

# nmi

2024 spring : lecture 02 : revisit rounding error wrt stability

# theorem 01 wrt bound, aggregate error

suppose  $i = 1, \dots, n$  and  $0 < \delta_i \leq \mu_M$  and  $e_i \in \{-1, +1\}$ . additionally suppose  $n\mu_M < 1$ . then

$$\prod^n (1 + \delta_i)^{e_i} = 1 + \Theta_n,$$

where  $|\Theta_n| \leq \Upsilon_n = n\mu_M / (1 - n\mu_M)$ . **ie,  $\Theta_n$  aggregates error and  $\Upsilon_n$  is its bound.**

note:  $\mu_M$  is rounding error. in FPS,  $\mu_M = \frac{1}{2} \epsilon_M$ , machine error.

proof-lite. (its just a sketch.)

$$\prod^n (1 + \delta_i)^{e_i} \leq \prod^n (1 + \delta_i) \leq \prod^n (1 + n\mu_M) = (1 + n\mu_M)^n.$$

then by binomial theorem  $\Rightarrow$

$$(1 + n\mu_M)^n \leq n\mu_M / (1 - n\mu_M).$$

more intensely:

[Formal Proofs of Rounding Error Bounds](#)

# eg 01

forward stability

forward stability. ie,  $\eta > 0$  such that  $\|\Delta y\| \leq \eta \|y\|$ .

$$\Delta y = \hat{y} - y$$

$$\phi : \mathbb{R} \rightarrow \mathbb{R}, x \mapsto x^3$$

$$\hat{\phi} : \mathbb{F} \rightarrow \mathbb{F}, \hat{x} \mapsto \hat{x}^3 \Rightarrow \hat{y} = \hat{\phi}(\hat{x}) = \hat{x}^3$$

$$\begin{aligned}\hat{y} = \hat{x}^3 &= (\hat{x} \otimes \hat{x}) \otimes \hat{x} \\ &= (x(1 + \delta_x) \otimes x(1 + \delta_x)) \otimes (x(1 + \delta_x)) \\ &= x^2(1 + \delta_x)^2(1 + \delta_{\otimes}) \otimes (x(1 + \delta_x)) \\ &= x^3(1 + \delta_x)^3(1 + \delta_{\otimes})^1(1 + \delta_{\otimes})^1 \\ &= x^3(1 + \theta_5) \text{ [theorem 01]}\end{aligned}$$

$$\hat{y} - y = x^3(1 + \theta_5) - x^3 = x^3\theta_5$$

$$|\hat{y} - y| = |x^3||\theta_5| \leq |y|\gamma_5 \text{ [theorem 01]}$$

$$\gamma_5 = \frac{n\mu_M}{1-n\mu_M} = \frac{5\mu_M}{1-5\mu_M}.$$

let  $\eta = \gamma_5$ , then  $|\hat{y} - y| \leq \eta |y|$ . ✓

# eg 02 numerical stability

numerical stability. ie,  $\eta > 0$  such that  $|\Delta y| \leq \eta|y|$  and  $\epsilon > 0$  such that  $|\Delta x| \leq \epsilon|x|$ .

$y = \phi(x) = x^3$ . need  $\Delta x, \Delta y$ .

note:  $\hat{y} = x^3(1 + \theta_5)$ ,  $\gamma_5$  from eg 01.

$$\begin{aligned}\hat{y} + \Delta y &= \phi(x + \Delta x) \\ x^3(1 + \theta_5) + \Delta y &= (x + \Delta x)^3 = x^3 + 3x^2(\Delta x) + 3x(\Delta x)^2 + (\Delta x)^3 \\ \Rightarrow \Delta y &= 3x^2(\Delta x) + 3x(\Delta x)^2 + (\Delta x)^3 - x^3\theta_5\end{aligned}$$

what  $\eta$  for  $|\Delta y| \leq \eta|y|$ ?

$\Rightarrow$  what  $\epsilon$  for  $|\Delta x| \leq \epsilon|x|$ ?

$$\begin{aligned}|\Delta y| &\leq 3|x|^2|\Delta x| + 3|x||\Delta x|^2 + |\Delta x|^3 + |x|^3\theta_5 \\ &\leq 3|x|^2\epsilon|x| + 3|x|\epsilon^2|x|^2 + \epsilon^3|x|^3 + |x|^3\gamma_5 \\ &\Rightarrow 3\epsilon|x|^3 + 3\epsilon^2|x|^3 + \epsilon^3|x|^3 + |x|^3\gamma_5 \Rightarrow (3\epsilon + 3\epsilon^2 + \epsilon^3 + \gamma_5)|y|\end{aligned}$$

$$\Rightarrow |\Delta y| \leq \eta|y| \Rightarrow \eta = 3\epsilon + 3\epsilon^2 + \epsilon^3 + \gamma_5$$

ie, take any  $\epsilon > 0$ , any  $\Delta x$  such that  $|\Delta x| \leq \epsilon|x|$ . then let  $\eta = 3\epsilon + 3\epsilon^2 + \epsilon^3 + \gamma_5$ ,  $\Delta y = 3x^2(\Delta x) + 3x(\Delta x)^2 + (\Delta x)^3 - x^3\theta_5$ .

by construction (ie, by choice), we know  $|\Delta y| \leq \eta|y|$ ,  $\hat{y} + \Delta y = \phi(x + \Delta x)$ . ✓

ie,  $\hat{\phi}(x) = \hat{x}^3$  is numerically stable.

also, "any"  $\epsilon$  vs "the best, smallest"  $\epsilon$ ,  $\eta$ . note that the only part of the derived  $\eta$  that you cannot control is  $\gamma_5$ , so choose  $\epsilon \ll 1$  and  $\eta$  will barely above  $\gamma_5$ . ie,  $\eta = \gamma_5 + c_0$  where  $c_0 \geq 1$ .

# nmi

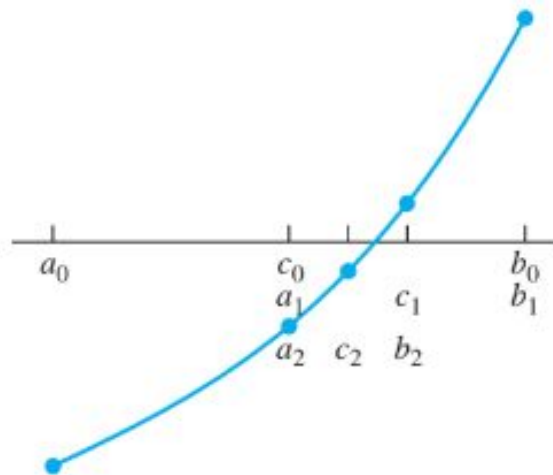
2024 spring : lecture 03 : bisection, fixed-point iteration

# root finding, bisection

given  $f(x) = x^3 + x - 1$ , find its root on the interval  $[0, 1]$ . let  $a, b$  represent the endpoints of the interval considered and  $c$  as its bisector,  $\frac{a+b}{2}$ .

a	f(a)	b	f(b)	c	f(c)
0	-1	1	1	$\frac{1}{2}$	$-\frac{3}{8}$
$\frac{1}{2}$	$-\frac{3}{8}$	1	1	$\frac{3}{4}$	+
$\frac{1}{2}$	$-\frac{3}{8}$	$\frac{3}{4}$	+	...	...

how will you stop this function? when is  $x$  good enough? or is when  $f(x)$  is good enough?  
thats on you! an error less than  $0.5 \times 10^{-p}$  for precision of  $p$  decimal places.



# root finding, fixed-point iteration, usw

$$x_{n+1} = f(x_n) \text{ with } n=0,1,2,\dots$$

$$\Rightarrow x_0, x_1, x_2, \dots = x_0, f(x_0), f(f(x_0)), \dots$$

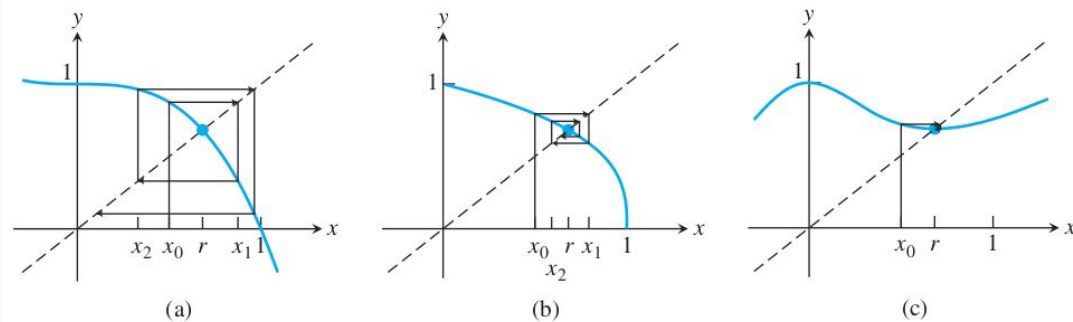
$$\Rightarrow f(x_{\text{fix}}) = x_{\text{fix}}.$$

eg, rearrange the previous  $f(x) = x^3 + x - 1$

$$\Rightarrow x = 1 - x^3 = g(x)$$

$$\Rightarrow x_1 = g(x_0), x_2 = g(x_1), x_3 = g(x_2), \dots$$

image right shows cobweb pattern of convergence.



**Figure 1.3 Geometric view of FPI.** The fixed point is the intersection of  $g(x)$  and the diagonal line. Three examples of  $g(x)$  are shown together with the first few steps of FPI. (a)  $g(x) = 1 - x^3$  (b)  $g(x) = (1 - x)^{1/3}$  (c)  $g(x) = (1 + 2x^3)/(1 + 3x^2)$

and heres another fpi, newton-raphson.

$$x_{n+1} = x_n - f(x_n)/f'(x_n) = g(x_n).$$

# fpi, extended

for the first part of this revisitation, lets re-rig that failed example.

start by considering that  $f(x) = x^3 + x - 1$  has three roots. (thats three roots - not necessarily three distinct roots.) so

$$f(x) = 0 \Rightarrow x = 1 - x^3 = g(x) \Rightarrow x_{n+1} = 1 - x_n^3.$$

thats the one that barfs, so lets look at what happens when  $x^3$  is on the left side of that equation - in colab... *refer to blackboard for python notebook, demo\_02\_rootfinding.*

what makes those last two methods increasingly better? #2 addresses that there are three roots and #3 reduces the order of  $g(x)$ .



next time

newtons method

secant method