

## 0/1 Knapsack in C++

```
#include <iostream>
#include <vector>

using namespace std;

class ZeroOneKnapsack {
public:
    int knapsack(int n, vector<int>& vals,
vector<int>& wts, int cap) {
        vector<vector<int>>> dp(n + 1, vector<int>(cap +
1, 0));

        for (int i = 1; i <= n; i++) {
            for (int j = 1; j <= cap; j++) {
                if (j >= wts[i - 1]) {
                    int remainingCap = j - wts[i - 1];

                    if (dp[i - 1][remainingCap] + vals[i - 1] >
dp[i - 1][j]) {
                        dp[i][j] = dp[i - 1][remainingCap] +
vals[i - 1];
                    } else {
                        dp[i][j] = dp[i - 1][j];
                    }
                } else {
                    dp[i][j] = dp[i - 1][j];
                }
            }
        }

        return dp[n][cap];
    }
};

int main() {
    ZeroOneKnapsack solution;

    // Input parameters
    int n = 5;
    vector<int> vals = {15, 14, 10, 45, 30};
    vector<int> wts = {2, 5, 1, 3, 4};
    int cap = 7;

    // Compute maximum value using knapsack
    function
    int maxVal = solution.knapsack(n, vals, wts, cap);

    // Output the maximum value
    cout << "Maximum value that can be obtained: " <<
maxVal << endl;

    return 0;
}
```

### Dry Run

#### Input:

- Number of items:  $n = 5$
- Values:  $\text{vals} = \{15, 14, 10, 45, 30\}$
- Weights:  $\text{wts} = \{2, 5, 1, 3, 4\}$
- Capacity:  $\text{cap} = 7$

#### Steps:

1. **Initialize the DP Table:**
  - $\text{dp}$  is a 2D table of size  $(n+1) \times (\text{cap}+1)$  (i.e.,  $6 \times 8$ ).
  - Initially, all entries are 0.

### DP Table Construction

#### Base Case:

$\text{dp}[0][j] = 0$  for all  $j$   
 $\text{dp}[i][0] = 0$  for all  $i$

#### DP Transitions:

- **Row 1 ( $i = 1$ , item with value = 15, weight = 2):**
  - For  $j = 1$ :  $\text{dp}[1][1] = 0$  (weight exceeds capacity).
  - For  $j = 2$ :  $\text{dp}[1][2] = 15$  (item included).
  - For  $j = 3$  to  $7$ :  $\text{dp}[1][j] = 15$  (item included).
- **Row 2 ( $i = 2$ , item with value = 14, weight = 5):**
  - For  $j = 1$  to  $4$ :  $\text{dp}[2][j] = \text{dp}[1][j]$ .
  - For  $j = 5$ :  $\text{dp}[2][5] = \max(\text{dp}[1][5], \text{vals}[1] + \text{dp}[1][5 - \text{wts}[1]]) = \max(15, 14) = 15$ .
  - For  $j = 6$ :  $\text{dp}[2][6] = \max(15, 14) = 15$ .
  - For  $j = 7$ :  $\text{dp}[2][7] = \max(15, 15 + 14) = 29$ .
- **Row 3 ( $i = 3$ , item with value = 10, weight = 1):**
  - Updates based on the new item's inclusion.
- **Row 4 ( $i = 4$ , item with value = 45, weight = 3):**
  - Updates based on the new item's inclusion.
- **Row 5 ( $i = 5$ , item with value = 30, weight = 4):**
  - Updates based on the new item's inclusion.

	<p><b>Final DP Table:</b></p> <pre> dp = {     { 0, 0, 0, 0, 0, 0, 0, 0 },     { 0, 0, 15, 15, 15, 15, 15, 15 },     { 0, 0, 15, 15, 15, 15, 29, 29 },     { 0, 10, 15, 25, 25, 25, 29, 40 },     { 0, 10, 15, 45, 55, 55, 70, 70 },     { 0, 10, 15, 45, 55, 55, 70, 75 }, }</pre>
<p>Output:  Maximum value that can be obtained: 75  The maximum value that can be obtained is stored in <code>dp[5][7] = 75</code>.</p>	