

## PART A

1. Prove:

$$(\text{append} (\text{append } xs \ ys) \ zs) = (\text{append } xs \ (\text{append } ys \ zs))$$

### PART A : BASE CASE $\{ xs \rightarrow '() \}$

Assume  $xs = '()$

$$(\text{append} (\text{append } xs \ ys) \ zs)$$

$$= \{ \text{by assumption, } xs = '() \}$$

$$(\text{append} (\text{append } '() \ ys) \ zs)$$

$$= \{ \text{by append-nil law} \}$$

$$(\text{append } ys \ zs)$$

$$= \{ \text{by append-nil law from right to left} \}$$

$$(\text{append } '() \ (\text{append } ys \ zs))$$

$$= \{ \text{by assumption } xs = '() \}$$

$$(\text{append } xs \ (\text{append } ys \ zs))$$

### PART B : INDUCTIVE HYPOTHESIS

Given  $xs = (\text{cons } a \ as)$

Then inductive hypothesis:

$$(\text{append} (\text{append } as \ ys) \ zs) = (\text{append } as \ (\text{append } ys \ zs))$$

### PART C: INDUCTIVE CASE ( $xs \mapsto (\text{cons } a \ as)$ )

$$\begin{aligned} & (\text{append} (\text{append } xs \ ys) \ zs) \\ &= \{ \text{by assumption } xs = (\text{cons } a \ as) \} \\ & (\text{append} (\text{append} (\text{cons } a \ as) \ ys) \ zs) \\ &= \{ \text{by append-cons law} \} \\ & (\text{append} (\text{cons } a \ (\text{append } as \ ys)) \ zs) \\ &= \{ \text{by append-cons law} \} \\ & (\text{cons } a \ (\text{append} (\text{append } as \ ys) \ zs)) \\ &= \{ \text{by inductive-hypothesis} \} \\ & (\text{cons } a \ (\text{append } as \ (\text{append } ys \ zs))) \\ &= \{ \text{by append-cons law from right to left} \} \\ & (\text{append} (\text{cons } a \ as) \ (\text{append } ys \ zs)) \\ &= \{ \text{by assumption, } xs = (\text{cons } a \ as) \} \\ & (\text{append } xs \ (\text{append } ys \ zs)) \end{aligned}$$

### PART D: CONCLUSION

Thus from the above inductive proof, we can conclude that appending lists is an associative operation,

ie

$$(\text{append} (\text{append } xs \ ys) \ zs) = (\text{append } xs \ (\text{append } ys \ zs))$$