

* **Quick Sort →**

Algorithm →

Step 1: Write partition logic to find left and right part of the array. Pivot index divides the array into two.

Step 2: Make a recursive call for left part. Left part is from $s \rightarrow \text{pivot} - 1$.

Step 3: Make a recursive call for right part. Right part is from $\text{pivot} + 1 \rightarrow e$.

quickSort(arr, s, e)

- 1. // Base case
- 2. Break into left and right using pivot. Let pivot be starting index.
- 3. Make recursive call for left and right.

partition (arr, s, e)

1. Take starting index as pivot.
2. Count all elements which are less than pivot element.
3. Place pivot element at the position it must be.
4. Update pivot index.
5. Place elements less than pivot element to left of pivot index and elements greater than pivot element to right of pivot index.
6. Return new pivot index.

Dry Run →

array = {40, 20, 10, 30, 50},

0	1	2	3	4
40	20	10	30	50

↑ ↑
s p e

- $p = 0$

- $\text{count} = 3$ (Count of elements less than pivot element.)

- $\text{rightIndex} = s + \text{count}$ (rightIndex indicates right position of pivot element)
 $= 0 + 3$
 $= 3$

- $\text{swap}(\text{arr}[p], \text{arr}[\text{rightIndex}])$

30	20	10	40	50
↑	↑	↑	↑	

i p j

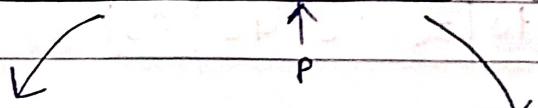
(new pivot index = rightIndex = 3)

- All elements less than pivot element are placed left to pivot element and elements greater than pivot element are placed right to pivot element. So, no need to swap here.

- $\text{pivotIndex} = 3$

Now, the array is →

0	1	2	3	4
30	20	10	40	50



Call for left →

$$s=0, e=p-1=2$$

0	1	2
30	20	10

s p

e

50

↓

Base case
So, return

Call for right →

$$s=p+1=4, e=4$$

- $P = 0$
- Count = 2
- rightIndex = 2
- swap(arr[P], arr[rightIndex])

0	1	2	3	4
10	20	30	40	50

J P

(new pivot index =
right index = 2)

Call for left →
 $s=0, e=p-1=1$

Call for right →
 $s=p+1, e=e$

0	1
10	20

s p

e

There is no such element
So, $s > e$ (Base case)

$P=0, \text{count}=0$
rightIndex = 0

swap

0	1
10	20

↓ P

call for left → $s=P, e=p-1$
($s > e$) return.

call for right → $s=p+1=1, e=1$

[20] → $s==e$, return

So, the final array we got after sorting is →

10	20	30	40	50
----	----	----	----	----

Code →

```
#include <bits/stdc++.h>
using namespace std;

// this function divides the array into left and
// right using pivot

int partition(int arr[], int s, int e) {
    // let's take starting index as pivot
    int pivotIndex = s;
    int pivotElement = arr[s];

    // count all elements less than pivot element
    int count = 0;
    for (int i = s + 1; i <= e; i++) {
        if (arr[i] <= pivotElement) {
            count++;
        }
    }

    // place pivot element at the position it must be have
    int rightIndex = s + count;
    swap(arr[pivotIndex], arr[rightIndex]);
    // update pivot index
    pivotIndex = rightIndex;
```

// pivot element se chotte uske left mai and
// bade right mai kardo

```
int J=s;
int J=e;
while (i<pivotIndex && J>pivotIndex){
    while (arr[i]<=pivotElement){
        i++;
    }
    while (arr[J]>pivotElement){
        J--;
    }
}
```

```
if (i<pivotIndex && J>pivotIndex){
    swap (arr[i], arr[J]);
}
```

```
return pivotIndex;
```

```
void quickSort (int arr[], int s, int e){
```

// base case

```
if (s >= e) {
```

// in case of empty array & single element

```
return;
```

// partition logic

```
int pivot = partition (arr, s, e);
```

// recursive call for left →

```
quickSort (arr, s, pivot - 1);
```

// recursive call for right →

```
quickSort (arr, pivot + 1, e);
```

```

int main() {
    int arr[] = {40, 20, 10, 30, 50};
    int size = 5;
    int s = 0;
    int e = size - 1;
    quicksort(arr, s, e);
    for (int i = 0; i < size; i++) {
        cout << arr[i] << " ";
    }
}

```

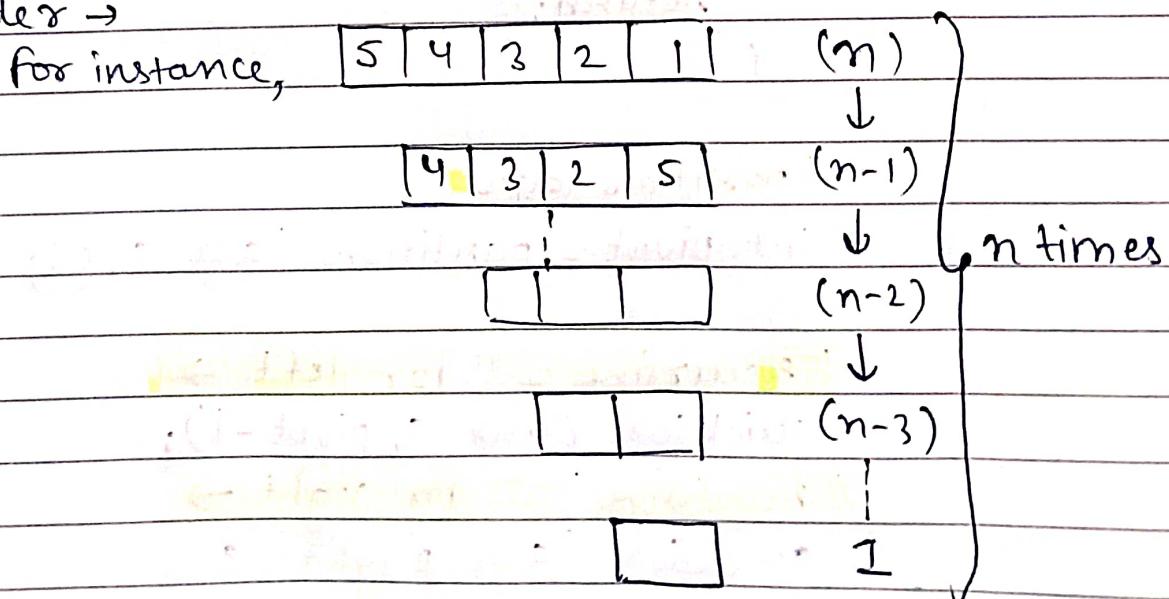
Time complexity → (worst case)

for partitioning logic →

$O(n)$ → Comparing whole elements in array with pivot elements. That's why traversing whole array.

$O(n)$ → for arranging elements to the left and right of the pivot element.

In worst case, elements are arranged in descending order →



So, partitioning of array is occurring n times
and every partition takes n time.

That's why $T(n) = O(n^2)$

Space complexity $\rightarrow O(n)$ [stack space]

* Another approach \rightarrow Quick Sort

Code \rightarrow

```
#include <bits/stdc++.h>
using namespace std;

void quickSort(int arr[], int start, int end){
    // base case
    if (start >= end) return;

    int i = start - 1;
    int j = start;

    // let's take end as pivot
    int pivot = end;
    while (j < pivot){
        if (arr[j] < arr[pivot]){
            i++;
            swap(arr[i], arr[j]);
        }
        j++;
    }

    i++;
    swap(arr[i], arr[pivot]);
```

// recursive call for left half →

quicksort(arr, start, i-1);

// recursive call for right half →

quicksort(arr, i+1, end);

}

int main() {

int arr[] = {9, 2, 7, 4};

int size = 4;

int s = 0;

int e = size - 1;

quicksort(arr, s, e);

for (auto num : arr) {

cout << num << " ";

}

return 0;

partition of arr

partition of arr