Circulations with demands

Input

-directed graph G(V, E)

-each vertex v has
a demand dv

- capacity c of each edge

RT=0(E.D)

Flow-network

(standard form)

-directed graph G'(V', E')

• add super-source s* and

super-sink t*

• Keep the edges in G and

their capacities

• add edges from s* to each

source node v with capacity dy

• add edges from each sink

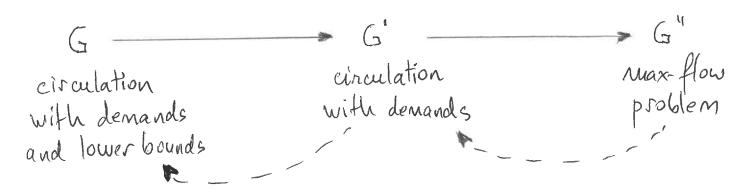
node v to t* with capacity dv

RT analysis

Ford-Fulkerson => RT=0(E'.|f*|) $|f^*| \leq D$ $|E'| \leq |E| + |V|$ $= |F| \leq |E| + |V| = |E| \approx |V| - 1 \Rightarrow |V| = 0(E)$ = |F'| = 0(E)

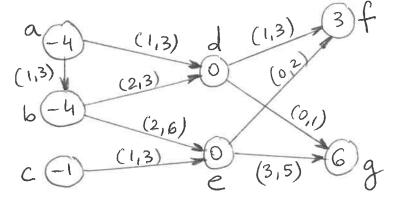
Consider the graph G(V, E) with the following capacity and demand constraints: Find a feasible circulation. Solution a, 6 - source nodes e, d- sink nodes construct grouph G compute max-flow of G' P=<5*, a, e, d, t*> Cp(p)=2 P= <s*, b, e, c, t*> Cf (p)=1 P = < 5*, a,c, t*> Cf(p)=L If = 4. Since If = D => G has a feasible circulation:

Circulations with demands and lower bounds



Example
Consider the graph G(V, E) below, with the following constraints
for demands, capacity and lower bounds:

graph G



Find a feasible circulation.

Solution

graphs

a = 4 (1,3) d (1,3) 3

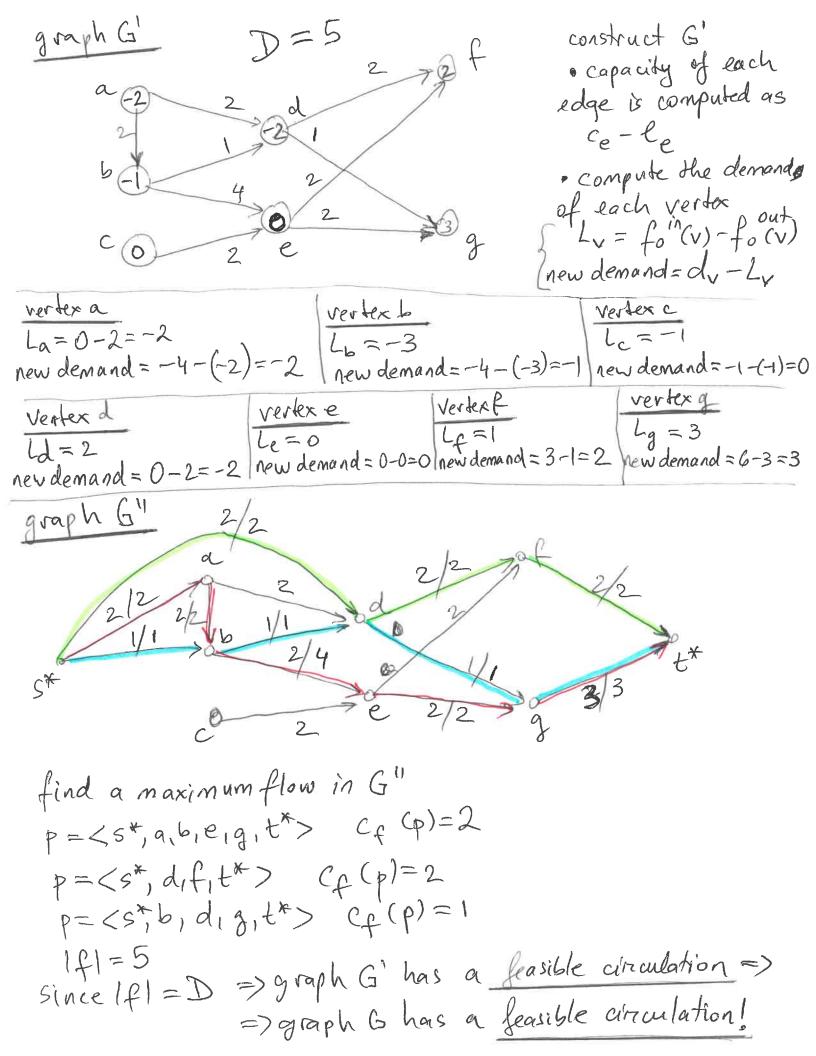
V(1,3) 2 (2,3) 0 0 (0,3)

b -4 2 (2,6) 0 (0,1)

c -1 1 (1,3) 0 3 (3,5) 6

flow fo

estarted with a circulation fo(e)= le



flow f graph G = circulation for graph G