

## Equivalent Subarrays in C++

```
#include <iostream>
#include <unordered_map>
#include <unordered_set>
#include <vector>

using namespace std;

int main() {
    int ans = 0;
    vector<int> arr = {2, 1, 3, 2, 3};
    unordered_set<int> set;

    // Insert unique elements into the set
    for (int i = 0; i < arr.size(); i++) {
        set.insert(arr[i]);
    }

    int k = set.size();
    int i = -1;
    int j = -1;
    unordered_map<int, int> map;

    while (true) {
        bool f1 = false;
        bool f2 = false;

        // Expand the window until all unique elements
        // are covered
        while (i < arr.size() - 1) {
            f1 = true;
            i++;
            map[arr[i]] = map[arr[i]] + 1; // Add current
            // element to the map
            if (map.size() == k) { // If all unique elements
            // are covered
                ans += arr.size() - i; // Add the number of
                // valid subarrays ending at index i
                break;
            }
        }

        // Slide the window to the right until the
        // uniqueness condition is violated
        while (j < i) {
            f2 = true;
            j++;
            if (map[arr[j]] == 1) {
                map.erase(arr[j]); // Remove element from
                // map if its count is reduced to 0
            } else {
                map[arr[j]] = map[arr[j]] - 1; // Decrease the
                // count of the element
            }
        }

        // If the map size matches k, add the number of
        // valid subarrays again
        if (map.size() == k) {
            ans += arr.size() - i;
        } else {
            break;
        }
    }
}
```

### Input:

arr = {2, 1, 3, 2, 3}

We are looking for subarrays with **all unique elements** in the array.

### Initial Setup:

- arr = {2, 1, 3, 2, 3}
- We initialize an unordered\_set called set to store the unique elements of the array.
- We calculate k = set.size(), which is the number of unique elements in the array.
  - set = {2, 1, 3} -> k = 3 (3 unique elements).

### Algorithm Steps:

1. **Initialization:**
  - i = -1, j = -1 (start indices for the sliding window).
  - map = {} (tracks the frequency of elements in the current window).
  - ans = 0 (tracks the number of valid subarrays).
2. **Start the outer while (true) loop:**
  - The outer loop keeps expanding and shrinking the window until we can no longer process subarrays.

### First pass through the outer loop:

- **Expand the window (moving i):**
  - Initially i = -1, we move i to 0 (i.e., arr[i] = 2).
  - Add arr[i] = 2 to the map:
    - map = {2: 1}
  - Now, i = 0, and map.size() = 1 (we have only one unique element, so we move on).
  - Move i to 1 (i.e., arr[i] = 1).
  - Add arr[i] = 1 to the map:
    - map = {2: 1, 1: 1}
  - Now, i = 1, and map.size() = 2 (we still have not covered all unique elements, so continue expanding).
  - Move i to 2 (i.e., arr[i] = 3).
  - Add arr[i] = 3 to the map:
    - map = {2: 1, 1: 1, 3: 1}
  - Now, i = 2, and map.size() = 3 (we have all 3 unique elements).
  - At this point, we have a valid subarray from arr[0] to arr[2]. All unique elements are covered.
  - **Count valid subarrays:**
    - Subarrays ending at i = 2:
      - The number of subarrays is

```

    }

    // If both windows cannot be expanded or
    contracted further, break the loop
    if (!f1 && !f2) {
        break;
    }
}

// Print the total number of equivalent subarrays
cout << ans << endl;

return 0;
}

```

- calculated as  
 $\text{arr.size()} - i = 5 - 2 = 3$ .
- $\text{ans} += 3 \rightarrow \text{ans} = 3$ .
  - **Shrink the window (moving j):**
    - Now we start shrinking the window by moving the left pointer (j).
    - Move j to 0 (i.e.,  $\text{arr}[j] = 2$ ):
      - Since  $\text{map}[\text{arr}[j]] = 1$ , we remove  $\text{arr}[j]$  from the map:
        - $\text{map} = \{1: 1, 3: 1\}$
    - Move j to 1 (i.e.,  $\text{arr}[j] = 1$ ):
      - Since  $\text{map}[\text{arr}[j]] = 1$ , we remove  $\text{arr}[j]$  from the map:
        - $\text{map} = \{3: 1\}$
    - Now,  $\text{map.size()} = 1$ , which is less than k. We stop shrinking.

#### Second pass through the outer loop:

- Now, we repeat the steps again by expanding and shrinking the window, but the  $\text{map.size()}$  will no longer match k (the number of unique elements).

Output:-  
0