#include <iostream> #include <vector> #include <deque> using namespace std; struct Pair { int i; int j; string psf; Pair(int i, int j, string psf) { this->i = i; this->j = j; this->psf = psf; **}**; void printPaths(vector<vector<int>>& dp, vector<int>& vals, vector<int>& wts, int i, int j, string psf, deque<Pair>& que) { while (!que.empty()) { Pair rem = que.front(); que.pop_front(); if (rem.i == 0 | rem.j == 0)cout << rem.psf << endl;</pre> } else { int exc = dp[rem.i - 1][rem.j];if $(rem.j \ge wts[rem.i - 1])$ { int inc = dp[rem.i - 1][rem.j wts[rem.i - 1]] + vals[rem.i - 1]; $if (dp[rem.i][rem.j] == inc) {$ que.push_back(Pair(rem.i - 1, rem.j - wts[rem.i - 1], to string(rem.i - 1) + " " + rem.psf)); if (dp[rem.i][rem.j] == exc) { que.push_back(Pair(rem.i - 1, rem.j, rem.psf)); } void knapsackPaths(vector<int>& vals, vector<int>& wts, int cap) { int n = vals.size();

vector < vector < int >> dp(n + 1,

Paths of 0-1 knapsack In C++

Dry Run Using a Table

Step 1: Initialize DP Table

We define a **DP table (dp[i][j])**, where:

 dp[i][j] = Maximum value that can be obtained using the first i items with a capacity j.

Step 1.1: Base Case

• If i = 0 (no items), or j = 0 (zero capacity), dp[i] [i] = 0.

Step 1.2: Fill the DP Table

If including the item **does not exceed capacity**, we check:

- Exclude item $i \rightarrow dp[i-1][j]$
- Include item $i \rightarrow dp[i-1][j wts[i-1]] + vals[i-1]$

i∖j	0	1	2	3	4	5	6	7
0	0	0	0	0	0	0	0	0
1 (val= 15, wt=2	0	0	15	15	15	15	15	15
2 (val= 14, wt=5)	0	0	15	15	15	15	15	15

```
vector < int > (cap + 1, 0));
  for (int i = 1; i \le n; i++) {
     for (int j = 1; j \le cap; j++) {
        dp[i][j] = dp[i - 1][j];
        if (j \ge wts[i - 1]) {
          dp[i][j] = max(dp[i][j], dp[i - 1][j -
wts[i-1]] + vals[i-1]);
  int ans = dp[n][cap];
  cout << "Maximum value: " << ans <<
endl;
  deque<Pair> que;
  que.push_back(Pair(n, cap, ""));
  printPaths(dp, vals, wts, n, cap, "", que);
int main() {
  vector<int> vals = \{15, 14, 10, 45, 30\};
  vector<int> wts = \{2, 5, 1, 3, 4\};
  int cap = 7;
  knapsackPaths(vals, wts, cap);
  return 0;
```

3 (val= 10, wt=1	1	10	15	25	25	25	25	25
4 (val= 45, wt=3)	0	10	15	45	55	60	70	70
5 (val= 30, wt=4		10	15	45	55	60	70	75

The **maximum value** obtained is 75 at dp[5][7].

Step 2: Print All Paths

Using **backtracking**, the function printPaths reconstructs paths that lead to dp[n][cap] = 75.

Backtracking Paths

- 1. Start at dp[5][7] = 75
 - o dp[4][3] = $45 \rightarrow$ Item 5 (index 4, value 30, weight 4) is included.
- 2. Now at dp[4][3] = 45
 - o $dp[3][0] = 0 \rightarrow Item 4 (index 3, value 45, weight 3) is included.$

Thus, one of the optimal selections is {30, 45}.

Final Output

Maximum value: 75 4 3

Output:-	
Output:- Maximum value: 75	
3 4	