Bus Routes in C++

```
#include <iostream>
#include <vector>
#include <unordered map>
#include <queue>
#include <unordered_set>
using namespace std;
int numBusesToDestination(vector<vector<int>>&
routes, int S, int T) {
  int n = routes.size();
  unordered_map<int, vector<int>> map;
  // Building a map of bus stops to their respective bus
routes
  for (int i = 0; i < n; ++i) {
    for (int j = 0; j < routes[i].size(); ++j) {
       int busStopNo = routes[i][j];
       map[busStopNo].push_back(i);
  }
  queue<int> q;
  unordered set<int> busStopVisited;
  unordered set<int> busVisited;
  int level = 0;
  q.push(S);
  busStopVisited.insert(S);
  // Performing BFS to find the minimum number of
buses
  while (!q.empty()) {
    int size = q.size();
    while (size-> 0) {
       int currentStop = q.front();
       q.pop();
       if (currentStop == T) {
         return level;
       if (map.find(currentStop) != map.end()) {
         vector<int>& buses = map[currentStop];
         for (int bus: buses) {
            if (busVisited.count(bus) > 0) {
              continue;
            }
            vector<int>& busRoute = routes[bus];
            for (int nextStop : busRoute) {
              if (busStopVisited.count(nextStop) > 0) {
                 continue:
              q.push(nextStop);
              busStopVisited.insert(nextStop);
            busVisited.insert(bus);
    ++level;
```

Input:

• Bus routes:

```
routes = {
     {1, 2, 7},  // Bus 0
     {3, 6, 7}  // Bus 1
}
```

- Source bus stop (S = 1)
- Destination bus stop (T = 6)

Step 1: Build the Map

The program constructs a map where each bus stop points to the buses that stop there. The map is:

```
map = {
    1: {0},
    2: {0},
    7: {0, 1},
    3: {1},
    6: {1}
}
```

Here:

- 1 is served by bus 0.
- 7 is served by buses 0 and 1.
- 6 is served by bus 1, etc.

Step 2: BFS Initialization

- Queue q is initialized with the source stop (s = 1): q = {1}.
- Visited sets:
 - o busStopVisited = {1} (to
 track visited bus stops).
 - o busVisited = {} (to track
 visited buses).
- level = 0 (tracks the number of buses taken).

Step 3: BFS Process

Level 0:

```
return -1; // If destination is not reachable
}

int main() {
    // Hardcoded input values
    vector<vector<int>> routes = {
        {1, 2, 7},
        {3, 6, 7}
    };
    int src = 1; // source bus stop
    int dest = 6; // destination bus stop

cout << numBusesToDestination(routes, src, dest)
<< endl;
    return 0;
}</pre>
```

- Queue size = 1 (contains 1).
- Process bus stop 1:
 - Stops at 1 are served by bus 0 (from map).
 - Bus 0 is not visited, so:
 - Add all stops from bus 0 ({1, 2, 7}) to the queue:
 - Add 2 to q.
 - Add 7 to q.
 - Mark stops 2 and7 as visited(busStopVisited= {1, 2, 7}).
 - Mark bus 0 as visited (busVisited = {0}).
- End of level 0:
 - o Queue: $q = \{2, 7\}.$
 - o Increment level = 1.

Level 1:

- Queue size = 2 (contains 2, 7).
- Process bus stop 2:
 - Stops at 2 are served by bus 0, which is already visited
 (busVisited = {0}).
 - Skip further processing for stop2.
- Process bus stop 7:
 - Stops at 7 are served by buses 0 and 1 (from map).
 - o Bus o is already visited.
 - o Bus 1 is not visited, so:
 - Add all stops from bus 1 ({3, 6, 7}) to the queue:
 - Add 3 to q.
 - Add 6 to q.
 - Mark stops 3 and 6 as visited (busStopVisited = {1, 2, 3, 6, 7}).
 - Mark bus 1 as visited (busVisited = {0, 1}).
- End of level 1:
 - o Queue: $q = \{3, 6\}.$

Level 2:

• Queue size = 2 (contains 3, 6).
• Process bus stop 3:
• Stops at 3 are served by bus 1, which is already visited (busVisited = {0, 1}).
• Skip further processing for stop 3.
• Process bus stop 6:
• 6 is the destination (T = 6).
• Return level = 2.

Output:

The minimum number of buses required to travel from stop 1 to stop 6 is:

Coloring Border in C++

```
#include <iostream>
#include <vector>
using namespace std;
vector<vector<int>> dirs = {{0, 1}, {1, 0}, {0, -1}, {-1, 0}};
void dfs(vector<vector<int>>& grid, int row, int col, int
clr) {
  grid[row][col] = -clr;
  int count = 0:
  for (auto dir : dirs) {
     int rowdash = row + dir[0];
     int coldash = col + dir[1];
     if (rowdash < 0 \mid | coldash < 0 \mid | rowdash >=
grid.size() \mid | coldash >= grid[0].size() \mid |
abs(grid[rowdash][coldash]) != clr) {
        continue;
     count++;
     if (grid[rowdash][coldash] == clr) {
        dfs(grid, rowdash, coldash, clr);
  }
  if (count == 4) {
     grid[row][col] = clr;
}
void coloring_border(vector<vector<int>>& grid, int row,
int col, int color) {
  dfs(grid, row, col, grid[row][col]);
  for (int i = 0; i < grid.size(); i++) {
     for (int j = 0; j < grid[0].size(); j++) {
        if (grid[i][j] < 0) {
          grid[i][j] = color;
  }
}
int main() {
  // Hardcoded input
  int m = 4;
  int n = 4:
  vector<vector<int>> arr = {
     \{2, 1, 3, 4\},\
     \{1, 2, 2, 2\},\
     {3, 2, 2, 2},
     \{1, 2, 2, 2\}
  };
  int row = 1;
  int col = 1;
  int color = 3;
```

Step-by-Step Dry Run:

Step 1: Call to coloring_border

- **Initial call:** coloring_border(arr, row=1, col=1, color=3).
- Call dfs(grid, row=1, col=1, clr=2) to mark the connected component and determine the border cells.

Step 2: DFS Traversal

The function dfs will:

- 1. Mark cells in the connected component as -2 (negate the value).
- 2. Identify border cells and keep them marked as -2. Interior cells are restored to clr (2) if they are surrounded by four other cells of the same value.

Traversal Process:

- 1. Start DFS at (1, 1):
 - o Mark grid[1][1] = -2.
 - Neighboring cells:
 - $(1, 2) \rightarrow \text{Proceed (same value 2)}.$
 - $(2, 1) \rightarrow \text{Proceed (same value 2)}.$
 - Other directions lead to invalid cells or cells with a different value.
 - Connected cells count so far:
- 2. Move to (1, 2):
 - $\text{o} \quad \text{Mark grid}[1][2] = -2.$
 - o Neighboring cells:
 - $(1, 3) \rightarrow \text{Proceed (same value 2)}.$
 - $(2, 2) \rightarrow \text{Proceed (same value 2)}.$
 - Connected cells count so far: 4.
- 3. Move to (1, 3):
 - o Mark grid[1][3] = -2.
 - Neighboring cells:
 - $(2, 3) \rightarrow \text{Proceed (same value 2)}.$
- 4. Continue this process until all connected cells are visited and marked as -2. For the input grid, all 2 values connected to (1, 1) are part of the connected component.

```
coloring_border(arr, row, col, color);

// Print the modified grid
for (int i = 0; i < m; i++) {
    for (int j = 0; j < n; j++) {
        cout << arr[i][j] << "\t";
    }
    cout << endl;
}

return 0;
}</pre>
```

• Marked Grid after DFS:

```
Copy code
2 1 3 4
1 -2 -2 -2
3 -2 -2 -2
1 -2 -2 -2
```

Step 3: Check for Border Cells

The DFS function identifies border cells:

- A cell is on the border if it:
 - Is not surrounded on all four sides by other cells of the same value.
- Border cells remain marked as -2.

Step 4: Update Colors

- Traverse the grid:
 - o Any cell marked as -2 is changed to the new color (3).
 - Interior cells (those surrounded by the same value) are reverted to clr (2).
- Final Grid:

```
2 1 3 4
1 3 3 3
3 3 3 3
1 3 3 3
```

Output:-

2	1	3	4
2 1 3	3	3	3
3	3	2	3
1	3	3	3

Min Cost to collect all cities in C++

```
#include <iostream>
#include <vector>
#include <queue>
using namespace std;
struct Edge {
  int v;
  int wt;
   Edge(int nbr, int weight) {
     this->v = nbr:
    this->wt = weight;
};
struct CompareEdge {
  bool operator()(const Edge& e1, const Edge& e2) {
    return e1.wt > e2.wt; // Min-Heap based on edge
weight
  }
};
int main() {
  // Hardcoded input
  int vtces = 7;
  int edges = 8;
  vector<vector<Edge>> graph(vtces);
  // Hardcoded edges
  vector<vector<int>> hardcoded_edges = {
     \{0, 1, 10\},\
    \{1, 2, 10\},\
    \{2, 3, 10\},\
     \{0, 3, 40\},\
     \{3, 4, 2\},\
    {4, 5, 3},
    \{5, 6, 3\},\
    \{4, 6, 8\}
  };
  // Populating the graph with hardcoded edges
  for (auto& edge : hardcoded_edges) {
    int v1 = edge[0];
    int v2 = edge[1];
    int wt = edge[2];
    graph[v1].emplace_back(v2, wt);
    graph[v2].emplace_back(v1, wt);
  }
  int ans = 0;
  priority queue<Edge, vector<Edge>, CompareEdge>
pq;
  vector<br/>bool> vis(vtces, false);
  pq.push(Edge(0, 0)); // Start with any vertex (0 in this
case) with 0 weight
  while (!pq.empty()) {
     Edge rem = pq.top();
    pq.pop();
    if (vis[rem.v]) {
```

Step-by-Step Execution

Step 1: Populate the Graph

The adjacency list (graph) for the given input will look like this:

```
0: (1, 10), (3, 40)

1: (0, 10), (2, 10)

2: (1, 10), (3, 10)

3: (2, 10), (0, 40), (4, 2)

4: (3, 2), (5, 3), (6, 8)

5: (4, 3), (6, 3)

6: (5, 3), (4, 8)
```

Step 2: Initialize Priority Queue and Visited Array

- Start with vertex 0 and push an edge (0,
 0) to the priority queue (pq).
- Initially:

```
pq: [(0, 0)] (min-heap: weight 0, vertex 0) vis: [false, false, false, false, false, false] ans: 0
```

Step 3: Prim's Algorithm

• Iteration 1:

- o Pop (0, 0) from pq.
- Add weight 0 to ans. Now ans =0.
- Mark vertex 0 as visited (vis[0] = true).
- Push neighbors (1, 10) and (3, 40) to pq.

```
pq: [(1, 10), (3, 40)] vis: [true, false, false, false, false, false] ans: 0
```

• Iteration 2:

- o Pop (1, 10) from pq.
- Add weight 10 to ans. Now ans = 10.
- Mark vertex 1 as visited (vis[1] = true).
- Push neighbors (2, 10) to pq (skip (0, 10) because 0 is already visited).

```
pq: [(2, 10), (3, 40)]
vis: [true, true, false, false, false, false]
```

```
continue;
}
vis[rem.v] = true;
ans += rem.wt;

for (Edge nbr : graph[rem.v]) {
    if (!vis[nbr.v]) {
        pq.push(nbr);
    }
}
cout << ans << endl;
return 0;
}</pre>
```

ans: 10

• Iteration 3:

- o Pop (2, 10) from pq.
- Add weight 10 to ans. Now ans = 20.
- Mark vertex 2 as visited (vis[2] = true).
- Push neighbor (3, 10) to pq (skip (1, 10) because 1 is already visited).

pq: [(3, 10), (3, 40)] vis: [true, true, true, false, false, false, false] ans: 20

• Iteration 4:

- o Pop (3, 10) from pq.
- Add weight 10 to ans. Now ans = 30.
- Mark vertex 3 as visited (vis[3] = true).
- Push neighbors (4, 2) to pq (skip (2, 10) and (0, 40) because 2 and 0 are already visited).

pq: [(3, 40), (4, 2)] vis: [true, true, true, true, false, false, false] ans: 30

• Iteration 5:

- o Pop (4, 2) from pq.
- Add weight 2 to ans. Now ans = 32.
- Mark vertex 4 as visited (vis[4] = true).
- O Push neighbors (5, 3) and (6, 8) to pq (skip (3, 2) because 3 is already visited).

pq: [(3, 40), (5, 3), (6, 8)] vis: [true, true, true, true, true, false, false] ans: 32

• Iteration 6:

- o Pop (5, 3) from pq.
- Add weight 3 to ans. Now ans = 35.
- Mark vertex 5 as visited (vis[5] = true).
- Push neighbor (6, 3) to pq (skip (4, 3) because 4 is already visited).

pq: [(3, 6), (6, 8), (3, 40)] vis: [true, true, true, true, true, false]

ans: 35 **Iteration 7:** o Pop (6, 3) from pq. Add weight 3 to ans. Now ans = Mark vertex 6 as visited (vis[6] = true). o Skip pushing neighbors because all are already visited. pq: [(3, 40), (6, 8)] vis: [true, true, true, true, true, true, true] ans: 38 Final MST Weight: All vertices are visited, and the MST weight is 38. Output:-

38

Negative Wt Cycle Detection in C++ #include <iostream> #include <vector> #include <climits> using namespace std; struct Edge { int u, v, weight; **}**; bool isNegativeWeightCycle(int n, vector<Edge>& edges) vector<int> dist(n, INT MAX); dist[0] = 0; // Starting from vertex 0 // Relaxation process for (int i = 0; i < n - 1; ++i) { for (const auto& edge : edges) { if (dist[edge.u] != INT_MAX && dist[edge.u] + edge.weight < dist[edge.v]) { dist[edge.v] = dist[edge.u] + edge.weight; } } } // Checking for negative weight cycles for (const auto& edge : edges) { if (dist[edge.u] != INT_MAX && dist[edge.u] + edge.weight < dist[edge.v]) { return true; // Negative weight cycle detected } return false; // No negative weight cycle found } int main() { // Hardcoded input int n = 3; // Number of vertices int m = 3; // Number of edges $vector < Edge > edges = \{\{0, 1, -1\}, \{1, 2, -4\}, \{2, 0, 3\}\}\}; //$ Edges with (u, v, weight) if (isNegativeWeightCycle(n, edges)) { cout << "1\n"; // Negative weight cycle detected } else { cout << "0\n"; // No negative weight cycle found return 0;

Step-by-Step Execution

Input Edges:

```
Edge 1: (0 \rightarrow 1, \text{ weight} = -1)
Edge 2: (1 \rightarrow 2, \text{ weight} = -4)
Edge 3: (2 \rightarrow 0, \text{ weight} = 3)
```

Initial State:

 $dist = [0, INT_MAX, INT_MAX]$

Relaxation Process (n-1 = 2 times):

Iteration 1:

```
Relax Edge (0 \rightarrow 1, \text{ weight = -1}):
dist[1] = min(INT MAX, dist[0])
+ (-1) = \min(INT\_MAX, 0 + (-1))
= -1
dist = [0, -1, INT\_MAX]
```

Relax Edge $(1 \rightarrow 2, \text{ weight = -4})$:

Relax Edge $(2 \rightarrow 0, \text{ weight = 3})$:

$$dist[0] = min(0, dist[2] + 3) = min(0, -5 + 3) = -2$$

 $dist = [-2, -1, -5]$

Iteration 2:

Relax Edge $(0 \rightarrow 1, \text{ weight = -1})$:

$$dist[1] = min(-1, dist[0] + (-1)) = min(-1, -2 + (-1)) = -3$$

 $dist = [-2, -3, -5]$

Relax Edge $(1 \rightarrow 2, \text{ weight} = -4)$:

$$dist[2] = min(-5, dist[1] + (-4)) = min(-5, -3 + (-4)) = -7$$

 $dist = [-2, -3, -7]$

Relax Edge $(2 \rightarrow 0, \text{ weight} = 3)$:

$$dist[0] = min(-2, dist[2] + 3) = min(-2, -7 + 3) = -4$$

 $dist = [-4, -3, -7]$

Check for Negative Weight Cycles:

Try relaxing all edges once more:

	$\circ \text{Relax Edge } (0 \to 1, \text{ weight = -1}):$	
	dist[1] = min(-3, dist[0] + (-1)) = min(-3, -4 + (-1)) = -5	
	 At this point, the distance to vertex 1 changes, which means a negative weight cycle exists. 	
	Output:	
	Since a negative weight cycle is detected:	
	1	
Output:-		
1		

```
No of Distinct Island in C++
#include <iostream>
#include <vector>
#include <unordered_set>
using namespace std;
// Function prototypes
void dfs(vector<vector<int>>& arr, int row, int col,
string& psf);
int numDistinctIslands(vector<vector<int>>& arr);
// Depth-first search to mark all connected land cells of
an island
void dfs(vector<vector<int>>& arr, int row, int col,
string& psf) {
       arr[row][col] = 0; // Marking current cell as visited
      int n = arr.size();
      int m = arr[0].size();
      // Directions: up, right, down, left
       vector<pair<int, int>> dirs = \{\{-1, 0\}, \{0, 1\}, \{1, 0\}, \{0, 1\}, \{1, 0\}, \{0, 1\}, \{1, 0\}, \{0, 1\}, \{1, 0\}, \{0, 1\}, \{1, 0\}, \{0, 1\}, \{1, 0\}, \{0, 1\}, \{1, 0\}, \{0, 1\}, \{1, 0\}, \{0, 1\}, \{1, 0\}, \{0, 1\}, \{1, 0\}, \{0, 1\}, \{1, 0\}, \{0, 1\}, \{1, 0\}, \{0, 1\}, \{1, 0\}, \{0, 1\}, \{1, 0\}, \{0, 1\}, \{1, 0\}, \{0, 1\}, \{1, 0\}, \{0, 1\}, \{1, 0\}, \{0, 1\}, \{1, 0\}, \{0, 1\}, \{1, 0\}, \{0, 1\}, \{1, 0\}, \{0, 1\}, \{1, 0\}, \{0, 1\}, \{1, 0\}, \{0, 1\}, \{1, 0\}, \{0, 1\}, \{1, 0\}, \{0, 1\}, \{1, 0\}, \{0, 1\}, \{1, 0\}, \{0, 1\}, \{1, 0\}, \{0, 1\}, \{1, 0\}, \{0, 1\}, \{1, 0\}, \{0, 1\}, \{1, 0\}, \{0, 1\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}, \{1, 0\}
       string dirStr = "urdl"; // Corresponding directions
characters
      for (int i = 0; i < 4; ++i) {
             int newRow = row + dirs[i].first;
             int newCol = col + dirs[i].second;
             if (\text{newRow} \ge 0 \&\& \text{newRow} \le n \&\& \text{newCol} \ge 0)
&& newCol < m && arr[newRow][newCol] == 1) {
                    psf += dirStr[i]; // Append direction character to
path string
                    dfs(arr, newRow, newCol, psf);
      psf += "a"; // Append anchor to indicate end of island
path
// Function to find number of distinct islands
int numDistinctIslands(vector<vector<int>>& arr) {
       int n = arr.size():
      if (n == 0) return 0;
      int m = arr[0].size();
      unordered_set<string> islands; // Set to store distinct
island paths
      for (int i = 0; i < n; ++i) {
             for (int j = 0; j < m; ++j) {
                   if (arr[i][j] == 1) {
                           string psf = "x"; // Starting character to
represent new island
                           dfs(arr, i, j, psf);
                          islands.insert(psf); // Insert island path into set
      }
      return islands.size(); // Return the number of distinct
islands
```

Input:

Grid: 100 $0\ 1\ 0$ 111

Execution Steps

Step 1: Initializing Variables

- The grid has 3 rows x 3 columns.
- An empty set islands is initialized to store distinct island shapes.

Step 2: Traversing the Grid

- 1. At (0, 0):
 - o Start a DFS and encode the path:

sql Copy code psf = "x" (start)Move down: psf = "xurda" (anchor added after exploring up, right, down, left)

- Add "xurda" to islands.
- 2. At (1, 1):
 - Start a DFS and encode the path:

arduino Copy code psf = "x" (start)No other cells connected to (1, 1): psf = "xurda" (isolated cell)

- o Add "xurda" to islands (already present).
- 3. At (2, 0):
 - Start a DFS and encode the path:

sql Copy code psf = "x" (start)Move right: psf = "xrd" (connects $(2, 0) \rightarrow (2, 1)$ Move right again: psf = "xrdrr" (connects $(2, 1) \rightarrow (2, 2)$) Move up: psf = "xrdrru" (connects $(2, 2) \rightarrow (1, 2)$) Add anchors: psf ="xrdrruarrarrarr"

Add "xrdrruarrarrarr" to

```
islands.
int main() {
  // Hardcoded input
                                                               Step 3: Count Distinct Islands
  vector<vector<int>> arr = {
    \{1, 0, 0\},\
                                                                       The set islands contains:
    \{0, 1, 0\},\
    \{1, 1, 1\}
                                                                       {"xurda", "xrdrruarrarrarr"}
  };
                                                                       The size of the set is 2.
  // Calculating number of distinct islands
  cout << numDistinctIslands(arr) << endl;</pre>
  return 0;
                                                               Output:
                                                               2
Output:-
```

#include <iostream> #include <vector> using namespace std; void dfs(vector<vector<int>>& arr, int i, int j) { if (i < 0 | | j < 0 | | i > = arr.size() | | j > = arr[0].size() $| | arr[i][j] == 0) {$ return; } arr[i][j] = 0;dfs(arr, i + 1, j);dfs(arr, i - 1, j); dfs(arr, i, j + 1);dfs(arr, i, j - 1); int numEnclaves(vector<vector<int>>& arr) { int m = arr.size();int n = arr[0].size();// Marking connected components touching the boundaries for (int i = 0; i < m; ++i) { for (int j = 0; j < n; ++j) { if $((i == 0 \mid | j == 0 \mid | i == m - 1 \mid | j == n - 1) &&$ $arr[i][j] == 1) {$ dfs(arr, i, j); // Counting remaining land cells int count = 0; for (int i = 0; i < m; ++i) { for (int j = 0; j < n; ++j) { $if (arr[i][j] == 1) {$ ++count; return count; } int main() { int m = 4, n = 4; vector<vector<int>> arr = { $\{0, 0, 0, 0\},\$ $\{1, 0, 1, 0\},\$ $\{0, 1, 1, 0\},\$ $\{0, 0, 0, 0\}$ **}**; int result = numEnclaves(arr); cout << result << endl; return 0; }

No of enclaves in C++

Dry Run

Input Grid:

```
\{0, 0, 0, 0\},\
\{1, 0, 1, 0\},\
\{0, 1, 1, 0\},\
\{0, 0, 0, 0\}
```

Step 1: DFS from Boundary Cells

- Boundary cells: We start by scanning the boundary cells (first and last rows, first and last columns). The boundary cells are:
 - o Row 0: {0, 0, 0, 0} Row 3: {0, 0, 0, 0}
 - Column 0: {1, 0, 0, 0} Column 3: {0, 0, 0, 0}
 - The boundary cells that are 1 (land) are:

Step 2: Marking Land Cells Connected to **Boundary**

1. DFS starting at (1, 0):

(1, 0)

- Mark arr[1][0] as 0.
- Explore its neighbors (down: (2, 0), left: out of bounds, right: (1, 1), up: (0, 0)).
- No other connected land cells.

Step 3: Count Remaining Land Cells

After marking the connected land cells to the boundary, the grid looks like this:

```
\{0, 0, 0, 0\},\
\{0, 0, 1, 0\},\
\{0, 1, 1, 0\},\
\{0, 0, 0, 0\}
```

Now, we count the remaining land cells (1) in the grid:

(1, 2), (2, 1),and (2, 2) are the remaining land cells.

Final Answer:

The number of enclosed land cells is 3.

Output:

	3
Output:-	
3	

Optimize water distribution in C++ #include <iostream> #include <vector> #include <queue> #include <utility> using namespace std; class Pair { public: int vtx; int wt: Pair(int vtx, int wt) { this->vtx = vtx; this->wt=wt;bool operator>(const Pair& other) const { return this->wt > other.wt; } **}**; int minCostToSupplyWater(int n, vector<int>& wells, vector<vector<int>>& pipes) { vector < Pair >> graph(n + 1);for (const auto& pipe : pipes) { int u = pipe[0]; int v = pipe[1]; int wt = pipe[2]; graph[u].emplace_back(v, wt); graph[v].emplace_back(u, wt); for (int i = 1; $i \le n$; ++i) { graph[i].emplace_back(0, wells[i - 1]); graph[0].emplace_back(i, wells[i - 1]); int ans = 0; priority_queue<Pair, vector<Pair>, greater<Pair>> pq; pq.emplace(0, 0);vector < bool > vis(n + 1, false);while (!pq.empty()) { Pair rem = pq.top(); pq.pop(); if (vis[rem.vtx]) continue; ans += rem.wt; vis[rem.vtx] = true; for (const Pair& nbr : graph[rem.vtx]) { if (!vis[nbr.vtx]) { pq.push(nbr); } return ans; int main() { int v = 3, e = 2; vector \leq int \geq wells = $\{1, 2, 2\}$; vector<vector<int>> pipes = {{1, 2, 1}, {2, 3, 1}};

cout << minCostToSupplyWater(v, wells, pipes) <<</pre>

```
int v = 3, e = 2;
vector<int> wells = {1, 2, 2};
vector<vector<int>> pipes = {{1, 2, 1}, {2, 3, 1}};
```

- v = 3: Number of houses (vertices).
- wells = {1, 2, 2}: The cost to build a well at house 1, 2, and 3.
- pipes = {{1, 2, 1}, {2, 3, 1}}: The pipes connecting houses, with respective costs.

Step 1: Construct the Graph

We begin by creating an adjacency list that represents the graph, including both the pipes and wells.

• Graph Construction:

- Create an adjacency list graph with v + 1 = 4 nodes (including the virtual node 0).
- Add edges for the pipes between houses:
 - Pipe from 1 to 2 with cost 1.
 - Pipe from 2 to 3 with cost 1.
- Add edges for the wells:
 - Well for house 1 (cost 1), connect node 0 to node 1.
 - Well for house 2 (cost 2), connect node 0 to node 2.
 - Well for house 3 (cost 2), connect node 0 to node 3.

• Graph Representation:

```
Node 0 (virtual node) \rightarrow {(1, 1), (2, 2), (3, 2)}
Node 1 \rightarrow {(2, 1), (0, 1)}
Node 2 \rightarrow {(1, 1), (3, 1), (0, 2)}
Node 3 \rightarrow {(2, 1), (0, 2)}
```

Step 2: Prim's Algorithm with Min-Heap

We will use **Prim's Algorithm** to find the Minimum Spanning Tree (MST) with a priority queue (min-heap).

• **Priority Queue Initialization**: Start with node 0 (virtual node), which has no cost yet, so we push (0, 0) into the priority queue.

```
endl;
                                                                       Priority Queue: [(0, 0)]
  return 0;
                                                                       Step 3: First Iteration (start with
                                                                       node 0)
                                                                              Pop from the priority queue:
                                                                           0
                                                                               (0, 0) is popped, meaning we're
                                                                               at the virtual node 0 with a cost
                                                                               Visit Node 0 and explore its
                                                                               neighbors (nodes 1, 2, 3):
                                                                                       Add the edges to the
                                                                                       priority queue:
                                                                                                Edge (0 \rightarrow 1,
                                                                                                cost 1)
                                                                                                Edge (0 \rightarrow 2,
                                                                                                cost 2)
                                                                                                Edge (0 \rightarrow 3,
                                                                                                cost 2)
                                                                       After this step:
                                                                       Priority Queue: [(1, 1), (2, 2), (2, 3)]
                                                                       Visited nodes: [0]
                                                                       Total Cost: 0
                                                                       Step 4: Second Iteration (pop node
                                                                              Pop from the priority queue:
                                                                               (1, 1) is popped, meaning we're
                                                                               now at node 1 with a cost of 1.
                                                                               Visit Node 1 and explore its
                                                                               neighbors:
                                                                                       Node 0 has already been
                                                                                       visited, so ignore.
                                                                                       Add edge (1 \rightarrow 2, \cos t 1)
                                                                                       to the priority queue.
                                                                       After this step:
                                                                       Priority Queue: [(1, 2), (2, 3), (1, 2)]
                                                                       Visited nodes: [0, 1]
                                                                       Total Cost: 1
                                                                       Step 5: Third Iteration (pop node 2)
                                                                               Pop from the priority queue:
                                                                               (1, 2) is popped, meaning we're
                                                                               now at node 2 with a cost of 1.
                                                                               Visit Node 2 and explore its
                                                                               neighbors:
                                                                                       Node 1 has already been
                                                                                       visited, so ignore.
                                                                                       Node 3 is unvisited, so
                                                                                       add edge (2 \rightarrow 3, \cos t 1)
                                                                                       to the priority queue.
                                                                                       Node 0 has already been
                                                                                       visited, so ignore.
                                                                       After this step:
```

Priority Queue: [(1, 3), (2, 3), (1, 2)]

Visited nodes: [0, 1, 2] Total Cost: 2 Step 6: Fourth Iteration (pop node Pop from the priority queue: (1, 3) is popped, meaning we're now at node 3 with a cost of 1. Visit Node 3 and explore its neighbors: Node 2 has already been visited, so ignore. Node 0 has already been visited, so ignore. After this step: Priority Queue: [(2, 3)] Visited nodes: [0, 1, 2, 3] Total Cost: 3 **Step 7: Termination** The priority queue is empty, and all nodes are visited. Final Total Cost: 3. **Final Output** The minimum cost to supply water is: 3Output:-

```
Redundant Connection in C++
#include <iostream>
#include <vector>
using namespace std;
class UnionFind {
public:
  vector<int> parent;
  vector<int> rank;
  UnionFind(int n) {
    parent.resize(n + 1);
    rank.resize(n + 1, 1);
    for (int i = 1; i \le n; ++i) {
       parent[i] = i;
  }
  int find(int x) {
    if (parent[x] != x) {
       parent[x] = find(parent[x]); // Path compression
    return parent[x];
  }
  void unionSets(int x, int y) {
    int rootX = find(x);
    int rootY = find(y);
    if (rootX != rootY) {
       if (rank[rootX] > rank[rootY]) {
          parent[rootY] = rootX;
       } else if (rank[rootX] < rank[rootY]) {</pre>
          parent[rootX] = rootY;
       } else {
          parent[rootY] = rootX;
          rank[rootX]++;
};
vector<int>
findRedundantConnection(vector<vector<int>>& edges) {
  int n = edges.size():
  UnionFind uf(n);
  for (auto& edge : edges) {
    int u = edge[0];
    int v = edge[1];
    if (uf.find(u) == uf.find(v)) {
       return edge; // This edge is a redundant
connection
    uf.unionSets(u, v);
  }
  return {};
int main() {
```

Dry Run

Input:

```
edges = {
   \{1, 2\},\
   {1, 3}.
   \{2, 3\}
```

Step-by-Step Execution:

- Initialization:
 - o UnionFind:
 - parent = [1, 2, 3] (each node is its own parent initially)
 - rank = [1, 1, 1] (all ranks start as 1)

Processing Edges:

- 1. **Edge (1, 2)**:
 - find(1) = 1, $find(2) = 2 \rightarrow$ Different components.
 - Union the components:
 - Set parent[2] = 1.
 - Update rank[1] = 2.
 - Updated state:
 - parent = [1, 1, 3]
 - rank = [1, 2, 1].
- 2. **Edge (1, 3)**:
 - find(1) = 1, $find(3) = 3 \rightarrow$ Different components.
 - Union the components:
 - Set parent[3] = 1.
 - Updated state:
 - parent = [1, 1, 1]
 - rank = [1, 2, 1].
- 3. **Edge (2, 3)**:
 - find(2) = 1, $find(3) = 1 \rightarrow Same$ component (cycle detected).
 - Return the edge (2, 3) as the redundant connection.