

* Balanced Binary Tree - (Leetcode - 110)

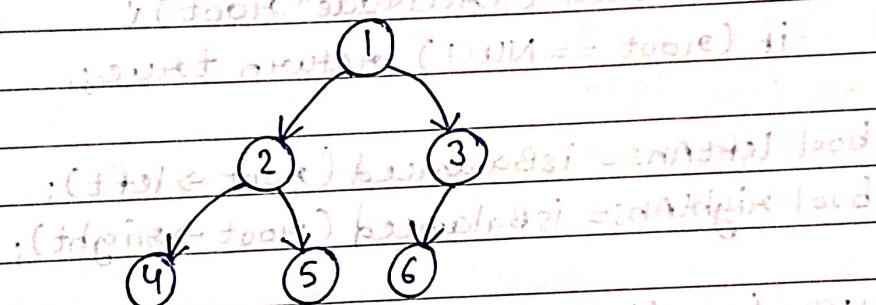
Tree is height balanced if tree ki har node ke left subtree aur right subtree ki height ka absolute difference atmost 1 hai.

If this condition is true at every node, then tree is height-balanced.

I/P -

$[1, 2, 3, 4, 5, 6]$

O/P - true



At every node, 3 things must be true -

- (1) Left subtree should be balanced.
- (2) Right subtree should be balanced.
- (3) $\text{abs}(\text{leftHeight} - \text{rightHeight}) \leq 1$

Code -

```
class Solution {
public:
```

```
    int height(TreeNode* root) {
        if (root == NULL) return 0;
        int leftHeight = height(root->left);
        int rightHeight = height(root->right);
        return max(leftHeight, rightHeight) + 1;
    }
```

```
bool isBalanced(TreeNode* root) {
    if (root == NULL) return true;
```

```
    bool leftAns = isBalanced(root->left);
    bool rightAns = isBalanced(root->right);
```

```
    bool diff = (abs(height(root->left) - height(root->right))  
                <= 1);
```

```
    return diff && leftAns && rightAns;
```

```
}
```

→ node which has both p and q as its descendants.

* Lowest Common Ancestor of a Binary Tree (LCA) (Leetcode - 236)

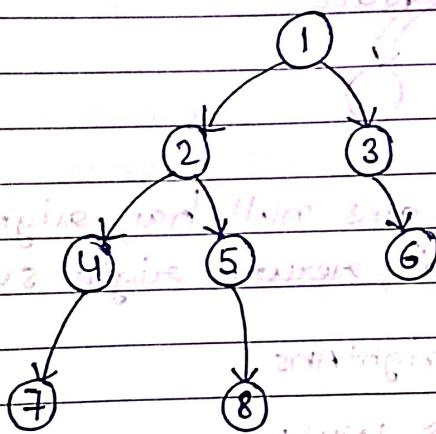
lowest common ancestor is the lowest node that has both p and q as its descendants.

I/P -

[1, 2, 3, 4, 5, null, 6, 7, null, null, 8]

O/P -

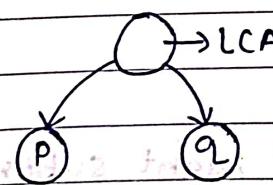
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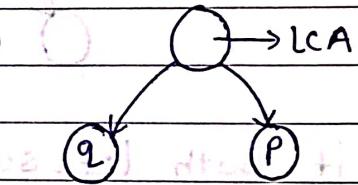
Here, if $p=2$ and $q=8$.
O/P will be 2 i.e. lowest common Ancestor of node

Cases -

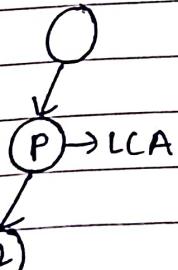
(1)



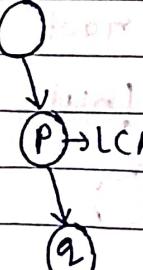
(2)



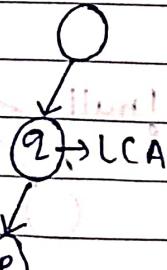
(3)



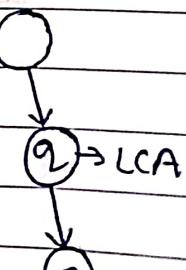
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(5)

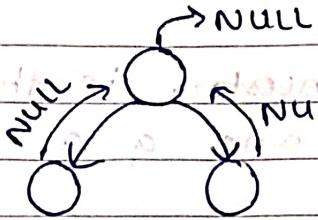


(6)

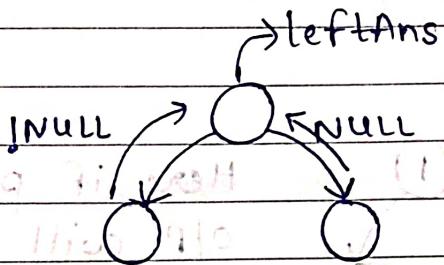


Answer below 4 ways se aa sakte hai -

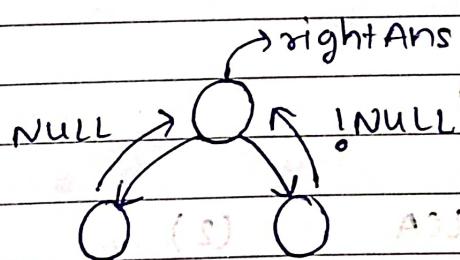
- (1) if left subtree ka ans null hai, right subtree ka ans null hai, return null.



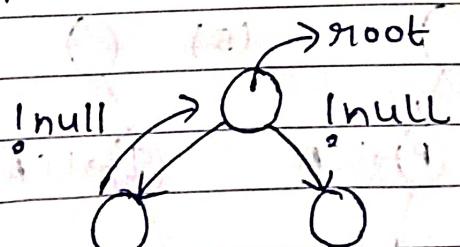
- (2) if left subtree ka ans (!null) hai, right subtree ka ans null hai, return left subtree ka ans.



- (3) if left subtree ka ans null hai, right subtree ka ans (!null) hai, return right subtree ka ans.



- (4) if both left subtree and right subtree ka ans (!null) hai, return root because that root is LCA.



Code -

```
class Solution {  
public:  
    TreeNode* lowestCommonAncestor(TreeNode* root,  
                                    TreeNode* p, TreeNode* q) {  
        if (root == NULL) return NULL;  
        if (root == p) return p;  
        if (root == q) return q;  
  
        TreeNode* leftAns = lowestCommonAncestor(root->left, p, q);  
        TreeNode* rightAns = lowestCommonAncestor(root->right, p, q);  
  
        if (leftAns == NULL && rightAns == NULL) return NULL;  
        else if (leftAns != NULL && rightAns == NULL) return leftAns;  
        else if (leftAns == NULL && rightAns != NULL) return rightAns;  
        else return root;  
    };
```

* - 100% AC

* - 2nd approach

1. idea: traverse and sum = from 31
- if node meets, store total w/ its node + 2
- if not, then one by one compare
- calculate difference

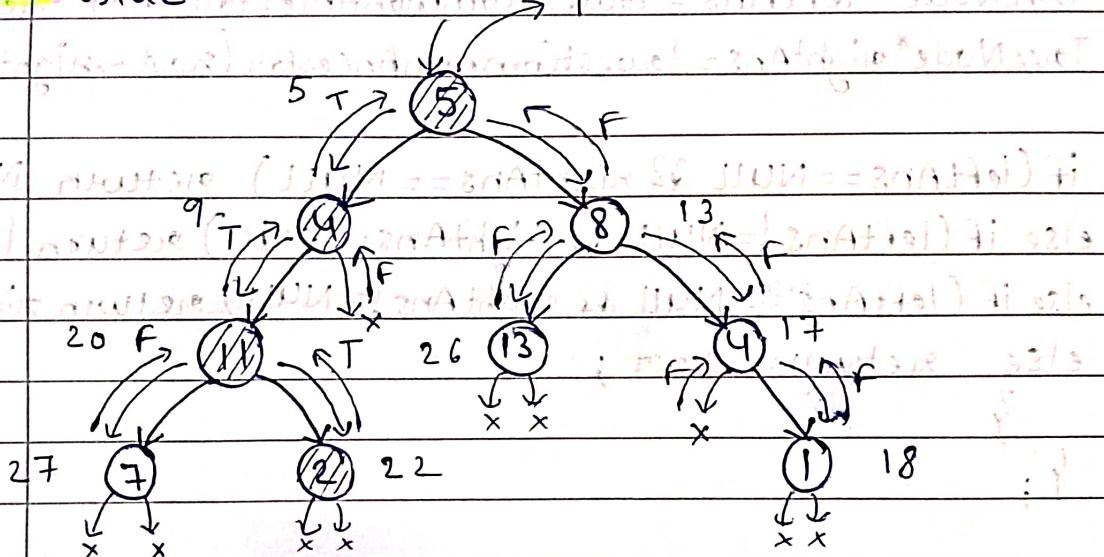
2. idea: trip up to last and all the previous steps
- first find LCA
- then find sum of both
- calculate difference

* Path sum (leetcode - 112) Traditional Recursion

Given root of binary tree and an integer targetSum, return true if tree has a root to leaf path such that adding up all the values along the path equals targetSum.

I/P - $\text{root} = [5, 4, 8, 11, \text{null}, 13, 7, 2, \text{null}, \text{null}, \text{null}, 1]$
 $\text{targetSum} = 22$

O/P - true



Approach -

= Base cases -

- If $\text{root} = \text{NULL}$ hai, return false.
- If root is a leaf node, then check if -
 - targetSum reached, return true
 - else return false.
- Make recursive calls for left & right subtree.
- return $\text{leftAns} \text{ || } \text{rightAns}$.

विद्युत वितरण के लिए जल संकट होता है।
इसके लिए जल की बचत करना चाहिए।
जल की बचत करने के लिए जल का उपयोग
उपयोग के लिए जल का उपयोग करना चाहिए।

Code -

```
class Solution {  
public:  
    bool solve(TreeNode* root, int &targetSum, int sum)  
    {  
        if (root == NULL) return false;  
        sum += root->val;  
        if (root->left == NULL && root->right == NULL) {  
            if (targetSum == sum) return true;  
            return false;  
        }  
        bool leftAns = solve(root->left, targetSum, sum);  
        bool rightAns = solve(root->right, targetSum, sum);  
        return leftAns || rightAns;  
    }
```

```
bool hasPathSum(TreeNode* root, int targetSum){  
    int sum = 0;  
    return solve(root, targetSum, sum);  
}
```

→ Code mei sum ko pass by value kiya hai, if
by reference berenge to backtracking wala
step abhi add karna pdega.

Note -

Base case brane ke liye 2 chize dhyam
mei rakhni hai, i.e. rukna kahai.

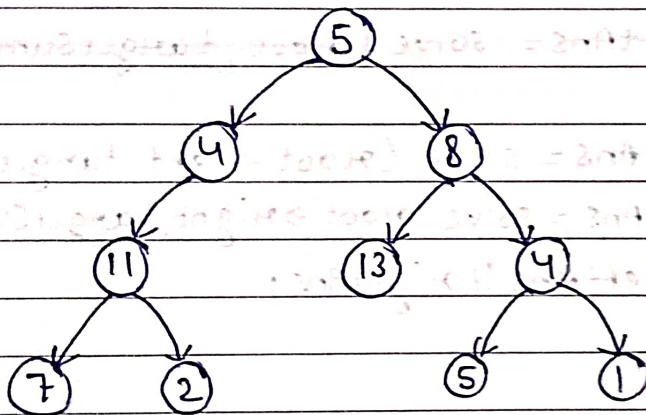
Other is kya extra kaam kar rhe hai jisme
return jana pad raha hai.

* Path Sum II (Leetcode - 113)

Given root of binary tree and targetSum,
return all root-to leaf path where sum of
node values in the path equals targetSum

I/P - root = [5, 4, 8, 11, null, 13, 4, 7, 2, null, null, 5, 1]
targetSum = 22

O/P - [[5, 4, 11, 2], [5, 8, 4, 5]]



Approach -

- Base cases - If root = null, return .
If root is a leaf node, store temp in ans.
- At every node → sum += root → val;
temp.push_back(root → val);
- Make recursive calls for left and right subtree.

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Code -

```
class Solution {
public:
    void solve(TreeNode* root, int &targetSum, int sum,
               vector<int> temp, vector<><int>>&ans) {
        if (root == NULL) return;
        sum += root->val;
        temp.push_back(root->val);
        if (root->left == NULL && root->right == NULL) {
            if (sum == targetSum) ans.push_back(temp);
            return;
        }
        solve(root->left, targetSum, sum, temp, ans);
        solve(root->right, targetSum, sum, temp, ans);
    }
}
```

```
vector<><int>> pathSum(TreeNode* root, int targetSum) {
    int sum = 0;
    vector<int> temp;
    vector<><int>> ans;
    solve(root, targetSum, sum, temp, ans);
    return ans;
}
```

→ Code mein sum and temp ko pass by value kr rhe
so that backtracking vala step add na karna
pde.

Important -

- * Construct Binary Tree from Preorder and Inorder Traversal (Leetcode - 105)

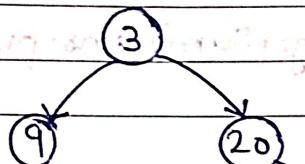
Given preorder and inorder of a binary tree, construct a tree from it.

I/P -

preorder = [3, 9, 20, 15, 7]

inorder = [9, 3, 15, 20, 7]

O/P - [3, 9, 20, null, null, 15, 7]



Approach -

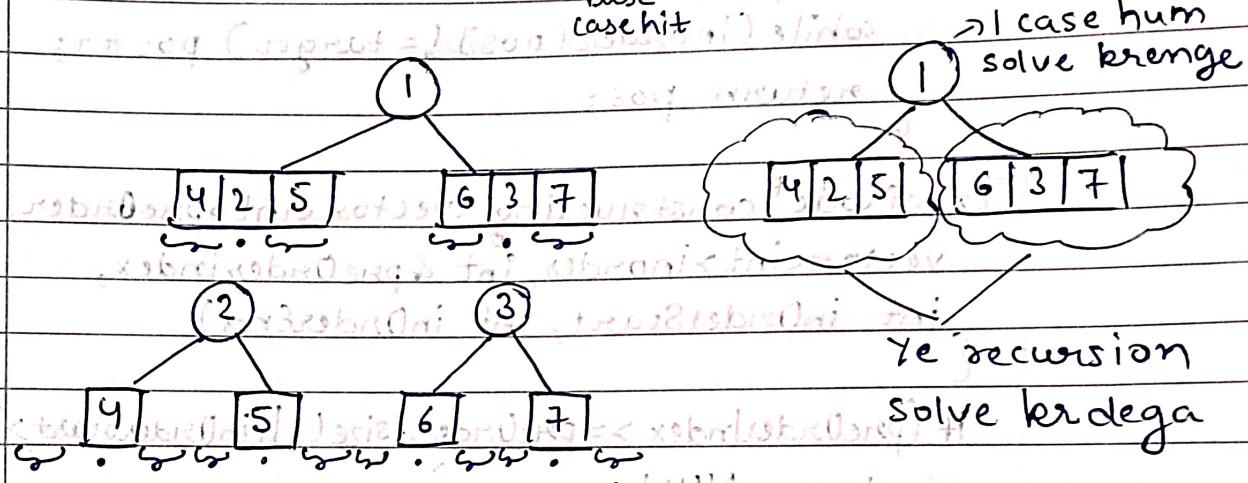
Preorder - NLR

- Preorder array ka leftmost element i.e. 0th index element chaimesha root node hoga tree ka.
- Inorder - LNR
Inorder array me wo element search kro. Us element ke left ke sare elements left subtree me aayenge and right ke sare elements right subtree me aayenge..
- Iske baad preorder array me index ko increment kro and repeat above process until preorderIndex < preorder.size()

Implementation of tree, so first we will implement
preorder traversal as this is the most common and
easiest approach for traversal. When $i=0$, $j=6$
and $k=0$ to $n-1$, then $i < j$ and $k < n$.
and $\text{inorderStartIndex} > \text{inorderEndIndex}$

let -

preorder = $[1|2|4|5|3|6|7]$ inorder = $[4|2|5|1|6|3|7]$



Nothing to

Nothing to

solve in

solve in $i < j$ = $i < j < k$

left &

left & right,

right,

return

return

$(\text{subtree1}, \text{subtree2})$ when $\text{subtree1} = \text{nothing}$ $\text{subtree2} = \text{nothing}$

So, tree is -

$(\text{subtree1}, \text{subtree2})$ when $\text{subtree1} = \text{nothing}$ $\text{subtree2} = \text{nothing}$

\Rightarrow nothing

$(\text{subtree1}, \text{subtree2})$ when $\text{subtree1} = \text{nothing}$ $\text{subtree2} = \text{nothing}$

\Rightarrow nothing

$(\text{subtree1}, \text{subtree2})$ when $\text{subtree1} = \text{nothing}$ $\text{subtree2} = \text{nothing}$

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\Rightarrow nothing

$(\text{subtree1}, \text{subtree2})$ when $\text{subtree1} = \text{nothing}$ $\text{subtree2} = \text{nothing}$

\Rightarrow nothing

Note - Isme preOrderIndex ko hamesha by reference bhejna hai ki to call se return karte time iska index change ho Jayega which we don't want i.e. preorder array mei ek index per ek baar hi jana hai isliye.

Code -

```
class Solution{
```

```
public:
```

```
int searchInOrder(vector<int> inorder, int target) {  
    int pos = 0;  
    while (inorder[pos] != target) pos++;  
    return pos;  
}
```

```
TreeNode* constructing(vector<int> preOrder,  
                      vector<int> inorder, int &preOrderIndex,  
                      int inOrderStart, int inOrderEnd)
```

```
{  
    if (preOrderIndex >= preOrder.size() || inOrderStart > inOrderEnd)  
        return NULL;
```

```
    int element = preOrder[preOrderIndex++];
```

```
    TreeNode* Root = new TreeNode(element);
```

```
    int position = searchInOrder(inorder, element);
```

```
    Root->left = constructing(preOrder, inorder, preOrderIndex,  
                               inOrderStart, position - 1);
```

```
    Root->right = constructing(preOrder, inorder, preOrderIndex,  
                               position + 1, inOrderEnd);
```

```
    return Root;
```

```
}
```

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```

12. TreeNode* buildTree (vector<int>& preorder,
13.                                     vector<int>& inorder)
14. {
15.     int preOrderIndex = 0; // A. First iteration = 1
16.     int inOrderStart = 0; // E. 0, 1, 2 = left subtree
17.     int inOrderEnd = inorder.size() - 1;
18.     TreeNode* root = constructing (preorder, inorder,
19.                                     preOrderIndex, inOrderStart, inOrderEnd);
20.     return root;
21. }

```

Explain above function. Inverse et abhi to
 agar ek node ko external form mai terminal form ke liye
 \rightarrow Is code mein inorder array mein element ko
 search karne ke har baar $O(n)$ time complexity
 jari hai.

So, instead of using this, we can use map here.
 Ye map har element ke corresponding jo index
 hogा vo lebtayega. Ek sort of hash table. Har
 ek node, joki usko process kiji get
 apna set sort ho.

Har node jise jaldi sort kiji jati h. So, if I want to
 sort ek node, kisi sort ke upar jao
 + sort ke saath sort ke index daal do.

*

Construct Binary Tree from Inorder and Postorder traversal (leetcode-106)

I/P -

inorder = [9, 3, 15, 20, 7]

postorder = [9, 15, 7, 20, 3]

O/P -

[3, 9, 20, null, null, 15, 7]

Approach -

- Postorder - LRN

postorder traversal mei root node always rightmost element i.e. last index of array hoga.
So,isme postOrderIndex ko postorder.size() - 1

start, i kina chahi. id and id. need array

last first

Postorder array mei right to left kaate time
pehle right subtree ke elements aayenge then
left subtree ke. That's why pehle right subtree
ke liye recursive call lgani hai, then left
subtree ke liye.

- Same idhr bhi, postOrderIndex ko by reference
pass karna pdega so that jis index pr ja
chuke hai wope fir se na jaye.

Code -

```
class Solution {
public:
    void createMapping(map<int, int> &valueToIndexMap,
                      vector<int> &inorder) {
        for (int i = 0; i < inorder.size(); i++) {
            int element = inorder[i];
            valueToIndexMap[element] = i;
        }
    }

    TreeNode* constructTree(map<int, int> &valueToIndexMap,
                           vector<int> &inorder, vector<int> &postorder,
                           int &postOrderIndex, int inOrderStart, int inOrderEnd)
    {
        if (postOrderIndex < 0 || inOrderStart > inOrderEnd)
            return NULL;

        int element = postorder[postOrderIndex--];
        TreeNode* Root = new TreeNode(element);

        int position = valueToIndexMap[element];

        Root->right = constructTree(valueToIndexMap, inorder,
                                      postorder, postOrderIndex, position + 1, inOrderEnd);
        Root->left = constructTree(valueToIndexMap, inorder,
                                   postorder, postOrderIndex, inOrderStart, position - 1);

        return Root;
    }
}
```

```
TreeNode* buildTree(vector<int>& inorder, vector<int>& postorder) {
```

```
    int postOrderIndex = postOrder.size() - 1; // postOrderIndex is index of current root node
    int inOrderStart = 0;
    int inOrderEnd = inorder.size() - 1;
    map<int, int> valueToIndexMap; // map to store index of each value
    CreateMapping(valueToIndexMap, inorder);
```

```
    TreeNode* root = constructing(valueToIndexMap, // constructing is helper function
        inorder, postorder, postOrderIndex, inOrderStart, // pass all required parameters
        inOrderEnd);
```

```
    return root;
```

```
}
```