PROGRAMMING ASSIGNMENT - 2

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QUESTION - 1:

❖ Assumptions:

- > Human player always moves first
- > Human has been allocated "x"
- ➤ Human player doesn't take <u>any</u> invalid moves, i.e., once a cell is occupied, the human player doesn't select that cell again.

Methodology for MINIMAX :

- ➤ After, the human takes the turn, the computer makes the moves using MINIMAX algorithm.
- > The algorithm is as follows:

```
 \begin{array}{l} \textbf{function } \ \mathsf{MINIMAX-DECISION}(state) \ \textbf{returns} \ an \ action \\ \textbf{return } \ \mathrm{arg\,max}_{a \ \in \ \mathsf{ACTIONS}(s)} \ \mathsf{MIN-VALUE}(\mathsf{RESULT}(state, a)) \end{array}
```

```
function MAX-VALUE(state) returns a utility value if Terminal-Test(state) then return Utility(state) v \leftarrow -\infty for each a in Actions(state) do v \leftarrow \text{Max}(v, \text{Min-Value}(\text{Result}(s, a))) return v
```

```
function MIN-VALUE(state) returns a utility value

if TERMINAL-TEST(state) then return UTILITY(state)

v \leftarrow \infty

for each a in ACTIONS(state) do

v \leftarrow \text{MIN}(v, \text{MAX-VALUE}(\text{RESULT}(s, a)))

return v
```

♦ Methodology for ALPHA-BETA:

- ➤ After, the human takes the turn, the computer makes the moves using MINIMAX algorithm.
- > The algorithm is as follows:

```
function ALPHA-BETA-SEARCH(state) returns an action
   v \leftarrow \text{MAX-VALUE}(state, -\infty, +\infty)
   return the action in ACTIONS(state) with value v
function MAX-VALUE(state, \alpha, \beta) returns a utility value
   if TERMINAL-TEST(state) then return UTILITY(state)
   v \leftarrow -\infty
   for each a in ACTIONS(state) do
      v \leftarrow \text{MAX}(v, \text{MIN-VALUE}(\text{RESULT}(s, a), \alpha, \beta))
     if v \geq \beta then return v
      \alpha \leftarrow \text{MAX}(\alpha, v)
   return v
function MIN-VALUE(state, \alpha, \beta) returns a utility value
   if TERMINAL-TEST(state) then return UTILITY(state)
   v \leftarrow +\infty
   for each a in ACTIONS(state) do
      v \leftarrow \text{MIN}(v, \text{MAX-VALUE}(\text{RESULT}(s, a), \alpha, \beta))
     if v \leq \alpha then return v
     \beta \leftarrow \text{MIN}(\beta, v)
   return v
```

TIC TAC TOE RESULTS

MINIMAX:

```
11 11 1
 11 11×1
Time taken to make this move = 0.12117362022399902 seconds
SOPHIA'S MOVE
  | | | | | |
  11 11x1
Play your move : 3
| || ||x|
 11011 1
  11 11×1
Time taken to make this move = 0.005606412887573242 seconds SOPHIA'S MOVE \mid \ \mid
Play your move : 4
|x||o||o|
  11 11×1
Time taken to make this move = 0.0004973411560058594 seconds SOPHIA'S MOVE |o|| ||x||
|x||0||0|
 11 11x1
Play your move : 7
|o|| ||x|
IXII IIXI
Time taken to make this move = 7.867813110351562e-05 seconds SOPHIA'S MOVE |O|| ||X|
|x||0||0|
```

```
Play your move : 2
|o||x||x|
|-------
|x||o||o|
|-------
|x||o||x|
|------
```

ALPHA-BETA:

```
Time taken to make this move = 0.014664411544799805 seconds
50PHIA'S MOVE
| || || |
 11011 1
 II IIxI
Play your move : 3
 11011 1
lime taken to make this move = 0.0024344921112060547 seconds
SOPHIA'S MOVE
| || ||x|
 llollol
 II II×I
Play your move : 4
| || ||x|
|x||o||o|
 II IIxI
Time taken to make this move = 0.00044608116149902344 seconds SOPHIA'S MOVE |0|| ||x||
 II IIxI
Play your move : 7
|o|| ||x|
Time taken to make this move = 9.1552734375e-05 seconds SOPHIA'S MOVE |o|| ||x||
|x||o||o|
|x||o||x|
```

```
Play your move : 2
|o||x||x|
------
|x||o||o|
------
|x||o||x|
------
```

♦ Observation:

- There's quite a visible difference in the time taken to make the first move between minimax and alpha-beta.
- > However, it gets decreased with the advancement of the game.

♦ Inference :

- The number of states that minimax visits for the initial moves is quite large as compared to alpha-beta, hence, the large amount of time for making the initial moves
- ➤ However, with the game's advancement, this number gets reduced and both perform almost equally towards the end.

GENETIC, MEMETIC, CSP

QUESTION - 2:

❖ Assumptions :

- > Professors have been mapped to courses uniformly.
- > The total number of slots to fill in a week is a hyper parameter.
- > Number of lecture halls = 4
- > Number of professors = 20
- > Number of courses = 40
- > Number of days/week = 5
- > Number of slots/day = 8

♦ Logical Constraints:

- A professor cannot be present in two lecture halls for a particular day during the same slot in a chromosome.
- Two courses cannot be allotted the same lecture hall for a particular day during the same slot.

GENETIC ALGORITHM

Methodology :

- > Firstly, I have mapped all the courses to professors uniformly.
- > Then, randomly created a gene having the given structure.

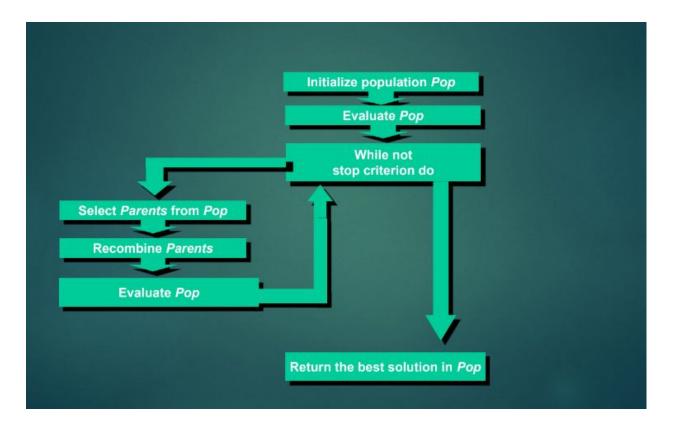


Fig. Structure of a single gene

- > Then created a chromosome of size = 8 * 5 * (number of lecture halls).
- > The chromosome is 1-D array of genes.
- Now, similarly, we create a random population of genes.
- ➤ Then calculate the fitness of each chromosome in the population, and sort in non-decreasing order.
- ➤ Then selected top k chromosomes for crossover and created another population.
- Now, merged both, the initial as well as the crossover population.
- > Then calculated its fitness and again sorted.
- > Then, cut the population size to 100.
- > Performed mutation on these 100 chromosomes.
- Now, added the top- k performing parents from the unmutated population to this new population.
- Make this the new population and repeat the process for it.

- ◆ FITNESS EVALUATION (same for both GA and MA):
 - ➤ Initially, the fitness score is zero.
 - > For each conflict in a chromosome the score is increased by 1.

❖ Algorithm:



Aspects of the process followed :

- ➤ Genetic Algorithm :

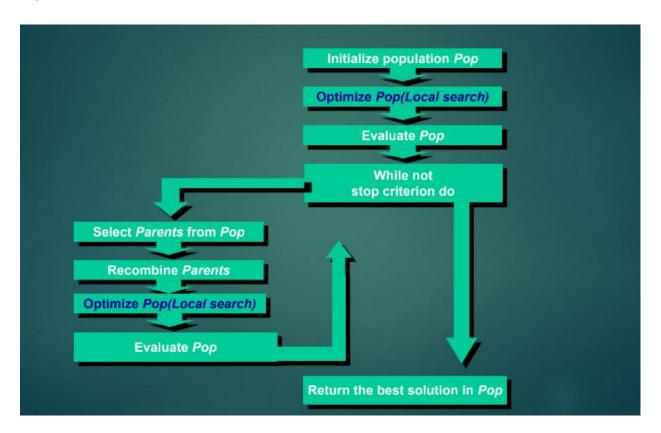
 - **Chromosome generation**: Whatever random genes are formed, they are added to form a chromosome.
 - **Population generation**: A fixed number of chromosomes are generated using the above method for creating a population.
 - **Crossover methods**: Single-point crossover is being used by randomly generating the crossover point.
 - Evaluation Function : A +1 is added to the fitness value for each conflict.
 - **Mutation method**: A fixed number of genes are chosen at random and replaced with new ones.

MEMETIC ALGORITHM

❖ Methodology :

- > The generation of gene and chromosome is the same as for genetic algorithm.
- ➤ However, initially, only 10 chromosomes are being generated.
- ➤ If there are any conflicting genes in these chromosomes, they are replaced with new ones.
- > The process doesn't move forward unless there's some improvement in its fitness.
- > After this, crossover is done to form a population.
- > This population is then sorted according to the fitness values of its chromosomes.
- ➤ And then again, top 10 chromosomes are selected for the process to repeat.

Algorithm:



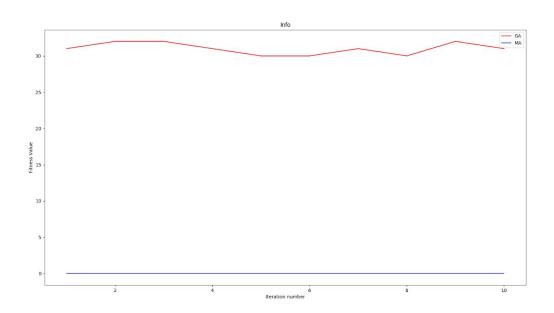
♦ Aspects of the process followed :

➤ Memetic Algorithm :

- **Chromosome generation**: Whatever random genes are formed, they are added to form a chromosome.
- **Population generation**: A fixed number of chromosomes are generated using the above method for creating a population.
- **Local Search**: A conflicting gene is replaced by a newly generated one, and the chromosome is checked for any improvement in its fitness.
- **Crossover methods**: Single-point crossover is being used by randomly generating the crossover point.
- Evaluation Function : A +1 is added to the fitness value for each conflict.

♦ Comparison between GA and MA:

> Running GA and MA for randomly generated datasets:



Observation :

- ➤ GA averages out at a worse fitness score than MA.
- > MA is able to completely solve the problem at hand.

♦ Inference :

- ➤ The reason for GA averaging at a worse score is because of crossover, there is some localisation of genes and hence, the entire population might settle at some average score, i.e., achieve convergence.
- ➤ MA is able to solve the problem completely because with the help of local search we are able to target the exact conflicting genes and deterministically improve the fitness score.

CSP

Assumptions :

- Professors have been mapped to courses uniformly.
- The problem is to fill in a given number of slots.
- > The CSP problem here has been solved for 150 slots.

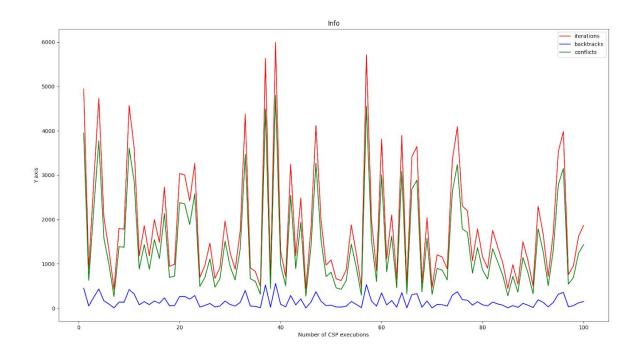
Methodology:

- > We have a list of the given number of slots.
- ➤ We randomly generate 10 tuples and place one of them into the schedule list.
- > Put the other 9 in a backstack.
- ➤ If the inserted tuple is valid, i.e, it satisfies the given constraints, we move on to the next slot and repeat the process.
- ➤ In case the tuple doesn't satisfy the constraints, we pop a new tuple from the backstack and insert it into the schedule and check for validity.
- > We repeat this process until all the slots are filled with valid tuples.

Algorithm :

```
function AC-3(csp) returns false if an inconsistency is found and true otherwise
  inputs: csp, a binary CSP with components (X, D, C)
  local variables: queue, a queue of arcs, initially all the arcs in csp
  while queue is not empty do
     (X_i, X_i) \leftarrow REMOVE-FIRST(queue)
     if REVISE(csp, X_i, X_j) then
       if size of D_i = 0 then return false
       for each X_k in X_i. NEIGHBORS - \{X_i\} do
          add (X_k, X_i) to queue
  return true
function REVISE(csp, X_i, X_j) returns true iff we revise the domain of X_i
  revised \leftarrow false
  for each x in D_i do
     if no value y in D_i allows (x,y) to satisfy the constraint between X_i and X_i then
       delete x from D_i
       revised \leftarrow true
  return revised
```

♦ NUMBER OF ITERATIONS OF CSP VS PARAMETERS PLOT



Observations:

- ➤ Here we observe that for each iteration of CSP, the number of conflicts are quite close to the total number of iterations performed to solve that particular problem.
- > Also, that the number of times that we have to backtrack is quite less as compared to the total number of conflicts encountered.

Inference:

- The total conflicts are close to the number of iterations as the genes are being generated randomly and the probability of a conflict is quite high.
- ➤ However, the times that we actually backtrack is quite less than the number of conflicts because some gene from the backstack will eventually be a valid one.

Result:

```
Run - Artificial_Intelligence_Assignment_2

-
Run: 🏺 csp
                                                                 [[list([19, 19]) 5 8 4 0]
[list([2, 22]) 2 7 3 1]
[list([4, 24]) 1 1 2 2]
[list([9, 9]) 2 5 1 3]
[list([11, 21]) 3 6 1 4]
[list([13, 13]) 1 5 3 5]
[list([14, 34]) 3 8 3 6]
[list([11, 31]) 4 5 4 7]
[list([11, 31]) 4 5 4 7]
[list([11, 11]) 4 6 1 8]
[list([15, 5]) 1 3 1 9]
[list([18, 38]) 5 8 2 10]
[list([2, 2]) 4 2 4 11]
[list([15, 15]) 5 4 4 13]
[list([15, 15]) 5 4 4 13]
[list([13, 33]) 3 5 4 14]
[list([14, 34]) 3 5 2 19]
[list([17, 37]) 4 3 1 17]
[list([17, 37]) 4 3 1 17]
[list([17, 37]) 4 5 3 18]
[list([17, 37]) 4 5 3 18]
[list([18, 18]) 3 7 2 20]
[list([18, 18]) 3 7 2 20]
[list([18, 18]) 3 7 2 20]
[list([18, 38]) 1 7 4 24]
[list([18, 38]) 1 7 4 24]
[list([10, 30]) 3 2 1 27]
[list([10, 30]) 3 2 1 27]
[list([14, 34]) 1 2 2 30]
[list([14, 34]) 1 2 2 30]
[list([14, 34]) 1 6 3 34]
[list([14, 34]) 1 6 3 34]
[list([14, 34]) 1 6 3 34]
[list([19, 10]) 4 7 2 38]
[list([14, 34]) 1 6 3 34]
[list([14, 34]) 2 7 4 40]
[list([15, 5]) 3 1 3 3 41]
[list([16, 16]) 1 7 2 42]
[list([18, 38]) 3 8 2 46]
[list([19, 20]) 1 4 4 49]
[list([10, 30]) 5 1 5 1 51]
[list([15, 5]) 5 7 1 51]
[list([17, 77]) 4 6 2 50]
[list([17, 77]) 4 6 2 50]
[list([17, 77]) 2 8 1 52]
[list([17, 77]) 4 6 2 50]
[list([17, 77]) 2 8 1 52]
[list([17, 77]) 2 8 1 52]
[list([17, 77]) 2 8 1 52]
[list([17, 77]) 3 8 3 53]
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