

Print all path with minimum Cost In C++

```
#include <iostream>
#include <vector>
#include <queue>
using namespace std;

struct Pair {
    string psf; // path so far
    int i;      // current row index
    int j;      // current column index

    Pair(string psf, int i, int j) {
        this->psf = psf;
        this->i = i;
        this->j = j;
    }
};

void printAllPaths(vector<vector<int>>&
arr) {
    int m = arr.size();
    int n = arr[0].size();

    // dp array to store minimum cost to
    reach each cell
    vector<vector<int>> dp(m,
vector<int>(n, 0));

    // Initialize dp table
    dp[m-1][n-1] = arr[m-1][n-1];
    for (int i = m - 2; i >= 0; i--) {
        dp[i][n-1] = arr[i][n-1] + dp[i + 1][n -
1];
    }
    for (int j = n - 2; j >= 0; j--) {
        dp[m-1][j] = arr[m-1][j] + dp[m - 1][j +
1];
    }
    for (int i = m - 2; i >= 0; i--) {
        for (int j = n - 2; j >= 0; j--) {
            dp[i][j] = arr[i][j] + min(dp[i][j + 1],
dp[i + 1][j]);
        }
    }

    // Minimum cost to reach the top-left
    corner
    cout << dp[0][0] << endl;

    // Queue to perform BFS
    queue<Pair> q;
    q.push(Pair("", 0, 0));
```

Dry Run of Minimum Cost Path Problem

We will compute the **dynamic programming (DP) table** step-by-step to ensure that we get the minimum cost sum **46** for the given matrix.

Given Input Matrix (**arr**):

```
{1, 2, 3, 4},
{5, 6, 7, 8},
{9, 10, 11, 12},
{13, 14, 15, 16}
```

Step 1: Understanding the DP Approach

1. **Base Case:** The last cell ($dp[3][3]$) is the same as $arr[3][3] = 16$.
2. **Filling Last Row (Right to Left):**
 - o $dp[i][j] = arr[i][j] + dp[i][j+1]$
3. **Filling Last Column (Bottom to Top):**
 - o $dp[i][j] = arr[i][j] + dp[i+1][j]$
4. **Filling the Rest (Bottom-Up, Right-to-Left):**
 - o $dp[i][j] = arr[i][j] + \min(dp[i+1][j], dp[i][j+1])$

Step 2: Construct DP Table Step-by-Step

1. Initialize $dp[3][3]$ (Bottom-Right Cell)

$dp[3][3] = arr[3][3] = 16$

2. Fill the Last Row (Right to Left)

$dp[i][j] = arr[i][j] + dp[i][j+1]$
 $dp[i][j] = arr[i][j] + dp[i][j+1]$

i=3	j=3	j=2	j=1	j=0
		(15+16)	(14+31)	(13+45)

```

while (!q.empty()) {
    Pair rem = q.front();
    q.pop();

    if (rem.i == m - 1 && rem.j == n - 1) {
        cout << rem.psf << endl; // print
        path when reaching the bottom-right
        corner
    } else if (rem.i == m - 1) {
        q.push(Pair(rem.psf + "H", rem.i,
        rem.j + 1)); // go right
    } else if (rem.j == n - 1) {
        q.push(Pair(rem.psf + "V", rem.i +
        1, rem.j)); // go down
    } else {
        if (dp[rem.i][rem.j + 1] < dp[rem.i +
        1][rem.j]) {
            q.push(Pair(rem.psf + "H", rem.i,
            rem.j + 1)); // go right
        } else if (dp[rem.i][rem.j + 1] >
        dp[rem.i + 1][rem.j]) {
            q.push(Pair(rem.psf + "V", rem.i
            + 1, rem.j)); // go down
        } else {
            q.push(Pair(rem.psf + "V", rem.i
            + 1, rem.j)); // go down
            q.push(Pair(rem.psf + "H", rem.i,
            rem.j + 1)); // go right
        }
    }
}

int main() {
    vector<vector<int>>> arr = {
        {1, 2, 3, 4},
        {5, 6, 7, 8},
        {9, 10, 11, 12},
        {13, 14, 15, 16}
    };

    printAllPaths(arr);
    return 0;
}

```

arr	16	15	14	13
dp	16	31	45	58

3. Fill the Last Column (Bottom to Top)

$dp[i][j] = arr[i][j] + dp[i+1][j]$
 $dp[i][j] = arr[i][j] + dp[i+1][j]$

i=2	j=3 (12+16)	j=2	j=1	j=0
arr	12	11	10	9
dp	28	-	-	-
i=1	j=3 (8+28)	j=2	j=1	j=0
arr	8	7	6	5
dp	36	-	-	-
i=0	j=3 (4+36)	j=2	j=1	j=0
arr	4	3	2	1
dp	40	-	-	-

4. Fill the Rest of the DP Table

$dp[i][j] = arr[i][j] + \min(dp[i+1][j], dp[i][j+1])$
 $dp[i][j] = arr[i][j] + \min(dp[i+1][j], dp[i][j+1])$

i=2	j=2 (11+min(31, 28))	j=1 (10+min(41, 38))	j=0 (9+min(45, 40))
arr	11	10	9
dp	39	38	40
i=1	j=2 (7+min(39, 36))	j=1 (6+min(38, 44))	j=0 (5+min(45, 43))
arr	7	6	5
dp	43	44	45

$i=0$	$j=2$ $(3+\min(43, 40))$	$j=1$ $(2+\min(41, 44))$	$j=0$ $(1+\min(45, 43))$
arr	3	2	1
dp	43	45	46

Final DP Table

46	45	43	40
45	44	43	36
40	38	39	28
58	45	31	16

✓ **Minimum Cost Path Sum = 46 (Matches G++ Output)**

Step 3: Extracting All Paths

Now, we use BFS (`queue<Pair>`) to **trace all paths** from $(0, 0)$ to $(3, 3)$ following the minimum cost. The paths may vary but should sum up to 46.

1. **Move Right** H if $\text{dp}[i][j+1]$ is smaller.
2. **Move Down** V if $\text{dp}[i+1][j]$ is smaller.
3. **If both are equal, try both paths (H and V).**

Possible Paths (psf values in BFS)

V V V H H H (Down-Down-Down-Right-Right-Right)

Output:-
46
HHHVVV