

Recommendations for Streaming Data

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Introduction

- Most current recommender systems are designed in the context of offline setting.
- It is desirable to provide real-time recommendations in large-scale scenarios.
- Some applications: social networks, movie/book suggestions, dating.

Challenges

- Changing user preferences:
 - New items keep appearing.
 - The underlying user patterns keep changing.
 - This causes the recommendations to vary with time.
- In-core memory for memory-resident operations is quite limited.
- Classical methods like neighborhood-based and latent factor models have shortcomings.
 - They require a computationally expensive offline phase.
 - Factorizing large matrices is cumbersome when the said matrices are rapidly changing with time.

The Setup

- The ratings are received in the format: <userID, itemID, rating>.
- If rating is drawn from $\{-1, +1\}$, then let users who have given a rating of $+1$ to item i at time t be represented by $P(i, t)$ and -1 by $N(i, t)$.
- Since exact similarity computation is intensive, we compute it probabilistically by imposing a sort order on the users with the help of hash functions.
 - Use d mutually independent hash functions.
 - Each hash function takes in an identifier of a user and outputs a random number uniformly distributed in $(0, 1)$.

Probabilistic Similarity

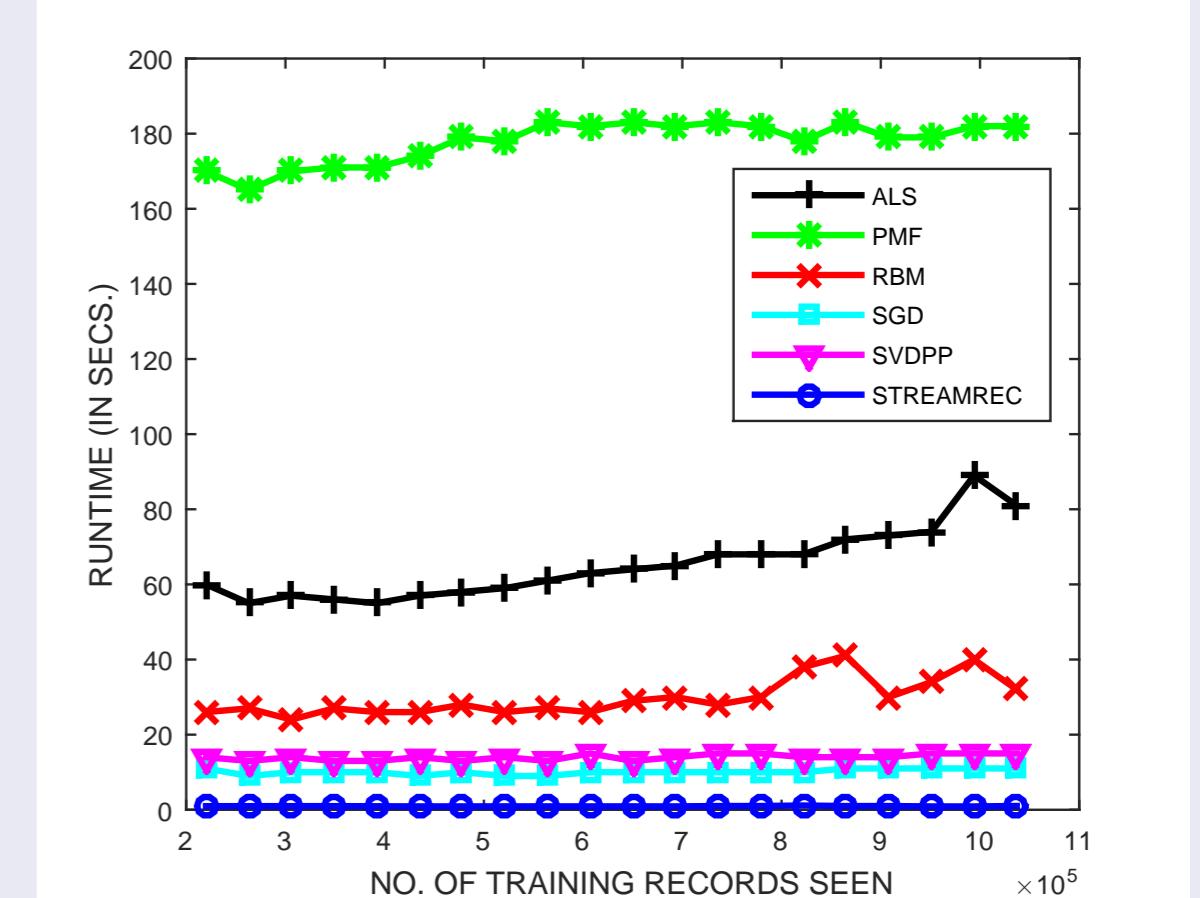
- Now, for a given sort order, what is the probability that the first user with a positive rating for item i is the same as the first user with a positive rating for item j ?
- This is the probability that both i and j take on the value $+1$ when at least one of them takes on the value of $+1$ given by:

$$\frac{P(i, t) \cap P(j, t)}{P(i, t) \cup P(j, t)}$$
- The above expression represents the similarity between items i and j with respect to positive ratings.

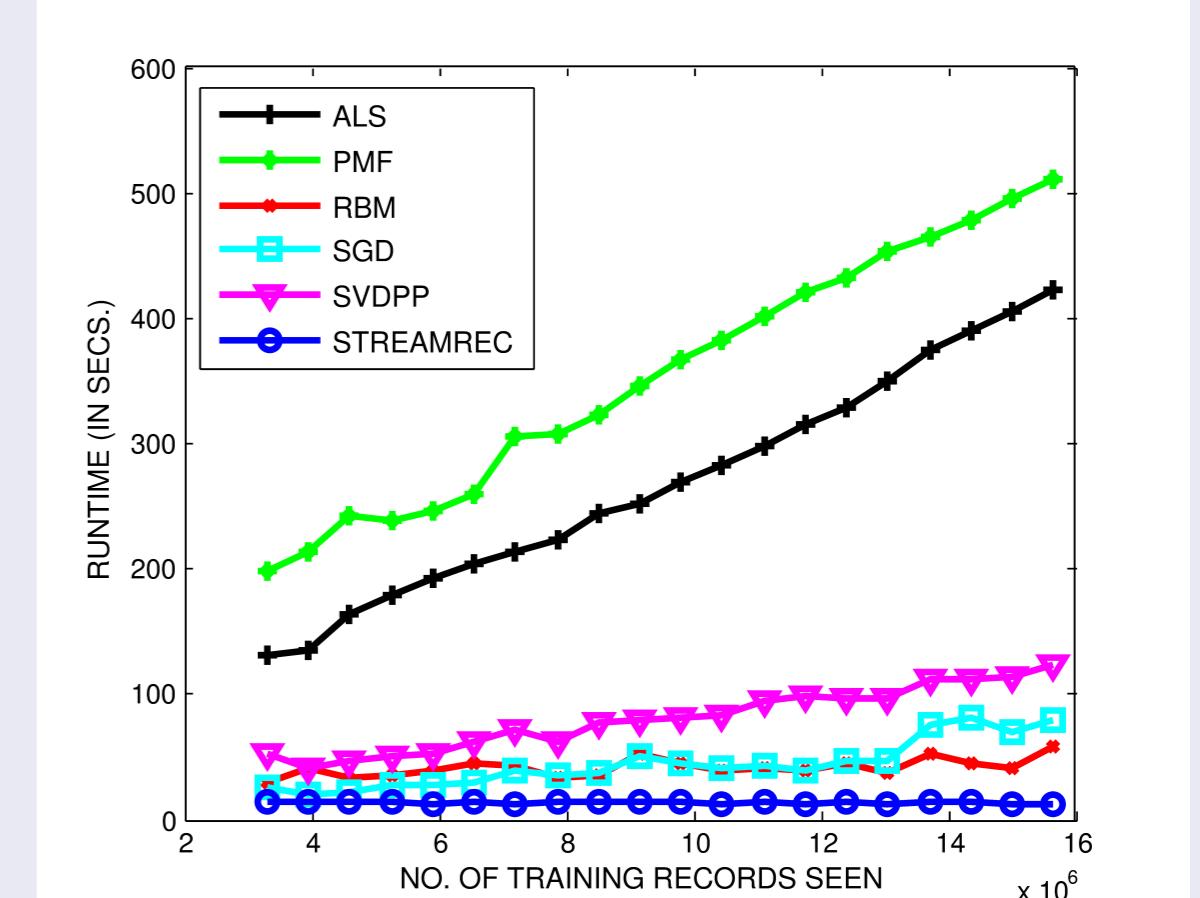
Probabilistic Similarity

- The d mutually independent hash functions are applied to the user indices that have rated item j positively.
- For each hash function, the least hash value (*min-hash value*) among these positive users and the corresponding user index (*min-hash index*) are maintained.
- d of such pairs are maintained for the n items seen which are easily updatable. This drastically reduces memory requirements.
- Similarly, the process is repeated for negatively rated items. The system can be extended for scenarios with multiple ratings as well.

Results



(a) Books Runtime



(b) Dating Runtime

Figure: Runtime

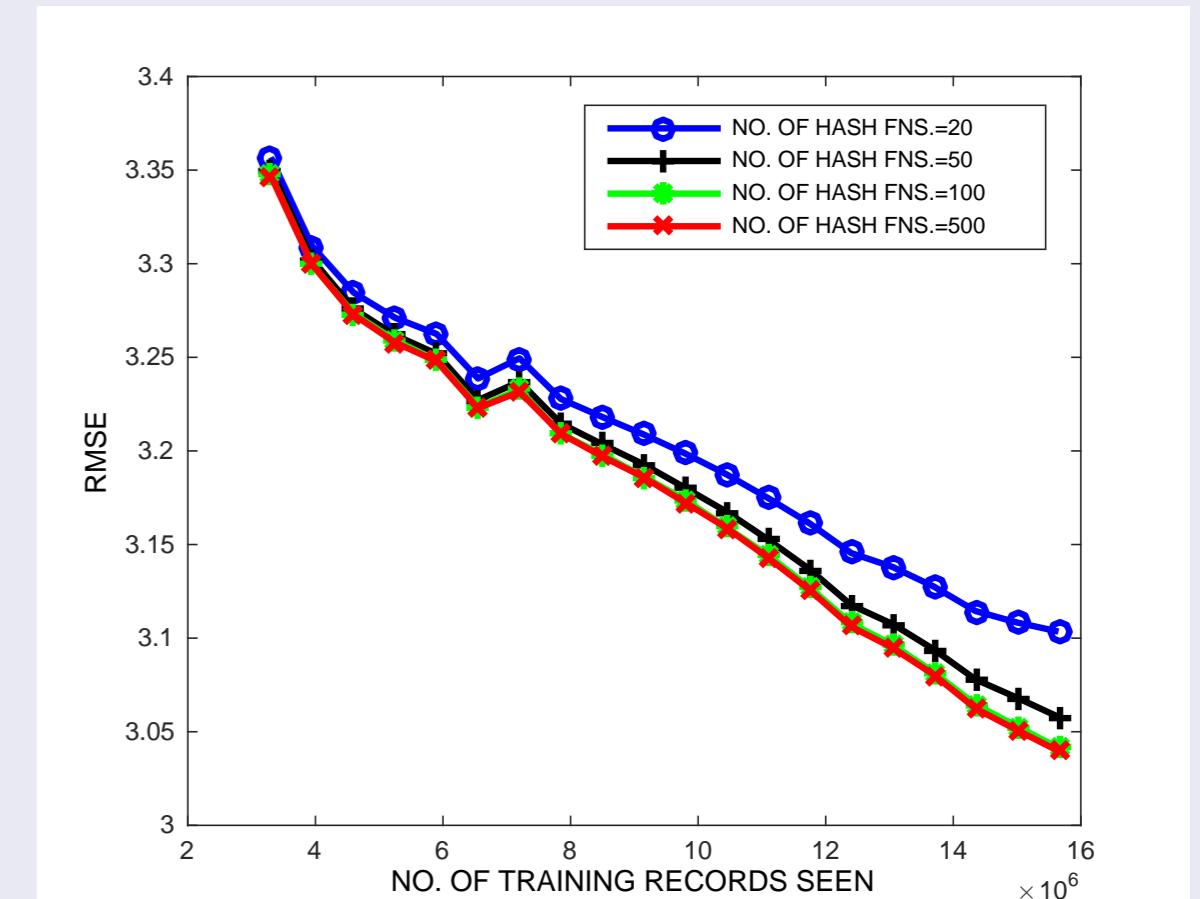
Probabilistic Similarity

- How good is the quality of the similarity measure computed in this manner?
- Let $R^+(i, j, t)$ be an approximation to the Jaccard coefficient computed by the min-hash approach and $S^+(i, j, t)$ be the actual value. Then, we prove the following:
 - Lower Tail Bound: For any $\epsilon \in (0, 1)$, $R^+(i, j, t)$ lies outside $S^+(i, j, t)$ by a factor of $(1 - \epsilon)$ with the probability:

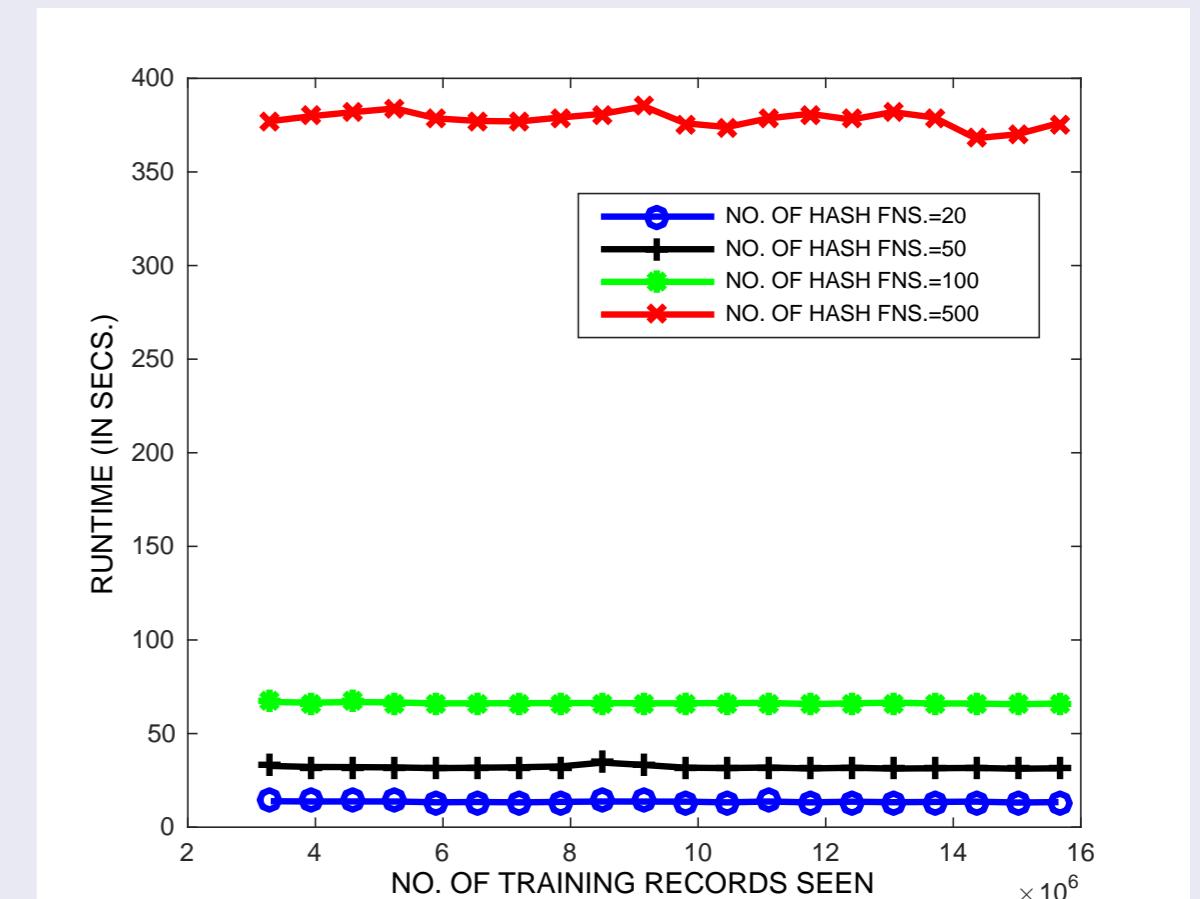
$$P(R^+(i, j, t) < (1 - \epsilon) \cdot S^+(i, j, t)) \leq \exp(-d \cdot S^+(i, j, t) \cdot \epsilon^2 / 2)$$
 - Upper Tail Bound: For any $\epsilon \in (0, 2 \cdot e - 1)$, $R^+(i, j, t)$ lies outside by a factor of $(1 + \epsilon)$ with the probability:

$$P(R^+(i, j, t) > (1 + \epsilon) \cdot S^+(i, j, t)) \leq \exp(-d \cdot S^+(i, j, t) \cdot \epsilon^2 / 4)$$

Results



(a) Sensitivity Dating RMSE



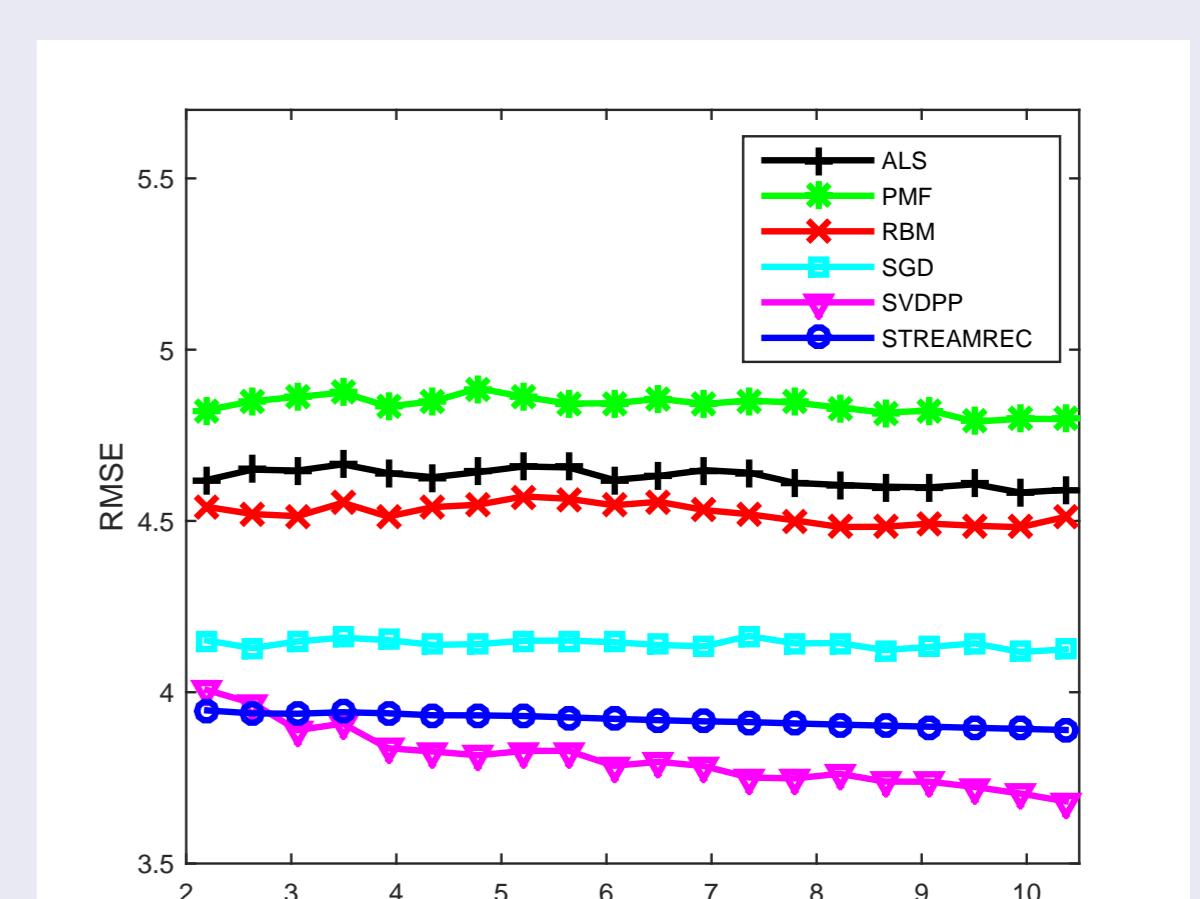
(b) Sensitivity Dating Runtime

Figure: Sensitivity

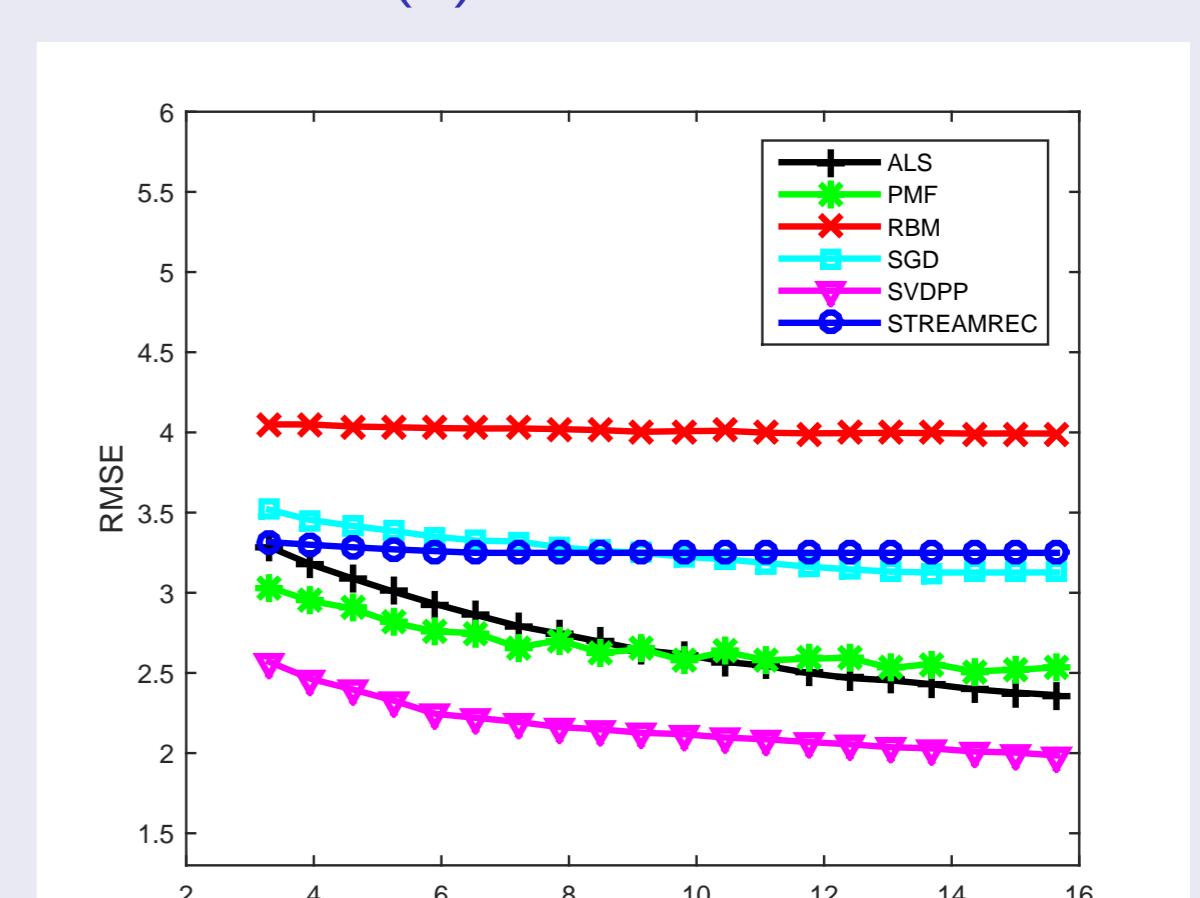
Summary

- We address the problem of providing recommendations in a streaming setting.
- We provide an efficient algorithm to perform streaming recommendations using a probabilistic model.
- The proposed algorithm stores the rating matrix compactly and hence memory requirements are low.
- Our thorough experiments show that the proposed method performs as good or better than the state-of-the-art.

Results



(a) Books RMSE



(b) Dating RMSE

Figure: Efficiency