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Optimal strategy for a game In C++
#include <iostream>
#include <algorithm>
using namespace std;
int main() {
  int arr[] = \{20, 30, 2, 10\};
  int n = sizeof(arr) / sizeof(arr[0]);
  int dp[n][n]; // Create a 2D array of size n x n
  for (int g = 0; g < n; g++) {
     for (int i = 0, j = g; j < n; i++, j++) {
        if (g == 0) {
          dp[i][j] = arr[i];
       else if (g == 1) {
          dp[i][j] = max(arr[i], arr[j]);
       } else {
          int val1 = arr[i] + min((i + 2 \le j ? dp[i
+2][j]:0), (i + 1 \le j - 1? dp[i + 1][j - 1]:0));
          int val2 = arr[j] + min((i + 1 \le j - 1))
dp[i + 1][j - 1] : 0), (i \le j - 2? dp[i][j - 2] : 0));
          dp[i][j] = max(val1, val2);
     }
  }
  cout \ll dp[0][n - 1] \ll endl; // Print the
maximum value that can be collected
  return 0;
```

Dry Run with the Input $arr[] = \{20, 30, 2, 10\}$

We need to compute dp[0][n-1], which gives the result for the entire array.

Initialization:

dp table (initial values):

```
dp | \Pi = \{
   \{0, 0, 0, 0\},\
   \{0, 0, 0, 0\},\
   \{0, 0, 0, 0\},\
   \{0, 0, 0, 0\}
```

Step-by-Step Iteration:

- 1. $\mathbf{g} = \mathbf{0}$ (Subarrays of size 1): \circ dp[0][0] = arr[0] = 20 dp[1][1] = arr[1] = 30dp[2][2] = arr[2] = 2dp[3][3] = arr[3] = 10
- 2. $\mathbf{g} = \mathbf{1}$ (Subarrays of size 2): \circ dp[0][1] = max(arr[0], arr[1]) =
 - max(20, 30) = 30dp[1][2] = max(arr[1], arr[2]) =
 - max(30, 2) = 30
 - dp[2][3] = max(arr[2], arr[3]) =max(2, 10) = 10
- 3. $\mathbf{g} = \mathbf{2}$ (Subarrays of size 3):
 - o For dp[0][2]: We compute two options:
 - val1 = arr[0] + min(dp[2])[2], dp[1][1] = 20 + min(2,30) = 22
 - val2 = arr[2] + min(dp[1])[1], dp[0][0]) = 2 + min(30, 20) = 22
 - dp[0][2] = max(22, 22) = 22
 - o For dp[1][3]: We compute two options:
 - val1 = arr[1] + min(dp[3])[3], dp[2][2]) = 30 +min(10, 2) = 32
 - val2 = arr[3] + min(dp[2])[2], dp[1][1] = 10 + min(2,30) = 12
 - dp[1][3] = max(32, 12) = 32
- 4. $\mathbf{g} = \mathbf{3}$ (Subarrays of size 4): For dp[0][3]: We compute two

options:

- val1 = arr[0] + min(dp[2][3], dp[1][2]) = 20 + min(10, 30) = 30
- val2 = arr[3] + min(dp[1]
 [2], dp[0][1]) = 10 + min(30, 30) = 40
- dp[0][3] = max(30, 40) = 40

Final DP Table:

After all iterations, the final dp table looks like this:

```
dp[[]] = \{ \\ \{20, 30, 22, 40\}, \\ \{0, 30, 30, 32\}, \\ \{0, 0, 2, 10\}, \\ \{0, 0, 0, 10\} \}
```

Result:

The final value dp[0][3] = 40 is the maximum sum that can be collected by the first player in this game.

Output:-

40