## **Assignment #2**

## RSA with Cipher Block Chaining (RSA-CBC)

- TCP Protocol
- Cipher Block Chaining
- RSA
- Start-up codes: TCP client-server programs

#### Main task

Implement a hybrid encryption algorithm called RSA with Cipher Block Chaining for securely communicating TCP client-server applications.

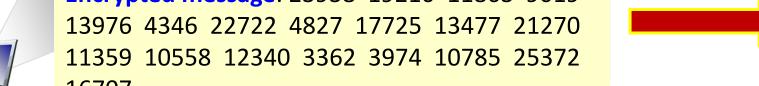
**Decrypt the encrypted** 

message: 23938 19210 11863 9019 13976 4346 22722 4827 17725 13477 21270 11359 10558 12340 3362 3974 10785 25372 16707

meet me at sunset.

plaintext message: meet me at sunset

**Encrypted message**: 23938 19210 11863 9019 16707





Server



## RSA\_CBC

#### Encryption

m – plaintext message

 $m = 'AAA \r\n'$ 

No regular pattern should emerge

RSA\_CBC (m) 
$$\rightarrow$$
 12589 50082 54789 13389 59510

ciphertext

Client (sender)

Server (receiver)

To reduce the complexity of the assignment, the scope is limited to one-way communication. The client is the sender of encrypted messages, and the server is the receiver

## Client (sender)

- has a public key (e,n), and a private key (d,n)
- has a certificate issued by a CA.
   It is denoted as dCA(e,n), it is an encrypted public key of the server

## Client (sender)

**eCA** – is an RSA public key

nonce – to be used as part of CipherBlock Chaining

- has a public key (e,n), and a private key (d,n)
- has a certificate issued by a CA.
   It is denoted as dCA(e,n), it is an encrypted public key of the server

## Client (sender)

eCA – is an RSA public key

nonce – to be used as part of CipherBlock Chaining

## Server (receiver)

- has a public key (e,n), and a private key (d,n)
- has a certificate issued by a CA.
   It is denoted as dCA(e,n), it is an encrypted public key of the server

Care must be taken when assigning values to keys and nonce:

Client's variable:

nonce

nCA > n

nonce < n

Certification Authority's keys:

eCA, dCA, nCA

Server's keys:

e, d, n

#### Sequence of steps:

## Client (sender)

eCA – is an RSA public keynonce – to be used as part of CipherBlock Chaining

TCP connection establishment

## Server (receiver)

- has a public key (e,n), and a private key (d,n)
- has a certificate issued by a CA.
   It is denoted as dCA(e,n), it is an encrypted public key of the server

1. send certificate dCA(e,n) to client

#### Sequence of steps:

## Client (sender)

eCA – is an RSA public keynonce – to be used as part of CipherBlock Chaining

2. extract the server's **public key (e,n)** using its copy of the Certification Authority's public key.

 $eCA (dCA(e,n)) \rightarrow (e,n)$ 

- has a public key (e,n), and a private key (d,n)
- has a certificate issued by a CA.
   It is denoted as dCA(e,n), it is an encrypted public key of the server
- 1. send dCA(e,n) to client

#### Sequence of steps:

## Client (sender)

eCA – is an RSA public keynonce – to be used as part of CipherBlock Chaining

2. extract the server's public key (e,n) using its copy of the Certification Authority's public key.

$$eCA (dCA(e,n)) \rightarrow (e,n)$$

3. Send "ACK 226 public key received".

- has a public key (e,n), and a private key (d,n)
- has a certificate issued by a CA.
   It is denoted as dCA(e,n), it is an encrypted public key of the server
- 1. send dCA(e,n) to client

#### Sequence of steps:

## Client (sender)

eCA – is an RSA public keynonce – to be used as part of CipherBlock Chaining

2. extract the server's public key **(e,n)** using its copy of the Certification Authority's public key.

$$eCA (dCA(e,n)) \rightarrow (e,n)$$

- 3. Send "ACK 226 public key received".
- 4. Send e(nonce).

- has a public key (e,n), and a private key (d,n)
- has a certificate issued by a CA.
   It is denoted as dCA(e,n), it is an encrypted public key of the server
- 1. send dCA(e,n) to client

#### Sequence of steps:

## Client (sender)

eCA – is an RSA public keynonce – to be used as part of CipherBlock Chaining

2. extract the server's public key (e,n) using its copy of the Certification Authority's public key.

$$eCA (dCA(e,n)) \rightarrow (e,n)$$

- 3. Send "ACK 226 public key received".
- 4. Send e(nonce).

## Server (receiver)

- has a public key (e,n), and a private key (d,n)
- has a certificate issued by a CA.
   It is denoted as dCA(e,n), it is an encrypted public key of the server
- 1. send dCA(e,n) to client

5. Extract the **nonce** using its private key, **(d,n)**. This involves the RSA decryption operation, **d(e(nonce))**. Subsequently, it sends "ACK 220 nonce ok".

Sequence of steps:

Client (sender)

Server (receiver)

#### **Communication Session**

Successive messages coming from the client will then be encrypted accordingly using RSA-CBC, and the server will decrypt the received messages using RSA-CBC.

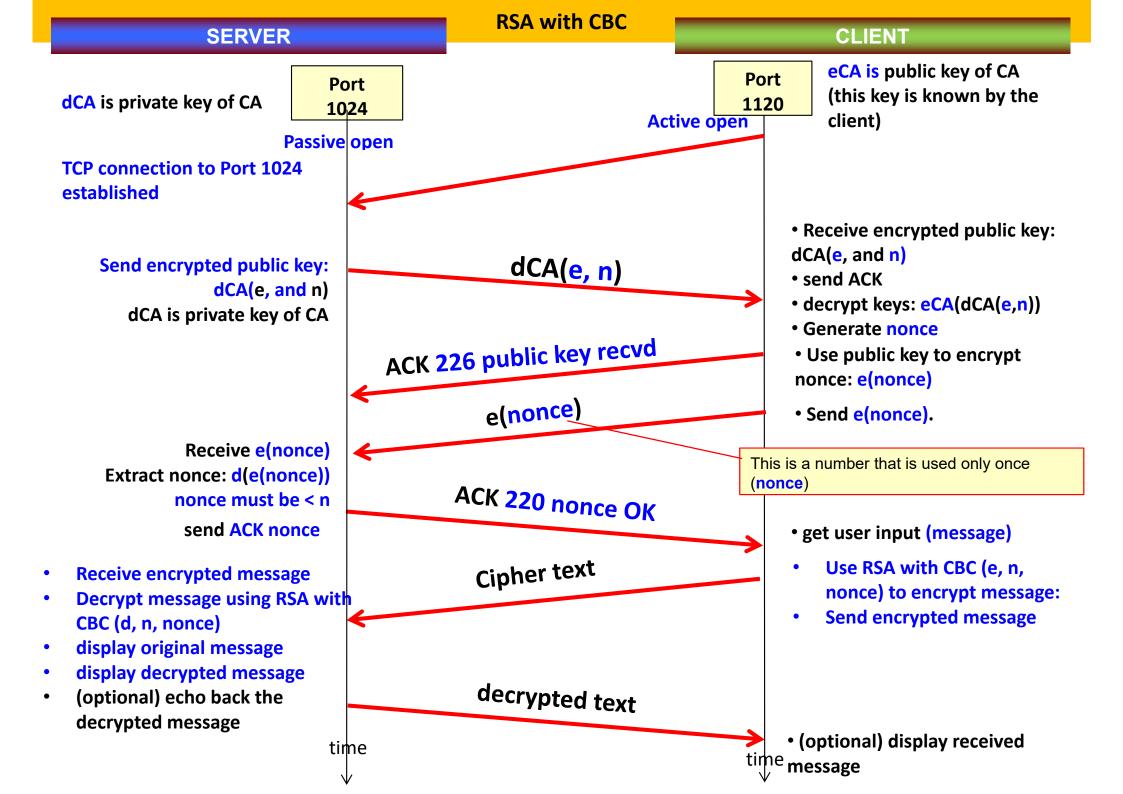
Sequence of steps:

Client (sender)

Server (receiver)

Successive messages coming from the client will then be encrypted accordingly using RSA-CBC, and the server will decrypt the received messages using RSA-CBC.

The client-server programs must be able to show the original, encrypted and decrypted messages for verification purposes.



## Demonstration

```
CLIENT
===<<< SECURE TCP CLIENT >>>===
===<< by: n.h.reyes@massey.ac.nz >>===
The Winsock 2.2 dll was initialised.
USAGE: ClientWindows IP-address [port]
Default portNum = 1234
Using default settings, IP:localhost, Port:1234
Connected to <<<SERVER>>> with IP address: localhost, IPv6 at port: 1234
Connected to <<<SERVER>>> extracted IP address: ::1, IPv6 at port: 1234
Received Server's Certificate: PUBLIC KEY 51177, 42697, scannedItems = 2
Decrypted Server's Public key: [e = 737, n = 67637]
---> Sending reply to SERVER: ACK 226 Public key received
 --> Sending Nonce to SERVER: NONCE 16150
Received packet: ACK 220 nonce ok, scannedItems = 1
nonce ACKed by Server
you may now start sending commands to the <<∢SERVER>>>.
 Enter message to send:
```

```
Decrypted server's public key (e,n)=(737, 67637)
```

```
SERVER
===<<< SECURE TCP SERVER >>>===
===<< by: n.h.reyes@massey.ac.nz >>===
The Winsock 2.2 dll was initialised.
Using DEFAULT PORT = 1234
getaddrinfo response 1
        Flags: 0x0
        Family: AF INET6 (IPv6)
        IPv6 address [::]:1234
        Socket type: SOCK STREAM (stream)
        Protocol: IPPROTO TCP (TCP)
        Length of this sockaddr: 28
        Canonical name: (null)
<><SERVER>>> is listening at PORT: 1234
A <<<CLIENT>>> has been accepted.
Connected to <<<Cli>Connected to <<<Cli>Connected to <<<Cli>Client>>> with IP address:::1, at Port:55362
Sending packet: PUBLIC KEY 51177, 42697
Received packet: ACK 226 Public key received, scannedItems = 1
Received packet: NONCE 16150 scannedItems = 1
After decryption, received nonce = 1234
nonce = 1234
Sending packet: ACK 220 nonce ok
the <<<SERVER>>> is waiting to receive commands.
```

Encrypted server's public key dCA(e,n)=(51177, 42697)

## Demonstration

```
CLIENT
===<<< SECURE TCP CLIENT >>>===
===<< by: n.h.reyes@massey.ac.nz >>===
The Winsock 2.2 dll was initialised.
USAGE: ClientWindows IP-address [port]
Default portNum = 1234
Using default settings, IP:localhost, Port:1234
Connected to <<<SERVER>>> with IP address: localhost, IPv6 at port: 1234
Connected to <<<SERVER>>> extracted IP address: ::1, IPv6 at port: 1234
Received Server's Certificate: PUBLIC KEY 51177, 42697, scannedItems = 2
Decrypted Server's Public key: [e = 737, n = 67637]
---> Sending reply to SERVER: ACK 226 Public key received
 --> Sending Nonce to SERVER: NONCE 16150
Received packet: ACK 220 nonce ok, scamedItems = 1
nonce ACKed by Server
you may now start sending commands to the <<<SERVER>>>.
-------
Enter message to send:
```

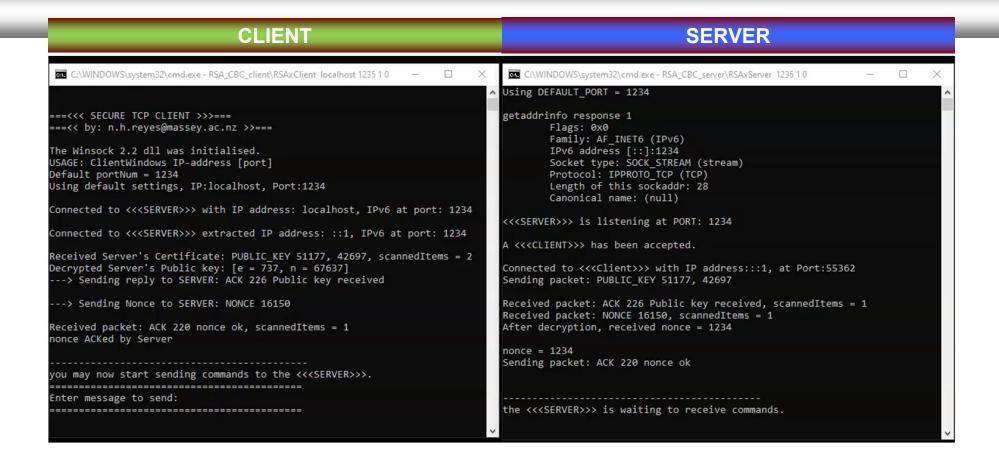
```
===<<< SECURE TCP SERVER >>>===
===<< by: n.h.reyes@massey.ac.nz >>===
The Winsock 2.2 dll was initialised.
Using DEFAULT PORT = 1234
getaddrinfo response 1
        Flags: 0x0
        Family: AF INET6 (IPv6)
        IPv6 address [::]:1234
        Socket type: SOCK STREAM (stream)
        Protocol: IPPROTO TCP (TCP)
        Length of this sockaddr: 28
        Canonical name: (null)
<<<SERVER>>> is listening at PORT: 1234
A <<<CLIENT>>> has been accepted.
Connected to <<<Cli>ent>>> with IP address:::1, at Port:55362
Sending packet: PUBLIC KEY 51177, 42697
Received packet: ACK 226 Public key received, scannedItems = 1
Received packet: NONCE 16150, scannedItems = 1
After decryption, received nonce = 1234
nonce = 1234
Sending packet: ACK 220 nonce ok
the <<<SERVER>>> is waiting to receive commands
```

**SERVER** 

Encrypted nonce: e(nonce) = 16150

decrypted nonce
nonce=1234

## **Demonstration**



## **RSA:** Choosing keys

- 1. Choose two large prime numbers p, q. (e.g., 1024 bits each)
- 2. Compute n = pq, z = (p 1)(q 1)
- 3. Choose e (with e < n) that has no common factors with z. (e, z are 'relatively prime').
- 4. Choose d such that ed -1 is exactly divisible by z. (in other words: ed mod z = 1).
- 5. Public key is (e, n). Private key is (d, n).



In mathematics, a prime number (or a prime) is a natural number that has exactly two (distinct) natural number divisors, which are 1 and the prime number itself. The first 30 prime numbers are 2, 3, 5, 7, 11, 13, 17, 19, 23, 29, 31, 37, 41, 43, 47, 53, 59, 61, 67, 71, 73, 79, 83, 89, 97, 101, 103, 107, 109, and 113

## **RSA:** Encryption, decryption

- 0. Given a public key (e, n) and a private key (d, n)
- 1. To encrypt bit pattern, m, compute

$$c = m^e \mod n$$
 (i.e., remainder when  $m^e$  is divided by  $n$ )

2. To decrypt received bit pattern, c, compute

$$m = c^d \mod n$$
 (i.e., remainder when  $c^d$  is divided by  $n$ )

Magic happens! 
$$m = (m^e \mod n)^d \mod n$$

## **RSA** example:

```
Bob chooses p=5, q=7. Then n=35, z=24.

e=5 (so e, z relatively prime).

d=29 (so ed-1 exactly divisible by z)

(equivalently, ed mod z=1).
```

encrypt:

letter

m

me

 $c = m^e \mod n$ 

12

248832

17

decrypt:

С

C

 $m = c^{d} \mod n$ 

letter

12

l

$$\underline{c}^{d}$$
 = 481968572106750915091411825223072000 - too big !! (int type)

## How to solve this problem:

## Repeated Squaring: calculate y = x e mod n

```
long repeatSquare(long x, long e, long n) {
y=1;//initialize y to 1, very important
while (e > 0) {
         if (( e % 2 ) == 0) {
                  x = (x*x) \% n;
                   e = e/2;
         else {
                   y = (x*y) \% n;
                   e = e-1;
return y; //the result is stored in y
```

```
c = m^e \mod n
```

$$m = c^d \mod n$$

Use repeatSquare for both encryption and decryption. You might want to change the datatypes used in this function to accommodate bigger numbers in the encryption scheme.



## Summary of Techniques

RSA public key (e, n) and a private key (d, n)

Check if e & z are coprimes.

Euclidean Algorithm gcd(e,z):

Solve for d

zx + ed = gcd(z,e)=1

Extended Euclidean Algorithm

Bézout's identity.

It's not very difficult to implement both Algorithms for automatically generating the RSA keys.

## Example of RSA keys

#### 1024-bit RSA encryption key (in hex format):

#### n=

A9E167983F39D55FF2A093415EA6798985C8355D9A915BFB1D01DA197026170F BDA522D035856D7A986614415CCFB7B7083B09C991B81969376DF9651E7BD9A9 3324A37F3BBBAF460186363432CB07035952FC858B3104B8CC18081448E64F1C FB5D60C4E05C1F53D37F53D86901F105F87A70D1BE83C65F38CF1C2CAA6AA7EB

**e**=010001

#### d=

67CD484C9A0D8F98C21B65FF22839C6DF0A6061DBCEDA7038894F21C6B0F8B35 DE0E827830CBE7BA6A56AD77C6EB517970790AA0F4FE45E0A9B2F419DA8798D6 308474E4FC596CC1C677DCA991D07C30A0A2C5085E217143FC0D073DF0FA6D14 9E4E63F01758791C4B981C3D3DB01BDFFA253BA3C02C9805F61009D887DB0319

https://www.boost.org/

http://www.di-mgt.com.au/rsa alg.html#realexample

You need a big number library, like **Boost** to implement values at this range



## Sample Small Keys

```
p=257 and q=293
e=529;
d=24305;
n=75301;
```

Use keys at least around the following range of values. Remember, the bigger the keys, the better! However, you need to make sure that the data type you are using can accommodate the range of values for which the algorithm will be operating on.

Given the binary representation of a plain text message (m), compute for its cipher text using CBC. Assume that the message is processed in blocks of 3 bits, and that the initial random number is equal to 010. Fill-up the following table with the details of your solution.

m = 110110110

#### 1) ENCRYPTION

i	1	2	3
m(i)			
r(i)	010 <		
c(i)			
INFO. TO SEND			

#### Table representing **K**<sub>s</sub>()

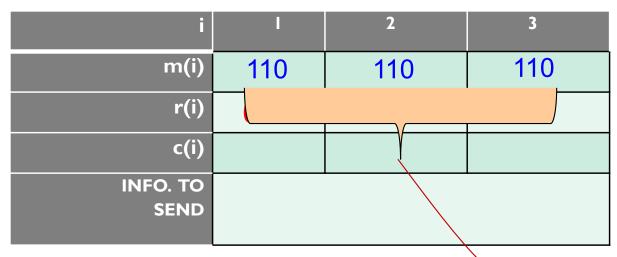
input	output	input	output
000	110	100	011
001	111	101	010
010	101	110	000
011	100	111	001

Put initial random number as the value for r(1)

Given the binary representation of a plain text message (m), compute for its cipher text using CBC. Assume that the message is processed in blocks of 3 bits, and that the initial random number is equal to 010. Fill-up the following table with the details of your solution.

m = 110110110

#### 1) ENCRYPTION



#### Table representing **K**<sub>s</sub>()

input	output	input	output
000	110	100	011
001	111	101	010
010	101	110	000
011	100	111	001

Split the original plaintext message into blocks of 3-bits

Given the binary representation of a plain text message (m), compute for its cipher text using CBC. Assume that the message is processed in blocks of 3 bits, and that the initial random number is equal to 010. Fill-up the following table with the details of your solution.

m = 110110110

#### 1) ENCRYPTION

i	1	2	3
m(i)	110	110	110
r(i)	010		
c(i)	011		
INFO. TO			
SEND			

#### Table representing **K**<sub>s</sub>()

input	output	input	output
000	110	/ 100	011
001	111	101	010
010	101	110	000
011	100	111	001

Solve for c(1):

$$K_s(m(l) \oplus r(l)) = K(110 \oplus 010) = K_s(100) = 011$$

Given the binary representation of a plain text message (m), compute for its cipher text using CBC. Assume that the message is processed in blocks of 3 bits, and that the initial random number is equal to 010. Fill-up the following table with the details of your solution.

m = 110110110

#### 1) ENCRYPTION

i	1	2	3
m(i)	110	110	110
r(i)	010	<b>011</b>	
c(i)	011		
INFO. TO			
SEND			

Table representing **K**<sub>s</sub>()

input	output	input	output
000	110	100	011
001	111	101	010
010	101	110	000
011	100	111	001

The next random number block is equal to the previous cipherblock, C(1). Therefore, copy c(1) onto r(2).

Given the binary representation of a plain text message (m), compute for its cipher text using CBC. Assume that the message is processed in blocks of 3 bits, and that the initial random number is equal to 010. Fill-up the following table with the details of your solution.

m = 110110110

#### 1) ENCRYPTION

i	1	2	3
m(i)	110	110	110
r(i)	010	011	
c(i)	011	010	
INFO. TO			
SEND			

#### Table representing **K**<sub>s</sub>()

input	output	input	output
000	110	100	011
001	111	, 101	010
010	101	110	000
011	100	111	001

Compute for c(2).

Given the binary representation of a plain text message (m), compute for its cipher text using CBC. Assume that the message is processed in blocks of 3 bits, and that the initial random number is equal to 010. Fill-up the following table with the details of your solution.

m = 110110110

#### 1) ENCRYPTION

i	1	2	3
m(i)	110	110	110
r(i)	010	011	010
c(i)	011	010	011
INFO. TO			
SEND			

Table representing **K**<sub>s</sub>()

input	output	input	output
000	110	100	011
001	111	101	010
010	101	110	000
011	100	111	001

Continue until we have computed all the cipher blocks.

Message to send:  $= 010 \ 011 \ 010 \ 011$ 

#### 2) DECRYPTION

Using your answer in Item#1 (information to send), decrypt the message using CBC. Fill-up the table with the details of your solution.

Encrypted message = 010 011 010 011

Table representing **K**<sub>s</sub>()

input	output	input	output
000	110	100	011
001	111	101	010
010	101	110	000
011	100	111	001

i	ı	2	3
c(i)	011	110	110
r(i)	010	011	110
Ks(c(i))			
m(i)			

#### 2) DECRYPTION

Using your answer in Item#1 (information to send), decrypt the message using CBC. Fill-up the table with the details of your solution.

Encrypted message = 010 011 010 011

Table representing **K**<sub>s</sub>()

input	output	input	output
000	110	100	011
001	111	101	010
010	101	110	000
011	100	111	001

i	' '	2	3
c(i)	011	110	110
r(i)	010	011	110
Ks(c(i))	100		
m(i)	110		

Compute for m(1).

$$K_{5}(011) = 100$$
 $K_{5}(011) \oplus 010 = 100 \oplus 010 = 110$ 

#### 2) DECRYPTION

Using your answer in Item#1 (information to send), decrypt the message using CBC. Fill-up the table with the details of your solution.

Encrypted message = 010 011 010 011

Table representing **K**<sub>s</sub>()

input	output	input	output
000	110	100	011
001	111	101	010
010	101	110	000
011	100	111	001

i	1	2	3
c(i)	011	110	110
r(i)	010	011	110
Ks(c(i))	100	101	100
m(i)	110	110	110

Compute for m(2) and m(3) following a similar strategy.

# Cryptography RSA with

# Cipher Block Chaining

- Combination of Public key and Symmetric key cryptography
- Avoid generating the same ciphertext for identical plaintext blocks
- Avoid sending twice the number of ciphertext bits by applying a formula to calculate the sequence of random numbers automatically, except for the very first random number.
- No need for any mini-table mapping of plaintext to ciphertext.

## **RSA with Cipher Block Chaining**

## **Encryption Example**

Plaintext message

 $m = 'AAA \n'$ 

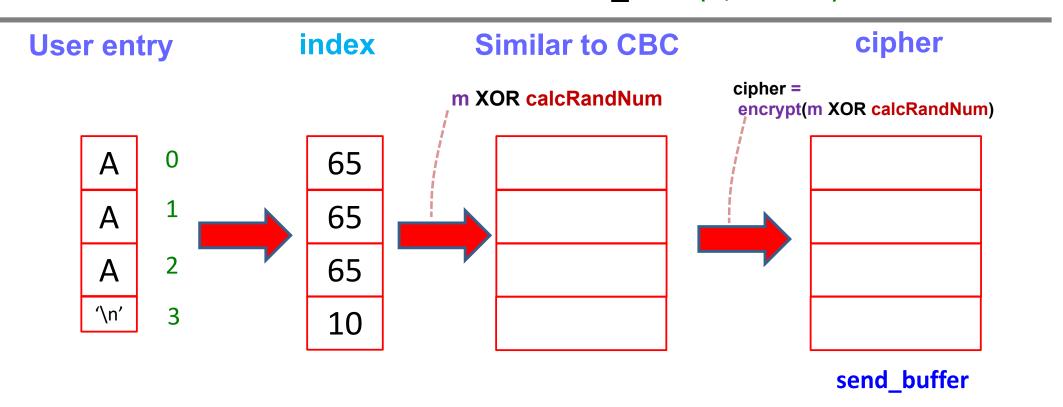
Server's PUBLIC\_KEY: (3, 25777)

Initial NONCE: 1234

Encryption is performed on the index number of the character XORed with a random number.

ENCRYPTION:  $c = m^e \mod n$ 

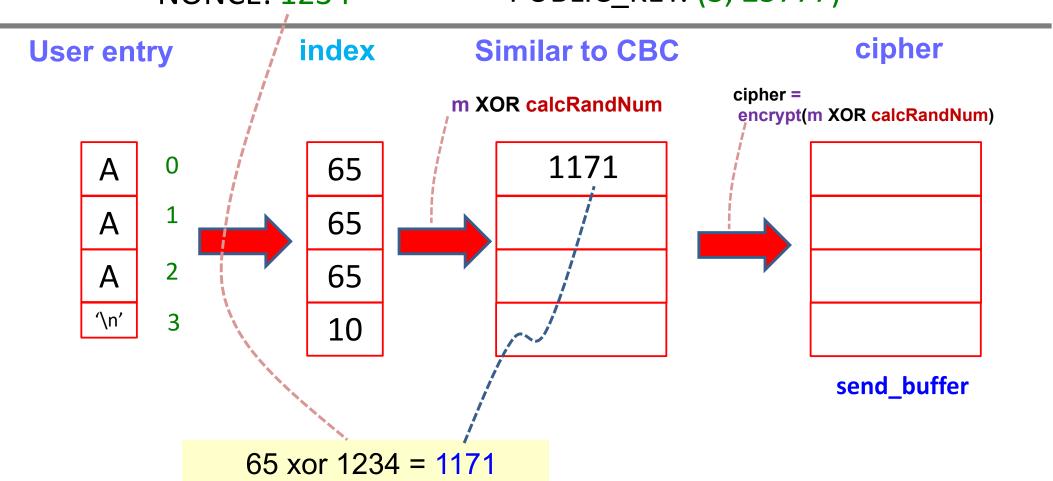
NONCE: 1234 PUBLIC\_KEY: (3, 25777)



Encryption is performed on the index number of the character XORed with a random number.

ENCRYPTION:  $c = m^e \mod n$ 

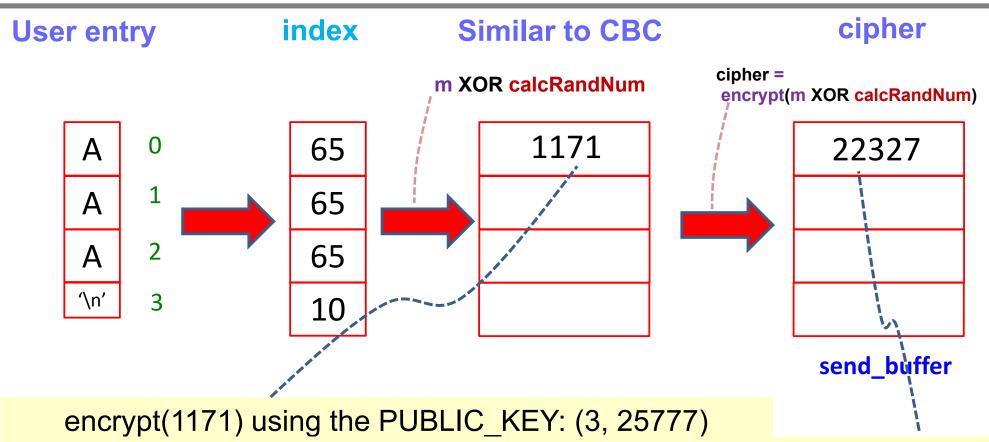
NONCE: 1234 PUBLIC\_KEY: (3, 25777)



Encryption is performed on the index number of the character XORed with a random number.

ENCRYPTION:  $c = m^e \mod n$ 

NONCE: 1234 PUBLIC KEY: (3, 25777)

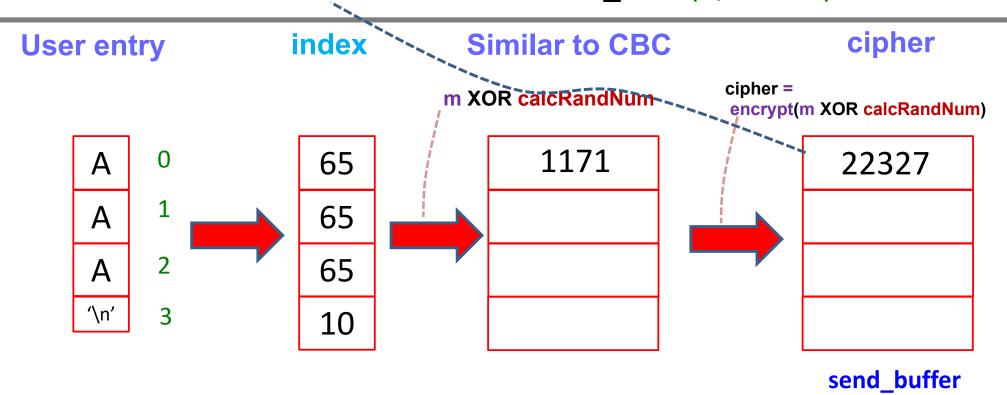


This is implemented by calling repeatSquare(1171, 3, 25777) = 22327

Encryption is performed on the index number of the character XORed with a random number.

ENCRYPTION:  $c = m^e \mod n$ 

NONCE: 22327 PUBLIC\_KEY: (3, 25777)

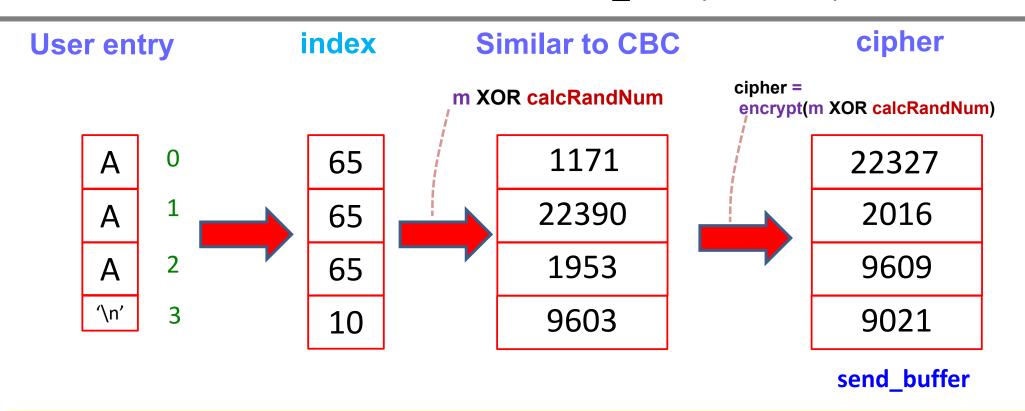


NONCE is updated to 22327 (the most recent cipher block computed).

Encryption is performed on the index number of the character XORed with a random number.

ENCRYPTION:  $c = m^e \mod n$ 

NONCE: 22327 PUBLIC\_KEY: (3, 25777)

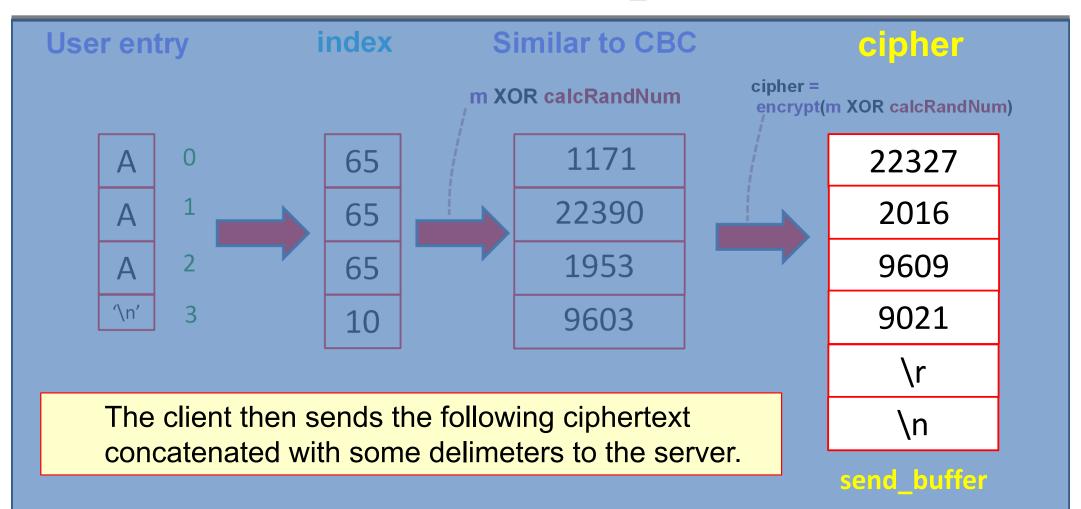


Continuing with the same sequence of operations, we get the complete ciphertext.

Encryption is performed on the index number of the character XORed with a random number.

ENCRYPTION:  $c = m^e \mod n$ 

NONCE: 22327 PUBLIC\_KEY: (3, 25777)



### **RSA with Cipher Block Chaining**

#### **Decryption Example**

Cipher

Server's PRIVATE\_KEY: (16971, 25777)

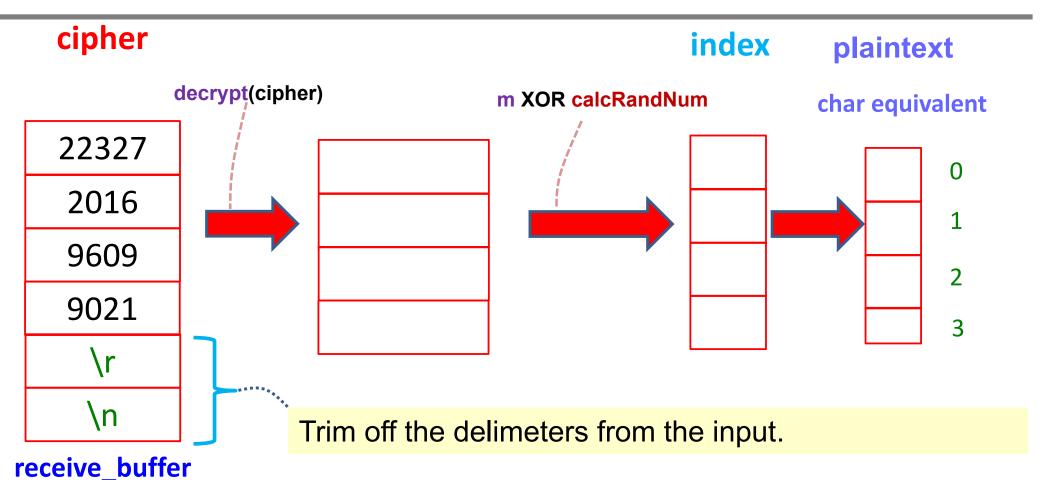
Initial NONCE: 1234

 $c = 22327 2016 9606 9021 \r\n$ 

Decryption is performed on the received cipher, then the result is XORed with a random number.

DECRYPTION:  $m = c^d \mod n$ 

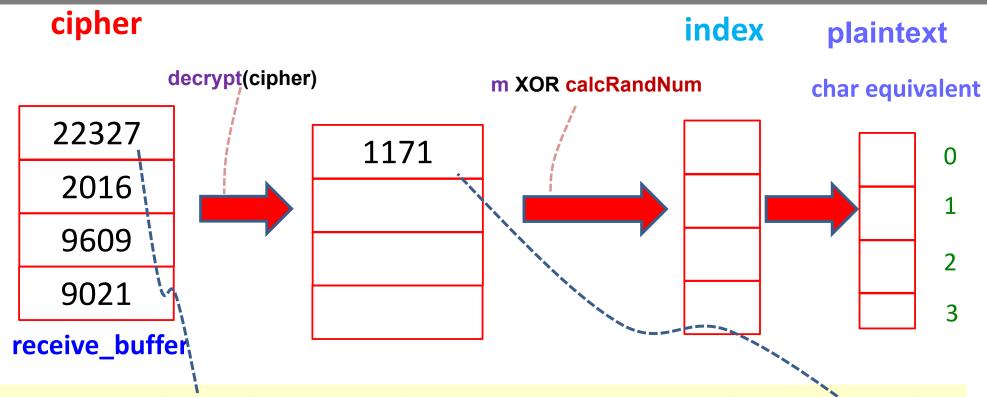
NONCE: 1234 PRIVATE\_KEY: (16971, 25777)



Decryption is performed on the received cipher, then the result is XORed with a random number.

DECRYPTION:  $m = c^d \mod n$ 

NONCE: 1234 PRIVATE\_KEY: (16971, 25777)



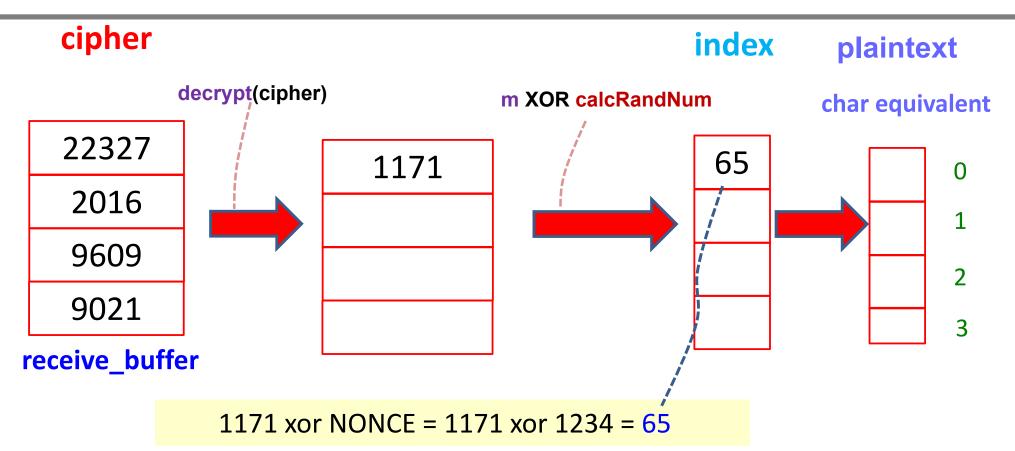
Decrypt(22327) using the Server's PRIVATE\_KEY: (16971, 25777)

This is implemented by calling repeatSquare(22327, 16971, 25777) = 1171

Decryption is performed on the received cipher, then the result is XORed with a random number.

DECRYPTION:  $m = c^d \mod n$ 

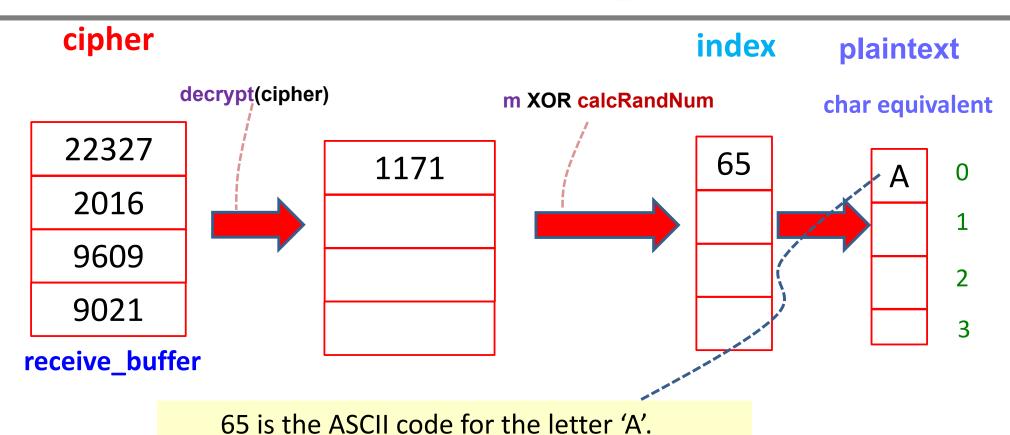
NONCE: 1234 PRIVATE\_KEY: (16971, 25777)



Decryption is performed on the received cipher, then the result is XORed with a random number.

DECRYPTION:  $m = c^d \mod n$ 

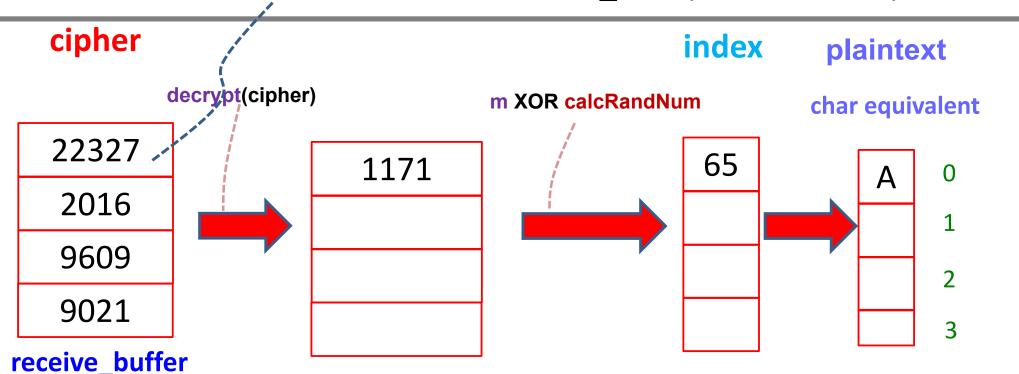
PRIVATE KEY: (16971, 25777) **NONCE: 1234** 



Decryption is performed on the received cipher, then the result is XORed with a random number.

DECRYPTION:  $m = c^d \mod n$ 

NONCE: 22327 PRIVATE\_KEY: (16971, 25777)



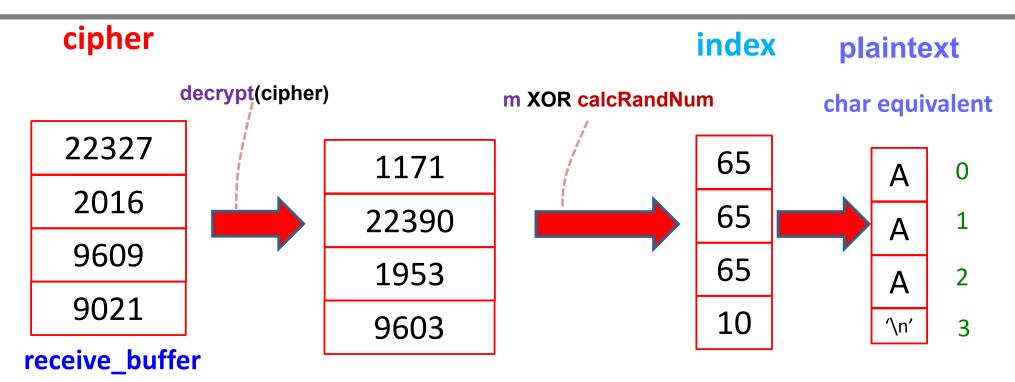
NONCE is then set to 22327 (the previous cipher block).

This will ensure that any repeating characters in the input will not generate any repeating patterns in the ciphertext.

Decryption is performed on the received cipher, then the result is XORed with a random number.

DECRYPTION:  $m = c^d \mod n$ 

NONCE: 1234 PRIVATE\_KEY: (16971, 25777)



Continue with the decryption process until all plaintext characters have been extracted.

### Checklist

Checklist (Include this in your submission)						
Please accomplish the following check list <u>in order to</u> allow for accurate marking of your assignment.						
Item		your assignment details		Comments		
1	Names and ID numbers of Group Members	your assignme	nt details	(Maximum of 3 members in a group)		
2	Operating System(s) used for testing your client and server codes	Windows 10				
3	Compiler used	g++ 11.2.0		g++ 11.2.0 is required		
4	IDE used					
5	Required Functionalities	Sending encrypted RSA public key to the client, and decryption of the public key by the client		Indicate 'full', if you have completed the implementation of the required command, 'partial', if you are only submitting a partial		
		Sending encrypted nonce to the server, and decryption of nonce by client		implementation, or ' <b>none</b> ', if not accomplished.		
		RSA with Cipher Block Chaining	full/partial/none			
6	Snap shots of sample interactions, encryption and decryption: one for each required test input.	Indicate 'full', if your codes can encrypt the following test inputs and then decrypt them correctly, snap shots must be included; 'partial', if your codes can only solve them partially (not all characters can be encrypted, then decrypted back to their original form), snap shots must be included; write 'none', if they are not solvable by your codes.				
	Input containing all letters of the alphabet	Test Input #1: the quick brown fox jumps over the lazy dog	full/partial/none	Must include snap shots		
	Input with repeating characters	Test Input #2: AAA	full/partial/none	The encrypted message must not exhibit any repeating pattern as cipher block chaining addresses this problem. Show snapshots.		
	Input with repeating characters	Test Input #3: 555	full/partial/none	The encrypted message must not exhibit any repeating pattern as cipher block chaining addresses this problem. Show snapshots.		

### Checklist

#### **Optional**

7 Extra work done (Max. bonus of 2 marks)	Program uses the Boost library's big number facility. Very large public key and	Yes/No
marks)	private keys are being used.	
	Include instructions on how to install the big compilation instructions.	number library you used and the
	Write CA's large Public key:	
	Write CA's large Private key:	
	Write Server's large Public key:	
	Write Server's large Private key:	

# Other useful notes

## Reading characters from stdin

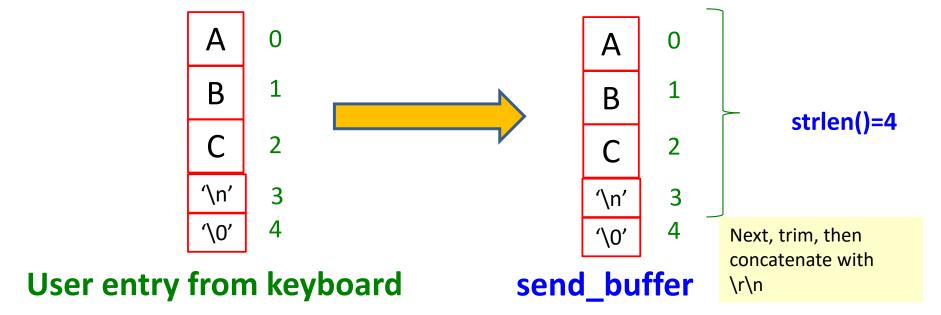
get string from stream

**CLIENT** 

#### char\* fgets(char\* send\_buffer, int num, FILE \*stream)

- reads characters from stream until (num-1) characters have been read or until it encounters:
  - a new line character (copied into send\_buffer)
  - a NULL-termination character ('\0') is automatically appended
  - if an error occurs, a NULL pointer is returned

strlen() - counts the number of characters excluding the NULL-character



### String Tokenizer

#### Ŷ

#### **Example**

```
1 /* strtok example */
2 #include <stdio.h>
3 #include <string.h>
 5 int main ()
6
    char str[] ="- This, a sample string.";
   char * pch;
   printf ("Splitting string \"%s\" into tokens:\n",str);
pch = strtok (str," ,.-");
    while (pch != NULL)
12
13
      printf ("%s\n",pch);
      pch = strtok (NULL, " ,.-");
14
15
16
    return 0;
17 }
```

## The End

#### **Advanced Reading**

- http://www.di-mgt.com.au/rsa\_alg.html#KALI93
- The Handbook of Applied Cryptography
  - http://cacr.uwaterloo.ca/hac/