

# Francis's QR $\equiv$ QR

Let  $A$  be properly upper Hessenberg.

## Shifted QR

Set  $A_0 = A$

for  $j = 0, 1, \dots$

(i) Compute  $p_j(A_j)$ .

(ii) Find reflectors  $Q_j^{(1)}, \dots, Q_j^{(n-1)}$  such that  $Q_j^{(n-1)} \dots Q_j^{(1)} p_j(A_j) = R_j$  is upper triangular

(iii) Compute

$$A_{j+1} := Q_j^{(n-1)} \dots Q_j^{(1)} A_j Q_j^{(1)} \dots Q_j^{(n-1)}.$$

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(i) Compute  $p_j(A_j)e_1$ .

(ii) Find reflector  $Q_j^{(1)}$  such that  $Q_j^{(1)} p_j(A_j)e_1 = \alpha_j e_1$

(iii) Compute  $Q_j^{(1)} A_j Q_j^{(1)}$ .

(iv) Find reflectors  $\hat{Q}_j^{(2)}, \dots, \hat{Q}_j^{(p)}$  such that  $A_{j+1} := \hat{Q}_j^{(p)} \dots \hat{Q}_j^{(2)} Q_j^{(1)} A_j Q_j^{(1)} \hat{Q}_j^{(2)} \dots \hat{Q}_j^{(p)}$  is upper Hessenberg.

Shifted QR finds  $Q_j := Q_j^{(1)} \dots Q_j^{(n-1)}$  such that  $p_j(A_j) = Q_j R_j$  is a QR decomposition of  $p_j(A_j)$  and sets  $A_{j+1} = Q_j^* A_j Q_j$ .

But Francis's Shifted QR **also** finds a QR decomposition  $p_j(A_j) = \tilde{Q}_j \tilde{R}_j$  and sets  $A_{j+1} = \tilde{Q}_j^* A_j \tilde{Q}_j$  where

$$\tilde{Q}_j := Q_j^{(1)} \hat{Q}_j^{(2)} \dots \hat{Q}_j^{(p)} \text{ and } \tilde{R}_j := \alpha_j K(A_{j+1}, e_1) [K(A_j, e_1)]^{-1}.$$

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The iterates  $A_j$  are symmetric tridiagonal if  $A$  is Hermitian or real symmetric and upper Hessenberg otherwise.



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↪ Finding  $v$  requires forming  $T$ !



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- ▶ If  $A$  is not Hermitian, solve upper triangular systems  $Tv = \lambda v$  for  $v$  and compute  $Qv$ .
- ▶ If  $A$  is real, arrange all computations to stay in real arithmetic till the very end.

# Implicit QR: One of the top 10 algorithms of the 20th century

