

A proof of the Hübsch conjecture on constructing 2D Adinkras from 1D Adinkras

Kevin Iga and Yan X Zhang

March 11, 2015

1 Preliminaries

1.1 1-d Adinkras

Adinkras in [references] will be referred to as 1-d Adinkras in this paper, since they relate to supersymmetry in 1 dimension. We will review a definition of 1-d Adinkras now.

Definition 1.1. Let n be a non-negative integer. A 1-d Adinkra with n colors is (V, E, c, μ, g) where

- (V, E) is a finite undirected graph (called the underlying graph of the Adinkra) with vertex set V and edge set $E \subset V \times V$,
- $c : E \rightarrow \{1, \dots, n\}$ is a map called the coloring,
- $\mu : E \rightarrow \mathbf{Z}_2 = \{0, 1\}$ is a map called the dashing,
- $h : V \rightarrow \mathbf{Z}$ is a map called the grading.

These are required to satisfy the following:

1. If $(v, w) \in E$, then $(w, v) \in E$. Furthermore, $c(v, w) = c(w, v)$ and $\mu(v, w) = \mu(w, v)$. Intuitively, the edges are undirected.
2. For every $v \in V$ and $c \in \{1, \dots, n\}$, there exist exactly one $w \in V$ so that $(v, w) \in E$ and $c(v, w) = c$.
3. If $c_1, c_2 \in \{1, \dots, n\}$ with $c_1 \neq c_2$, and $v \in V$, then there exist w, x , and $y \in V$ so that $(v, w), (w, x), (x, y)$, and $(y, v) \in E$, and $c(v, w) = c(x, y) = c_1$ and $c(w, x) = c(y, v) = c_2$ and $\mu(v, w) + \mu(w, x) + \mu(x, y) + \mu(y, v) \equiv 1 \pmod{2}$. Note that μ is called admissible if this last condition holds for every v, w, x and y that satisfy the above.

4. If $(v, w) \in E$, then $|h(v) - h(w)| = 1$.

Note that in [Reference], there is also a bipartition of the vertices, where some vertices are represented by open circles and called bosons, and other vertices are represented by filled circles and called fermions. This is not necessary to include in our definition, because a vertex v is a boson if and only if $h(v)$ is even.

1.2 The action of \mathbf{Z}_2^n and the code

Let A be a 1-d Adinkra with n colors, with vertex set V . For all $1 \leq i \leq n$, define

$$q_i : V \rightarrow V$$

so that for all $v \in V$, $q_i(v)$ is the unique vertex joined to v by an edge of color i .

Proposition 1.2. *The map q_i is a graph isomorphism from the underlying graph of A to itself which preserves colors. It is an involution and for all i, j , we have*

$$q_i \circ q_j = q_j \circ q_i. \quad (1.3)$$

Proof. The statement that q_i is a graph homomorphism means that if (v, w) is an edge in A , then so is $(q_i(v), q_i(w))$. This follows from items 2 and 3 in the definition of an Adinkra above, using $j = c(v, w)$.

The fact that $q_i(q_i(v)) = v$ for all $v \in V$ follows from item 2 in the definition. This means that q_i is an involution and in particular is an isomorphism.

The equation (1.3) follows from item 3 of the definition when $i \neq j$ and is trivial when $i = j$. \square

By combining the q_1, \dots, q_n maps, we can define an action of \mathbf{Z}_2^n on the graph (V, E) underlying the Adinkra in the following way:

Definition 1.4. The action of \mathbf{Z}_2^n on the graph (V, E) underlying the Adinkra is given on vertices by

$$(x_1, \dots, x_N)v = q_1^{x_1} \circ \dots \circ q_n^{x_n}(v)$$

and the fact that it acts on edges can be seen by the fact that each q_i is a graph homomorphism. It preserves colors since each q_i does.

Definition 1.5. A path is a finite sequence of edges of the form $((v_1, v_2), (v_2, v_3), \dots, (v_{k-1}, v_k))$. The *color sequence* of the path is the sequence $(c(v_1, v_2), c(v_2, v_3), \dots, c(v_{k-1}, v_k))$.

So if there is a path from v to w with color sequence (i_1, \dots, i_k) , we have $w = q_{i_k} \circ \dots \circ q_{i_1}(v)$.

Proposition 1.6. *Let A be an Adinkra and let v be a vertex of A . Let σ be a color sequence. There exists a unique path in A that starts at v and has color sequence σ .*

Proof. This can be proved by induction on the length of σ , and using the fact that given a vertex v_i of A , and a color c_i , there exists a unique vertex v_{i+1} so that (v_i, v_{i+1}) is an edge of A with color c_i . \square

Now, define a map s that takes a color sequence and returns an element of $\mathbf{Z}_2^n = \{0, 1\}^n$ where the i -th coordinate is the number of times (modulo 2) that color i appears in the sequence. For example, $s(3, 1, 2, 1) = 0110$. Note that $s(\sigma)$ does not depend on the ordering of the color sequence σ . This relates to paths in Adinkras because of the following:

Proposition 1.7. *Let A be an Adinkra. Let v be a vertex of A and let p be a path that begins at v . Let σ be the color sequence obtained from p . Then the path p ends at the vertex $s(\sigma)v$.*

Proof. If the color sequence is $\sigma = (i_1, \dots, i_k)$, then the path p ends at $q_{i_k} \circ \dots \circ q_{i_1}(v)$. By the commutativity of the q_i , we can order them in non-decreasing order of i_j . If any of the q_i appear more than once, we use the fact that q_i^2 is the identity to reduce the number of q_i modulo 2. The result is $s(\sigma)v$. \square

Corollary 1.8. *Let A be an Adinkra. Let v be a vertex of A and let p and p' be paths that begin at v . Let σ and σ' be the color sequences obtained from p and p' , respectively. If $s(\sigma) = s(\sigma')$, then p and p' both end at the same point.*

Definition 1.9. Pick a vertex $v \in A$. Define $C(A, v)$ to be the stabilizer of v under this action of \mathbf{Z}_2^n . Since it is a subgroup of \mathbf{Z}_2^n , $C(A, v)$ is a binary block code of length n .

Proposition 1.10. *The Adinkra A is connected if and only if the \mathbf{Z}_2^n action is transitive on the vertex set of A .*

Proof. Let v, w be vertices of A . If A is connected, then there is a path in A connecting v to w . Let σ be the color sequence obtained from this path. Then by Proposition 1.7, $s(\sigma)v = w$.

Conversely, suppose the action is transitive. Let v and w be vertices of A . Then there exists a $\mathbf{x} \in \mathbf{Z}_2^n$ so that $w = \mathbf{x}v$. Write $\mathbf{x} = (x_1, \dots, x_n)$ and construct a color sequence σ by taking the i for which $x_i = 1$. By Proposition 1.6, there is a path starting at v that has σ as its color sequence. By Proposition 1.7, this path ends at $s(\sigma)v = \mathbf{x}v = w$. \square

Proposition 1.11. *If A is connected, then the code $C(A, v)$ does not depend on v .*

Proof. Let $w \in V$. By Proposition 1.10, there exists a $\mathbf{y} \in \mathbf{Z}_2^n$ so that $\mathbf{y}v = w$.

Let $\mathbf{x} \in \mathbf{Z}_2^n$. The result follows from the sequence of equivalences:

$$\mathbf{x}w = w \Leftrightarrow \mathbf{x}\mathbf{y}v = \mathbf{y}v \Leftrightarrow \mathbf{y}\mathbf{x}v = \mathbf{y}v \Leftrightarrow \mathbf{x}v = v$$

\square

Definition 1.12. Given a connected Adinkra A , the code for A , called $C(A)$, is defined to be $C(A, v)$, where v is a vertex of A .

Proposition 1.13. *Let A be an Adinkra. Let v be a vertex of A and let p and p' be paths that begin at v . Let σ and σ' be the color sequences obtained from p and p' , respectively. The paths p and p' end at the same vertex if and only if*

$$s(\sigma) - s(\sigma') \in C(A).$$

Proof. Suppose p and p' end at the same vertex. Then by Proposition 1.7,

$$s(\sigma)v = s(\sigma')v.$$

Then¹

$$v = s(\sigma')^{-1}(s(\sigma)v) = (s(\sigma) - s(\sigma'))v.$$

Thus, $s(\sigma) - s(\sigma') \in C(A)$.

Conversely, suppose $s(\sigma) - s(\sigma') \in C(A)$. Then by reversing the above argument,

$$s(\sigma)v = s(\sigma')v$$

and thus, by Proposition 1.7, p and p' end at the same vertex. \square

The following 1-d Adinkra was defined in [reference].

Construction 1.14. *For any non-negative integer n , we have an Adinkra $I^n = (V, E, c, \mu, h)$ with*

- $V = \mathbf{Z}_2^n = \{0, 1\}^n$,
- $E = \{(v, w) \mid v \text{ and } w \text{ differ in precisely one coordinate}\}$,
- $c(v, w)$ is the coordinate where v and w differ,
- $\mu(v, w)$ is the number of 1s in v before the coordinate $c(v, w)$, modulo 2,
- $h(v) = \text{wt}(v)$ is the number of 1s in v .

Definition 1.15. Given C a linear block code of length n , we can define a graph with edge colors I^n/C as the orbit space of the action of $C \subset \mathbf{Z}_2^n$ on I^n as a graph.

The fact that C acts on I^n via graph homomorphisms results in a graph (see later section on Quotients).

Definition 1.16. A graph homomorphism from a graph (V_1, E_1) to a graph (V_2, E_2) is a map $\phi : V_1 \rightarrow V_2$ so that if $(v, w) \in E_1$ is an edge, then $(\phi(v), \phi(w)) \in E_2$ is an edge. If there is a coloring $c_1 : E_1 \rightarrow \{1, \dots, n\}$ and a coloring $c_2 : E_2 \rightarrow \{1, \dots, n\}$, we say that ϕ preserves colors if $c_1(v, w) = c_2(\phi(v), \phi(w))$. If ϕ is bijective, we say that it is a graph isomorphism.

¹Note that in this sequence of equations, \mathbf{Z}_2^n is written additively but the group action is written multiplicatively.

Theorem 1.17. *If A is a connected Adinkra, then there is a graph isomorphism from $I^n/C(A)$ to A that preserves colors.*

Proof. Choose a vertex $\bar{0}$ in A . Let

$$\phi : I^n \rightarrow A$$

$$\phi(x_1, \dots, x_n) = (x_1, \dots, x_n)\bar{0}$$

where we are using the action of \mathbf{Z}_2^n on A as described above.

To see that this is a graph homomorphism, let $(x_1, \dots, x_n) \in \mathbf{Z}_2^n$ and let (y_1, \dots, y_n) be another vertex connected to (x_1, \dots, x_n) with an edge of color i . Then $y_j = x_j$ for all $j \neq i$ and $y_i = 1 - x_i$. Then $(y_1, \dots, y_n)\bar{0} = q_1^{y_1} \cdots q_n^{y_n}(\bar{0}) = q_1^{x_1} \cdots q_i^{1-x_i} \cdots q_n^{x_n}(\bar{0}) = q_i(q_1^{x_1} \cdots q_n^{x_n}(\bar{0})) = q_i(x_1, \dots, x_n)$. So $\phi(x_1, \dots, x_n)$ and $\phi(y_1, \dots, y_n)$ are connected by an edge of color i . Note that this also shows that the graph homomorphism preserves colors.

We now prove that ϕ is surjective. Since A is connected, \mathbf{Z}_2^n acts transitively on the vertex set of A , and so for any vertex v of A , there exists an element $\mathbf{x} \in \mathbf{Z}_2^n$ so that $\mathbf{x}\bar{0} = v$. Then $v = \phi(\mathbf{x})$.

To prove the isomorphism from $I^n/C(A)$ to A , we consider the necessary and sufficient conditions for $\phi(\mathbf{x}) = \phi(\mathbf{y})$ for $\mathbf{x}, \mathbf{y} \in \mathbf{Z}_2^n$. The condition $\phi(\mathbf{x}) = \phi(\mathbf{y})$ is equivalent to $\phi(\mathbf{x} - \mathbf{y})\bar{0} = \bar{0}$. This is equivalent to saying $\mathbf{x} - \mathbf{y} \in C(A)$. Thus, the map ϕ descends to $I^n/C(A)$ and gives an isomorphism. \square

In Reference ..., it was proved that if A is a connected Adinkra, then $C(A)$ is a doubly even code. Furthermore, if C is any doubly even binary block code of length n , then there is an Adinkra with $C(A) = C$.

Lemma 1.18. *Given connected 1-d Adinkras A and B and vertices a of A and b of B , there is at most one graph isomorphism from A to B that preserves colors and sends a to b .*

Proof. Suppose $\phi : A \rightarrow B$ is a graph isomorphism that preserves colors and sends a to b .

Let v be a vertex of A . Since A is connected, there is a path from a to v . This produces a color sequence σ . Since ϕ is a graph isomorphism that preserves colors, ϕ sends this path to a path in B starting from b with the same color sequence. By Proposition 1.6, there is only one path in B with this property. The end of this path is $\phi(v)$. \square

1.3 2-d Adinkras

The notion of 2-d Adinkras is described in Ref...[fill in references to various things]. We use a definition here that is equivalent to the one found in [some reference]: the proof is found in Appendix...[maybe?]

A 2-d Adinkra is similar to a 1-d Adinkra except that some colors are called “left-moving” and the other colors called “right-moving”. Edges are called “left-moving” if they are colored by left-moving edges, and right-moving otherwise. Furthermore, there are two gradings, one that is affected by the left-moving edges and the other for the right-moving edges.

More formally,

Definition 1.19. Let p and q be non-negative integers. A 2-d Adinkra with (p, q) colors is a 1-d Adinkra (V, E, c, μ, h) with $p + q$ colors, and two grading functions $h_L : V \rightarrow \mathbf{Z}$ and $h_R : V \rightarrow \mathbf{Z}$ so that

- $h(v) = h_L(v) + h_R(v)$
- if (v, w) is a left-moving edge, then $|h_L(v) - h_L(w)| = 1$ and $h_R(v) = h_R(w)$.
If (v, w) is a right-moving edge, then $|h_R(v) - h_R(w)| = 1$ and $h_L(v) = h_L(w)$.

1.4 Product of Adinkras

One way to get 2-d Adinkras is to take a product of two 1-d Adinkras, where the first Adinkra uses only left-moving colors and the second Adinkra uses only right-moving colors.

Construction 1.20. Let p and q be non-negative integers. Let $A_1 = (V_1, E_1, c_1, \mu_1, h_1)$ be a 1-d Adinkra with p colors and let $A_2 = (V_2, E_2, c_2, \mu_2, h_2)$ be a 1-d Adinkra q colors. We can define the product of these Adinkras $A_1 \times A_2$ as the following 2-Adinkra with (p, q) colors:

$$A_1 \times A_2 = (V, E, c, \mu, h_L, h_R)$$

where

$$\begin{aligned} V &= V_1 \times V_2 \\ E &= E_1 \cup E_2 \text{ where} \\ E_1 &= \{((v_1, w), (v_2, w)) \mid (v_1, v_2) \in E_1, \text{ and } w \in V_2\} \\ E_2 &= \{((v, w_1), (v, w_2)) \mid v \in V, \text{ and } (w_1, w_2) \in E_2\} \\ c((v_1, w), (v_2, w)) &= c_1(v_1, v_2) \text{ for all } ((v_1, w), (v_2, w)) \in E_1 \\ c((v, w_1), (v, w_2)) &= p + c_2(w_1, w_2) \text{ for all } (v, w_1), (v, w_2) \in E_2 \\ h_L(v, w) &= h_1(v) \\ h_R(v, w) &= h_2(w) \\ \mu((v_1, w), (v_2, w)) &= \mu_1(v_1, v_2) \\ \mu((v, w_1), (v, w_2)) &= \mu_2(w_1, w_2) + h_1(v) \pmod{2} \end{aligned}$$

Proposition 1.21. Given Adinkras A_1 and A_2 , μ is an admissible dashing for $A_1 \times A_2$. That is, given colors c_1 and c_2 , and vertices v, w, x , and y in $A_1 \times A_2$, with edges (v, w) , (w, x) , (x, y) , and (y, v) and colors $c(v, w) = c(x, y) = c_1$ and $c(w, x) = c(y, v) = c_2$, we have $\mu(v, w) + \mu(w, x) + \mu(x, y) + \mu(y, v) \equiv 1 \pmod{2}$.

Proof. If c_1 and c_2 are both left-moving colors, then write $v = (v_1, z)$. Then $w = (w_1, z)$, $x = (x_1, z)$, and $y = (y_1, z)$. Then we have

$$\mu(v, w) + \mu(w, x) + \mu(x, y) + \mu(y, v) = \mu_1(v_1, w_1) + \mu_1(w_1, x_1) + \mu_1(x_1, y_1) + \mu_1(y_1, v_1)$$

and the result follows from the fact that μ_1 was an admissible dashing on A_1 .

If c_1 and c_2 are both right-moving colors, then write $v = (z, v_2)$. Then $w = (z, w_2)$, $x = (z, x_2)$, and $y = (z, y_2)$. Then

$$\begin{aligned} & \mu(v, w) + \mu(w, x) + \mu(x, y) + \mu(y, v) \\ &= \mu_2(v_1, w_1) + h_1(z) + \mu_2(w_1, x_1) + h_1(z) + \mu_2(x_1, y_1) + h_1(z) + \mu_2(y_1, v_1) + h_1(z) \end{aligned}$$

which modulo 2 is the same as

$$\mu_2(v_1, w_1) + \mu_2(w_1, x_1) + \mu_2(x_1, y_1) + \mu_2(y_1, v_1).$$

The result follows from the fact that μ_2 was an admissible dashing on A_2 .

Now if c_1 is left-moving and c_2 is right-moving, then write $v = (v_1, v_2)$. Then $w = (w_1, v_2)$, $x = (w_1, x_2)$, and $y = (v_1, x_2)$. Now note that v_1 and w_1 are connected by an edge in A_1 , so $h_1(v_1)$ and $h_1(w_1)$ differ by 1. In particular, one of these is even and the other is odd. Without loss of generality, suppose $h_1(v_1)$ is even. Then $\mu(v, w) = \mu_1(v_1, w_1)$, $\mu(w, x) = \mu_2(v_2, x_2) + 1$, $\mu(x, y) = \mu_1(w_1, v_1)$, and $\mu(y, v) = \mu_2(x_2, v_2)$, and we have

$$\begin{aligned} & \mu(v, w) + \mu(w, x) + \mu(x, y) + \mu(y, v) \\ &= \mu_1(v_1, w_1) + \mu_2(v_2, x_2) + 1 + \mu_1(w_1, v_1) + \mu_2(x_2, v_2) \equiv 1 \pmod{2}. \end{aligned}$$

□

Definition 1.22. Let p and q be non-negative integers and let $n = p + q$.

Given a binary block code C of length p , we can define a binary block code \hat{C} of length n by appending to the end of every code word in C a string of 0s of length q .

Likewise, given a binary block code C of length q , we can define a binary block code \check{C} of length n by prepending to the beginning of every code word in C a string of 0s of length p .

Proposition 1.23. Let A_1 and A_2 be as above. Then

$$C(A_1 \times A_2) = \hat{C}(A_1) \oplus \check{C}(A_2).$$

Proof. Let $(v_1, v_2) \in A_1 \times A_2$. Let $\mathbf{x} \in \mathbf{Z}_2^N$. We can write $\mathbf{x} = \mathbf{x}_1 + \mathbf{x}_2$ where \mathbf{x}_1 is zero in the last q bits and \mathbf{x}_2 is zero in the first p bits. Now

$$\mathbf{x}(v_1, v_2) = (\mathbf{x}_1 + \mathbf{x}_2)(v_1, v_2) = (\mathbf{x}_1 v_1, \mathbf{x}_2 v_2).$$

This means that $\mathbf{x}(v_1, v_2) = (v_1, v_2)$ if and only if $\mathbf{x}_1 v_1 = v_1$ and $\mathbf{x}_2 v_2 = v_2$.

So $\mathbf{x} \in C(A_1 \times A_2)$ if and only if $\mathbf{x}_1 \in \hat{C}(A_1)$ and $\mathbf{x}_2 \in \check{C}(A_2)$. □

2 Structural Theorems

In this section, we show that the coherence conditions of 2-d adinkras force a lot of structure onto them. In particular, we can think of the vertices of 2-d adinkras as arranged in a rectangle, with the structure of the entire adinkra basically determined by a horizontal and a vertical “slice” of the picture.

2.1 A 2-d Adinkra Fits in a Rectangle

Let the *support* of a 2-d adinkra (and/or its bigrading function (h_L, h_R)) be defined as the range of (h_L, h_R) , its bigrading function. Now, we show that the support of a connected 2-d adinkra must form a rectangle in \mathbf{Z}^2 .

Let σ be a color sequence. Let $l(\sigma)$ be a rearrangement of σ that moves all the left-moving colors to the beginning, and let $r(\sigma)$ be a rearrangement of σ that moves all the right-moving colors to the beginning. Thus, suppose $p = q = 2$, we have $l((3, 1, 2, 1)) = (2, 1, 1, 3)$ and $r((3, 1, 2, 1)) = (3, 2, 1, 1)$. We always have $s(l(\sigma)) = s(r(\sigma))$ in general, since $l(\sigma)$ and $r(\sigma)$ are just permutations of each other.

Proposition 2.1. *Let A be a connected Adinkra. Suppose (x_1, y_1) and (x_2, y_2) are in the support of A . Then (x_1, y_2) and (x_2, y_1) are also in the support of A .*

Proof. The statement that (x_1, y_1) and (x_2, y_2) is in the support of A means that there exist vertices v_1 and v_2 of A with $(h_L(v_1), h_R(v_1)) = (x_1, y_1)$ and $(h_L(v_2), h_R(v_2)) = (x_2, y_2)$, respectively. Since A is connected, there exists a path from v_1 to v_2 .

Let σ be the color sequence of this path. Now $l(\sigma)$ is a color sequence that can be written as a concatenation $\alpha\beta$ where α is a color sequence only involving left-moving colors, and β is a color sequence only involving right-moving colors. Consider the vertex $u = s(\alpha)v_1$. The color sequence α describes a path from v_1 to u involving only left-moving colors, and so $h_R(u) = h_R(v_1) = y_1$. The color sequence β describes a path from u to v_2 involving only right-moving colors, and so $h_L(u) = h_L(v_2) = x_2$. To summarize, $(h_L(u), h_R(u)) = (x_2, y_1)$.

Repeating this procedure with $r(\sigma)$ likewise provides a vertex w with $(h_L(w), h_R(w)) = (x_1, y_2)$. \square

Corollary 2.2. *The support of a connected 2-d adinkra is a rectangle. That is, there exist integers x_0, x_1, y_0 , and y_1 so that the support is*

$$\{(x, y) \in \mathbf{Z}^2 \mid x_0 \leq x \leq x_1 \text{ and } y_0 \leq y \leq y_1\}$$

Proof. Let x_0 and y_0 be the minima of the x and y coordinates, respectively, of the support of the Adinkra. By Proposition 2.1, (x_0, y_0) is in the support as well.

Likewise, if x_1 and y_1 are the maxima of the x and y coordinates, respectively, of the support of the Adinkra, then (x_1, y_1) is in the support. By Proposition 2.1, (x_1, y_0) and (x_0, y_1) are also in the support.

Since the Adinkra is connected, there must be paths from vertices with bigrading (x_0, y_0) to vertices with bigrading (x_1, y_1) . Since h_L and h_R can change by at most 1 along these paths, we see that for all $x_0 \leq x \leq x_1$, there must exist y_x so that (x, y_x) is in the support. Likewise for all $y_0 \leq y \leq y_1$, there must exist x_y so that (x_y, y) is in the support. By application of Proposition 2.1 again, we get that (x, y) is in the support for all $x_0 \leq x \leq x_1$ and $y_0 \leq y \leq y_1$. \square

While it is neat that the vertices of a 2-d adinkra A line up nicely in a rectangle, we now show that there is even more regularity in its structure. Let A_L (resp. A_R) be the subgraphs of A induced by left-moving (resp. right-moving) edges of A .

Lemma 2.3. *Let A be a 2-d Adinkra. If X is a connected component of A_L and i is a right-moving color, then there is a graph isomorphism between X and $q_i(X)$ that preserves colors and h_L . The analogous statement for A_R also holds.*

Proof. Proposition 1.2 states that q_i is a graph isomorphism from the underlying graph of A to itself that preserves colors. If we restrict q_i to a connected component X of A_L , the restricted map is an isomorphism from X to $q_i(X)$ that preserves colors.

Since i is a right-moving color, then for all vertices $v \in X$, $h_L(v) = h_L(q_i(v))$. \square

Lemma 2.4. *Let A be a 2-d Adinkra. If X is a connected component of A_L and i is a right-moving color, then $q_i(X)$ is the vertex set of a connected component of A_L . The analogous statement for A_R also holds.*

Proof. Because the property of connectedness is preserved under graph isomorphism, we know that $q_i(X)$ is connected. If we let X' be the connected component of A_L that contains $q_i(X)$, then the same argument proves that $q_i(X')$ is connected as well. Since X was assumed to be a connected component of A_L , we have that $q_i(X') \subseteq X$. But since q_i^2 is the identity, $X = q_i^2(X) \subseteq q_i(X') \subseteq X$. This means $q_i(X') = X$. By the fact that q_i^2 is the identity, we also have $q_i(X) = X'$. \square

Proposition 2.5. *Let A be a connected 2-d Adinkra. All connected components of A_L (and respectively A_R) are isomorphic as graded posets.*

Proof. Let X and Y be two connected components of A_L . Pick vertices $x \in X$ and $y \in Y$. Since A is connected, there is a path from x to y in A . Reorder the path so that the right-moving edges occur before the left-moving edges. Since the left-moving edges stay in Y , the right-moving edges alone take x to a vertex $y' \in Y$. The sequence of right-moving edges provides a color sequence i_1, \dots, i_k , and thus, a sequence of compositions $q_{i_k} \circ \dots \circ q_{i_1}$. Now $q_{i_k} \circ \dots \circ q_{i_1}(x) = y'$. By repeated application of Lemma 2.4, we have that $q_{i_k} \circ \dots \circ q_{i_1}(X)$ is a connected component of A_L that contains y' , which is Y . By repeated application of Lemma 2.3, we have an isomorphism of graphs that preserves colors and the grading h_L . \square

With all this redundancy, what is the minimal amount of information required for us to understand a 2-d adinkra? Proposition 2.5 suggests we just need a single connected component for each direction to give us all the data; this turns out to basically be true, as we see in the next section.

3 Quotienting

To understand quotients, we first define the homomorphism from one 2-d Adinkra to another. This will be similar to the definition of homomorphism of graphs [give some standard reference to this terminology].

Definition 3.1. Let $A_1 = (V_1, E_1, c_1, \mu_1, h_{L1}, h_{R1})$ and $A_2 = (V_2, E_2, c_2, \mu_2, h_{L2}, h_{R2})$ be 2-Adinkras with (p, q) colors. A homomorphism from A_1 to A_2 is a graph homomorphism

$$\phi : (V_1, E_1) \rightarrow (V_2, E_2)$$

that preserves colors and so that:

- If $v \in V_1$ then $h_{1L}(v) = h_{2L}(\phi(v))$.
- If $v \in V_1$ then $h_{1R}(v) = h_{2R}(\phi(v))$.

Note that there is no condition on the dashings μ_1 and μ_2 .

Proposition 3.2. *If a group C acts on a graph $\Gamma = (V, E)$ via graph isomorphisms, then C acts on the set of edges in the following manner: if $c \in C$ and $(v, w) \in E$, then $c(v, w) = (cv, cw)$.*

Proof. The fact that (cv, cw) is an edge follows directly from the definition of graph homomorphism.

The fact that this is a group action follows from the fact that $c(c'(v, w)) = (c(c'(v)), c(c'(w))) = ((cc')v, (cc')w) = (cc')(v, w)$. \square

Definition 3.3. Suppose C is a group that acts on a graph (V, E) via graph isomorphisms. We say that this group action is distant if for every non-identity group element $c \in C$, and every vertex $v \in V$, every path from v to cv is of length greater than 2.

Construction 3.4. *If a group C acts on a graph $\Gamma = (V, E)$ via graph isomorphisms, then we create a graph Γ/C as follows: let the set of vertices be the orbit space V/C . Edges in Γ/C are pairs (Cv, Cw) where there exist $c_1, c_2 \in C$ so that (c_1v, c_2w) is an edge in E .*

Proposition 3.5. *If, in the previous construction, the group action is distant, then there is a bijection from E/C to the set of edges in Γ/C .*

Proof. For every orbit $C(v, w)$ of C on the set of edges of Γ , there is an edge (Cv, Cw) . This is an edge in Γ/C because (v, w) is an edge in Γ . This assignment is independent of the choice of representative (v, w) from $C(v, w)$ above since if (v', w') is another representative, then there is a group element $c \in C$ so that $v' = cv$ and $w' = cw$. Then $v' \in Cv$ and $w' \in Cw$ so that $Cv' = Cv$ and $Cw' = Cw$.

To see that this assignment is surjective follows directly from the definition: for every edge in Γ/C , there is an edge (v, w) in Γ that gets sent to it.

Now to prove that this is injective, suppose (v, w) and (v', w') are edges in Γ and suppose $(Cv, Cw) = (Cv', Cw')$. Then there are group elements c_1 and $c_2 \in C$ so that $c_1v = v'$ and $c_2w = w'$. On the other hand, there is an edge $(c_1v, c_1w) = (v', c_1w)$. So (v', c_1w) and $(v', c_2w) = (v', w')$ are edges in the graph. Then $c_2c_1^{-1}$ is a group element that sends c_1w to c_2w but we have just constructed a path of length 2 from c_1w to c_2w . Since the group action is distant, we know that $c_2c_1^{-1} = 1$ so that $c_1 = c_2$. Therefore $(v', w') = c_1(v, w)$. \square

Proposition 3.6. *If a group C acts on an Adinkra A using graph isomorphisms that preserve colors, and if the action is distant, then the coloring descends to a coloring on A/C that satisfies those properties of Adinkras that relate to colors and not to dashings or grading. That is, let E/C denote the set of edges of A/C as defined in the construction. Then*

1. *If (Cv, Cw) is an edge in A/C , then (Cw, Cv) is also an edge. Furthermore, $c(Cv, Cw) = c(Cw, Cv)$.*
2. *For every $Cv \in V/C$ and $c \in \{1, \dots, n\}$, there exist exactly one $Cw \in V/C$ so that $(Cv, Cw) \in E/C$ and $c(Cv, Cw) = c$.*
3. *If $c_1, c_2 \in \{1, \dots, n\}$ with $c_1 \neq c_2$, and $Cv \in V/C$, then there exist Cw, Cx , and $Cy \in V/C$ so that $(Cv, Cw), (Cw, Cx), (Cx, Cy)$, and $(Cy, Cv) \in E/C$, and $c(Cv, Cw) = c(Cx, Cy) = c_1$ and $c(Cw, Cx) = c(Cy, Cv) = c_2$*

Proof. Since Proposition ?? says that the edges of A/C can be identified with E/C , and since the action of C preserves colors, we can define the coloring on E/C , and thus, on the edges of A/C .

Property 1 follows from the construction of the edges of A/C .

Now we prove Property 2. Let $Cv \in V/C$ and $c \in \{1, \dots, n\}$. Since A is an Adinkra, there is exactly one vertex $w \in V$ so that $(v, w) \in E$ and $c(v, w) = c$. There is thus an edge (Cv, Cw) in A/C that has color c . To show uniqueness, suppose there is an $x \in V$ so that (Cv, Cx) is in A/C with color c . This means there are group elements $\alpha, \beta \in C$ so that $(\alpha v, \beta x)$ is an edge of A of color c . By acting by α^{-1} , we see that without loss of generality, $\alpha = 1$. Because A is an Adinkra, there is a unique edge from v that has color c . So $\beta x = w$. Therefore $x \in Cw$.

Now we prove Property 3. Let c_1 and c_2 be two different colors, and $Cv \in V/C$. Since A is an Adinkra, there are vertices w, x , and y in V so that $(v, w), (w, x), (x, y)$ and $(y, v) \in E$ and $c(v, w) = c(x, y) = c_1$ and $c(w, x) = c(y, v) = c_2$. Then Cw, Cx , and Cy are the requisite vertices in V/C with the given property. \square

[Change notation? Colors and codes are both C and c ... Graphs are often G but here they are Γ because groups are often G .]

4 Main Theorem

The main result of this section is the following theorem:

Theorem 4.1. *Let A be a connected 2-d Adinkra. Fix a vertex $\bar{0}$ in A and let A_L^0 (resp. A_R^0) be the connected component of A_L (resp. A_R) containing $\bar{0}$.*

There is a binary block code K of length n so that $K \cap \hat{C}(A_L^0) = 0$ and $K \cap \check{C}(A_R^0) = 0$, and there is a graph isomorphism

$$A \cong (A_L^0 \times A_R^0)/K \quad (4.2)$$

of the underlying graphs that preserves colors and the bigrading.

Before we begin it will help to understand the relationship of the codes.

Lemma 4.3.

$$\hat{C}(A_L^0) \oplus \check{C}(A_R^0) \subseteq C(A).$$

Proof. Let $g \in \hat{C}(A_L^0)$ and $h \in \check{C}(A_R^0)$. Then $g\bar{0} = \bar{0}$ because applying g to $\bar{0}$ results in a path that lies completely inside A_L^0 , and so the fact that $g\bar{0} = \bar{0}$ in A_L^0 (since $g \in \hat{C}(A_L^0)$) results in $g\bar{0} = \bar{0}$ in A . Likewise $h\bar{0} = \bar{0}$. So $(g+h)\bar{0} = g(h\bar{0}) = \bar{0}$ and $g+h \in C(A)$. \square

Proof of Theorem 3.2. From Lemma 3.4 and basic linear algebra, there exists a vector subspace K of \mathbf{Z}_2^N that is a vector space complement of $\hat{C}(A_L^0) \oplus \check{C}(A_R^0)$ in $C(A)$. That is,

$$C(A) = \hat{C}(A_L^0) \oplus \check{C}(A_R^0) \oplus K.$$

Then by Theorem 1.17, there is an isomorphism

$$i_1 : A \cong I^n/C(A)$$

that preserves colors. The proof of this began with choosing a special vertex in A , and we specify here that we should use our chosen vertex $\bar{0}$ of A . This isomorphism then sends $\bar{0}$ to the coset containing $(0, \dots, 0)$ in $I^n/C(A)$.

We now note that

$$\begin{aligned} I^n/C(A) &= I^n/(\hat{C}(A_L^0) \oplus \check{C}(A_R^0) \oplus K) \\ &= (I^n/(\hat{C}(A_L^0) \oplus \check{C}(A_R^0)))/K \\ &= (I^n/C(A_L^0 \times A_R^0))/K \end{aligned}$$

and again apply Theorem 1.17 to establish an isomorphism

$$i_2 : (I^n/C(A_L^0 \times A_R^0))/K \cong (A_L^0 \times A_R^0)/K$$

that preserves colors and sends the coset of $(0, \dots, 0)$ to $(\bar{0}, \bar{0})$.

We compose these isomorphisms to define

$$i_3 : A \rightarrow (A_L^0 \times A_R^0)/K$$

an isomorphism of graphs that preserves colors.

We now restrict i_3 to A_L^0 . It turns out that this sends every vertex v of A_L^0 to $(v, \bar{0})$ in $(A_L^0 \times A_R^0)/K$. This can be seen by noting that both the restriction of i_3 to A_L^0 and the map sending v to $(v, \bar{0})$ are graph isomorphisms that preserve colors and send $\bar{0}$ to $(\bar{0}, \bar{0})$. By Lemma 1.18, such a map is unique. Therefore, $i_3(v) = (v, \bar{0})$ for all vertices v in A_L^0 . Likewise, a similar proof shows that if $v \in A_R^0$, then $i_3(v) = (\bar{0}, v)$.

Now let v be any vertex of A . There is a path from $\bar{0}$ to v . This gives a color sequence σ . Consider the reordered color sequence $l(\sigma)$. Write this as $\alpha\beta$ where α only involves left-moving colors and β only uses right-moving colors. Then define $x = s(\alpha)\bar{0}$. Because α only involves left-moving colors, x is in A_L^0 .

Now since β consists of right-moving colors, and connects $i_3(x)$ with $i_3(v)$, we have $h_L(i_3(v)) = h_L(i_3(x)) = h_L(x, \bar{0}) = h_L(x)$. This is $h_L(v)$ because β only involves right-moving colors and connects x with v . Thus, $h_L(i_3(v)) = h_L(v)$.

A similar proof shows that $h_R(i_3(v)) = h_R(v)$. \square

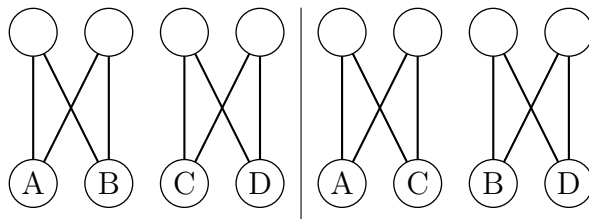


Figure 1: Taking the product of the two adinkras here with the following identification gives a non-disconnected adinkra with 16 vertices.

4.1 Dashing

We now consider the dashing. It is too much to expect the dashing on $A_L^0 \times A_R^0$ to agree with the dashing on A . It is even generally untrue that the dashing on $A_L^0 \times A_R^0$ is invariant under the quotienting action of K . But if we allow vertex flips, then we can accomplish this.

Theorem 4.4. *Let A be a connected 2-d Adinkra. Let $\bar{0}$, A_L^0 , A_R^0 , and K be as in Theorem 3.2.*

Then there is a sequence of vertex flips on $A_L^0 \times A_R^0$ so that the action of K preserves the dashing, and the quotient by K gives a dashing that agrees with the dashing on A when compared via the isomorphism given in Theorem 3.2.

Before proving this theorem, it is helpful to recognize that this theorem can be viewed as comparing two different dashings on $A_L^0 \times A_R^0$.

Consider the dashing μ on A . This restricts to A_L^0 and A_R^0 , and via Construction 1.20, this gives a dashing μ_1 on $A_L^0 \times A_R^0$.

At the same time, under the isomorphism (3.3), the dashing μ on A gets carried to $(A_L^0 \times A_R^0)/K$. The graph homomorphism

$$\pi : A_L^0 \times A_R^0 \rightarrow (A_L^0 \times A_R^0)/K$$

pulls back the dashing to μ_2 on $A_L^0 \times A_R^0$.

Proposition 4.5. *The map π above is a graph homomorphism.*

Proof. □

Proposition 4.6. *μ_2 is an admissible dashing. That is, given colors c_1 and c_2 , and vertices v, w, x , and y , with edges (v, w) , (w, x) , (x, y) , and (y, v) and colors $c(v, w) = c(x, y) = c_1$ and $c(w, x) = c(y, v) = c_2$, we have $\mu_2(v, w) + \mu_2(w, x) + \mu_2(x, y) + \mu_2(y, v) \equiv 1 \pmod{2}$.*

Proof. Since π is a graph homomorphism, and c_1, c_2, v, w, x , and y are as above, then $\pi(v), \pi(w), \pi(x)$, and $\pi(y)$ satisfy the same condition. If $\mu_2 = \pi^*(\mu)$, then $\mu_2(v, w) + \mu_2(w, x) + \mu_2(x, y) + \mu_2(y, v) = \mu(\pi(v, w)) + \mu(\pi(w, x)) + \mu(\pi(x, y)) + \mu(\pi(y, v))$ and the result follows from the fact that μ is an admissible dashing. □

In the language of cubical cohomology in [reference], what this theorem says is that $\mu_1 - \mu_2$ is 0 in cohomology.

Definition 4.7. Given a path given by vertices (v_0, \dots, v_k) , the parity of a dashing μ on the path is the following sum modulo 2:

$$\sum_{i=0}^{k-1} \mu(v_i, v_{i+1}) \pmod{2}.$$

Lemma 4.8. *The dashings μ_1 and μ_2 agree on $A_L^0 \times \{\bar{0}\}$ and on $\{\bar{0}\} \times A_R^0$.*

Proof. The construction of μ_1 gives all edges in $A_L^0 \times \{\bar{0}\}$ as the same as in A_L^0 under the association of every edge (v, w) with $((v, \bar{0}), (w, \bar{0}))$. For μ_2 , consider that i_3 sends every vector $v \in A_L^0$ to $(v, \bar{0})$, and likewise any edges in A_L^0 will be sent to the corresponding edge in $A_L^0 \times \{\bar{0}\}$. Therefore μ_1 and μ_2 agree on $A_L^0 \times \{\bar{0}\}$. Likewise for $\{\bar{0}\} \times A_R^0$. □

Lemma 4.9. *Given any cycle (v_0, v_1, \dots, v_k) with $v_0 = v_k$, the parities of μ_1 and μ_2 on the cycle are equal.*

Proof. We first consider the case where $v_0 = \bar{0}$. Then there is a color sequence σ from this. We modify the path (and thereby modify the color sequence) by “bubble sorting” the colors. Earlier, we noted that we could rearrange σ to $l(\sigma)$ so that the left-moving colors were before the right-moving colors. This can be more specifically done by a series of adjacent swaps of the type that sends the color sequence

$$(i_1, \dots, i_{j-1}, i_j, i_{j+1}, i_{j+2}, \dots, i_k)$$

to

$$(i_1, \dots, i_{j-1}, i_{j+1}, i_j, i_{j+2}, \dots, i_k).$$

A color swap of this type leads to a new path from $\bar{0}$, via Proposition 1.6. The path is unchanged up to v_{j-1} , but by the definition of Adinkras, property 3, the path returns to v_{j+1} so it is only v_j that has changed. The effect on the parity of any dashing is, by property 3, to add 1 modulo 2. In particular, μ_1 and μ_2 are both affected in the same way.

After a series of such color swaps, we rearrange σ to $l(\sigma)$, and the path in question consists of left-moving edges, which remain in $A_L^0 \times \{\bar{0}\}$, and then right-moving edges, which ends in $\bar{0}$, and so must be in $\{\bar{0}\} \times A_R^0$. By the previous lemma, μ_1 and μ_2 are equal here, and so their parity on this modified path is the same. Therefore, their parity on the original loop was the same.

Now we consider loops p where $v_0 \neq \bar{0}$. Since $A_L^0 \times A_R^0$ is connected, there is a path p_0 from $\bar{0}$ to v_0 . Take the path p_0 , followed by p , then followed by p_0^{-1} (meaning p_0 traversed in the opposite sense). This is a loop starting and ending in $\bar{0}$, but the parity of a dashing is the same as that of p , since every new edge in p_0 is counterbalanced by a new edge in p_0^{-1} . Therefore the parity of μ_1 and μ_2 agree on all loops. \square

Definition 4.10. Let $V' \subset V$ and let μ be a dashing. Then we can define a vertex flip of μ by V' to be a dashing $\psi_{V'}(\mu)$ so that

$$\psi_{V'}(\mu)(e) = \mu(e) + i(e, V') \pmod{2},$$

where $i(e, V')$ is the number of elements of V' that are incident with the edge e .

Note that the property 3 of Adinkras is unchanged if we use $\psi_{V'}(\mu)$ instead of μ , so the result of replacing μ with $\psi_{V'}(\mu)$ is still an Adinkra.

Lemma 4.11. *If μ_1 and μ_2 have the same parity on all cycles, then there is a set of vertices $V' \subset V$ so that $\mu_2 = \psi_{V'}(\mu_1)$.*

Two proofs: one using cohomology, one using spanning trees.

5 ESDE Codes

Let p and q be non-negative integers so that $p + q = n$, and let C be a linear block code of length n . For every element $\mathbf{x} = (x_1, \dots, x_n) \in \mathbf{Z}_2^n$, we define $\text{wt}_L(\mathbf{x})$ to be the number of 1s of (x_1, \dots, x_p) in the first p components, and $\text{wt}_R(\mathbf{x})$ to be the number of 1s of (x_{p+1}, \dots, x_n) in the last q components. Then $\text{wt}(\mathbf{x}) = \text{wt}_L(\mathbf{x}) + \text{wt}_R(\mathbf{x})$.

Define a *even-split doubly-even (ESDE) code* of length (p, q) to be a doubly-even code of length n so that every word \mathbf{x} in the code has $\text{wt}_L(\mathbf{x})$ and $\text{wt}_R(\mathbf{x})$ even.

Recall that 1-d chromotopologies are in bijection with quotients of the hamming cube I^n by a doubly-even code C , so any adinkra A has a well-defined associated code $C(A)$ that is uniquely determined by just the graph structure of the adinkra. Our goal is to show that ESDE codes are exactly the codes that appear for 2-d adinkras. To do this, we introduce a special family of 2-d adinkras.

In [references], there was a class of 1-d Adinkras that were particularly easy to construct and study, called *Valise Adinkras*, whose height function had values in two adjacent integers (usually 0 and 1). By analogy, given a ESDE code, we construct a 2-d Adinkra called the *Valise 2-d Adinkra* that has support in a 2×2 square.

Construction 5.1. *Let C be an ESDE code. We will describe a construction that provides a 2-d Adinkra with code C , called the Valise 2-d Adinkra.*

C is doubly-even, so there exists a connected 1-d Adinkra A with code $C(A) = C$. [reference: this was a construction somewhere else]

We now wish to define a bigrading (h_L, h_R) . To do this we do not use the grading h that comes from the 1-d Adinkra A . Instead, fix a vertex $\bar{0}$ of A . For every vertex v of A , pick a path from $\bar{0}$ to v . This produces a color sequence d . Define $h_L(v)$ to be the number modulo 2 of left-moving colors in d and likewise $h_R(v)$ is the number modulo 2 of right-moving colors in d . These functions are well-defined if every path from $\bar{0}$ to v has the same parity in number of left-moving colors and in number of right-moving colors. This occurs if and only if every loop starting from $\bar{0}$ has an even number of left-moving colors and an odd number of right-moving colors. If C is ESDE, then this is the case.

Theorem 5.2. *This construction gives a 2-d Adinkra with code C .*

Proof. By construction, A is a 1-d Adinkra with code C . The only thing left to check is whether the h_L and h_R that are defined satisfy the correct properties on edges.

Let (v, w) be a left-moving edge in the Adinkra A . Then the color of the edge is left-moving and so $h_L(v) \neq h_L(w)$ and $h_R(v) = h_R(w)$. Since the possible values for h_L and h_R are 0 and 1, it follows that $|h_R(v) - h_R(w)| = 1$.

The proof is similar for right-moving edges. □

Theorem 5.3. *For a code $C \subset \mathbf{Z}_2^n$, there exists a 2-d adinkra A with $C(A) = C$ if and only if C is a ESDE code.*

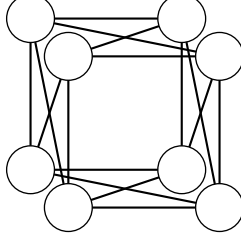


Figure 2: A valise 2-d adinkra that cannot be put into non-valise form.

Proof. Suppose $C(A) = C$ for some 2-d Adinkra A . We know that C is doubly-even. Consider any codeword $\alpha \in C$. Starting at $\bar{0} \in A$, moving by a path corresponding to α must end up back at $\bar{0}$. In particular, it must use an even number of left-moving (resp. right-moving) edges since each of such edge changes the h_L (resp. h_R) by 1 in absolute value. Thus, C must be ESDE.

Conversely, given an ESDE code C , we can use Construction 4.1 to get a 2-d Adinkra with code C . \square

5.1 Finding K

Given a 2-d Adinkra A , with special basepoint $\bar{0}$ and corresponding A_L^0 and A_R^0 , there remains the question of finding the code K . Recall that

$$C(A) = \hat{C}(A_L^0) \oplus \check{C}(A_R^0) \oplus K.$$

There is some choice involved in defining K ; it merely has to be a vector space complement to

$$C' = \hat{C}(A_L^0) \oplus \check{C}(A_R^0).$$

K will more naturally be defined as $C(A)/C'$, though computationally we can choose a representative to be K .

Let $V^0 = A_L^0 \cap A_R^0$. Now for every $v \in V^0$, we have $h_L(v) = h_L(\bar{0})$ and $h_R(v) = h_R(\bar{0})$ so in particular, V^0 has no edges; only vertices.

We construct a bijection between V^0 and $C(A)/C'$.

Construction 5.4. For every $v \in V^0$, there is a path p_L from $\bar{0}$ to v in A_L^0 and a path p_R from v to $\bar{0}$ in A_R^0 . These paths give color sequences σ_L and σ_R with $\beta_L = s(\sigma_L)$ and $\mathbf{x}_R = s(\sigma_R)$ so that \mathbf{x}_L is 0 in the last q coordinates and \mathbf{x}_R is 0 in the first p coordinates. Since the paths p_L and p_R form a loop in A , we know that $\mathbf{x}_L + \mathbf{x}_R \in C(A)$.

The paths p_L and p_R are not canonical, but if p'_L and p'_R are other paths with those properties, then $p_L(p'_L)^{-1}$ is a loop in A_L^0 from $\bar{0}$. Let σ'_L and σ'_R be the corresponding color sequences, then $s(\sigma_L) + s(\sigma'_L) \in C(A_L^0)$. Likewise, $s(\sigma_R) + s(\sigma'_R) \in C(A_R^0)$.

Thus, the choice of $\mathbf{x}_L + \mathbf{x}_R$ is well-defined modulo $C' = \hat{C}(A_L^0) \oplus \check{C}(A_R^0)$. This provides a map from V^0 to $K = C(A)/C'$.

Theorem 5.5. *The above construction provides a bijective map*

$$\Phi : V^0 \rightarrow C(A)/C'$$

Proof. We now prove this map is bijective by providing its inverse. Suppose $\mathbf{x} \in C(A)$. Then by Proposition 1.6, there is a loop in A starting at $\bar{0}$ that follows the colors indicated by \mathbf{x} . This gives a color sequence σ so that $s(\sigma) = \mathbf{x}$. Write $l(\sigma) = \alpha\beta$ where α only involves left-moving colors and β only involves right-moving colors. Let $v = s(\alpha)\bar{0}$. Since α only involves left-moving colors, we have $v \in A_L^0$. Since $s(\beta)v = \bar{0}$, and β only involves right-moving colors, we have that $v \in A_R^0$. Therefore $v \in V^0$. Note that v does not depend on the order of the left-moving colors in the color sequence, as long as they all come before the right-moving colors. Thus we have a map

$$\Psi : C(A) \rightarrow V^0.$$

Now suppose $\mathbf{y} \in \hat{C}(A_L^0)$ and $\mathbf{z} \in \check{C}(A_R^0)$. We consider how the loop in the previous paragraph changes if we replace \mathbf{x} with $\mathbf{x} + \mathbf{y} + \mathbf{z}$. We use Proposition 1.6 to create a loop p_1 in A_L^0 that starts at $\bar{0}$ and follows the colors indicated by \mathbf{y} . Then let p_2 be the path from $\bar{0}$ following the colors indicated by $s(\alpha)$ and p_3 the path that continues this using $s(\beta)$. Finally, let p_4 be the path from there using the colors of \mathbf{z} . By Proposition 1.6, since p_2 starts from $\bar{0}$, it must still end at the same point v . Therefore the map Ψ descends to a map

$$\tilde{\Psi} : C(A)/C' \rightarrow V^0.$$

We now show that this function is the inverse of Construction 4.4. Let $v \in V^0$. The construction gives a path p_L of left-moving edges starting from $\bar{0}$ and ending in v and a path p_R of right-moving edges starting from v to $\bar{0}$. Composing these paths gives the result of the Construction. Then we apply $\tilde{\Psi}$. This takes the code word and finds the loop, which by Proposition 1.6 must be $p_L p_R$. Then Ψ of the code word is the vertex at the end of p_L , which is v .

[show the other composition is the identity]

□

To do:

See if there is a place for a Lemma that says you can reorder a path so left colors come first.

A Equivalence with other notions of 2-d Adinkras

If I read Tristan's stuff right, we can completely translate the combinatorial rules to: a *2-d adinkra* (of dimension n) is a finite simple connected graph A such that:

- It is an 1-d adinkra (with the associated ranking, dashing, etc.).

- It has $p + q = n$ colors, where the first p -colors are called “left-moving” and the second q -colors are called “right-moving.”
- A coherence condition: for any cycle, we imagine the following sum: going up (here “up” comes from the grading we have from the engineering dimension in our ranking for the 1-d adinkra) a left-handed edge adds -1 , and going up a right-handed edge adds 1 ; going down the edges give contributions with opposite signs. The sum of this around any cycle must be 0 . (in particular, this rules out things like ambidextrous bow-ties)

Assuming I interpreted these rules correctly, now I can do combinatorics without needing any physics.

The first structural fact we can impose is a bi-grading that is compatible with the grading we already have from the 1-d adinkra structure, in the sense that the 1-d grading is simply one of the coordinates of our bi-grading.

Proposition A.1. *A 1-d adinkra can be extended to a 2-d adinkra if and only if the 1-d adinkra has a bigrading to \mathbf{Z}^2 . This is a map $g : V \rightarrow \mathbf{Z}^2$, such that all left-moving edges correspond to displacements of $(0, 1)$ and right-moving edges correspond to displacements of $(1, 0)$.*

Proof. Proof delayed until talking more with Kevin and Tristan about the easiest way to write things up to avoid reinventing wheels. \square

B Misc. (unorganized)

Corollary B.1. *Consider the vertex $\bar{0} \in A$. Let the connected component of A_L (resp. A_R) that $\bar{0}$ belongs to be labeled A_L^0 (resp. A_R^0). The adinkra A is uniquely determined² by A_L^0 and A_R^0 .*

Proof. Consider the color sequence d of any path from $\bar{0}$ to a vertex v . We can permute the sequence so that the left-moving colors all occur before the right-moving colors. Thus, we first make some moves in A_L^0 , then by Proposition 2.5 we make the remaining moves in a copy of A_R^0 . \square

Corollary B.2. *A valise 2-d adinkra A of type (n, k) is uniquely determined by $B_L(A)$, $B_R(A)$, and an identification of $V(B_L(A), 0)$ and $V(B_R(A), 0)$. Furthermore, $|B_L(A)| = |B_R(A)| = 2^{n-k-1}$.*

Problem 1. What are all the 2-d adinkras A with the same valise adinkra $\text{Val}(A)$? Not all lifts are possible. For example, the adinkra in Figure 2 cannot be lifted to any non-valise form!

Here are some other problems:

²This is fairly nuanced; we need to know not just the shape of A_L^0 and A_R^0 , but also their vertices

Problem 2. Given two valise 1-d adinkras $B_L(A)$ and $B_R(A)$ of equal size, what identifications of $V(B_L(A), 0)$ and $V(B_R(A), 0)$ are possible?

Proof. Data: if we have $\{0, 1\} \cup \{2, 3\}$ on one side, the other side must be $\{0, 2\} \cup \{1, 3\}$. \square

[*** There are two kinds of quotienting that we can think of: one quotient is directly quotienting the mega hypercube adinkra by a ESDE code; one quotient is given the valise adinkra with associated $B_L(A)$, $B_R(A)$ each with 2^{d+1} vertices, the necessary d -dimensional quotienting that occurs when we naively tensor the two parts (which gives 2^{2d} vertices in each corner, for 2^{2d+2} total vertices, when in the end we just want 2^{d+2} vertices. ***]