LIS in C++ #include <iostream> #include <vector> #include <algorithm> // For std::max using namespace std; void LIS(const vector<int>& arr) { int n = arr.size();vector<int> dp(n, 1); // dp[i] will store the length of LIS ending at index i int omax = 1; // To store the overall maximum length of LIS // Compute the length of the Longest Increasing Subsequence for (int i = 1; i < n; i++) { $int max_len = 0;$ for (int j = 0; j < i; j++) { if (arr[i] > arr[j]) { $if (dp[j] > max_len)$ {

cout << omax << " "; // Print the length of the LIS

 $max_len = dp[j];$

dp[i] = max len + 1;if(dp[i] > omax) {

omax = dp[i];

}

return 0;

```
// Printing the LIS length values (optional)
 for (int i = 0; i < n; i++) {
   cout << dp[i] << " ";
 cout << endl;
}
int main() {
  3};
 LIS(arr);
```

Dry Run of the Program

Input:

Array arr = {10, 22, 9, 33, 21, 50, 41, 60, 80,

Initializations:

- n = arr.size() = 10
- dp is initialized to {1, 1, 1, 1, 1, 1, 1, 1, 1, 1} because each element starts with a subsequence length of 1.
- omax = 1, which stores the overall maximum length of the LIS.

Steps:

We iterate through the array and compute the LIS length for each index:

Iteration 1 (i = 1):

- arr[1] = 22
- For j = 0: arr[1] > arr[0] (22 > 10), so we check dp[0] = 1.
 - o $\max_{\text{len}} = \max(0, dp[0]) = 1$
- dp[1] = max len + 1 = 1 + 1 = 2
- omax = $\max(\text{omax}, dp[1]) = \max(1, 2) = 2$

Iteration 2 (i = 2):

- arr[2] = 9
- For j = 0: arr[2] > arr[0] (9 > 10) is false.
- For j = 1: arr[2] > arr[1] (9 > 22) is false.
- No update in dp[2], it remains 1.
- omax = $\max(\text{omax}, dp[2]) = \max(2, 1) = 2$

Iteration 3 (i = 3):

- arr[3] = 33
- For j = 0: arr[3] > arr[0] (33 > 10), so we check dp[0] = 1.
 - o $\max_{len} = \max(0, dp[0]) = 1$
- For j = 1: arr[3] > arr[1] (33 > 22), so we check dp[1] = 2.
 - o $\max_{\text{len}} = \max(1, dp[1]) = 2$
- For j = 2: arr[3] > arr[2] (33 > 9), so we check dp[2] = 1.
 - o $\max_{l} = \max(2, dp[2]) = 2$
- $dp[3] = max_len + 1 = 2 + 1 = 3$
- omax = $\max(\text{omax}, dp[3]) = \max(2, 3) = 3$

Iteration 4 (i = 4):

- arr[4] = 21
- For j = 0: arr[4] > arr[0] (21 > 10), so we check dp[0] = 1.

- o $\max_{\text{len}} = \max(0, dp[0]) = 1$
- For j = 1: arr[4] > arr[1] (21 > 22) is false.
- For j = 2: arr[4] > arr[2] (21 > 9), so we check dp[2] = 1.
 - o $\max_{e} = \max(1, dp[2]) = 1$
- For j = 3: arr[4] > arr[3] (21 > 33) is false.
- $dp[4] = max_len + 1 = 1 + 1 = 2$
- omax = max(omax, dp[4]) = max(3, 2) = 3

Iteration 5 (i = 5):

- arr[5] = 50
- For j = 0: arr[5] > arr[0] (50 > 10), so we check dp[0] = 1.
 - o $\max_{l} = \max(0, dp[0]) = 1$
- For j = 1: arr[5] > arr[1] (50 > 22), so we check dp[1] = 2.
 - \circ max_len = max(1, dp[1]) = 2
- For j = 2: arr[5] > arr[2] (50 > 9), so we check dp[2] = 1.
 - o $\max_{len} = \max(2, dp[2]) = 2$
- For j = 3: arr[5] > arr[3] (50 > 33), so we check dp[3] = 3.
 - o $\max_{e} = \max(2, dp[3]) = 3$
- For j = 4: arr[5] > arr[4] (50 > 21), so we check dp[4] = 2.
 - \circ max_len = max(3, dp[4]) = 3
- $dp[5] = max_{len} + 1 = 3 + 1 = 4$
- omax = $\max(\text{omax}, dp[5]) = \max(3, 4) = 4$

Iteration 6 (i = 6):

- arr[6] = 41
- For j = 0: arr[6] > arr[0] (41 > 10), so we check dp[0] = 1.
 - o max len = $\max(0, dp[0]) = 1$
- For j = 1: arr[6] > arr[1] (41 > 22), so we check dp[1] = 2.
 - o $\max_{e} = \max(1, dp[1]) = 2$
- For j = 2: arr[6] > arr[2] (41 > 9), so we check dp[2] = 1.
 - o \max_{e} len = $\max(2, dp[2]) = 2$
- For j = 3: arr[6] > arr[3] (41 > 33), so we check dp[3] = 3.
 - \circ max_len = max(2, dp[3]) = 3
- For j = 4: arr[6] > arr[4] (41 > 21), so we check dp[4] = 2.
 - \circ max_len = max(3, dp[4]) = 3
- For j = 5: arr[6] > arr[5] (41 > 50) is false.
- $dp[6] = max_{len} + 1 = 3 + 1 = 4$
- omax = max(omax, dp[6]) = max(4, 4) = 4

Iteration 7 (i = 7):

- arr[7] = 60
- For j = 0: arr[7] > arr[0] (60 > 10), so we check dp[0] = 1.
 - o $\max_{l} = \max(0, dp[0]) = 1$
- For j = 1: arr[7] > arr[1] (60 > 22), so we check dp[1] = 2.

 \circ max_len = max(1, dp[1]) = 2

- For j = 2: arr[7] > arr[2] (60 > 9), so we check dp[2] = 1.
 - o $\max_{e} = \max(2, dp[2]) = 2$
- For j = 3: arr[7] > arr[3] (60 > 33), so we check dp[3] = 3.
 - \circ max_len = max(2, dp[3]) = 3
- For j = 4: arr[7] > arr[4] (60 > 21), so we check dp[4] = 2.
 - o $\max_{e} = \max(3, dp[4]) = 3$
- For j = 5: arr[7] > arr[5] (60 > 50), so we check dp[5] = 4.
 - o $\max_{e} = \max(3, dp[5]) = 4$
- For j = 6: arr[7] > arr[6] (60 > 41), so we check dp[6] = 4.
 - \circ max_len = max(4, dp[6]) = 4
- $dp[7] = max_len + 1 = 4 + 1 = 5$
- omax = $\max(\text{omax}, dp[7]) = \max(4, 5) = 5$

Further iterations will follow the same process, updating dp[i] and omax as needed.

Finally, after processing all elements, the length of the longest increasing subsequence will be omax =

Output:-

6

 $1\; 2\; 1\; 2\; 4\; 4\; 5\; 6\; 1$