

Binary Tree 2 LL in C++

```
#include <iostream>
using namespace std;
class Node {
public:
    int key;
    Node* left;
    Node* right;

    Node(int value) {
        key = value;
        left = nullptr;
        right = nullptr;
    }
};

class BinTree2LL {
private:
    static Node* prev;

public:
    static void flatten(Node* root) {
        if (root == nullptr) return;

        flatten(root->right);
        flatten(root->left);

        root->right = prev;
        root->left = nullptr;
        prev = root;
    }

    static void printList(Node* root) {
        while (root->right != nullptr) {
            cout << root->key << "->";
            root = root->right;
        }
        cout << root->key;
    }
};

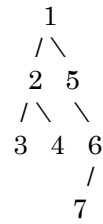
Node* BinTree2LL::prev = nullptr;

int main() {
    Node* root = new Node(1);
    root->left = new Node(2);
    root->left->left = new Node(3);
    root->left->right = new Node(4);
    root->right = new Node(5);
    root->right->right = new Node(6);
    root->right->right->left = new Node(7);

    BinTree2LL::flatten(root);
    BinTree2LL::printList(root);

    // Clean up allocated memory (not present in Java
    version)
    while (root != nullptr) {
        Node* temp = root;
        root = root->right;
        delete temp;
    }
}
```

Original Binary Tree Structure



✳ Flattening Logic: Reverse Postorder (Right → Left → Node)

The algorithm works like this:



- Traverse the tree in **reverse postorder**.
- Use a static prev pointer to keep track of the previously processed node.
- Set the current node's right to prev, and its left to nullptr.

📊 Step-by-Step Tabular Dry Run

We will track:

- The current node being visited
- The state of prev
- Links updated

Step	Node Visited	Previous (prev)	Action	Updated Links
1	7	nullptr	Set 7.right = nullptr, 7.left = nullptr, prev = 7	7 → nullptr
2	6	7	Set 6.right = 7, 6.left = nullptr, prev = 6	6 → 7
3	5	6	Set 5.right = 6, 5.left = nullptr, prev = 5	5 → 6 → 7
4	4	5	Set 4.right = 5, 4.left = nullptr, prev = 4	4 → 5 → 6 → 7
5	3	4	Set 3.right = 4, 3.left = nullptr, prev = 3	3 → 4 → ...
6	2	3	Set 2.right = 3, 2.left = nullptr,	2 → 3 → ...

<pre> } return 0; }</pre>			prev = 2	
	7	1	2	Set 1.right = 2, 1.left = nullptr, prev = 1
<p> Final Flattened Linked List (Right Pointers)</p> <p>1 -> 2 -> 3 -> 4 -> 5 -> 6 -> 7</p> <p>All left pointers are nullptr, forming a single right-skewed list.</p> <p> Output</p> <p>1->2->3->4->5->6->7</p>				
1->2->3->4->5->6->7				