

Number Theory Lecture 7

Saturday, 13 July 2024

8:15 PM

Problem:- Find a^b given a, b

Brute Force:-

```
ans = 1;
for (i = 1 → b) {
    ans = ans * a;
}
```

Time = $O(b)$

$$a^b = \underbrace{a \times a \times a \dots \times a}_{b \text{ times}}$$

eg: $2^{10} = ?$

$$\begin{aligned} 2^{10} &= 2^5 \times 2^5 \leftarrow O(1) \\ 2^5 &= 2^2 \times 2^2 \times 2 \leftarrow O(1) \\ 2^2 &= 2^1 \times 2^1 \leftarrow O(1) \end{aligned} \left. \vphantom{\begin{aligned} 2^{10} \\ 2^5 \\ 2^2 \end{aligned}} \right\} \text{3 operations}$$

$$a^b$$

$$\underline{\text{pow}}(a, b) = \begin{cases} \text{pow}(a, b/2) * \text{pow}(a, b/2) & \text{if } b \text{ is even} \\ \text{pow}(a, b/2) * \text{pow}(a, b/2) * a & \text{if } b \text{ is odd} \\ 1, & \text{if } b \text{ is } 0 \end{cases}$$

```

int pow(a, b) {
    if (b == 0) return 1;
    if (b % 2 == 0)
        → return pow(a, b/2) * pow(a, b/2);
    else
        → return pow(a, b/2) * pow(a, b/2) * a;
}

```

$$\rightarrow T(b) = 2 * T(b/2) + c$$

$$\boxed{T(b) = O(b)}$$



① Will this code work? Yes ✓

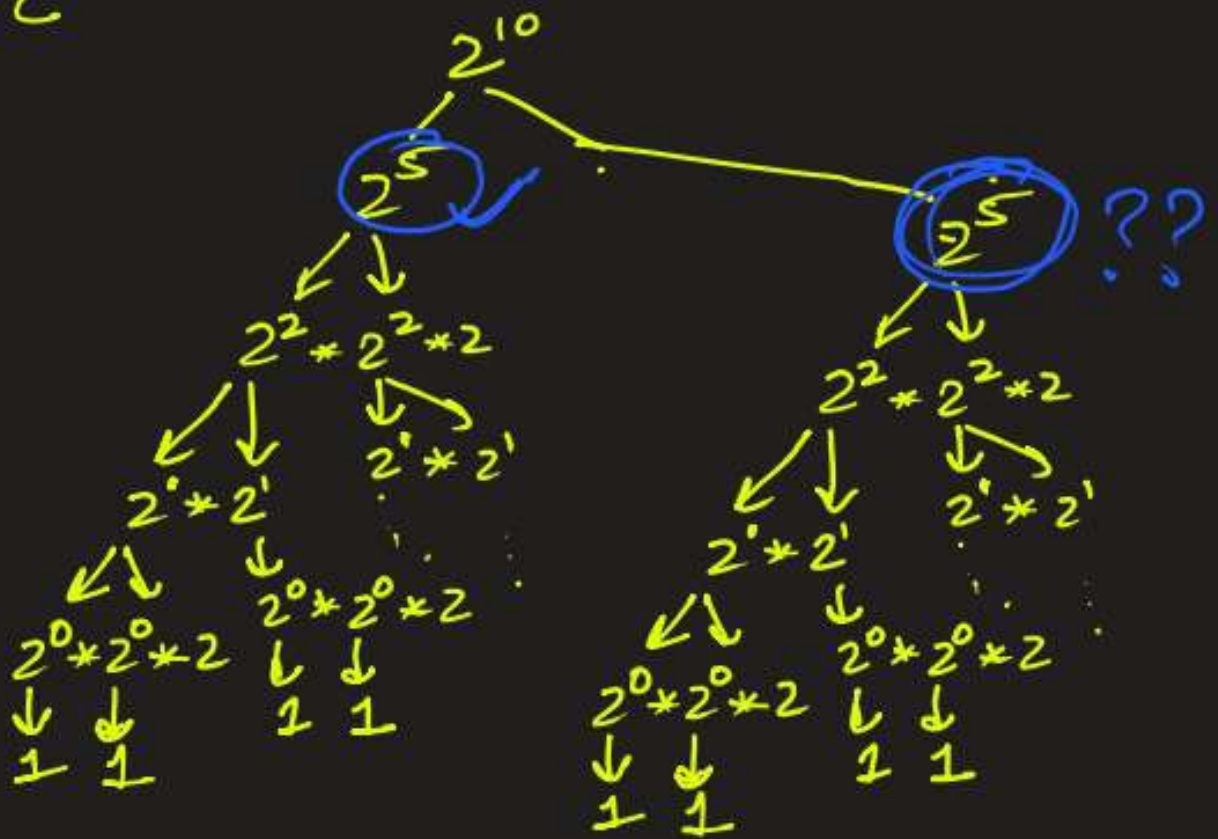
② Time complexity = $O(b)$

$$T(n) = 2$$

$$= 4$$

$$T(b/2) + c$$

$$O(b)$$



$$T(n) = 2T(n/2) + c$$

$$= 4T(n/4) + c$$

$$= 8T(n/8) + c$$

$$\vdots$$

$$= n \cdot T(1) + c$$

$$= O(n)$$

$$= 8T(1)$$

$$= n \cdot T(1)$$

$$= O(n)$$

Optimal:-

```
int pow( int a, int b) {
    if (b==0) return 1;
    int ans = pow(a, b/2);
    if (b%2==0)
        ans = ans * ans;
    else
        ans = ans * ans * a;
    return ans;
}
```

$$T(b) = T(b/2) + c$$

$$T = O(\log b)$$

$$S = O(\log b)$$

```
int ans = 1;
while(b) {
    if(b%2 == 1) ans = ans*a;
    b >>= 1;
    a *= a;
}
cout << ans << endl;
```

$T = O(\log b)$
 $S = O(1)$

2^{10}

ans=1, b=10, a=2
 ans=1, b=5, a=4
 ans=4, b=2, a=16
 ans=4, b=1, a=256
 ans=1024, b=0, a=256*256;

$$1007^{1007} \% \underline{p} \xrightarrow{\text{prime number}} < \textcircled{p}$$

$$\searrow a \% b < b$$

$$\text{int} \sim 10^9 + 7$$

$$\underline{1000000007}$$

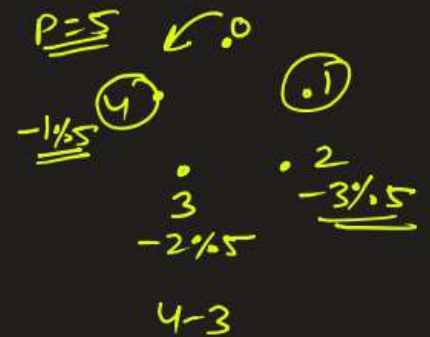
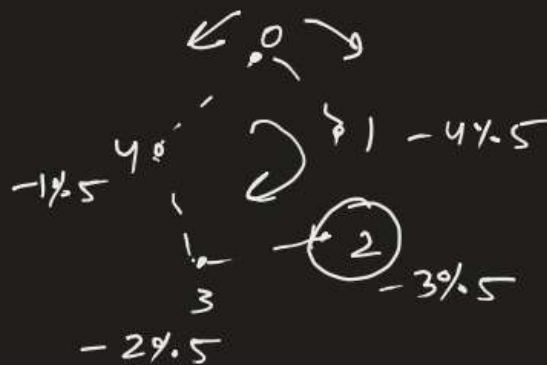
Modular Arithmetic

- $(a+b) \% p = ((a \% p) + (b \% p)) \% p$ ✓
- $(a-b) \% p = ((a \% p) - (b \% p)) \% p$ ✓
- $(a \times b) \% \underline{p} = ((a \% p) \times (b \% p)) \% p$ ✓
- ~~$(a/b) \% p = ((a \% p) / (b \% p)) \% p$~~ Doesn't work

$$p = \underline{\underline{5}}$$

$$\begin{aligned}(13 + 24) \% 5 &= (13 \% 5 + 24 \% 5) \% 5 \\ &= (3 + 4) \% 5 \\ &= 7 \% 5 \\ &= \boxed{2}\end{aligned}$$

$$-3 \% 5 = \boxed{2}$$



$$\underbrace{(a * a * a * a \dots * a)}_{b \text{ times}} \% p$$

$$= (((a * a) \% p * a) \% p) \% p \dots * a) \% p$$

$$(p + a) \% p = a \% p$$

$$(7 + 3) \% 7 = 3 \% 7$$



$$p = 7$$

$$\begin{aligned}
 &(-3) \% 7 \\
 &= (-3 + 7) \% 7 \\
 &= 4 \% 7 = \textcircled{4}
 \end{aligned}$$

$=x=$

How to define $\left(\frac{a}{b}\right) \bmod p$?

$$\frac{1}{b} \bmod p \quad ??$$

modular inverse
of $b \bmod p$

$$\frac{1}{5} \bmod 7$$

$$= 5^{-1} \bmod 7$$

Modular inverse of $5 \bmod 7$

$$b \times \frac{1}{b} = 1$$

$$(b * b^{-1}) \bmod p \equiv 1$$

$$(5 * 5^{-1}) \bmod 7 = 1$$

$$5^{-1} \bmod 7 = 3$$

$$(5 * \textcircled{3}) \bmod 7 = \textcircled{1}$$

$$\frac{1}{5} \bmod 7 = 3$$

$$\left. \begin{aligned}
 (5 \times 1) \bmod 7 &= 5 \\
 (5 \times 2) \bmod 7 &= 3 \\
 (5 \times \underline{3}) \bmod 7 &= \underline{1} \\
 (5 \times 4) \bmod 7 &= 6 \\
 (5 \times 5) \bmod 7 &= 4 \\
 (5 \times 6) \bmod 7 &= 2 \\
 (5 \times 7) \bmod 7 &= 0
 \end{aligned} \right\} \downarrow$$

Fermat's Little Theorem

$$\bullet (a^{p-1}) \bmod p = 1$$

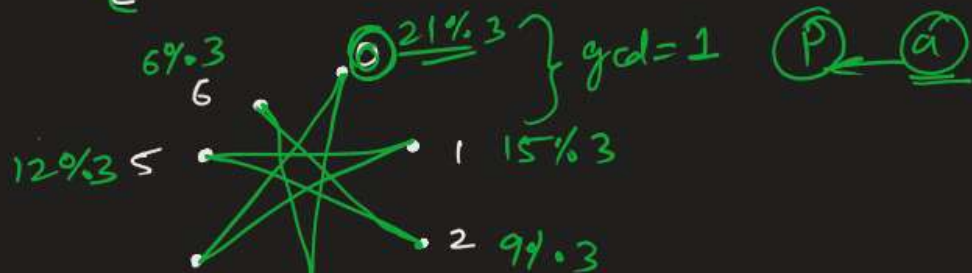
if p is prime
 a is not a multiple of p

$$p=7 \quad a=4$$

$$\left. \begin{array}{l} 5^6 \% 7 = 1 \\ 4^6 \% 7 = 1 \\ 3^6 \% 7 = 1 \\ 2^6 \% 7 = 1 \\ 1^6 \% 7 = 1 \\ 6^6 \% 7 = 1 \end{array} \right\}$$

Proof: $\rightarrow a \% p, 2a \% p, 3a \% p, 4a \% p \dots (p-1)a \% p$
 $\{1, 2, 3 \dots, p-1\}$

$$\rightarrow \begin{array}{cccccc} 5 \% 7, & 10 \% 7, & 15 \% 7, & 20 \% 7, & 25 \% 7, & 30 \% 7 \\ \{5 & 3 & 1 & 6 & 4 & 2\} \end{array}$$



$$\begin{aligned}
 (a * 2a * 3a * 4a * \dots * (p-1)a) \% p &= (1 * 2 * 3 * \dots * (p-1)) \% p \\
 &= (a^{p-1} (p-1)!) \% p \\
 &= \underline{[a^{p-1} \% p] [(p-1)! \% p]} \\
 &= \underline{(p-1)! \% p}
 \end{aligned}$$

$$\begin{aligned}
 (a^{p-1} \% p) \cdot ((p-1)! \% p) &= (p-1)! \% p \\
 \Rightarrow \boxed{a^{p-1} \% p = 1}
 \end{aligned}$$

Given that $a^{p-1} \% p = 1$

Can you calculate the modulo inverse of a mod p ?

$$a^{-1} \bmod p$$

$$x = \frac{1}{a} \bmod p$$

$$\Rightarrow (x \cdot a) \bmod p = 1$$

$$\boxed{a^{p-1} \% p = 1}$$

$$\Rightarrow a^{p-2} \% p = \frac{1}{a} \bmod p$$

$$\Rightarrow a^{-1} \bmod p = (a^{p-2}) \bmod p$$

$$\star \underline{a^{-1} \bmod p} = \boxed{a^{p-2} \bmod p} \quad \checkmark$$

$$5^{-1} \bmod 7 = 3$$

$$5^{-1} \bmod 7 = 5^5 \bmod 7 = \textcircled{3}$$

$$3^{-1} \bmod 7 = \textcircled{5}$$

$$(3 \times 5) \bmod 7 = 1$$

```
#include<bits/stdc++.h>
using namespace std;
```

```
const int MOD = 7;
```

```
ll pow(int a, int b) {
    if(b==0) return 1;
    ll ans = pow(a, b/2);
    ans = (ans*ans)%MOD;
    if(b%2 == 1) ans = (ans*a)%MOD;
    return ans;
}
```

```
ll mod_inv(int a) {
    return pow(a, MOD-2);
}
```

```
void solve() {
    for(int i=1; i<7; i++) {
        cout << mod_inv(i) << endl;
    }
}
```

H.w:- Find ${}^nC_r \% p$ given n and r , ✓ $p = 10^9 + 7$

limits:- T testcases

$$\boxed{\begin{array}{l} 1 \leq T \leq 10^5 \\ 1 \leq r \leq n \leq 10^5 \end{array}}$$

Hint:- ${}^nC_r = \frac{n!}{\underline{(n-r)!} \cdot \underline{r!}}$

$$\Rightarrow {}^nC_r \% p = (n! \% p) \cdot \text{mod_inv}((n-r)! \% p) \cdot \text{mod_inv}(r! \% p)$$

$$\cdot \left[\frac{a}{b} \% p = ((a \% p) \cdot (b^{-1} \% p)) \% p \right]$$