

0/1 KnapSack in C++

```
#include <iostream>
#include <vector>

using namespace std;

class ZeroOneKnapsack {
public:
    int knapsack(int n, vector<int>& vals,
vector<int>& wts, int cap) {
        vector<vector<int>>> dp(n + 1, vector<int>(cap +
1, 0));

        for (int i = 1; i <= n; i++) {
            for (int j = 1; j <= cap; j++) {
                if (j >= wts[i - 1]) {
                    int remainingCap = j - wts[i - 1];

                    if (dp[i - 1][remainingCap] + vals[i - 1] >
dp[i - 1][j]) {
                        dp[i][j] = dp[i - 1][remainingCap] +
vals[i - 1];
                    } else {
                        dp[i][j] = dp[i - 1][j];
                    }
                } else {
                    dp[i][j] = dp[i - 1][j];
                }
            }
        }

        return dp[n][cap];
    }
};

int main() {
    ZeroOneKnapsack solution;

    // Input parameters
    int n = 5;
    vector<int> vals = {15, 14, 10, 45, 30};
    vector<int> wts = {2, 5, 1, 3, 4};
    int cap = 7;

    // Compute maximum value using knapsack
function
    int maxVal = solution.knapsack(n, vals, wts, cap);

    // Output the maximum value
    cout << "Maximum value that can be obtained: " <<
maxVal << endl;

    return 0;
}
```

Dry Run of the ZeroOneKnapsack Problem:

Input:

```
n = 5;
vals = {15, 14, 10, 45, 30};
wts = {2, 5, 1, 3, 4};
cap = 7;
```

Step 1: Initialize the DP Table

We initialize a 2D DP table of size $(n + 1) \times (cap + 1)$ to store the maximum values obtainable for each subproblem. Each cell $dp[i][j]$ will represent the maximum value achievable with the first i items and a knapsack capacity of j .

Initially, the DP table is filled with zeros:

i\j	0	1	2	3	4	5	6	7
0	0	0	0	0	0	0	0	0
1	0	0	0	0	0	0	0	0
2	0	0	0	0	0	0	0	0
3	0	0	0	0	0	0	0	0
4	0	0	0	0	0	0	0	0
5	0	0	0	0	0	0	0	0

Step 2: Fill the DP Table

We iterate through each item ($i = 1$ to n) and each knapsack capacity ($j = 1$ to cap). The idea is to decide whether to include the current item or not.

Item 1 (Value = 15, Weight = 2)

- **Capacity 1:** $dp[1][1] = 0$ (Cannot include this item as the weight is greater than the capacity)
- **Capacity 2:** $dp[1][2] = \max(dp[0][2], dp[0][0] + 15) = \max(0, 15) = 15$
- **Capacity 3:** $dp[1][3] = \max(dp[0][3], dp[0][1] + 15) = \max(0, 15) = 15$
- **Capacity 4:** $dp[1][4] = \max(dp[0][4], dp[0][2] + 15) = \max(0, 15) = 15$
- **Capacity 5:** $dp[1][5] = \max(dp[0][5], dp[0][3] + 15) = \max(0, 15) = 15$
- **Capacity 6:** $dp[1][6] = \max(dp[0][6], dp[0][4] + 15) = \max(0, 15) = 15$
- **Capacity 7:** $dp[1][7] = \max(dp[0][7], dp[0][5] + 15) = \max(0, 15) = 15$

i\j	0	1	2	3	4	5	6	7
0	0	0	0	0	0	0	0	0

i\j	0	1	2	3	4	5	6	7
1	0	0	15	15	15	15	15	15
2	0	0	0	0	0	0	0	0
3	0	0	0	0	0	0	0	0
4	0	0	0	0	0	0	0	0
5	0	0	0	0	0	0	0	0

Item 2 (Value = 14, Weight = 5)

- **Capacity 1 to 4:** The weight is greater than the capacity, so we can't include this item.
- **Capacity 5:** $dp[2][5] = \max(dp[1][5], dp[1][0] + 14) = \max(15, 14) = 15$
- **Capacity 6:** $dp[2][6] = \max(dp[1][6], dp[1][1] + 14) = \max(15, 14) = 15$
- **Capacity 7:** $dp[2][7] = \max(dp[1][7], dp[1][2] + 14) = \max(15, 29) = 29$

i\j	0	1	2	3	4	5	6	7
0	0	0	0	0	0	0	0	0
1	0	0	15	15	15	15	15	15
2	0	0	15	15	15	15	15	29
3	0	0	0	0	0	0	0	0
4	0	0	0	0	0	0	0	0
5	0	0	0	0	0	0	0	0

Item 3 (Value = 10, Weight = 1)

- **Capacity 1:** $dp[3][1] = \max(dp[2][1], dp[2][0] + 10) = \max(0, 10) = 10$
- **Capacity 2:** $dp[3][2] = \max(dp[2][2], dp[2][1] + 10) = \max(15, 10) = 15$
- **Capacity 3:** $dp[3][3] = \max(dp[2][3], dp[2][2] + 10) = \max(15, 25) = 25$
- **Capacity 4:** $dp[3][4] = \max(dp[2][4], dp[2][3] + 10) = \max(15, 25) = 25$
- **Capacity 5:** $dp[3][5] = \max(dp[2][5], dp[2][4] + 10) = \max(15, 25) = 25$
- **Capacity 6:** $dp[3][6] = \max(dp[2][6], dp[2][5] + 10) = \max(15, 25) = 25$
- **Capacity 7:** $dp[3][7] = \max(dp[2][7], dp[2][6] + 10) = \max(29, 25) = 29$

i\j	0	1	2	3	4	5	6	7
0	0	0	0	0	0	0	0	0
1	0	0	15	15	15	15	15	15
2	0	0	15	15	15	15	15	29
3	0	10	15	25	25	25	25	29
4	0	0	0	0	0	0	0	0
5	0	0	0	0	0	0	0	0

Item 4 (Value = 45, Weight = 3)

- **Capacity 1 to 2:** Cannot include this item.
- **Capacity 3:** $dp[4][3] = \max(dp[3][3], dp[3][0] + 45) = \max(25, 45) = 45$
- **Capacity 4:** $dp[4][4] = \max(dp[3][4], dp[3][1] + 45) = \max(25, 55) = 55$
- **Capacity 5:** $dp[4][5] = \max(dp[3][5], dp[3][2] + 45) = \max(25, 55) = 55$
- **Capacity 6:** $dp[4][6] = \max(dp[3][6], dp[3][3] + 45) = \max(25, 70) = 70$
- **Capacity 7:** $dp[4][7] = \max(dp[3][7], dp[3][4] + 45) = \max(29, 70) = 70$

Item 5 (Value = 30, Weight = 4)

- **Capacity 1 to 3:** Cannot include this item.
- **Capacity 4:** $dp[5][4] = \max(dp[4][4], dp[4][0] + 30) = \max(55, 30) = 55$
- **Capacity 5:** $dp[5][5] = \max(dp[4][5], dp[4][1] + 30) = \max(55, 30) = 55$
- **Capacity 6:** $dp[5][6] = \max(dp[4][6], dp[4][2] + 30) = \max(70, 30) = 70$
- **Capacity 7:** $dp[5][7] = \max(dp[4][7], dp[4][3] + 30) = \max(70, 75) = 75$

Step 3: Final DP Table

i \ j	0	1	2	3	4	5	6	7
0	0	0	0	0	0	0	0	0
1	0	0	15	15	15	15	15	15
2	0	0	15	15	15	15	15	29
3	0	10	15	25	25	25	25	29
4	0	10	15	45	55	55	70	70
5	0	10	15	45	55	55	70	75

Result:

The maximum value that can be obtained with a knapsack capacity of 7 is 75.

Output:

Maximum value that can be obtained: 75

The maximum value that can be obtained is stored in $dp[5][7] = 75$.