COT 6405 ANALYSIS OF ALGORITHMS

Traveling Salesman Problem Branch-and-Bound Algorithm

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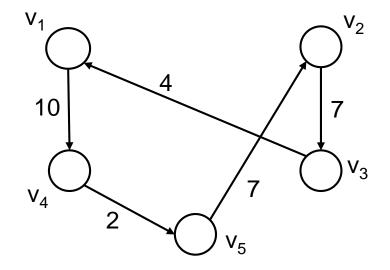
Traveling Salesman Problem (TSP)

Given *n* cities with known distances between each pair, find the shortest tour that passes through all the cities exactly once before returning to the starting city

Example:

Adjacency matrix:

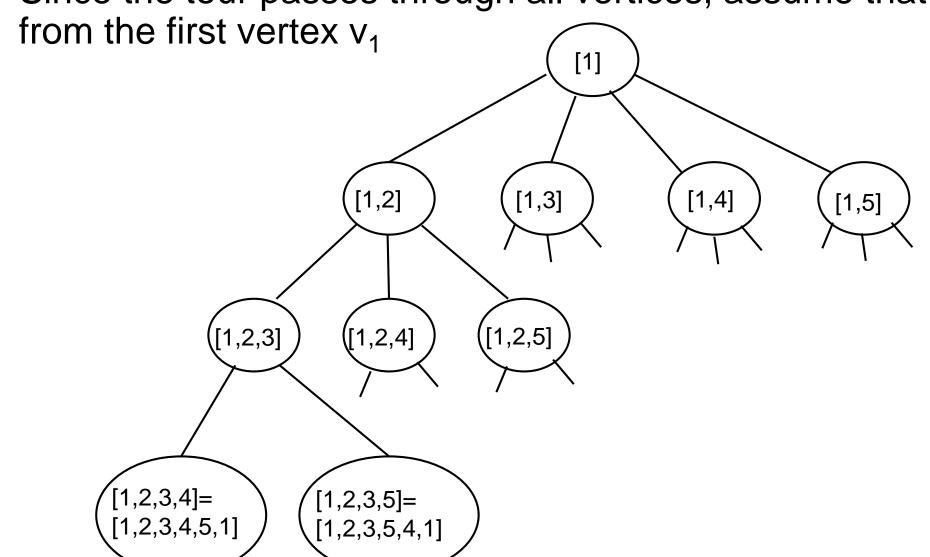
$$\begin{bmatrix} 0 & 14 & 4 & 10 & 20 \\ 14 & 0 & 7 & 8 & 7 \\ 4 & 5 & 0 & 7 & 16 \\ 11 & 7 & 9 & 0 & 2 \\ 18 & 7 & 17 & 4 & 0 \end{bmatrix}$$



min tour length = 30

TSP – search tree example

Since the tour passes through all vertices, assume that it starts



TSP with BestFS with Branch-and-Bound pruning

- compute a bound for each node
- lower bound on the length of any tour that can be obtained by expanding beyond a given node
- any tour must leave each vertex exactly once, then a lowerbound is to take minimum edge leaving every vertex:

```
v_1 minimum(14,4,10,20) = 4

v_2 minimum(14,7,8,7) = 7

v_3 minimum(4,5,7,16) = 4

v_4 minimum(11,7,9,2) = 2

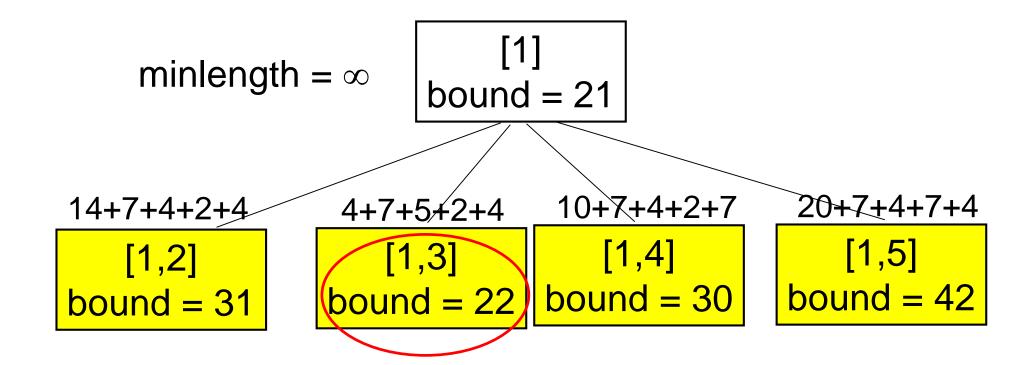
v_5 minimum(18,7,17,4) = 4
```

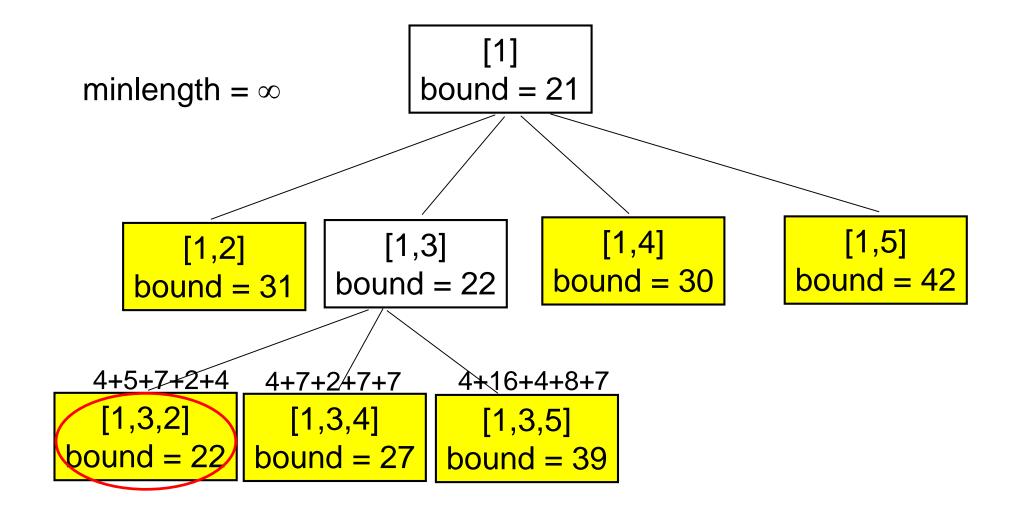
- lower-bound on the length of a tour is 4+7+4+2+4=21
- observation: this does not mean there is a tour with this length; it means there is no tour with a shorter length

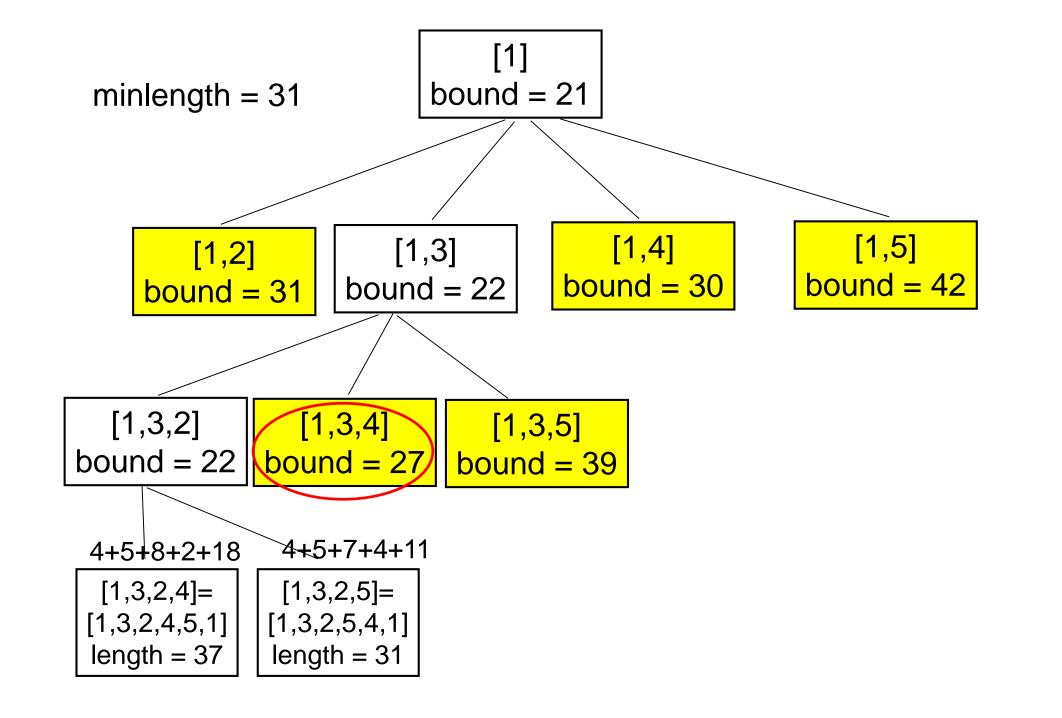
TSP Example

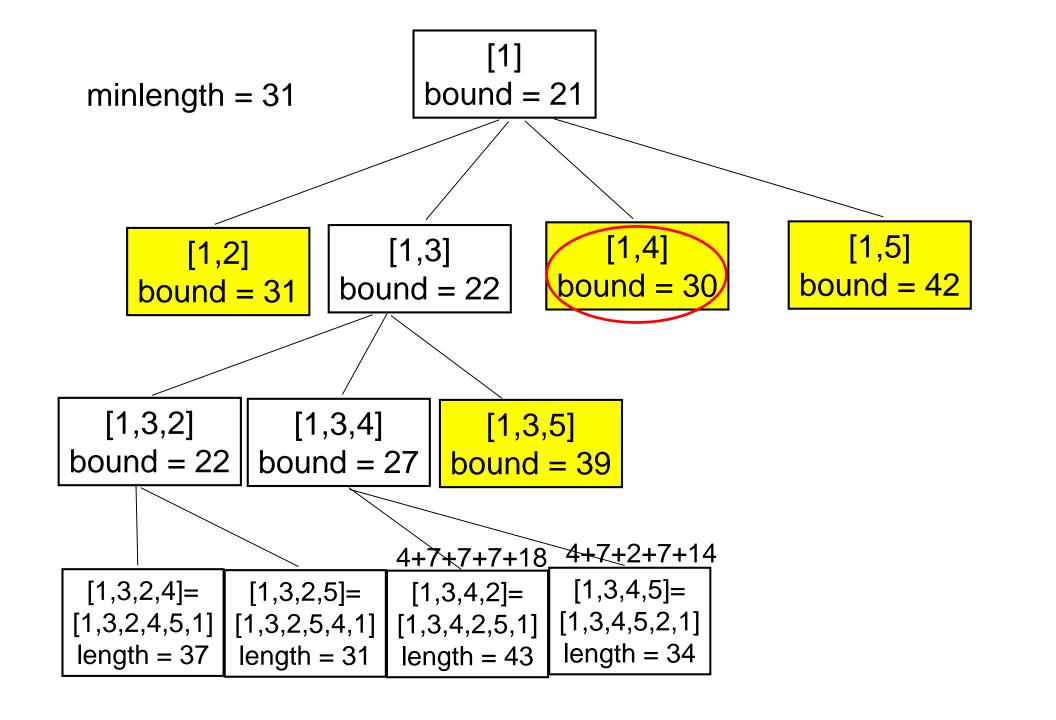
remove from PQ

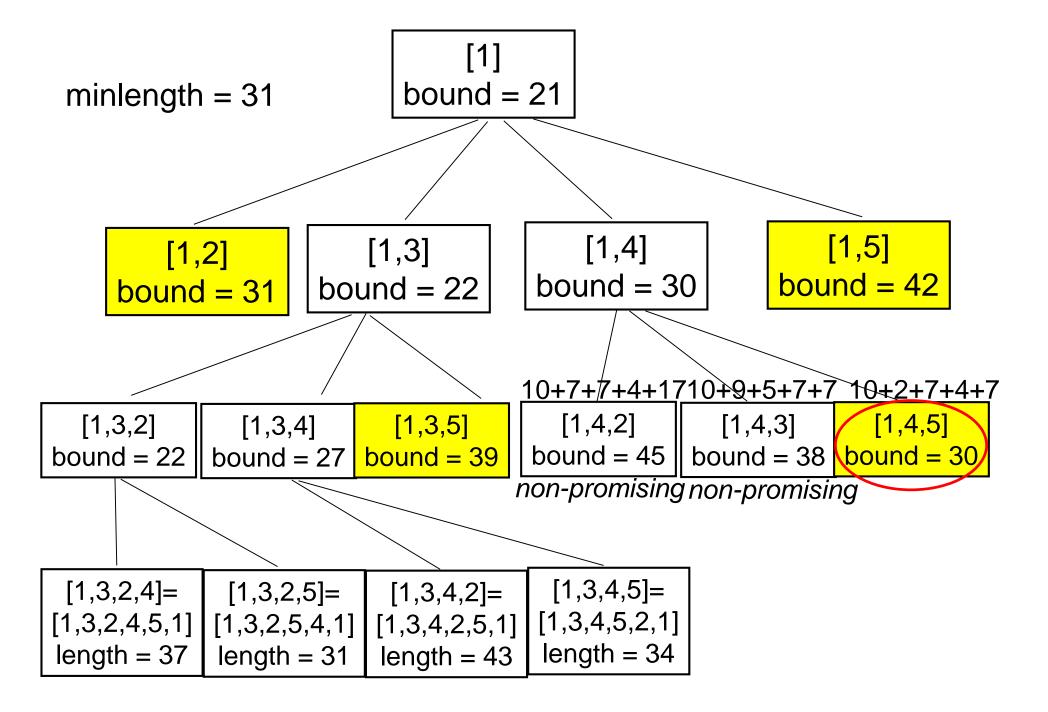
minlength =
$$\infty$$
 [1] bound = 21

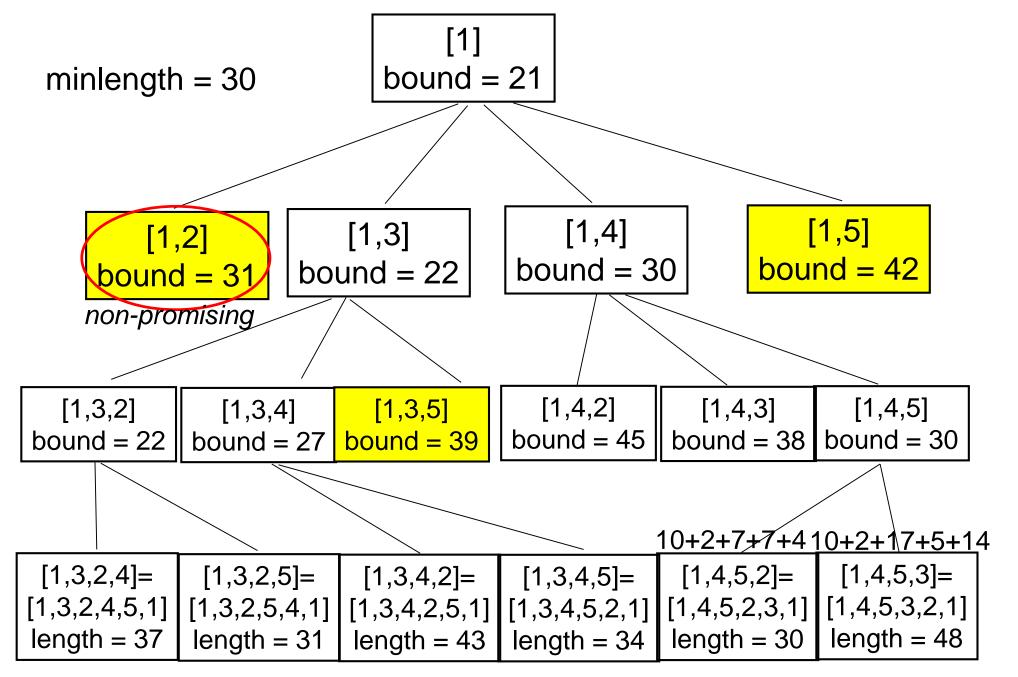


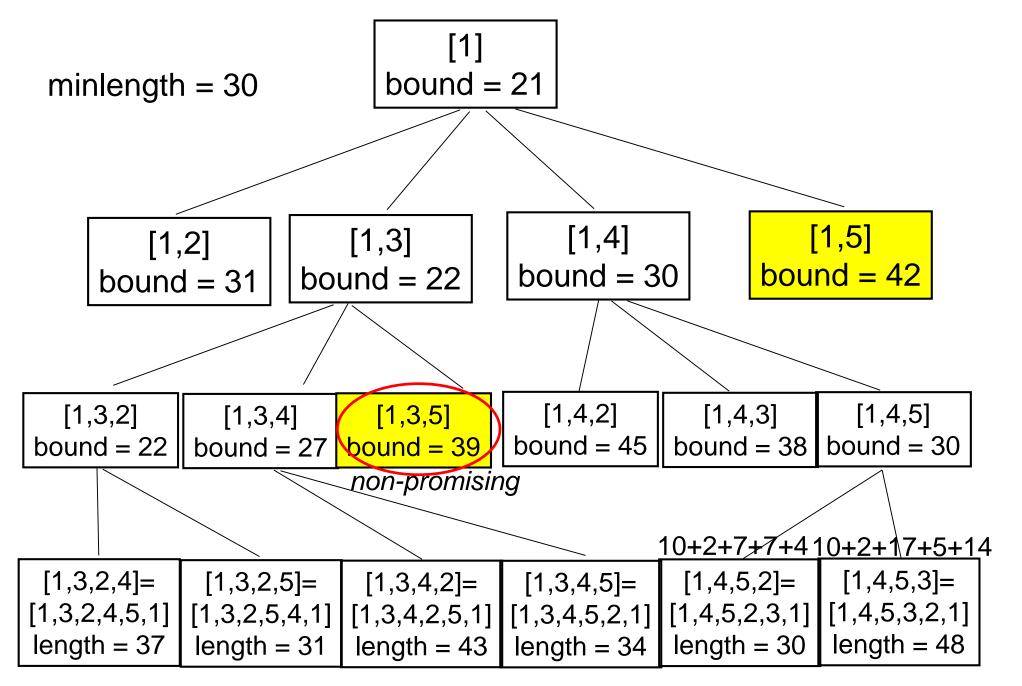


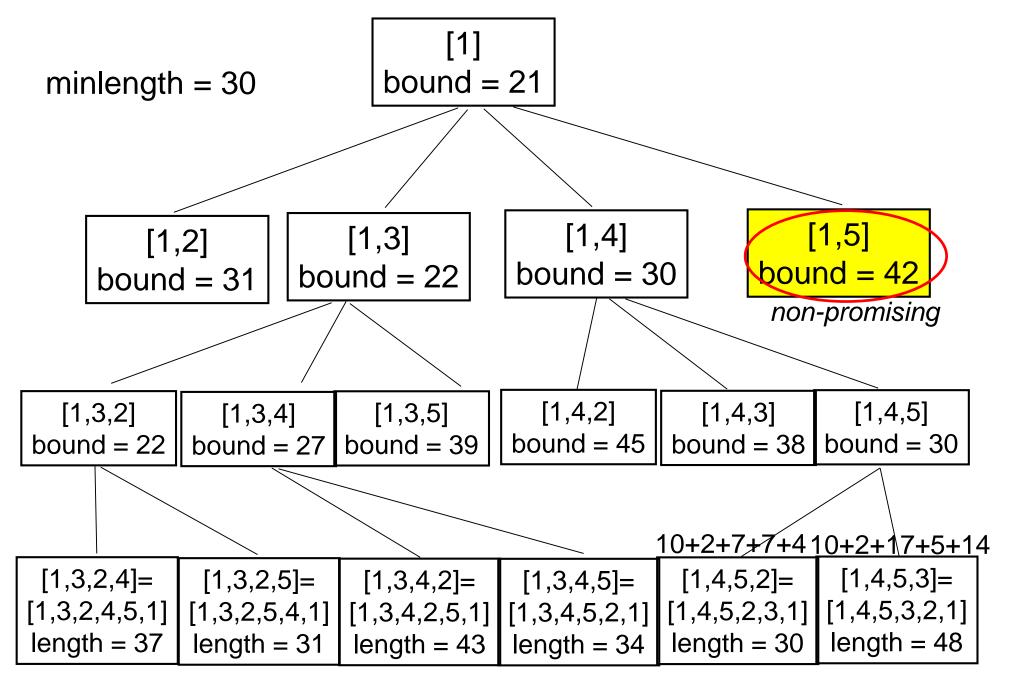


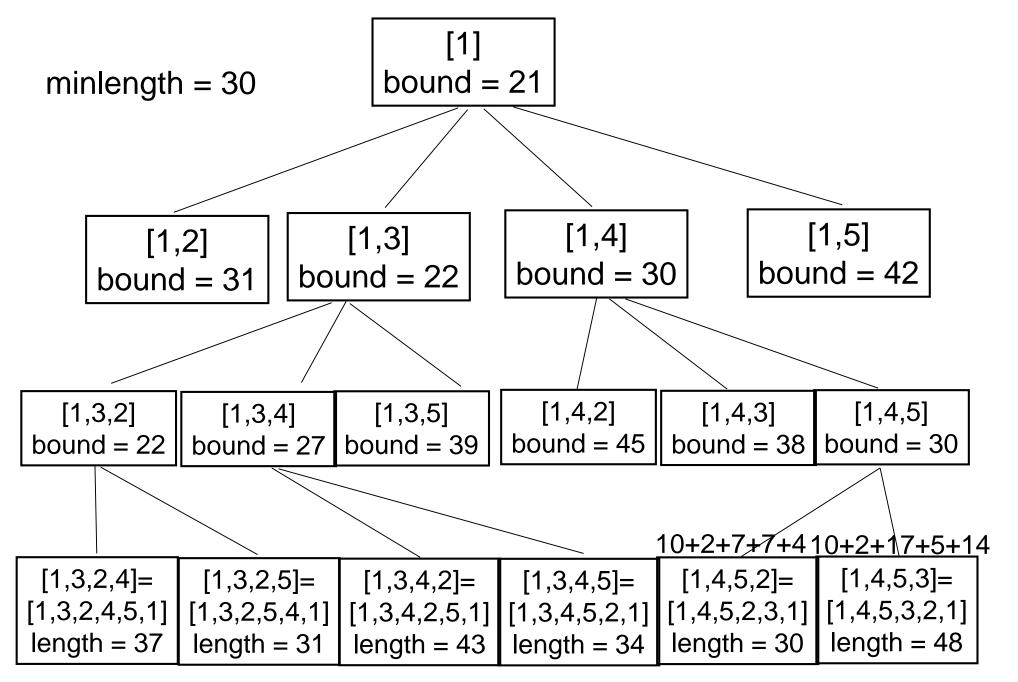












TSP with BestFS with Branch-and-bound pruning

- Each node is an object with fields:
 - v.level node's level in the tree
 - v.path path of the partial tour up to node v
 - v.bound lower bound on the length of any tour expanding beyond this partial tour

TSP with BestFS w/ Branch-and-bound pruning(n,W[][],opttour, minlength)

```
PQ = \emptyset
r.level = 0
r.path = [1]
r.bound = bound(r)
minlength = \infty
insert(PQ,r)
while PQ \neq \emptyset
   v = remove(PQ)
   if v.bound < minlength</pre>
              u.level = v.level+1
              for all i such that 2 \le i \le n and i is not in v.path
                           u.path = v.path
                           add i at the end of u.path
                           if u.level == n-2 // check if next vertex completes the tour
                                         put index of only vertex not in u.path at the end of u.path
                                         put 1 at the end of u.path
                                         if length(u) < minlength</pre>
                                                       minlength = length(u)
                                                       opttour = u.path
                           else
                                         u.bound = bound(u)
                                         if u.bound < minlength</pre>
                                                       insert(PQ,u)
```