

Semantic Enhancements for IPDL

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1 Omitting Arguments

The problem has been solved by adapting the way the rules were written. We explain this with the help of the substitution rule. The source of the problem is Maude's treatment of variables.

The original rule for substitutions of families is (some details omitted):

```
cr1 [subst-families-gen] :
  pConfig(Sigma, Delta,
    (family (fns2[blist2]) nlist2 blist2 ::= cases) ||
    (family (fns1[blist1]) nlist1 blist1 ::=
      nf((x : T <- read (fns2[tlist]) ) BRL, R)
    ), I, O, A
  )
=>
  pConfig(Sigma, Delta,
    (family (fns2[blist2]) nlist2 blist2 ::= cases) ||
    (family (fns1[blist1]) nlist1 blist1 ::=
      preNF((x : T <- R2) BRL, R)
    ), I, O, A
  )
if
  (project2Index
    (family (fns2[blist2]) nlist2 blist2 ::= cases)
    nlist2 tlist A empty)
  ==
  (fns2[tlist] ::= R2)
[nonexec]
.
```

and since R2 does not occur in the term on the left of the equality sign nor in the configuration that we are rewriting, the rule must be marked as nonexecutable and R2 must be explicitly provided when the rule is applied.

The strategy

$$\text{substNFFamiliesGen}((\text{fam}(\text{ns1}[\text{bdlist1}]), \text{fam}(\text{ns2}[\text{bdlist2}])), R)$$

that calls this rule has thus three arguments, so we can apply it in the strategy as

```
subst-families-gen [
  fns2 : NameWithScripts <- ns1 ,
  fns1 : NameWithScripts <- ns2 ,
  R2 : Reaction <- R
]
```

Instead of using a comparison, we have written a helper **getReaction** that gives us the reaction assigned to **fns2[tlist]** when doing the projection, and we can use this expression instead of R2 in the resulting pConfig. The rule is no longer nonexecutable and we can remove the third argument of **substNFFamiliesGen**.

With this change, all substitution strategies have two arguments and we can write a meta-strategy that handles all possible combinations:

```
strat subst : CNameBound CNameBound @ ProtocolConfig .
sd subst(chn C1, chn C2) :=
  substNFRead(C1, C2)
  or-else substNF(C1, C2)
.
sd subst(fam (fns1 [bdlist1]), fam (fns2 [bdlist2])) :=
  substNFFamiliesGen(
    fam (fns1 [bdlist1]), fam (fns2 [bdlist2])
  )
.
sd subst(fam (fns1 [bdlist1]), chn C2) :=
  substNFFamilyOne(fam (fns1 [bdlist1]), C2)
.
sd subst(chn C1, fam (fns2 [bdlist2])) :=
  substChannelFamilyOne(chn C1, fam (fns2 [bdlist2]))
.
```

which was our original goal.

We had a similar issue with the type of an expression/reaction. Just like above, it can be computed and thus eliminated as an argument of the strategy.

2 IPDL-specific Error Handling

We have added type checking steps in the analysis of protocols in the SpeX implementation of IPDL. We have attempted to do this in the original implemen-

tation, but Maude provides limited capabilities for printing values of any type. SpeX provides an appropriate solution to this, and we can simply reuse it. We have refactored the analysis of protocol declarations by adding a pre-processing step that removes the token constructors (SpeX-specific implementation detail), followed by a type checking step that reuses the type checking already implemented and throws an error if some condition fails to hold (this code is again written at a SpeX-specific level). In Fig. 1 we see how IPDL displays various error messages. Notice that now the types can be displayed instead of getting just a type assignment error, and we get a pointer to where the error happened (provided by SpeX).

Figure 1: Type-checking errors.

```

o2
P1
Warning: testErrors.ipdl, line 12: type mismatch for assignment : channel has type bool and reaction has type unit
|
| protocol E1 = o ::= return () . <--- here
|
P2
Warning: testErrors.ipdl, line 15: type mismatch for assignment : channel has type bool and reaction has type unit
|
| protocol E2 = o ::= read j <--- here
|
P3
P4
Warning: testErrors.ipdl, line 20: expecting an expression of type bool and got unit
|
| protocol E3 = o ::= if () then return True else return False . <--- here
|
Warning: testErrors.ipdl, line 21: reaction types do not match first type is bool second type is unit
|
| protocol E4 = o ::= if True then return True else return () . <--- here
|
P5
Warning: testErrors.ipdl, line 24: type mismatch for reaction expecting bool and got unit
|
| protocol E5 = o ::= x : bool <- return () ; return True . <--- here
|
Warning: testErrors.ipdl, line 25: already declared
|
| protocol E6 = new o <--- here
|
IPDL >

```

We have also added some support for the case that a strategy does not apply to the current protocol. Maude only answers this with "No solution.". While the problem of why a rule cannot be applied to a protocol is too complex to be solved in general, we can at least provide the left-hand side of the rule that we are trying to apply and the protocol fragment that we are trying to apply it to, and the user can compare the two. In Fig. 2 we see the result of failing to apply a strategy to a protocol.

Figure 2: Rule cannot be applied.



```
mcodescu@precision: ~/repos/SpexIonut/Spex/src
mcodescu@precision: ~/repos/IPDL-Maude/doc
mcodescu@precision: ~/repos/SpexIonut/Spex/src

IPDL > fold o1 into o2
the rule cannot be applied
to
o1 ::= nf(nil, return True)
||
o2 ::= nf(nil, return False)
the LHS of the rule is
(c ::= R) || (o ::= nf((x : T <- read c) ..., S))
IPDL >
```