

Distributed Systems

CS 380D

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Types of knowledge

- Common knowledge:
 - known by everyone in group
 - each node **can** assume others know this
- Distributed knowledge:
 - known by some members of group
 - a node **cannot** assume others know this
- Simultaneous actions requires common knowledge

Common Knowledge

- **Impossible** to obtain if communication is over unreliable channels
- Demonstrated in the Coordinated Attack Problem
- Internal Common Knowledge:
 - assume something is common knowledge
 - hope no node encounters state that disproves assumption

Muddy Children Puzzle

- n children play, k get muddy
- Each can observe all others, don't know their own state
- Dad says "at least one of you is muddy"
- Dad asks each of them: "do you know if you are muddy"?
- claim: After $k-1$ rounds, all children will answer yes

Muddy Children Puzzle

- Children get information from:
 - Observation of other children
 - Hearing what other children say
 - Inferences based on previous rounds
- Common knowledge: father says at start "at least one of you is muddy"

Proof by induction

- $k = 1$:

- Muddy child observes all others are clean
- But father said someone is muddy
- Hence child realizes they are muddy, answers yes
- Once other children hear muddy child answer yes, they also answer yes

- $k = 2$:

- Each muddy child observes one other muddy child
- in first round, $k = 1$, all answer no as they are unsure of their own state
- Muddy child realizes they are muddy, since other muddy child answered no in first round (hence other child must see someone muddy)
- In second round, all answer "yes"

Proof by induction

- $k = 3$
 - Say muddy children are a, b, c
 - if a is clean, b and c would have answered yes in second round
 - Hence a is not clean; b and c do similar reasoning
 - All answer yes on third round

Does father need to provide common knowledge?

- One might think no: for $k > 1$, seems like children get the information from direct observation
- However, it is not common knowledge
- For $k = 2$, muddy child a observes muddy child b. But **does not know** if b observes a, and therefore knows $k \geq 1$

Does father need to provide common knowledge?

- Showing it does not work for $k = 2$:
 - Muddy children are A and B
 - In first round, even if A had seen all clean kids, they would have still answered "no" (because they do not know $k \geq 1$)
 - In second round, A and B realizing they are muddy depends on muddy child saying yes in round 1
 - A saying "no" in round 1 does not provide B with any information
 - B still thinks $k = 1$ or $k = 2$

Does father need to provide common knowledge?

- Valid sequence if $k = 1$ from B's viewpoint:
 - A is only muddy child
 - A does not realize $k \geq 1$, cannot decide between $k = 0$ and $k = 1$
 - A says "no" in first round
- B still cannot decide between $k = 1$ or $k = 2$ (both can happen with prior seq)

Common knowledge

- $k \geq 1$ is distributed knowledge, not common knowledge
- This case clearly shows the difference between the two

Hierarchy of States of Knowledge

- Agent's knowledge depends on:
 - Starting knowledge
 - Observed history since start
- If agent i knows P then $K_i(P)$
- Agents know only true things

Hierarchy of States of Knowledge

- $D(G, P)$ = group G has distributed knowledge of P (union of knowledge of G members = P)
- $S(G, P)$ = someone in G knows P
- $E(G, P)$ = everyone in G knows P
- $E(G, K, P) = E(E(E.. E(G,P))))$ k times
- $E(E(G,P))$ = everyone in G knows that everyone in G knows P
- Common knowledge: $E(G, K, P)$ for all $K \geq 1$

Muddy Children Puzzle

- m = "at least one child is muddy"
- Without father speaking,
 - $E(G, K-1, m) = \text{true}$
 - $E(G, K+1, m) = \text{false}$

Muddy Children Puzzle

- $E(G, m) = \text{true}$
- In $k = 2$,
 - $E(G, 1, m) = \text{true}$, everyone knows m
 - $E(G, 2, m) = \text{false}$, everyone does not know everyone knows m
 - Specifically, with muddy children A and B, A does not know B knows $K \geq 1$
- Father's statement makes $E(G, 2, m) = \text{true}$

Knowledge in distributed systems

- Communication in a distributed systems seeks to move up the hierarchy of knowledge:
 - changing $S(G, P) = E(G, P) = C(G, P)$
- Fact discovery:
 - Changing D to S to E to C
 - Example: finding deadlock in a set of distributed locks
- Fact publication:
 - Changing S to C
 - Example: new protocol for communication

Common Knowledge

- How does one establish it?
 - By being part of a community
 - Membership procedure imparts common knowledge
 - Example: community of licensed drivers knows what signs mean
 - By being co-present at knowledge creation
 - Example: children being in same room as father when he makes announcement

Coordinated Attack Problem

- General A sends time in message to General B
- A will not attack without ack from B
- B sends ack to A
- But B will not attack without ack from A
- A sends $\text{Ack}(\text{Ack}(A))$ to B
- A will not attack without ack of this message
- And so it goes.. A and B cannot agree with finite messages
- Can use induction to prove no set of K messages is enough

Coordinated Attack Problem

- Generals A and B need common knowledge of the attack time
- After A sends the first message to B, $E(G) = \text{true}$, but $E(E(G))$ is not
- After A sends the ack to B's ack, $E(E(G)) = \text{true}$, but $E(3, G) = \text{false}$
- Coordinated attack requires $C(G) = E(k, G) = \text{true}$ for any k