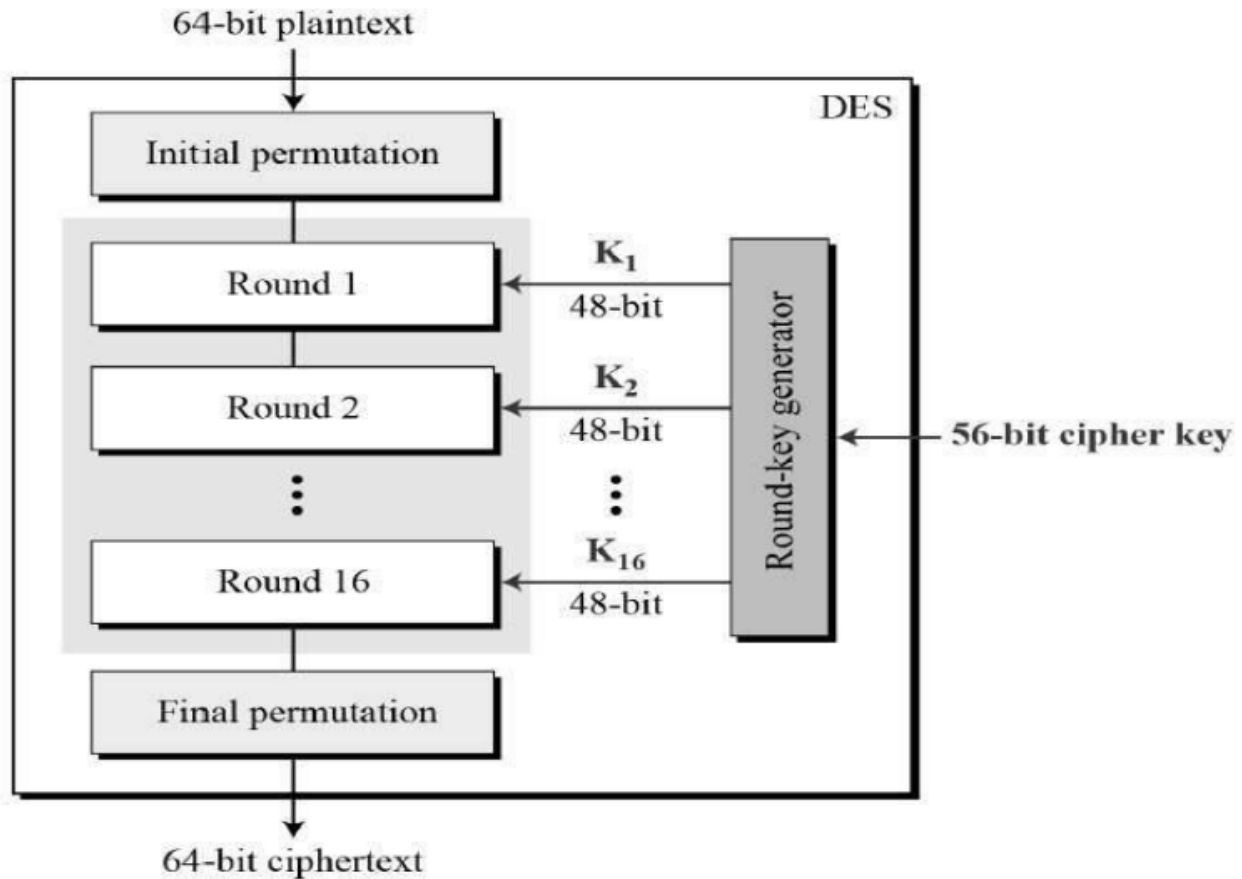


## ASSIGNMENT 6

Write a program to implement Data Encryption Standard for encryption and decryption.



### 1. Process the key.

i. Get a 64-bit key from the user.

ii. Calculate the key schedule.

1. Perform the following permutation on the 64-bit key. The parity bits are discarded, reducing the key to 56 bits. Bit 1 of the permuted block is bit 57 of the original key, bit 2 is bit 49, and so on with bit 56 being bit 4 of the original key.
2. Split the permuted key into two halves. The first 28 bits are called C[0] and the last 28 bits are called D[0].

3. Calculate the 16 subkeys. Start with  $i = 1$ .

1. Perform one or two circular left shifts on both C[i-1] and D[i-1] to get C[i] and D[i], respectively. The number of shifts per iteration are given in the table below.

Iteration #	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
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Left Shifts	1	1	2	2	2	2	2	2	1	2	2	2	2	2	2	1
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2. Permute the concatenation C[i]D[i] as indicated below. This will yield K[i], which is 48 bits long.

3. Loop back to 1.ii.c.1 until K[16] has been calculated.

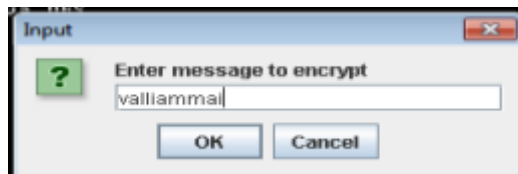
### 2. Process a 64-bit data block.

i. Get a 64-bit data block. If the block is shorter than 64 bits, it should be padded as appropriate for the application.

ii. Perform the initial permutation on the data block.

- iii. Split the block into two halves. The first 32 bits are called  $L[0]$ , and the last 32 bits are called  $R[0]$ .
- iv. Apply the 16 subkeys to the data block. Start with  $i = 1$ .
  - a. Expand the 32-bit  $R[i-1]$  into 48 bits according to the bit-selection function Expansion (E)
  - b. Exclusive-or  $E(R[i-1])$  with  $K[i]$ .
  - c. Break  $E(R[i-1]) \text{ xor } K[i]$  into eight 6-bit blocks. Bits 1-6 are  $B[1]$ , bits 7-12 are  $B[2]$ , and so on with bits 43-48 being  $B[8]$ .
  - d. Substitute the values found in the S-boxes for all  $B[j]$ . Start with  $j = 1$ . All values in the S-boxes should be considered 4 bits wide.
    - i. Take the 1st and 6th bits of  $B[j]$  together as a 2-bit value (call it  $m$ ) indicating the row in  $S[j]$  to look in for the substitution.
    - ii. Take the 2nd through 5th bits of  $B[j]$  together as a 4-bit value (call it  $n$ ) indicating the column in  $S[j]$  to find the substitution.
    - iii. Replace  $B[j]$  with  $S[j][m][n]$ .
    - iv. Loop back to 2.iv.d.i until all 8 blocks have been replaced.
  - e. Permute the concatenation of  $B[1]$  through  $B[8]$
  - f. Exclusive-or the resulting value with  $L[i-1]$ . Thus, all together, your  $R[i] = L[i-1] \text{ xor } P(S[1](B[1])...S[8](B[8]))$ , where  $B[j]$  is a 6-bit block of  $E(R[i-1]) \text{ xor } K[i]$ . (The function for  $R[i]$  is written as,  $R[i] = L[i-1] \text{ xor } f(R[i-1], K[i])$ .)
  - g.  $L[i] = R[i-1]$ .
  - h. Loop back to 2.iv.a until  $K[16]$  has been applied.
  - v. Perform the final permutation on the block  $R[16]L[16]$ .
3. Decryption : Use the keys  $K[i]$  in reverse order. That is, instead of applying  $K[1]$  for the first iteration, apply  $K[16]$ , and then  $K[15]$  for the second, on down to  $K[1]$

Input:



Output:

