

Meeting : 30 May 2019

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5/30/2019

Criteria

Tuning

Fast Algorithm (2018)

One-at-a-Time Algorithm

Criteria

Criteria

1. The total potential energy, which both algorithms aim to minimize:

$$\sum_{i \neq j} \frac{q(\mathbf{x}_i)q(\mathbf{x}_j)}{d(\mathbf{x}_i, \mathbf{x}_j)}$$

2. One-at-a-Time Algorithm criterion tries to minimize:

$$\left\{ \sum_{i \neq j} \left(\frac{q(\mathbf{x}_i)q(\mathbf{x}_j)}{d(\mathbf{x}_i, \mathbf{x}_j)} \right)^k \right\}^{1/k}$$

3. Fast Algorithm tries to minimize:

$$\max_{i \neq j} \frac{q(\mathbf{x}_i)q(\mathbf{x}_j)}{d(\mathbf{x}_i, \mathbf{x}_j)}$$

Tuning

Different Powers k and Stages K

- ▶ What happens when k gets large in the one-at-a-time algorithm's criterion:

$$\left\{ \sum_{i \neq j} \left(\frac{q(\mathbf{x}_i)q(\mathbf{x}_j)}{d(\mathbf{x}_i, \mathbf{x}_j)} \right)^k \right\}^{1/k}$$

- ▶ the criterion itself is changed
- ▶ For the fast algorithm, for $k = 1, \dots, K - 1$ steps, we obtain each design point $\mathbf{x}_j^{(k+1)}$ of the next stage \mathbf{D}_{k+1} by:

$$\mathbf{x}_j^{k+1} = \arg \min_{\mathbf{x} \in \mathbf{C}_j^{k+1}} \max_{i=1:(j-1)} \frac{q^{\gamma_k}(\mathbf{x}_i)q^{\gamma_k}(\mathbf{x})}{d(\mathbf{x}_i, \mathbf{x})}$$

What happens when we include more stages (larger K)?

- ▶ this isn't about a change in the criterion, but a change in the method for minimizing that criterion

In particular, we're interested in:

- ▶ the asymptotic behavior of the algorithms
- ▶ the limiting distributions of the design
- ▶ whether it improves the design's ability to minimize the criteria (its own criterion as well as the total potential energy criterion)

Parameters

► Parameters

```
mean_beta0 = 1 # slope of null model
mean_beta1 = 1 / 2 # slope of alternative model
var_mean = 0.001 # variance on beta
var_e = 0.01 # variance on error
```

► Settings for Fast and One-at-a-Time algorithms

```
N = 67
# for fast algorithm:
K = 20 # ceiling(4* sqrt(p))
numParameters = 1 # number of parameters (just slope!)
p = numParameters
# for one-at-a-time algorithm:
numCandidates = 105 # suggested 105
k = 4
```


Fast Algorithm (2018)

Fast Algorithm (2018)

In each of K stages, create a new design to iteratively minimize

$$\max_{i \neq j} \frac{q(\mathbf{x}_i)q(\mathbf{x}_j)}{d(\mathbf{x}_i, \mathbf{x}_j)}$$

1. Initialize space-filling design $\mathbf{D}_1 = \{\mathbf{x}_1^{(1)} \dots \mathbf{x}_N^{(1)}\}$
2. For $k = 1, \dots, K - 1$ steps, obtain each design point $\mathbf{x}_j^{(k+1)}$ of the next stage \mathbf{D}_{k+1} by:

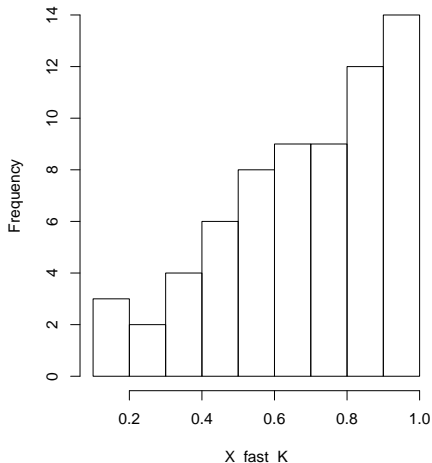
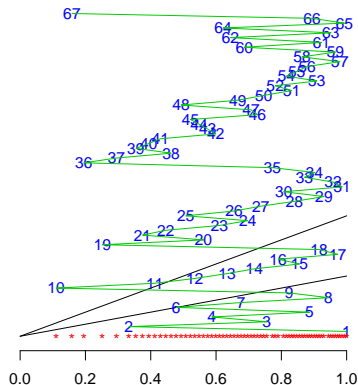
$$\begin{aligned}\mathbf{x}_j^{k+1} &= \arg \min_{\mathbf{x} \in \mathbf{C}_j^{k+1}} \max_{i=1:(j-1)} \frac{1}{f^{\gamma_k}(\mathbf{x}_i) f^{\gamma_k}(\mathbf{x}) d^{(2p)}(\mathbf{x}_i, \mathbf{x})} \\ &= \arg \min_{\mathbf{x} \in \mathbf{C}_j^{k+1}} \max_{i=1:(j-1)} \frac{q^{\gamma_k}(\mathbf{x}_i) q^{\gamma_k}(\mathbf{x})}{d(\mathbf{x}_i, \mathbf{x})}\end{aligned}$$

where $\gamma_k = k/(K - 1)$ and \mathbf{C}_j^{k+1} is the candidate set for design point \mathbf{x}_j at stage $k + 1$.

- ▶ points are no longer picked sequentially
- ▶ candidates are different for each design point

Design generated by Fast Algorithm ($K = 20$)

Histogram of X_{fast_K}

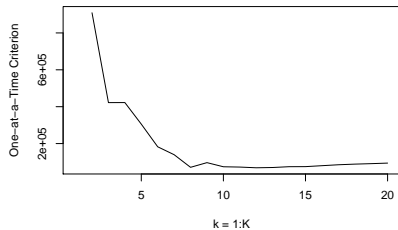
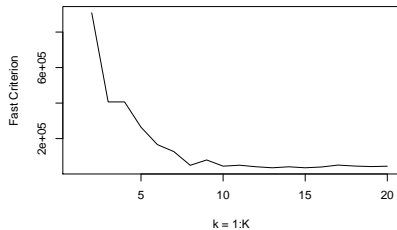
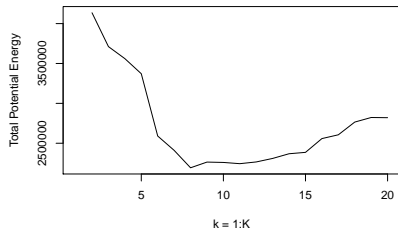
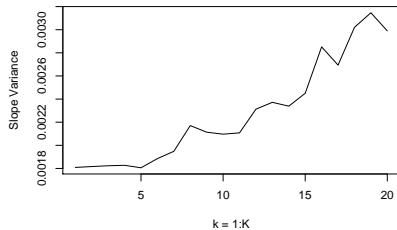


```
mean(X_fast_K); sd(X_fast_K)
```

```
## [1] 0.6836593
```

```
## [1] 0.2296987
```

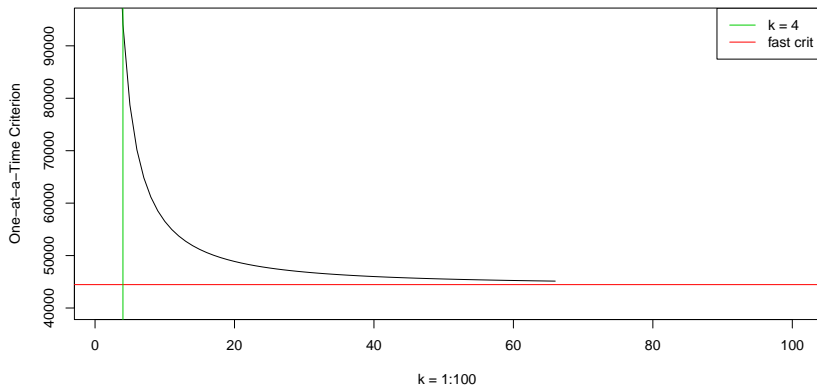
Over the K Designs



- ▶ Total Potential Energy criterion starts to increase as the design approaches $K = 20$. It could be that $K = 20$ is too large
- ▶ Variance on Slope ($\hat{\beta}$) also increases as the stages move away from a space-filling design and toward one that minimizes the fast algorithm's criterion. Why?

One-at-a-Time Criterion, Different k Powers

- ▶ To see if it decreases as $k \rightarrow \infty$, since this algorithm is supposed to be the asymptotic result of the one-at-a-time algorithm, here are the results for $k = 1:100$.
- ▶ For the fast algorithm's design, the one-at-a-time algorithm's criterion approaches the fast algorithm's criterion as k gets large (until it becomes Inf - likely due to a computational issue.)



One-at-a-Time Algorithm

One-at-a-Time Algorithm (2015)

1. Obtain *numCandidates* candidate points, \mathbf{x} , in $[0, 1]$.
2. Initialize D_N by choosing \mathbf{x}_j to be the candidate \mathbf{x} which optimizes f , where $f(\mathbf{x}) = \text{Wasserstein}(\phi_{0,\mathbf{x}}, \phi_{1,\mathbf{x}})$ and

$$\begin{aligned}\phi_{0,\mathbf{x}} &= N(\tilde{\beta}_0 \mathbf{x}, \sigma_{\epsilon_0}^2 + \mathbf{x}^2 \sigma_{\beta_0}^2), \\ \phi_{1,\mathbf{x}} &= N(\tilde{\beta}_1 \mathbf{x}, \sigma_{\epsilon_1}^2 + \mathbf{x}^2 \sigma_{\beta_1}^2)\end{aligned}$$

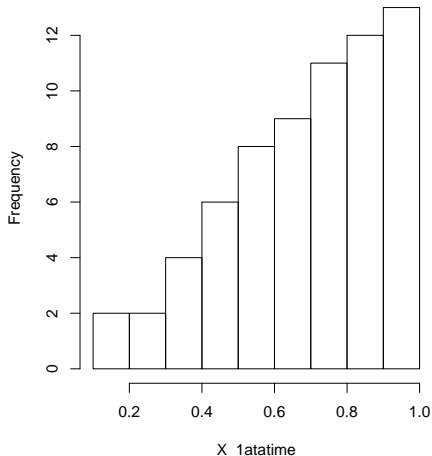
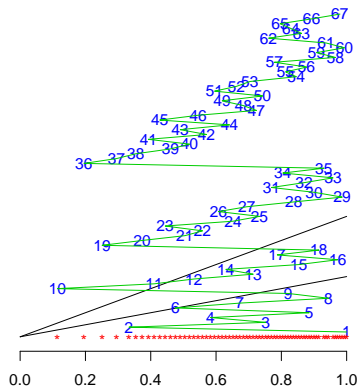
3. Choose the next point \mathbf{x}_{j+1} by:

$$\mathbf{x}_{j+1} = \arg \min_{\mathbf{x}} \sum_{i=1}^j \left(\frac{q(\mathbf{x}_i)q(\mathbf{x})}{d(\mathbf{x}_i, \mathbf{x})} \right)^k$$

where $q = 1/f^{(1/2p)}$, $d(x, y)$ is Euclidean distance and (suggested) $k = 4p$.

Design generated by One-at-a-Time Algorithm ($k = 4$)

Histogram of X_1atatime



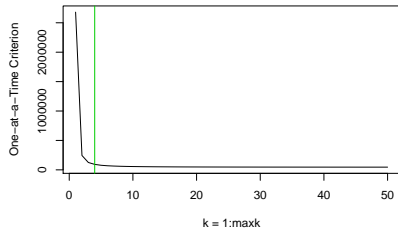
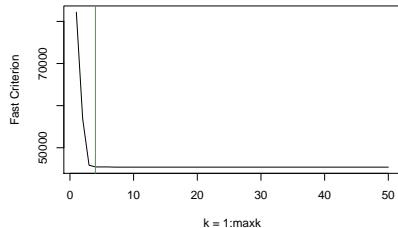
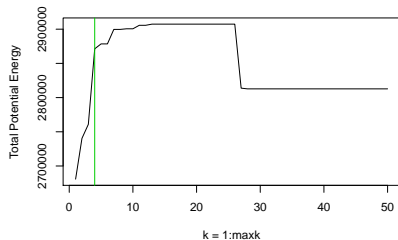
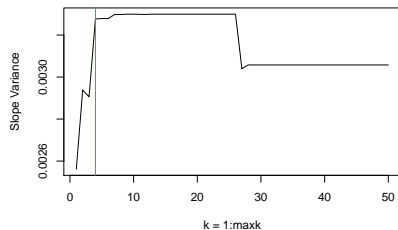
```
mean(X_1atatime)
```

```
## [1] 0.6886419
```

```
sd(X_1atatime)
```

```
## [1] 0.2187448
```

Robustness Across k Power



Comparing Evaluations for Criterion

- ▶ Each One-at-a-Time design from $k = 1:50$, evaluated at its corresponding one-at-a-time criterion.
- ▶ They appear to be approaching the $k = 4$ design's evaluation of the fast algorithm's criterion.
- ▶ Also recall that at $k = 4$ and after, the designs' evaluations of the fast algorithm's criterion are roughly the same.

