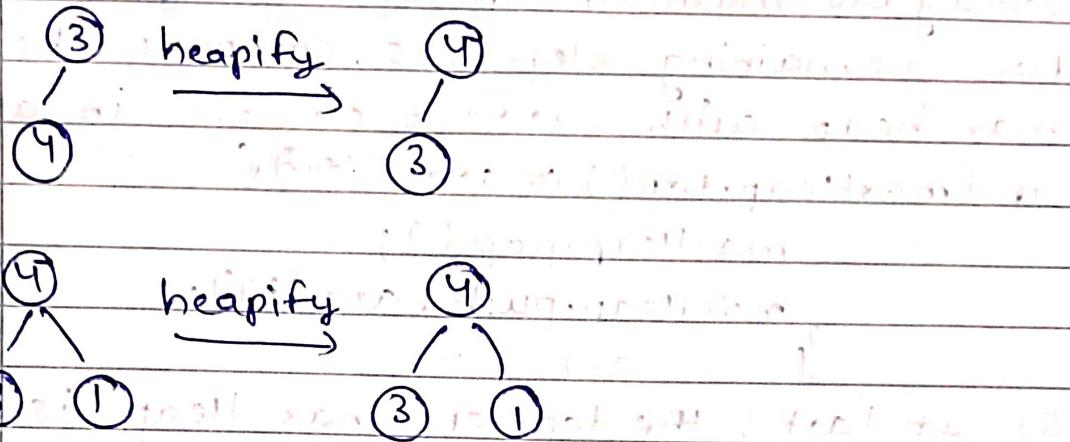
creating max-heap of first k elements -



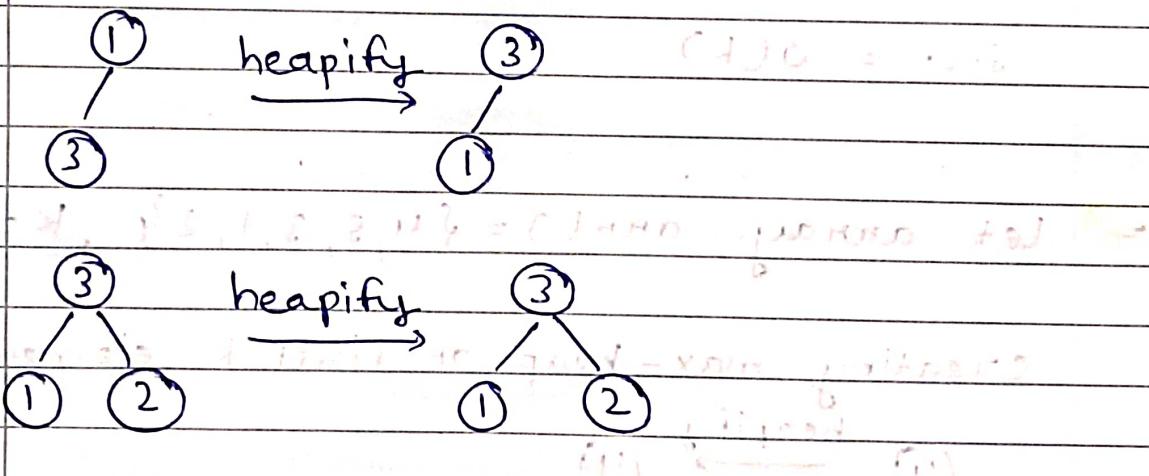
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Now, rest of the steps -

$5 > 1 \leftarrow$



$\leftarrow 4 > 2 \leftarrow$ heapify of (4, 3, 1)



Now, the element at top is the kth smallest element.

Code -

```
class Solution{
```

```
public:
```

```
int kthSmallest(int arr[], int l, int r, int k){
```

```
    priority_queue<int> maxHeap;
```

```
    for(int i=l; i<k; i++) {
```

```
        maxHeap.push(arr[i]);
```

```
    }
```

```
    for(int i=k; i<=r; i++) {
```

```
        if(maxHeap.top() > arr[i]) {
```

```
            maxHeap.pop();
```

```
            maxHeap.push(arr[i]);
```

```
}
```

```
    if(k == r) return maxHeap.top();
```

```
    else return maxHeap.top();
```

Time complexity of this solution

is O(n log n)

Reason is that we are doing a heap push operation.

Time complexity of this solution

is O(n log k)

Reason is that we are doing a heap pop operation.

Time complexity of this solution

is O(n log k)

* Kth largest Element in an Array (leetcode - 215)

I/P - nums = [3, 2, 1, 5, 6, 4], k = 2

O/P - 5

Approach - Create min heap of first k elements.

for remaining elements, check -

```
if (minHeap.top() < nums[i]) {  
    minHeap.pop();  
    minHeap.push(arr[i]);  
}
```

Now, the top() of the minHeap is kth largest element.

Code -

```
class Solution {  
public:
```

```
int findKthLargest(vector<int>& nums, int k) {  
    priority_queue<int, vector<int>, greater<int>>  
        minHeap;
```

```
    for (int i=0; i<k; i++) {  
        minHeap.push(nums[i]);  
    }
```

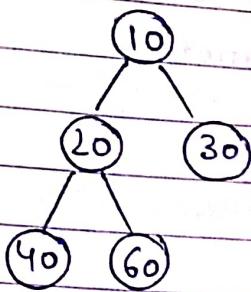
```
    for (int i=k; i<nums.size(); i++) {  
        if (minHeap.top() < nums[i]) {  
            minHeap.pop();  
            minHeap.push(arr[i]);  
        }  
    }
```

```
    return minHeap.top();  
}
```

```
};
```

* Is Binary Tree Heap (gfg) -
... kya hai, just go and check.

I/P -



O/P -

false

for a Binary tree to be heap -

It must follows = CBT + Heap property (Max heap)

Approach -

logic for complete Binary Tree -

Agr null ke baad koi non-null node mil jaye while traversing level order, then CBT nahi hai otherwise CBT hai.

logic for checking max heap -

for every node, agr parent node ki value children nodes ki value se greater hai, then max heap hai.

Code -

```
class Solution {
public: // for checking CBT
    bool isCompleteTree(Node* root) {
        queue<Node*> q;
        q.push(root);
        bool nullfound = 0;

        while (!q.empty()) {
            Node* front = q.front();
            q.pop();

            if (!front) {
                nullfound = 1;
            } else {
                if (nullfound) return 0;

                q.push(front->left);
                q.push(front->right);
            }
        }

        return 1;
    }
}
```

// this class will return two values for every node

```
class Info {
public:
    int maxVal;
    bool isHeap;
};

Info() {
    this->maxVal = INT_MIN;
    this->isHeap = true;
}
```

// for checking, whether max heap or not
Info checkMaxHeap (Node *root) {

Info ans;

if (!root) return ans;

if (!root->left && !root->right) {

ans.maxVal = root->data;

return ans;

}

Info leftAns = checkMaxHeap (root->left);

Info rightAns = checkMaxHeap (root->right);

if (root->data > leftAns.maxVal && root->data >

rightAns.maxVal && leftAns.isHeap && rightAns.isHeap) {

ans.maxVal = root->data;

return ans;

}

else {

ans.maxVal = max (root->data, max (leftAns.maxVal,
rightAns.maxVal));

ans.isHeap = false;

return ans;

}

bool isHeap (struct Node *tree) {

Info ans = checkMaxHeap (tree);

return ans.isHeap && isCompleteTree (tree);

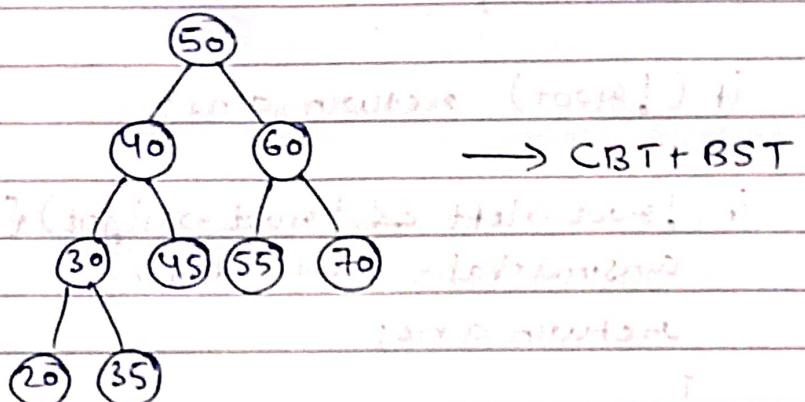
}

};

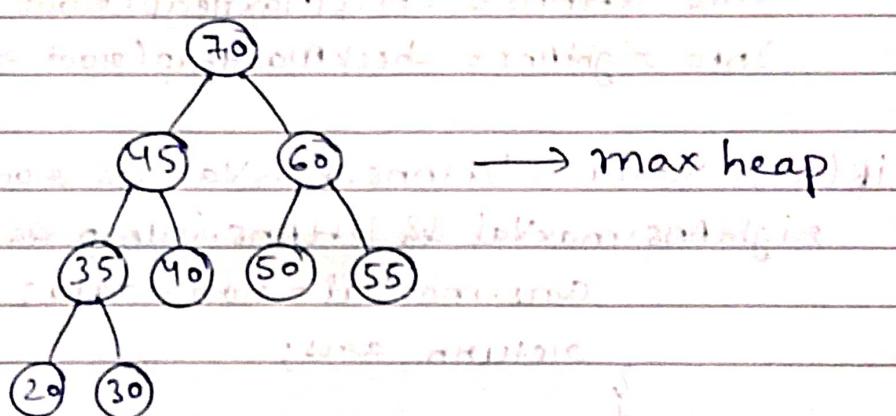
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* Convert CBT + BST into max heap -.

I/P -



O/P -



Approach -

- = find inorder for the given CBT + BST.
- = Replace the node's data using inorder array and by doing postorder traversal on tree.

- / -

Code -

```
class Solution {
public:
    void findInorder(Node* root, vector<int>& inorder) {
        if (!root) return;
        findInorder(root->left, inorder);
        inorder.push_back(root->data);
        findInorder(root->right, inorder);
    }

    void replaceUsingPostorder(Node* root, vector<int>& inorder,
                               int& index) {
        if (!root) return;
        replaceUsingPostorder(root->left, inorder, index);
        replaceUsingPostorder(root->right, inorder, index);
        root->data = inorder[index++];
    }

    void convertToMaxHeap(Node* root) {
        vector<int> inorder;
        findInorder(root, inorder);

        int index = 0;
        replaceUsingPostorder(root, inorder, index);
    }
};
```