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Cycle detection in undirected graph using Breadth First Search in C++
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#include <bits/stdc++.h>
using namespace std;
class Solution {
public:
  // Function to detect cycle in a directed graph.
  bool\ isCyclic(int\ V,\ vector{<}int{>}\ adj[])\ \{
     int indegree [V] = \{0\};
     for (int i = 0; i < V; i++) {
        for (auto it : adj[i]) {
          indegree[it]++;
     queue<int> q;
     for (int i = 0; i < V; i++) {
        if (indegree[i] == 0) {
           q.push(i);
     int cnt = 0;
     // o(v + e)
     while (!q.empty()) {
        int node = q.front();
        q.pop();
        cnt++;
        // node is in your topo sort
        // so please remove it from the indegree
        for (auto it : adj[node]) {
           indegree[it]--;
           if (indegree[it] == 0) q.push(it);
     if (cnt == V) return false;
     return true;
  }
};
int main() {
  //V = 6;
  vector<int> adj[6] = {{}, {2}, {3}, {4, 5}, {2}, {}};
  int V = 6;
  Solution obj;
  bool ans = obj.isCyclic(V, adj);
  if (ans) cout << "True";
  else cout << "Flase";
  cout << endl;
  return 0;
```

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Graph looks like:

1 \rightarrow 2 \rightarrow 3 \rightarrow 4

\uparrow \qquad \downarrow

\downarrow \rightarrow 5

Adjacency list looks like:-

adj[0] = {}

adj[1] = {2}

adj[2] = {3}

adj[3] = {4, 5}

adj[4] = {2}

adj[5] = {}
```

Step 1: Calculate Indegree

- Initialize indegree $[] = \{0, 0, 0, 0, 0, 0, 0\}.$
- Traverse adjacency list to calculate indegree:
 - $1 \rightarrow 2$: indegree[2]++ \rightarrow indegree[] = {0, 0, 1, 0, 0, 0}
 - $2 \rightarrow 3$: indegree[3]++ \rightarrow indegree[] = {0, 0, 1, 1, 0, 0}
 - $3 \rightarrow 4$: indegree[4]++ \rightarrow indegree[] = {0, 0, 1, 1, 1, 0}
 - $3 \rightarrow 5$: indegree[5]++ \rightarrow indegree[] = {0, 0, 1, 1, 1, 1}
 - o $4 \rightarrow 2$: indegree[2]++ \rightarrow indegree[] = {0, 0, 2, 1, 1, 1}
- Final indegree[]: {0, 0, 2, 1, 1, 1}.

Step 2: Add Nodes with indegree == 0 to Queue

- Nodes with indegree == 0: 0, 1.
- Initialize queue = $\{0, 1\}$.

Step 3: Process Queue (Topological Sort)

1. Process Node 0:

- Dequeue 0, $cnt++ \rightarrow cnt = 1$.
- Node 0 has no outgoing edges; no changes to indegree[].
- o queue = $\{1\}$.

2. Process Node 1:

- o Dequeue 1, $cnt++ \rightarrow cnt = 2$.
- Node 1 → Node 2: Decrease indegree[2]-- → indegree[] = $\{0, 0, 1, 1, 1, 1\}$.
- Node 2 has indegree != 0, so it is not added to the queue.
- \circ queue = $\{\}$.

Step 4: Check for Remaining Nodes

• Cycle Exists:

- Processed nodes (cnt = 2) < Total nodes (V = 6).
- A cycle exists, as some nodes (like 2, 3, 4, 5) were never processed.

True The graph contains a cycle